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AIX 4.3 System Management Guide Communications and Networks

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Software

October 1999

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About This Book

This book is for AIX system administrators who maintain the system's network connections. Familiarity with the Base Operating System and the material covered in AIX 4.3 System Management Guide: Operating System and Devices and AIX 4.3 System User's Guide: Communications and Networks is necessary.

Who Should Use This Book

This book is intended for system administrators who perform system management tasks that involve communication within a network.

How to Use This Book

Each chapter provides overview, installation and configuration, maintenance, troubleshooting, and reference information in roughly that order. Procedural information is located near the relevant conceptual information for easier access.

Highlighting

The following highlighting conventions are used in this book:

Bold Identifies commands, keywords, files, directories, and other

items whose names are predefined by the system.

Identifies parameters whose actual names or values are to

be supplied by the user.

Monospace Identifies examples of specific data values, examples of text

similar to what you might see displayed, examples of portions of program code similar to what you might write as a programmer, messages from the system, or information

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AIX Support for the X/Open UNIX95 Specification

The AIX operating system is designed to support the X/Open UNIX95 Specification for portability of UNIX-based operating systems. Many new interfaces, and some current ones, have been added or enhanced to meet this specification. AIX 4.3 is even more open and portable for applications.

At the same time, compatibility with previous AIX releases is preserved. This is accomplished by the creation of a new environment variable that can be used to set the system environment on a per–system, per–user, or per–process basis.

To determine the proper way to develop a UNIX95–portable application, you may need to refer to the X/Open UNIX95 Specification.

Related Publications

The following book contains information about or related to communications:

AIX and Related Products Documentation Overview, order number 86 A2 71WE.

AIX 4.3 System Management Guide: Operating System and Devices, order number 86 A2 99HX.

AIX 4.3 System User's Guide: Communications and Networks, order number 86 A2 98HX.

AIX General Programming Concepts: Writing and Debugging Programs, order number 86 A2 34JX.

AIX Commands Reference, order number 86 A2 38JX to 86 A2 43JX.

AIX 4.3 Installation Guide, order number 86 A2 43GX.

The TTY Subsystem Overview in AIX General Programming Concepts: Writing and Debugging Programs provides general information on line disciplines and the tty code.

See also the following industry publications:

Token–Ring Network Architecture Reference, order number SC30–3374.

Albitz, Paul and Liu, Cricket, [1992], *DNS and BIND in a Nutshell*, O'Reilly & Associates, Inc., Sebastopol, CA, ISBN 0–56592–010–4.

Comer, Douglas E., [1991], *Internetworking with TCP/IP, Volume I: Principles, Protocols, and Architectures*, Prentice Hall, Englewood Cliffs, NJ, ISBN 0–13–468505–9.

Comer, Douglas E. and Stevens, David L., [1991], *Internetworking with TCP/IP, Volume II: Design, Implementation, and Internals*, Prentice Hall, Englewood Cliffs, NJ, ISBN 0–13–472242–6.

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Hunt, Craig, [1992], *TCP/IP Network Administration*, O'Reilly & Associates, Inc., Sebastopol, CA, ISBN 0-93-717582-X.

Stern, Hal, [1991], Managing NFS and NIS, O'Reilly & Associates, Inc., Sebastopol, CA, ISBN 0-93-717575-7.

Stevens, Richard W., [1990], *UNIX Network Programming*, Prentice Hall, Englewood Cliffs, NJ, ISBN 0–13–949876–1.

Tanenbaum, Andrew S., *Computer Networks*, Prentice Hall, Englewood Cliffs, NJ, ISBN 0–13–165183–8.

Vanel, Laurent, Steve Gardner, Praben Prima, Simon Robertson, and Oreste Villari, *AIX and Windows NT: Solutions for Interoperability*, International Business Machines, Inc. http://www.redbooks.ibm.com

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X-1

Chapter 1. Communications and Networks Overview

This chapter presents the conceptual foundation for understanding computer networking in general. System administrators unfamiliar with general networking principles should read this chapter. Those familiar with UNIX networking may safely skip this chapter.

A network is the combination of two or more computers and the connecting links between them. A physical network is the hardware (equipment such as adapter cards, cables, concentrators, and telephone lines) that makes up the network. The software and the conceptual model make up the logical network.

This overview provides the following information on networks:

- Communications Functions Introduction, on page 1-2
- Network Introduction, on page 1-3
- Physical Networks Introduction, on page 1-5
- System Communications Support, on page 1-6
- Communicating with Other Operating Systems, on page 1-8

Communications Functions Introduction

Networks allow for several user and application communications functions. They enable a user to:

- · Send electronic mail
- Emulate another terminal or log in to another computer
- Transfer data
- Run programs residing on a remote node

One of the most popular applications for computer networks is electronic mail. Electronic mail, or e-mail, allows one user to send a message to another user. The two users may be on the same system (in which case a communications network is not needed), across the hall, or across the world.

A communications network also can allow one computer to mimic another, and access information as if it were another type of computer or terminal. Remote login capabilities allow users to log in to a remote system and access the same programs and files as if they were using the machine without a communications network.

Networks also allow for the transfer of data from one system to another. Files, directories, and entire file systems can be migrated from one machine to another across a network, enabling remote backup of data, as well as assuring redundancy in case of machine failure.

Several different protocols have been devised to allow users and applications on one system to invoke procedures and applications on other systems. This can be useful for a number of environments, including the offloading of many computer-intensive routines in engineering and scientific applications.

Network Introduction

The complexity of modern computer networks has given rise to several conceptual models for explaining how networks work. One of the most common of these models is the International Standards Organization's Open Systems Interconnection (OSI) Reference Model, also referred to as the OSI seven—layer model. The seven layers of the OSI model are numbered beginning at the lowest (physical) layer.

- 7 Application6 Presentation5 Session4 Transport
- 3 Network

2

4 Dhusiasi

Data Link

1 Physical

Levels 1 through 3 are network specific, and will differ depending on what physical network you are using. Levels 4 through 7 comprise network—independent, higher—level functions. Each layer describes a particular function (instead of a specific protocol) that occurs in data communications. The seven layers function as follows:

Physical Describes the physical media of the network. For example, the fiber

optic cable required for a Fiber Distributed Data Interface (FDDI)

network is part of the physical layer.

Data Link Provides reliable delivery of data across the physical layer (which is

usually inherently unreliable).

Network Manages the connections to other machines on the network.

Transport Ensures error–free data transmission.

Session Manages the connections between applications.

Presentation Ensures that data is presented to the applications in a consistent

fashion.

Application Comprises the applications that use the network.

Note that while the OSI Reference Model is useful for discussing networking concepts, many networking protocols do not closely follow the OSI model. For example, when discussing Transmission Control Protocol/Internet Protocol (TCP/IP), the Application and Presentation Layer functions can be combined into a single level, as can the Session and Transport Layers, as well as the Data Link and Physical Layers.

Each layer in the OSI model communicates with the corresponding layer on the remote machine as shown in the OSI Reference Model figure. The layers pass data only to the layers immediately above and below. Each layer adds its own header information (and, in the case of the Data Link layer, footer), effectively encapsulating the information received from the higher layers.

7	Application							Data	
6	Presentation	1				[Head	ler Data	
5	Session					Header		Data	
ļ	Transport				Header		Data		
3	Network		[Header		Data			
2	Data Link		Header		Da	ıta			Footer
ı	Physical	Header			Da	ıta			

Individual users as well as organizations use networks for many reasons. A few possibilities are:

- Data entry
- Data queries
- Remote batch entry
- Resource sharing
- Data sharing
- Electronic mail

Data entry consists of entering data directly into either local or remote data files, reducing the need for such intermediate steps as posting, recording, or punching. Increased accuracy and efficiency are natural by-products of a one-step data transfer. Data queries entail searching data files for specified information. Data updating involves altering, adding, or deleting data stored in local or remote files. Remote batch entry consists of entering batches of data from a remote location, an activity often performed at night or during periods of low system usage. Because of such diverse capabilities, communications and networks are not only desirable but necessary.

Sharing resources is another function of networks. Users can share data as well as programs, file-storage space, and peripheral devices (for example, printers, modems, terminals, and fixed disks). Such sharing of system resources is cost effective (in the case of peripheral sharing), while eliminating the problems of keeping multiple copies of programs and data consistent (in the case of program and file sharing).

A communications network also provides a means for users to communicate using electronic mail. Electronic mail (e-mail) can be used between users on the same system or between users across the world.

