

# **Five Powerful Chapel Idioms**

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#### What is Chapel?

- A new parallel programming language
  - Under development at Cray Inc.
  - Supported through the DARPA HPCS program
- Availability
  - Version 1.1 release April 15, 2010
  - Open source via BSD license

http://chapel.cray.com/

http://sourceforge.net/projects/chapel/



# **Chapel Productivity Goals**

- Improve programmability over current languages
  - Writing parallel codes
  - Reading, changing, porting, tuning, maintaining, ...
- Support performance at least as good as MPI
  - Competitive with MPI on generic clusters
  - Better than MPI on more capable architectures
- Improve portability over current languages
  - As ubiquitous as MPI
  - More portable than OpenMP, UPC, CAF, ...
- Improve robustness via improved semantics
  - Eliminate common error cases
  - Provide better abstractions to help avoid other errors



## Outline

- What is Chapel
- The Five Idioms
  - Data distributions
  - Data-parallel loops
  - [Asynchronous] [remote] tasks
  - Nested parallelism
  - [Remote] transactions
- Performance Study



# **Idiom 1: Data Distributions**

```
const D = [1..n, 1..n]; // domain - index set
var A: [D] real; // array - data values
const DD = D dmapped X(...); // distributed domain
var DA: [DD] real; // distributed array
```

Syntax

domain-expr dmapped distribution-expr

- Semantics
  - Index set of *domain-expr* is partitioned via *distribution-expr*
  - Partitioned across 'locales' of a system
  - Locale abstraction of memory and processing capability



Standard Block distribution

```
const D = [1..n, 1..m];
var A: [D] real;
const DD = D dmapped Block(boundingBox=D);
var DA: [DD] real;
```





Standard Cyclic distribution

```
const D = [1..n, 1..m];
var A: [D] real;
const DD = D dmapped Cyclic(startIdx=D.low);
var DA: [DD] real;
```





User-defined MyBanded distribution

```
const D = [1..n, 1..m];
var A: [D] real;
const DD = D dmapped MyBanded(startIdx=D.low);
var DA: [DD] real;
```



#### Idiom 2: Data-Parallel Loops





Syntax

forall ( index-exprs ) in ( iterable-exprs ) do
 loop-body-stmts

Semantics

- Zipped (element-wise) iteration
- Shapes of iterable expressions must match







• Example 1: Non-distributed arrays





#### Idiom 2: Data-Parallel Loops

forall (a, b, c) in (A, B, C) do a = b + alpha \* c; Example 2: Block-distributed arrays = Δ + B α• ( Locales 2 3 1 0







• Example 3: Unaligned block-distributed arrays







forall (a, b, c) in (A, B, C) do
a = b + alpha \* c;

• Example 4: 2D Block-distributed arrays









```
forall (a, b, c) in (A, B, C) do
a = b + alpha * c;
```

- Other possibilities
  - Associative, sparse, and unstructured arrays
  - Domains and iterators with no associated data
  - A distributed tree or graph that supports iteration
- Preferred way of writing simple computations:

A = B + alpha \* C;

## **Chapel View of Compiler Transform**



Initial Code:

A = B + alpha \* C;

1. Promotion of scalar multiplication:

A = B + [c in C] alpha\*c;

2. Promotion of scalar addition:

A = [(b,f) in (B, [c in C] alpha\*c)] b+f;

3. Collapse of foralls:

A = [(b,c) in (B,C)] b+alpha\*c;

4. Expansion of assignment:

forall (a,f) in (A,[(b,c) in (B,C)] b+alpha\*c) do
 a=f;

5. Collapse of foralls:

forall (a,b,c) in (A,B,C) do
 a = b + alpha \* c;



on loc do begin f();

#### • Syntax

on expr do stmt begin stmt

- Semantics
  - On-statement evaluates locale of *expr* Then executes *stmt* on that locale
  - Begin-statement creates a new task to execute *stmt* Original task continues with the next statement



#### Idiom 3: Asynchronous Remote Tasks

on loc do begin f();

• Picture







#### Locales

- Abstraction of memory and processing capability
- Architecture-dependent definition optimizes local accesses

#### Tasks

- Abstraction of computation or thread
- Execution is on a locale
- Programming model support

Chapel	OpenMP	MPI	UPC	CAF	Titanium
Locales		Processes	Threads	Images	Demesnes
Tasks	Threads				



# **Idiom 4: Nested Parallelism**

• Task parallelism of data parallelism



Data parallelism of task parallelism



### **Idiom 5: Remote Transactions**



on A(i) do atomic A(i) = A(i) ^ i;

• Syntax

atomic stmt

Semantics

 Executes stmt with transaction semantics so that stmt appears to take effect atomically

Note: atomic statements are not implemented



#### Outline

- What is Chapel
- The Five Idioms
- Performance Study
  - HPCC Global Stream
  - HPCC EP Stream



# **Global STREAM Triad in Chapel (Excerpts)**





### **EP STREAM Triad in Chapel (Excerpts)**





local {

**var** A, B, C: [1..m] real;

forall (a,b,c) in (A,B,C) do

a = b + alpha \* c;



Machine Characteristics			
Model	Cray XT4		
Location	ORNL		
Nodes	7832		
Processor	2.1 GHz Quadcore AMD Opteron		
Memory	8 GB per node		

Benchmark Parameters				
STREAM Triad Memory	Least value greater than 25% of memory			
Random Access Memory	Least power of two greater than 25% of memory			
Random Access Updates	2 <sup>n-10</sup> for memory equal to 2 <sup>n</sup>			



# **STREAM Triad Performance**



#### CUG '10<sup>.</sup> Five Powerful Chapel Idioms

#### **Thank You**



# Chapel URL: <u>http://chapel.cray.com/</u> Chapel Source: <u>http://sourceforge.net/projects/chapel</u> Contact: <u>chapel\_info@cray.com</u>