Physical Networks Introduction

The physical network consists of the cable (coaxial cable, twisted pair, fiber optic, and telephone lines) that connects the different hardware residing on the network, the adapter cards used on the attached hosts, and any concentrators, repeaters, routers or bridges used in the network. (A *host*, in this sense, is a computer attached to the network.)

Physical networks vary both in size and in the type of hardware used. The two common kinds of networks are *local area networks* (LANs) and *wide area networks* (WANs). A LAN is a network where communications are limited to a moderately sized geographic area of 1 to 10 km (1 to 6 miles), such as a single office building, warehouse, or campus. A WAN is a network providing data communications capability throughout geographic areas larger than those serviced by LANs; for example, across a country or across continents. A newer, intermediate class of networks is beginning to emerge, called *metropolitan area networks* (MANs). This guide does not generally distinguish MANs; they are grouped with WANs.

LANs commonly use Standard Ethernet, IEEE 802.3 Ethernet, or token—ring hardware for the physical network, while WANs and asynchronous networks use communications networks provided by common carrier companies. Operation of the physical network in both cases is usually controlled by networking standards; for example, EIA 232D or CCITT V.35.

System Communications Support

All network communications involve use of hardware and software. Hardware consists of the physical equipment connected to the physical network. Software consists of the programs, and device drivers pertaining to the operation of a particular system.

The system hardware consists of adapter cards or other devices that provide a path or interface between the system software and the physical network. An adapter card requires an input/output (I/O) card slot in the system. Other devices, such as modems, may attach to one of the standard ports on the computer.

Adapter cards support the standards required by the physical network (for example, EIA 232D, Smartmodem, V.25 bis, EIA 422A, X.21, or V.35) and may, at the same time, support software protocols (for example, synchronous data link control or SDLC, high-level data link control or HDLC, and bisynchronous protocols). If the adapter does not contain software support, then this support must be provided by the adapter's device driver.

Protocols

All communications software uses protocols, sets of semantic and syntactic rules that determine the behavior of functional units in achieving communication. Protocols define how information is delivered, how it is enclosed to reach its destination safely, and what path it should follow. Protocols also coordinate the flow of messages and their acknowledgments.

Protocols exist at different levels within the kernel and cannot be manipulated directly. However, they are manipulated indirectly by what the user chooses to do at the application programming interface (API) level. The choices a user makes when invoking file transfer, remote login, or terminal emulation programs define the protocols used in the execution of those programs.

Addresses

A third feature common to communications networks is addresses. Addresses are associated with both software and hardware. The address is the means by which the sending or control station selects the station to which it sends data. In effect, addresses are a means of identifying receiving or storage locations. A physical address is a unique code assigned to each device or workstation connected to a network.

For example, on a token-ring network, the netstat -iv command displays the token-ring card address. This is the physical network address. The netstat -iv command also displays class-level and user-level address information. Addresses are often defined by software but can be created by the user as well.

Domains

An aspect of addresses common to many communications networks is the concept of domains. For example, the structure of the Internet illustrates how domains define the Internet Protocol (IP) address. The Internet is an extensive network that comprises many different smaller networks. To facilitate routing and addressing, Internet addresses are hierarchically structured in domains, with very broad categories at the top such as com for commercial users, edu for educational users, and gov for government users.

Within the com domain are many smaller domains corresponding to individual businesses: for example, ibm. Within the ibm.com domain are even smaller domains corresponding to the Internet addresses for various locations, such as austin.ibm.com or raleigh.ibm.com. At this level, we start seeing names of hosts. A host, in this context, is any computer connected to the network. Within austin.ibm.com, there may be hosts with the names hamlet and lear, which are addressed hamlet.austin.ibm.com and lear.austin.ibm.com.

Gateways and Bridges

A wide variety of networks reside on the Internet, often using different hardware and running different software. Gateways and bridges enable these different networks to communicate with each other. A bridge is a functional unit that connects two LANs that possibly use the same logical link control (LLC) procedure, such as Ethernet, but different medium access control (MAC) procedures. A gateway has a broader range than a bridge. It operates above the link layer and translates, when required, the interface and protocol used by one network into those used by another distinct network. Gateways allow data transfers across the various networks that constitute the Internet.

Routing

Using domain names for addressing and gateways for translation greatly facilitates the routing of the data being transferred. Routing is the assignment of a path by which a message reaches its destination. The domain name effectively defines the message destination. In a large network like the Internet, information is routed from one communications network to the next until that information reaches its destination. Each communications network checks the domain name and, based on the domains with which that network is familiar, then routes the information on to the next logical stop. In this way, each communications network that receives the data contributes to the routing process.

Local and Remote Nodes

A physical network is used by the hosts that reside on that network. Each host can be said to be a node on the network. A node is an addressable location in a communications network that provides host–processing services. The intercommunication of these various nodes gives rise to the concepts of local and remote. Local pertains to a device, file, or system accessed directly from your system, without the use of a communications line. Remote pertains to a device, file, or system accessed by your system over a communications line. Local files reside on your system, while remote files reside on a file server or at another node with which you communicate using a physical network, for example, Ethernet, token-ring, or phone lines.

Server and Client

Related to the concepts of local and remote are those of *client* and *server*. A server is a computer that contains data or provides facilities to be accessed by other computers on the network. Common server types are file servers, which store files; name servers, which store names and addresses; and code servers, which store programs and applications.

A client is a computer requesting services or data from a server. A client, for example, could request updated program code or the use of applications from a code server. To obtain a name or address, a client contacts a name server. A client could also request files and data for data entry, inquiry, or record updating from a file server.

Communicating with Other Operating Systems

Different types of computers can be connected on a network. That is, the computers can be from different manufacturers or be different models from the same manufacturer. The differences in operating systems of two or more types of computers can be bridged with communications programs.

Sometimes these programs require that another program previously be installed on the network. Other programs may require that such communications connectivity protocols as TCP/IP or Systems Network Architecture (SNA) exist on the network.

For example, with AIX versions 4.3.2 and later, AIX Fast Connect for Windows lets PC clients access AIX files and printers using native PC networking client software. PC users can use remote AIX file systems directly from their machines as if they were locally stored, and they can print jobs on printers using the AIX spool system, view available printers, and map a printer as a network printer.

Chapter 2. Mail

The mail facility provides a method for exchanging electronic mail with users on the same system or on multiple systems connected by a network. This section documents the mail system and the standard mail user interface.

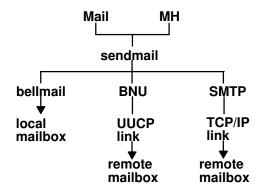
The mail facility provides a method for exchanging electronic mail (e-mail) with users on the same system or on multiple systems connected by a network. This section documents the mail system, the standard mail user interface, the Internet Message Access Protocol (IMAP), and the Post Office Protocol (POP).

The mail system is an internetwork mail delivery facility that consists of a user interface, a message routing program, and a message delivery program (or mailer). The mail system relays messages from one user to another on the same host, between hosts, and across network boundaries. It also performs a limited amount of message—header editing to put the message into a format that is appropriate for the receiving host.

A mail user interface enables users to create and send messages to, and receive messages from, other users. The mail system provides two user interfaces, mail and mhmail. The mail command is the standard mail user interface available on all UNIX systems. The mhmail command is the Message Handler (MH) user interface, an enhanced mail user interface designed for experienced users.

A message routing program routes messages to their destinations. The mail system's message routing program is the **sendmail** command, which is part of the Base Operating System (BOS) and is installed with BOS. The **sendmail** command is a daemon that uses information in the /etc/sendmail.cf file, the /etc/aliases file, and the /etc/sendmail.nl file to perform the necessary routing.

Depending on the type of route to the destination, the **sendmail** command uses different *mailers* to deliver messages (see figure).



As the figure illustrates:

- To deliver local mail, the sendmail program routes messages to the bellmail program.
 The bellmail program delivers all local mail by appending messages to the user's system mailbox, which is in the /varspool/mail directory.
- To deliver mail over a UNIX-to-UNIX Copy Program (UUCP) link, the sendmail program routes messages using Basic Network Utilities (BNU).
- To deliver Transmission Control Protocol/Internet Protocol (TCP/IP)—routed mail, the sendmail command establishes a TCP/IP connection to the remote system then uses Simple Mail Transfer Protocol (SMTP) to transfer the message to the remote system.

Mail Management Tasks

The following is a list of the tasks for which you, the mail manager, are responsible.

- 1. Configure the /etc/rc.tcpip file so that the sendmail daemon will be started at system boot time. See the instructions immediately following this list.
- Customize the configuration file /etc/sendmail.cf. The default /etc/sendmail.cf file is configured so that local mail can be delivered. In order to deliver mail through TCP/IP or BNU, you must customize and compile the /etc/sendmail.cf file. See the "sendmail.cf File" in AIX Files Reference for more information.
- 3. Define systemwide and domainwide mail aliases in the /etc/aliases file. See "Managing Mail Aliases", on page 2-3 for more information.
- 4. Manage the Mail Queue. See "Managing the Mail Queue Files and Directories", on page 2-6 for more information.
- 5. Manage the Mail Log. See "Managing Mail Logging", on page 2-11 for more information.

Configuring the /etc/rc.tcpip File to Start the sendmail Daemon

To configure the /etc/rc.tcpip file so that the sendmail daemon will be started at system boot time:

- 1. Edit the /etc/rc.tcpip file with your favorite editor.
- Find the line that begins with start /usr/lib/sendmail. By default, this line should be uncommented, that is, there is no # (pound sign) at the beginning of the line. However, if it is commented, delete the pound sign.
- 3. Save the file.

With this change, the system will start the **sendmail** daemon at boot time.

Managing Mail Aliases

Aliases map names to address lists using personal, systemwide, and domainwide alias files. You can define three types of aliases:

personal local system

Defined by individual users in the user's \$HOME/.mailrc file.

Defined by the mail system administrator in the /etc/aliases file. These aliases apply to mail handled by the sendmail command on the local system. Local system aliases rarely need to be changed.

domainwide

Aliases are found using the Network Information Service (NIS). (NIS must be configured and the \bigcirc AliasFile option must be set in the **sendmail.cf** file.) To enable use of the NIS aliases mapping:

- 1. Uncomment O AliasFile in the **sendmail.cf** file and specify the map name for NIS aliases.
- Recompile the sendmail.cf file with the command sendmail –bz.
- Recompile the alias database with the command sendmail -hi

If you use a domain hierarchy for NIS+, you may need one mail host for each domain in the namespace because names in separate domains may no longer be unique. For more information about NIS and NIS+ requirements, see *AIX 4.3 NIS/NIS+ Guide*.

/etc/aliases File

The /etc/aliases file consists of a series of entries in the following format:

```
Alias: Namel, Name2, ... NameX
```

where *Alias* can be any alphanumeric string that you choose (not including special characters, such as @ or !). *Name1* through *NameX* is a series of one or more recipient names. The list of names can span one or more lines. Each continued line begins with a space or a tab. Blank lines and lines beginning with a # (pound sign) are comment lines.

The /etc/aliases file must contain the following three aliases:

MAILER-DAEMON

The ID of the user who is to receive messages addressed to the mailer daemon. This name is initially assigned to the root user:

MAILER-DAEMON: root

postmaster

The ID of the user responsible for the operation of the local mail system. The **postmaster** alias defines a single mailbox address that is valid at each system in a network. This address enables users to send inquiries to the **postmaster** alias at any system, without knowing the correct address of any user at that system. This name is initially assigned to the root user:

postmaster: root

nobody

The ID that is to receive messages directed to programs such as **news** and **msgs**. This name is initially assigned to /dev/null:

```
nobody: /dev/null
```

To receive these messages, define this alias to be a valid user.

Any time you change this file, you must recompile it into a database format that the **sendmail** command can use. See *Building the Alias Database*, on page 2-4.

Creating Local System Aliases for Mail

To create or modify local system aliases:

- Edit the /etc/aliases file using your favorite editor.
- 2. On a blank line, add an alias, followed by a colon (:), followed by a list of comma—separated recipients. For example, the following entry defines the writers alias to be the names of people in that group:

```
writers: geo, mark@zeus, ctw@athena, brian
```

This definition could also be contained on several lines, as long as each added line begins with a space or a tab, for example:

3. Create an owner for any distribution list aliases. If the **sendmail** command has trouble sending mail to the distribution list, it sends an error message to the owner of that list. For example, the following set of entries in the /etc/aliases file defines a distribution list named editors, whose owner is glenda@hera:

```
editors: glenda@hera, davidm@kronos, perryw@athena owner-editors: glenda@hera
```

 Recompile the /etc/aliases file by following the instructions in the section, "Building the Alias Database".

Building the Alias Database

The **sendmail** command does not use directly the alias definitions in the local system /**etc/aliases** file. Instead, the **sendmail** command reads a processed database manager (dbm) version of the /**etc/aliases** file. You can compile the alias database using one of the following methods:

- Run the /usr/sbin/sendmail command using the -bi flag.
- Run newaliases. This command causes the sendmail command to read the local system /etc/aliases file and create two additional files containing the alias database information:

/etc/aliases.dir

/etc/aliases.pag

Run the sendmail command using the Rebuild Aliases flag. This rebuilds the alias
database automatically when it is out—of—date. Auto—rebuild can be dangerous on heavily
loaded machines with large alias files. If it might take more than the rebuild time—out
(normally five minutes) to rebuild the database, there is a chance that several processes
will start the rebuild process simultaneously.

Note:

- 1. If these files do not exist, the **sendmail** command cannot process mail and will generate an error message.
- 2. If you have multiple alias databases specified, the **-bi** flag rebuilds all the database types it understands (for example, it can rebuild Network Database Management (NDBM) databases but not NIS databases).

The /etc/netsvc.conf file contains the ordering of system services. To specify the service ordering of aliases, add the following line:

```
aliases=service, service
```

where service can be either files or nis. For example:

```
aliases=files, nis
```

tells the **sendmail** command to try the local alias file first; and if that fails, try nis. If nis is defined as a service, it should be running.

For further information on the /etc/netsvc.conf file, see AIX Files Reference.

Managing the Mail Queue Files and Directories

The mail queue is a directory that stores data and controls files for mail messages that the **sendmail** command delivers. By default, the mail queue is /**var/spool/mqueue**.

Mail messages may be queued for several reasons. First, the **sendmail** command can be configured to process the queue at certain intervals, rather than immediately. If this is so, mail messages must be stored temporarily. Second, if a remote host does not answer a request for a mail connection, the mail system queues the message and tries again later.

Printing the Mail Queue

The contents of the queue can be printed using the **mailq** command (or by specifying the **-bp** flag with the **sendmail** command).

This produces a listing of the queue IDs, the size of the message, the date the message entered the queue, and the sender and recipients.

Mail Queue Files

Each message in the queue has a number of files associated with it. The files are named according to the following conventions:

Typef ID

where *ID* is a unique message queue ID, and *Type* is one of the following letters indicating the type of file:

- **d** The data file containing the message body without the heading information.
- **q** The queue–control file. This file contains the information necessary to process the job.
- t A temporary file. This file is an image of the **q** file when it is being rebuilt. It is quickly renamed to the **q** file.
- A transcript file that exists during the life of a session and shows everything that happens during that session.

For example, if a message has a queue ID of AA00269, the following files are created and deleted in the mail queue directory while the **sendmail** command tries to deliver the message:

dfAA00269Data fileqfAA00269Control filetfAA00269Temporary filexfAA00269Transcript file

q Control File

The **q** control file contains a series of lines, each beginning with a code letter:

- B Specifies the body type. The remainder of the line is a text string defining the body type. If this field is missing, the body type is assumed to be undefined, and no special processing is attempted. Legal values are 7BIT and 8BITMIME.
- Contains the controlling address. The syntax is *localuser:aliasname*.

 Recipient addresses following this line are flagged so that deliveries will run as the *localuser* (a user name from the /etc/passwd file); *aliasname* is the name of the alias that expanded to this address (used for printing messages).

- **F** Contains flag bits represented as one letter per flag. Defined flag bits are **r** indicating that this is a response message and **w** indicating that a warning message has been sent announcing that mail has been delayed.
- H Contains a heading definition. There may be any number of these lines. The order in which the H lines appear determines their order in the final message. These lines use the same syntax as heading definitions in the /etc/sendmail.cf configuration file.
- Specifies the i-number of the data file; this can be used to recover your mail queue after a disk crash.
- **K** Specifies the time (as seconds) of the last delivery attempt.
- **M** Contains a message printed by the **mailq** command. It is mainly used to store status information.
- **N** Specifies the total number of delivery attempts.
- O Specifies the original message transfer system (MTS) value from the Ethernet Simple Mail Transfer Protocol (SMTP) transaction. Used for Delivery Status Notifications only.
- P Contains the priority of the current message. The priority is used to order the queue. Higher numbers mean lower priorities. The priority increases as the message sits in the queue. The initial priority depends on the message class and the size of the message.
- Q Contains the original recipient as specified by the ORCPT= field in an SMTP transaction. Used exclusively for Delivery Status Notifications. It applies only to the immediately following "R" line.
- **R** Contains a recipient address. There is one line for each recipient.
- **S** Contains the sender address. There is only one of these lines.
- **T** Contains the message creation time used to compute when to time out the message.
- V Specifies the version number of the queue file format used to allow new **sendmail** binaries to read queue files created by older versions. Defaults to version **zero**. Must be the first line of the file, if present.
- **Z** Specifies the original envelope ID (from the SMTP transaction). Used for Delivery Status Notifications only.
- \$ Contains a macro definition. The values of certain macros (**\$r** and **\$s**) are passed through to the queue run phase.

The q file for a message sent to amy@zeus would look similar to:

```
P217031
T566755281
MDeferred: Connection timed out during user open with zeus Sgeo
Ramy@zeus
H?P?return-path: <geo>
Hreceived: by george (0.13 (NL support)/0.01)
            id AA00269; Thu, 17 Dec 87 10:01:21 CST
H?D?date: Thu, 17 Dec 87 10:01:21 CST
H?F?From: geo
Hmessage-id: <8712171601.AA00269@george>
HTo: amy@zeus
Hsubject: test
```

P217031 Priority of the message

Where:

T566755281 Submission time in seconds

MDeferred: Connection timed out during user open with zeus

Status message

Sgeo ID of the sender Ramy@zeus ID of the receiver

HLines Header information for the message

Specifying Time Values in sendmail

To set the message time—out and queue processing interval, you must use a specific format for the time value. The format of a time value is:

```
-qNumberUnit
```

where *Number* is an integer value and *Unit* is the unit letter. *Unit* may have one of the following values:

S	Seconds
m	Minutes
h	Hours
d	Days
w	Weeks

If *Unit* is not specified, the **sendmail** daemon uses minutes (**m**) as the default. Here are three examples illustrating time—value specification:

```
/usr/sbin/sendmail -q15d
```

This command tells the **sendmail** daemon to process the queue every 15 days.

```
/usr/sbin/sendmail -q15h
```

This command tells the **sendmail** daemon to process the queue every 15 hours.

```
/usr/sbin/sendmail -q15
```

This command tells the **sendmail** daemon to process the queue every 15 minutes.

Forcing the Mail Queue

In some cases, you may find that the queue is clogged for some reason. You can force a queue to run using the $-\mathbf{q}$ flag (with no value). You can also use the $-\mathbf{v}$ flag (verbose) to watch what happens:

```
/usr/sbin/sendmail -q -v
```

You can also limit the jobs to those with a particular queue identifier, sender, or recipient using one of the queue modifiers. For example, **–qRsally** restricts the queue run to jobs that have the string **sally** in one of the recipient addresses. Similarly, **–qS** *string* limits the run to particular senders, and **–ql** *string* limits it to particular queue identifiers.

Setting the Queue Processing Interval

The interval at which the **sendmail** daemon processes the mail queue is determined by the value of the **-q** flag when the daemon starts.

The **sendmail** daemon is usually started by the /**etc/rc.tcpip** file, at system startup. The /**etc/rc.tcpip** file contains a variable called the queue processing interval (QPI), which it uses to specify the value of the **-q** flag when it starts the **sendmail** daemon. By default, the value of **qpi** is 30 minutes. To specify a different queue processing interval:

- 1. Edit the /etc/rc.tcpip file with your favorite editor.
- 2. Find the line that assigns a value to the *qpi* variable, such as:

```
qpi=30m
```

3. Change the value assigned to the *qpi* variable to the time value you prefer.

These changes will take effect at the next system restart. If you want the changes to take effect immediately, stop and restart the **sendmail** daemon, specifying the new **-q** flag value. See *Stopping the sendmail Daemon*, on page 2-10 and *Starting the sendmail Daemon*, on page 2-10 for more information.

Moving the Mail Queue

When a host goes down for an extended period, many messages routed to (or through) that host may be stored in your mail queue. As a result, the **sendmail** command spends a long time sorting the queue, severely degrading your system's performance. If you move the queue to a temporary place and create a new queue, the old queue can be run later when the host returns to service. To move the queue to a temporary place and create a new queue:

- 1. Stop the **sendmail** daemon by following the instructions in *Stopping the sendmail Daemon*, on page 2-10.
- 2. Move the entire queue directory by entering:

```
cd /var/spool
mv mqueue omqueue
```

- 3. Restart the **sendmail** daemon by following the instructions in *Starting the sendmail Daemon*, on page 2-10.
- 4. Process the old mail queue by entering:

```
/usr/sbin/sendmail -oQ/var/spool/omqueue -q
```

The **-oQ** flag specifies an alternate queue directory. The **-q** flag specifies to run every job in the queue. To get a report about the progress of the operation, use the **-v** flag.

Note: This operation can take a long time.

5. Remove the log files and the temporary directory when the queue is empty by entering:

```
rm /var/spool/omqueue/*
rmdir /var/spool/omqueue
```

Starting the sendmail Daemon

To start the **sendmail** daemon, enter either of the following commands:

```
startsrc -s sendmail -a "-bd -q15"
OR
  /usr/lib/sendmail -bd -q15
```

If the sendmail daemon is already active when you enter one of these commands, you will see the following message on the screen:

```
The sendmail subsystem is already active. Multiple instances are
not supported.
```

If the sendmail daemon is not already active, then you will see a message indicating that the sendmail daemon has been started.

Stopping the sendmail Daemon

To stop the **sendmail** daemon, execute the **stopsrc –s sendmail** command. If the sendmail daemon was not started with the startsrc command:

- Find the sendmail pid.
- Issue kill -9 sendmail_pid.

Managing Mail Logging

The **sendmail** command logs mail system activity through the **syslogd** daemon. The **syslogd** daemon must be configured and running for logging to occur. Specifically, the /etc/syslog.conf file may contain the uncommented line:

```
mail.debug /var/spool/mqueue/log
```

If it does not, use your favorite editor to make this change; be certain that the path name is correct. If you change the /etc/syslog.conf file while the syslogd daemon is running, refresh the syslogd daemon by entering the command:

```
refresh -s syslogd
```

If the /var/spool/mqueue/log file does not exist, you must create it by entering the command:

```
touch /var/spool/mqueue/log
```

Messages in the log file appear in the following format:

Each line in the system log consists of a time stamp, the name of the machine that generated it (for logging from several machines over the local area network), the word "sendmail:," and a message. Most messages are a sequence of *name=value* pairs.

The two most common lines are logged when a message is processed. The first logs the receipt of a message; there will be one of these per message. Some fields may be omitted. These message fields are:

from	Specifies the envelope sender address.
size	Specifies the size of the message in bytes.
class	Indicates the class (numeric precedence) of the message.
pri	Specifies the initial message priority (used for queue sorting).
nrcpts	Indicates the number of envelope recipients for this message (after aliasing and forwarding).
proto	Specifies the protocol used to receive the message, for example Ethernet Simple Mail Transfer Protocol (SMTP) or UNIX-to-UNIX Copy Program (UUCP).
relay	Specifies the machine from which it was received.

There is also one line logged per delivery attempt (so there can be several per message if delivery is deferred or there are multiple recipients). These fields are:

to	Contains a comma-separated list of the recipients to this mailer.
ctladdr	Specifies the <i>controlling user</i> , that is, the name of the user whose credentials are used for delivery.
delay	Specifies the total delay between the time this message was received and the time it was delivered.
xdelay	Specifies the amount of time needed in this delivery attempt.
mailer	Specifies the name of the mailer used to deliver to this recipient.
relay	Specifies the name of the host that actually accepted (or rejected) this recipient.
stat	Specifies the delivery status.

There is a large amount of information that can be logged. The log is arranged as a succession of levels. At the lowest level, only very unusual situations are logged. At the highest level, even the insignificant events are logged. As a convention, log levels under ten are considered "useful." Log levels above 64 are reserved for debugging purposes. Levels from 11-64 are reserved for verbose information.

The types of activities that the **sendmail** command puts into the log file are specified by the L option in the /etc/sendmail.cf file.

Managing the Log

Because information is continually appended to the end of the log file, it can become very large. Also, error conditions can cause unexpected entries to the mail gueue. To keep the mail queue and log from growing too large, execute the /usr/lib/smdemon.cleanu shell script. This script forces the sendmail command to process the queue and maintains four progressively older copies of log files, named log.0, log.1, log.2, and log.3. Each time the script runs it moves:

- log.2 to log.3
- log.1 to log.2
- log.0 to log.1
- log to log.0

This allows logging to start over with a new file. Run this script either manually or at a specified interval with the cron daemon.

Logging Traffic

Many SMTP implementations do not fully implement the protocol. For example, some personal computer-based Simple Mail Transfer Protocols (SMTPs) do not understand continuation lines in reply codes. These can be very hard to trace. If you suspect such a problem, you can set traffic logging using the **-X** flag. For example:

```
/usr/sbin/sendmail -X /tmp/traffic -bd
```

Using this command logs all traffic in the /tmp/traffic file.

This logs a lot of data very quickly and should never be used during normal operations. After starting such a daemon, force the **errant** implementation to send a message to your host. All message traffic in and out of sendmail, including the incoming SMTP traffic, will be logged in this file.

Using sendmail, you can log a dump of the open files and the connection cache by send it a **SIGUSR1** signal. The results are logged at **LOG DEBUG** priority.

Logging Mailer Statistics

The **sendmail** command tracks the volume of mail being handled by each of the mailer programs that interface with it (those mailers defined in the /etc/sendmail.cf file).

To start the accumulation of mailer statistics, create the /etc/sendmail.st file by entering:

```
touch /etc/sendmail.st
```

If the **sendmail** command encounters errors when trying to record statistics information, the command writes a message through the syslog subroutine. These errors do not affect other operations of the sendmail command.

The sendmail command updates the information in the file each time it processes mail. The size of the file does not grow, but the numbers in the file do. They represent the mail volume since the time you created or reset the /etc/sendmail.st file.

Displaying Mailer Information

The statistics kept in the /etc/sendmail.st file are in a database format that cannot be read as a text file. To display the mailer statistics, enter the command:

```
/usr/sbin/mailstats
```

This reads the information in the /etc/sendmail.st file, formats it, and writes it to standard output in the following format:

msgs_from by	tes_from	msgs_to	bytes_to	Mailer
1	2.	1	2.01	local

The fields in the report have the following meanings:

msgs_from	Contains the number of messages received by the local machine from the indicated mailer.
bytes_from	Contains the number of bytes in the messages received by the local machine from the indicated mailer.
msgs_to	Contains the number of messages sent from the local machine using the indicated mailer.
bytes_to	Contains the number of bytes in the messages sent from the local machine using the indicated mailer.

If the **sendmail** command transmits mail directly to a file, such as the **\$HOME**/**dead.letter** file or an alias target, the message and byte counts are credited to the prog mailer.

Debugging sendmail

There are a large number of debug flags built into the **sendmail** command. Each debug flag has a number and level, where higher levels mean print more information. The convention is that levels greater than nine print out so much information that you would not want to see them except for debugging a particular piece of code. Debug flags are set using the -d flag as shown in the following example:

```
debug-flag: -d debug-list
debug-list: debug-flag[.debug-flag]*
debug-flag: debug-range[.debug-level]
debug-range:integer|integer-integer
debug-level:integer
```

For example:

```
-d12Set flag 12 to level 1
-d12.3Set flag 12 to level 3
-d3-17Set flags 3 through 17 to level 1
-d3-17.4Set flags 3 through 17 to level 4
```

Available debug flags are as follows:

-d0	General debugging.
-d1	Show send information.
-d2	End with finis().
-d3	Print the load average.
-d4	Enough disk space.
-d5	Show events.
-d6	Show failed mail.
-d7	The queue file name.
-d8	DNS name resolution.
-d9	Trace RFC1413 queries.
-d9.1	Make host name canonical.
-d10	Show recipient delivery.
-d11	Trace delivery.
-d12	Show mapping of relative host.
-d13	Show delivery.
-d14	Show header field commas.
-d15	Show network get request activity.
-d16	Outgoing connections.
-d17	List MX hosts.

Internet Message Access Protocol (IMAP) and Post Office Protocol (POP)

AIX provides two Internet-based electronic mail protocol server implementations for accessing mail remotely:

- Post Office Protocol (POP)
- Internet Message Access Protocol (IMAP)

Both the POP and IMAP servers store and provide access to electronic messages. Using these mail access protocols on a server eliminates the requirement that, to receive mail, a computer must always be up and running.

This implementation includes both a POP server and an IMAP server.

The POP server provides an off–line mail system, whereby a client, using POP client software, can remotely access a mail server to retrieve electronic mail messages. The client can either download the mail messages and immediately delete the messages from the server, or download the messages and leave the messages resident on the POP server. After the mail is downloaded to the client machine, all mail processing is local to the client machine. The POP server allows access to a user's mailbox by only one client at a time.

The IMAP server provides a superset of POP functionality but has a different interface. (Thus, there are IMAP–specific mail clients and POP–specific mail clients.) The IMAP server provides an off–line service, as well as an on–line service and a disconnected service. The IMAP protocol is designed to permit manipulation of remote mailboxes as if they were local. For example, clients can perform searches and mark messages with status flags such as "deleted" or "answered." In addition, messages can remain in the server's database until explicitly removed. The IMAP server also allows simultaneous interactive access to user mailboxes by multiple clients.

Both the IMAP and POP servers are used for mail access only. These servers rely on the Simple Mail Transfer Protocol (SMTP) for sending mail.

Both IMAP and POP are open protocols, based on standards described in RFCs. The IMAP server is based on RFC 1730, and the POP server is based on RFC 1725. Both are connection—oriented using TCP sockets. The IMAP server listens on well—known port 143, and the POP server listens on well—known port 110. Both servers are handled by the **inetd** daemon.

Configuring IMAP and POP Servers

Prerequisites

You must have root authority.

Procedure

- 1. Uncomment the **imapd** and **pop3d** entries in the /etc/inetd.conf file.
- 2. Refresh the **inetd** daemon by running the following command:

refresh -s inetd

Running Configuration Tests

You may want to run a few tests to verify your imapd and pop3d servers are ready for operation.

First, verify the servers are listening on their well-known ports. To do this, enter:

```
netstat -a | grep imap
netstat -a | grep pop
```

You should receive the following output from the **netstat** command:

```
*.* LISTEN
*.* LISTEN
                      *.imap2
                0
                      *.pop3
tcp
```

If you do not receive this output, recheck the entries in the /etc/inetd.conf file and rerun the refresh -s inetd command.

To test the configuration of the imapd server, you can telnet into the imap2 port, 143. When you telnet in, you should get the imapd prompt. You can enter the IMAP Version 4 commands as defined in RFC 1730. To run one of the commands, type a period (.), followed by a space, and then the command name. For example:

CommandName

Note that passwords are echoed when you telnet into the imapd server.

In the following telnet example, you must provide your own password where id_password is indicated in the **login** command.

```
telnet e-xbelize 143
Trying...
Connected to e-xbelize.austin.ibm.com.
Escape character is '^]'.
* OK e-xbelize.austin.ibm.com IMAP4 server ready
. login id id_password
. OK
. examine /usr/spool/mail/root
* FLAGS (\Answered \Flagged \Draft \Deleted \Seen)
* OK [PERMANENTFLAGS (\Answered \Flagged \Draft \Deleted \Seen \*)]
* 0 EXISTS
* 0 RECENT
* OK [UIDVALIDITY 823888143]
. OK [READ-ONLY] Examine completed
. logout
* BYE Server terminating connection
. OK Logout completed
Connection closed.
```

To test the configuration of the pop3d server, you can telnet into the Post Office Protocol Version 3 (POP3) port, 110. When you telnet in, you should get the pop3d prompt. You can enter the POP commands that are defined in RFC 1725. To run one of the commands, type a period (.), followed by a space, and then the command name. For example:

CommandName

Note that passwords are echoed when you telnet into the pop3d server.

In the following telnet example, you must provide your own password where id password is indicated in the **pass** command.

```
telnet e-xbelize 110
Trying...
Connected to e-xbelize.austin.ibm.com.
Escape character is '^]'.
+OK e-xbelize.austin.ibm.com POP3 server ready
user id
+OK Name is a valid mailbox
pass id_password
+OK Maildrop locked and ready
list
+OK scan listing follows
.
stat
+OK 0 0
quit
+OK
Connection closed.
```

syslog Facility

The IMAP and POP server software sends log messages to the **syslog** facility.

To configure your system for IMAP and POP logging through the **syslog** facility, you must be the root user. Edit the /**etc/syslog.conf** file, and add an entry for *.**debug** as follows:

```
*.debug /usr/adm/imapd.log
```

The usr/adm/imapd.log file must exist before the syslogd daemon re—reads the /etc/syslog.conf configuration file. To create this file, enter the following command:

```
touch /usr/adm/imapd.log
```

Then, refresh the **syslogd** daemon to re–read its configuration file. Enter the following command:

```
refresh -s syslogd
```

Mail Reference

This section provides a quick reference to the various Mail commands, files, and directories.

List of Mail Commands

This list includes commands for using and managing the mail program.

bugfiler Stores bug reports in specific mail directories. comsat Notifies users of incoming mail (daemon). mailq Prints the contents of the mail queue. mailstats Displays statistics about mail traffic.

newaliases Builds a new copy of the alias database from the

/etc/aliases file.

Handles remote mail received through the uucp command rmail

of the Basic Networking Utilities (BNU).

sendbug Mails a system bug report to a specific address.

sendmail Routes mail for local or network delivery.

smdemon.cleanu Cleans up the **sendmail** queue for periodic housekeeping.

List of Mail Files and Directories

This list of files and directories is arranged by function.

Using the Mail Program

/usr/share/lib/Mail.rc Sets local system defaults for all users of the mail program.

A text file you can modify to set the default characteristics

of the mail command.

\$HOME/.mailrc Enables the user to change the local system defaults for the

mail program.

\$HOME/mbox Stores processed mail for the individual user.

/usr/bin/Mail. /usr/bin/mail.

or /usr/bin/mailx

Specifies three names linked to the same program. The

mail program is one of the user interfaces to the mail

system.

/var/spool/mail Specifies the default mail drop directory. By default, all mail

is delivered to the /var/spool/mail/UserName file.

/usr/bin/bellmail Performs local mail delivery.

/usr/bin/rmail Performs remote mail interface for BNU.

Contains the log file and temporary files associated with the /var/spool/mqueue

messages in the mail queue.

Using the sendmail Command

/usr/sbin/sendmail The **sendmail** command.

/usr/ucb/mailq Links to the /usr/sbin/sendmail. Using mailq is equivalent

to using the /usr/sbin/sendmail -bp command.

/usr/ucb/newaliases Links to the /usr/sbin/sendmail file. Using newaliases is

equivalent to using the /usr/sbin/sendmail -bi command.

/etc/netsvc.conf Specifies the ordering of certain name resolution services. /usr/sbin/mailstats Formats and prints the sendmail statistics as found in the

/etc/sendmail.st file if it exists. The /etc/sendmail.st file is

the default, but you can specify an alternative file.

/etc/aliases Describes a text version of the aliases file for the sendmail

command. You can edit this file to create, modify, or delete

aliases for your system.

/etc/aliasesDB Describes a directory containing the aliases database files,

DB.dir and DB.pag, that are created from the /etc/aliases

file when you run the sendmail -bi command.

/etc/aliasesDBI Identifies a lock file for the aliasesDB database.

/etc/sendmail.cf Contains the sendmail configuration information in text

form. Edit the file to change this information.

/etc/sendmail.cfDB Contains the processed version of the /etc/sendmail.cf

configuration file. This file is created from the

/etc/sendmail.cf file when you run the /usr/sbin/sendmail

-bz command.

/etc/sendmail.cfDBI Identifies a lock file for the /etc/sendmail.cfDB database.

/etc/sendmail.nl Contains the sendmail National Language Support (NLS)

configuration information in text form. Edit the file to change

this information.

/etc/sendmail.nIDB Contains the processed version of the /etc/sendmail.nl

configuration file. This file is created from the

/etc/sendmail.nl file when you run the /usr/sbin/sendmail

-bn command.

/etc/sendmail.nIDBI Identifies a lock file for the /etc/sendmail.nIDB database.

/usr/lib/smdemon.cleanu Specifies a shell file that runs the mail queue and maintains

the sendmail log files in the /var/spool/mqueue directory.

/etc/sendmail.st Collects statistics about mail traffic. This file does not grow.

Use the /usr/sbin/mailstats command to display the contents of this file. Delete this file if you do not want to

collect this information.

/var/spool/mqueue Describes a directory containing the temporary files

associated with each message in the queue. The directory

can contain the log file.

/var/spool/cron/crontabs Describes a directory containing files that the cron daemon

reads to determine which jobs to start. The **root** file contains a line to start the **smdemon.cleanu** shell script.

List of Internet Message Access Protocol and Post Office Protocol Commands

/usr/sbin/imapd The Internet Message Access Protocol (IMAP) server

process.

/usr/sbin/pop3d The Post Office Protocol Version 3 (POP3) server process.

Chapter 3. Transmission Control Protocol/Internet Protocol

This chapter describes the Transmission Control Protocol/Internet Protocol (TCP/IP) suite of networking software. TCP/IP is a powerful and flexible industry–standard way of connecting multiple computers to other machines.

The topics discussed in this chapter are:

- Planning Your TCP/IP Network, on page 3-2
- Installation and Configuration for TCP/IP, on page 3-3
- TCP/IP Protocols, on page 3-6
- TCP/IP Network Adapter Cards, on page 3-36
- TCP/IP Network Interfaces, on page 3-47
- TCP/IP Addressing, on page 3-52
- TCP/IP Address and Parameter Assignment Dynamic Host Configuration Protocol (DHCP), on page 3-58
- Configuring TCP/IP, on page 3-90
- TCP/IP Daemons, on page 3-91
- TCP/IP Name Resolution, on page 3-97
- TCP/IP Routing, on page 3-121
- Serial Line Interface Protocol (SLIP), on page 3-131
- Asynchronous Point-to-Point Protocol (PPP) Subsystem, on page 3-136
- TCP/IP Quality of Service (QoS), on page 3-139
- TCP/IP Security, on page 3-145
- TCP/IP Problem Determination, on page 3-152
- TCP/IP Reference, on page 3-161

Note: Most of the tasks discussed in this chapter require root authority.

Planning Your TCP/IP Network

Because TCP/IP is such a flexible networking tool, there are many ways you can customize TCP/IP to fit the specific needs of your organization. Listed below are the major issues you need to consider when planning your network. (The details of these issues are discussed at length later. This list is intended only to introduce you to the issues.)

- Decide which type of network hardware you want to use: token-ring, Ethernet Version 2, IEEE 802.3, Fiber Distributed Data Interface (FDDI), Serial Optical Channel, or Serial Line Interface Protocol (SLIP).
- 2. Plan the physical layout of the network.
 - Consider which functions each host machine will serve. For example, you need to decide which machine or machines will serve as gateways before you cable the network.
- 3. Decide whether a *flat* network or a *hierarchical* network organization best fits your needs.
 - If your network is fairly small, at a single site, and consists of one physical network, then a flat network will probably suit your needs. If your network is very large or complex with multiple sites or multiple physical networks, a hierarchical network may be a more efficient network organization for you.
- 4. If your network will be connected to other networks, plan how your gateways need to be set up and configured. You need to:
 - a. Decide which machine or machines will serve as gateways.
 - b. Decide whether you need to use static or dynamic routing, or a combination of the two. If you choose dynamic routing, decide which routing daemons each gateway will use in light of the types of communications protocols you need to support.
- 5. Decide on an addressing scheme.
 - If your network will not be part of a larger internetwork, choose the addressing scheme that best fits your needs. If you want your network to be connected to a larger internetwork such as the Internet, you will need to obtain an official set of addresses.
- 6. Decide whether your system needs to be divided into subnets. If so, decide how you will assign subnet masks.
- 7. Decide on a naming scheme. Each machine on the network needs its own unique host name.
- 8. Decide whether your network needs a name server for name resolution or if using the /etc/hosts file will be sufficient.
 - If you choose to use name servers, consider the type of name servers you need and how many you need to serve your network efficiently.
- 9. Decide the types of services your network will provide to remote users; for example, mail, remote login, remote command execution, file sharing, and other services.

Installation and Configuration for TCP/IP

For information on installing Transmission Control Protocol/Internet Protocol (TCP/IP), see the AIX Installation Guide.

Configuring TCP/IP

Once the TCP/IP software is installed on your system, you are ready to begin configuring your system.

Many TCP/IP configuration tasks can be performed in more than one way, either by:

- Using the Web-based System Manager wsm network application (fast path wsm network)
- Using the System Management Interface Tool (SMIT)
- · Editing a file format
- · Issuing a command at the shell prompt

For example, the **rc.net** shell script performs required minimum host configuration for TCP/IP during the system startup process (the **rc.net** script is run by the configuration manager program during the second boot phase). By using Web-based System Manager or SMIT to perform the host configuration, the **rc.net** file is configured automatically.

Alternatively, you can configure the **rc.net** file using a standard editor. With this method, you can specify the traditional UNIX TCP/IP configuration commands such as **ifconfig**, **hostname**, and **route**. (Refer to the "List of TCP/IP Commands", on page 3-161 for further information.)

A few tasks, such as configuring a name server, cannot be done using Web-based System Manager or SMIT. (Refer to Fast Paths for TCP/IP in AIX 4.3 System Management Guide: Operating System and Devices for a list of tasks in TCP/IP that can be performed using SMIT.)

Configuring Hosts

Each host machine on your network will need to be configured to function according to the needs of the end–users and the network as a whole. For each host on the network, you must configure the network interface, set the Internet address, and set the host name. You also need to set up static routes to gateways or other hosts, specify daemons to be started by default, and set up the /etc/hosts file for name resolution (or set up the host to use a name server for name resolution).

Configuring Hosts as Servers

If the host machine is to have a specific function, for example, if it is to serve as a gateway, file server, or name server, you must perform the necessary configuration tasks after the basic configuration is complete.

For example, if your network is organized hierarchically and you want to use the **DOMAIN** protocol to resolve names into Internet addresses, you will need to configure at least one name server to provide this function for your network.

Remember, a server host does not have to be a dedicated machine: it can be used for other things as well. For example, if the name server function for your network is fairly small, the machine might also be used as a workstation or as a file server for your network.

Note: If your system has either NIS or NIS+ installed, these services can also provide name resolution. For more information, see *AIX 4.3 NIS/NIS+ Guide*.

Configuring Gateways

If your network is going to communicate with other networks, you will need to configure at least one gateway host. You will need to consider which communications protocols you will need to support, and then use whichever routing daemon (the **routed** or **gated** daemon) supports those protocols.

TCP/IP System Manager Commands

The following list contains the commands you will use to configure and manage your TCP/IP network:

arp Displays or changes the Internet address to hardware address translation

tables used by the Address Resolution protocol.

finger Returns information about users on a specified host.

host Shows the Internet address of a specified host or the host name of a

specified Internet address.

hostname Shows or sets the Internet name and address of the local host.

ifconfig Configures network interfaces and their characteristics.

netstat Shows local and foreign addresses, routing tables, hardware statistics, and

a summary of packets transferred.

no Sets or shows current network kernel options.

ping Determines whether a host is reachable.

route Permits you to manipulate the routing tables manually.

ruptime Shows status information on hosts that are connected to local physical

networks and are running the **rwhod** server.

rwho Shows status information for users on hosts that are connected to local

physical networks and running the **rwhod** server.

setclock Reads the network time service and sets the time and date of the local

host accordingly.

timedc Returns information about the timed daemon.trpt Reports protocol tracing on TCP sockets.whois Provides the Internet name directory service.

Configuring a TCP/IP Network Checklist

Use the following procedure as a guide for configuring your network, both initially and when you make changes over time. Be certain that you have read and understood the appropriate material. Taking extra time now to understand what you are doing will save you time and frustration later, by avoiding having to correct an earlier mistake.

After you have your network up and running properly, you may find it useful to refer to this checklist for the purpose of debugging the inevitable problems that occur in every network.

Prerequisites

- Network hardware is installed and cabled (refer to "TCP/IP Network Adapter Cards", on page 3-36).
- 2. TCP/IP software is installed (see the AIX Installation Guide).

Procedure

- 1. Read "TCP/IP Protocols", on page 3-6 for the basic organization of TCP/IP. You should understand:
 - the layered nature of TCP/IP (that is, different protocols reside at different layers)
 - how data flows through the layers
- 2. Minimally configure each host machine on the network. Specifically, this means adding a network interface, assigning an IP address, and assigning a hostname to each host, as well as defining a default route to your network. First, read "TCP/IP Network Interfaces", on page 3-47, "TCP/IP Addressing", on page 3-52, and "Choosing Names for Hosts on Your Network", on page 3-99. Then, follow the instructions in "Configuring TCP/IP", on page 3-90.

Note: Each machine on the network needs this basic configuration whether it will be an end—user host, a file server, a gateway, or a name server.

- 3. Configure and start the **inetd** daemon on each host machine on the network. Read "TCP/IP Daemons", on page 3-91 and then follow the instructions in "Configuring the inetd Daemon", on page 3-94.
- 4. Configure each host machine either to perform local name resolution or to use a name server. If you are setting up a hierarchical DOMAIN network, configure at least one host to function as a name server. Read and follow the instructions in "TCP/IP Name Resolution", on page 3-97.
- 5. If your network will communicate with any remote networks, configure at least one host to function as a gateway. The gateway can use static routes or a routing daemon to perform internetwork routing. Read and follow the instructions in "TCP/IP Routing", on page 3-121.
- 6. Decide which services each host machine on the network will use. By default, all services are available. Follow the instructions in "Client Network Services", on page 3-95 if you wish to make a particular service unavailable.
- 7. Decide which hosts on the network will be servers, and which service a particular server will provide. Follow the instructions in "Server Network Services", on page 3-96 to start the server daemons you wish to run.
- 8. Configure any remote print servers you will need. See "Printer Overview" in *AIX Guide to Printers and Printing* for more information.
- 9. If desired, configure a host to use or serve as master time server for the network. See the **timed** daemon in the *AIX Commands Reference* for more information.

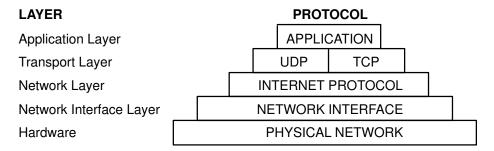
TCP/IP Protocols

The topics discussed in this section are:

- IP6 Overview, on page 3-9
- IPv6 in AIX: Supplementary Information, on page 3-16
- Packet Tracing, on page 3-20
- Network Interface Packet Headers, on page 3-20
- Internet Network–Level Protocols, on page 3-23
- Internet Transport-Level Protocols, on page 3-28
- Internet Application–Level Protocols, on page 3-31
- Assigned Numbers, on page 3-35

Protocols are sets of rules for message formats and procedures that allow machines and application programs to exchange information. These rules must be followed by each machine involved in the communication in order for the receiving host to be able to understand the message.

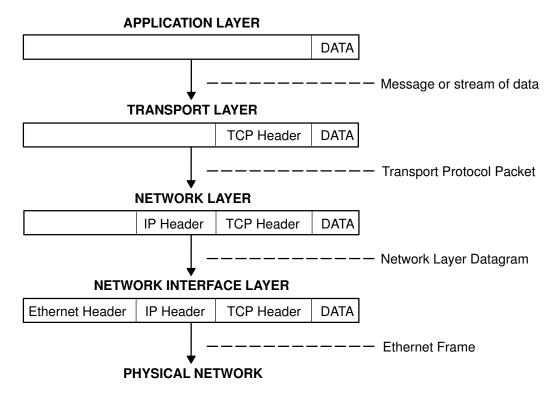
The TCP/IP suite of protocols (see figure) can be understood in terms of layers (or levels).



TCP/IP carefully defines how information moves from sender to receiver. First, application programs send messages or streams of data to one of the Internet Transport Layer Protocols, either the User Datagram Protocol (UDP) or the Transmission Control Protocol (TCP). These protocols receive the data from the application, divide it into smaller pieces called *packets*, add a destination address, and then pass the packets along to the next protocol layer, the Internet Network layer.

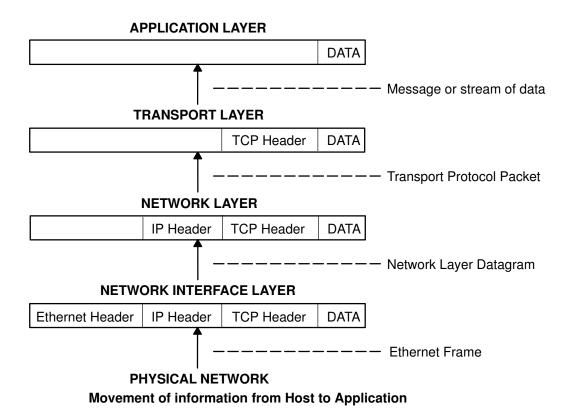
The Internet Network layer encloses the packet in an Internet Protocol (IP) datagram, puts in the datagram header and trailer, decides where to send the datagram (either directly to a destination or else to a gateway), and passes the datagram on to the Network Interface layer.

The Network Interface layer accepts IP datagrams and transmits them as *frames* over a specific network hardware, such as Ethernet or Token–Ring networks (see figure).

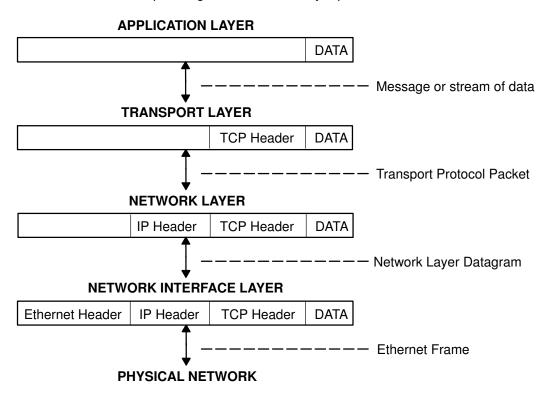


Movement of information from Sender Application to Receiver Host

Frames received by a host go through the protocol layers in reverse. Each layer strips off the corresponding header information, until the data is back at the application layer (see figure). Frames are received by the Network Interface layer (in this case, an Ethernet adapter). The Network Interface layer strips off the Ethernet header, and sends the datagram up to the Network layer. In the Network layer, the Internet Protocol strips off the IP header and sends the packet up to the Transport layer. In the Transport layer, the Transmission Control Protocol (in this case) strips off the TCP header and sends the data up to the Application layer.



Since hosts on a network both send and receive information simultaneously, the "Host Data Transmission and Reception" figure more accurately represents a host as it communicates.



Note: Headers are added and stripped in each protocol layer as data is transmitted and received by a host.

Host Data Transmission and Reception

Internet Protocol (IP) Version 6 Overview

Internet Protocol (IP) Version 6 (IPv6 or IPng) is the next generation of the Internet Protocol and has been designed to be an evolutionary step from IP Version 4 (IPv4). While IPv4 has allowed the development of a global Internet, it is not capable of carrying much farther into the future because of two fundamental factors: limited address space and routing complexity. IPv4's 32–bit addresses do not provide enough flexibility for global Internet routing. The deployment of Classless InterDomain Routing (CIDR) has extended the lifetime of IPv4 routing by a number of years, but the effort to manage the routing, however, continues to increase. Even if IPv4 routing could be scaled up, the Internet will eventually run out of network numbers.

The Internet Engineering Task Force (IETF) recognized that IPv4 would not be able to support the phenomenal growth of the Internet, so the IETF IPng working group was formed. Of the proposals that were made, Simple Internet Protocol Plus (SIPP) was chosen as an evolutionary step in the development of IP. This was renamed to IPng, and RFC1883 was finalized in December of 1995.

IPv6 extends the maximum number of Internet addresses to handle the ever increasing Internet user population. As an evolutionary change from IPv4, IPv6 has the advantage of allowing the new and the old to coexist on the same network. This coexistence enables an orderly migration from IPv4 (32 bit addressing) to IPv6 (128 bit addressing) on an operational network.

This overview is intended to give the reader a general understanding of the IP*ng* protocol. For detailed information, please refer to RFCs 1883, 1884, 1885, 1886, 1970, 1971, and 2133.

Expanded Routing and Addressing

IPv6 increases the IP address size from 32 bits to 128 bits, thereby supporting more levels of addressing hierarchy, a much greater number of addressable nodes, and simpler autoconfiguration of addresses.

In IPv6, there are three types of addresses:

unicast

A packet sent to a unicast address is delivered to the interface identified by that address. A unicast address has a particular scope: link-local, site-local, global. There are also two special unicast addresses:

- ::/128 (unspecified address)
- ::1/128 (loopback address)

In IPv6, this is only one address, not an entire network.

multicast

A packet sent to a multicast address is delivered to all interfaces identified by that address. A multicast address is identified by the prefix ff::/8. As with unicast addresses, multicast addresses have a similar scope: node—local, link—local, site—local, and organization—local.

anycast

An anycast address is an address that has a single sender, multiple listeners, and only one responder (normally the "nearest" one, according to the routing protocols' measure of distance). An example may be several web servers listening on an anycast address. When a request is sent to the anycast address, only one responds.

An anycast address is indistinguishable from a unicast address. A unicast address becomes an anycast address when more than one interface is configured with that address.

Note: There are no broadcast addresses in IPv6. Their function has been superseded by the multicast address.

Auto Configuration

The primary mechanisms available that enable a node to boot up and start communicating with other nodes over an IPv4 network are hard–coding, BOOTP, and DHCP. These work fairly well, but each has its own difficulties: hard–coding makes renumbering inordinately difficult; BOOTP and DHCP require that the host broadcast and are dependent on a remote server.

IPv6 introduces the concept of *scope* to IP addresses, one of which is link–local. This allows a host to construct a valid address from the predefined link–local prefix and its local identifier, which is typically the medium access control (MAC) address of the interface to be configured. Using this address, the node can multicast to a server, rather than broadcast and, for a fully–isolated subnet, may not need any other address configuration.

Meaningful Addresses

With IPv4, the only generally recognizable meaning in addresses are broadcast (typically all 1 or all 0), and classes (for example, a class D is multicast). With IPv6, the prefix can be quickly examined to determine *scope* (for example, link–local), multicast versus unicast, and a mechanism of assignment (provider–based, geography–based, etc.).

Routing information may be explicitly loaded into the upper bits of addresses as well, but this has not yet been finalized by the IETF (for provider–based addresses, routing information is implicitly present in the address).

Duplicate Address Detection

When an interface is initialized or reinitialized, it uses autoconfiguration to tentatively associate a link–local address with that interface (the address is not yet assigned to that interface in the traditional sense). At this point, the interface joins the all–nodes and solicited–nodes multicast groups, and sends a neighbor discovery message to these groups. By using the multicast address, the node can determine whether that particular link–local address has been previously assigned, and choose an alternate address. This reduces a common error in network management, namely assigning the same address to two different interfaces on the same link. (It is still possible to create duplicate global–scope addresses for nodes that are not on the same link.)

Neighbor Discovery/Stateless Address Autoconfiguration

Neighbor Discovery (ND) protocol for IPv6 is used by nodes (hosts and routers) to determine the link–layer addresses for neighbors known to reside on attached links, and maintain per–destination routing tables for active connections. Hosts also use Neighbor Discovery to find neighboring routers that are willing to forward packets on their behalf and detect changed link–layer addresses. Neighbor Discovery Protocol (NDP) uses the Internet Control Message Protocol (ICMP) Version 6 with its own unique message types. In general terms, the IPv6 Neighbor Discovery protocol corresponds to a combination of the IPv4 Address Resolution Protocol (ARP), ICMP Router Discovery (RDISC), and ICMP Redirect (ICMPv4), but with many improvements over these IPv4 protocols.

IPv6 defines both a stateful and stateless address autoconfiguration mechanism. *Stateless autoconfiguration* requires no manual configuration of hosts; minimal, if any, configuration of routers; and no additional servers. The stateless mechanism allows a host to generate its own addresses using a combination of locally available information and information advertised by routers. Routers advertise prefixes that identify the subnets associated with a link, while hosts generate an interface token that uniquely identifies an interface on a subnet. An address is formed by combining the two. In the absence of routers, a host can only generate link–local addresses. However, link–local addresses are sufficient for allowing communication among nodes attached to the same link.

Routing Simplification

To simplify routing issues, IPv6 addresses are considered in two parts: a prefix and an ID. This may seem no different from the IPv4 net–host address breakdown, but it has two advantages:

no class No fixed number of bits for prefix or ID, which allows for a

reduction in loss due to over-allocation

nesting An arbitrary number of divisions can be employed by

considering different numbers of bits as the prefix.

Case 1:

128 bits	
Node Address	

Case 2:

n bits	128– <i>n</i> bits
Subnet prefix	Interface ID

Case 3:

n bits	80- <i>n</i> bits	48 bits
Subscriber prefix	Subnet ID	Interface ID

Case 4:

s bits	n bits	<i>m</i> bits	128– <i>s–n–m</i> bits
Subscriber prefix	Area ID	Subnet ID	Interface ID

Generally, IPv4 cannot go beyond Case 3, even with VLSM. (This is as much an artifact of the shorter address length as the definition of variable length prefixes, but is worth noting nonetheless).

Header Format Simplification

IPv6 simplifies the IP header, either by removing entirely or by moving to an extension header, some of the fields found in the IPv4 header and it defines a more flexible format for optional information (the extension headers). Specifically, note the absence of:

- header length (length is constant)
- · identification
- flags
- fragment offset (moved into fragmentation extension headers)
- header checksum (upper-layer protocol or security extension header handles data integrity)

IPv4 Header:

Version	IHL	Type of Service	Total Len	gth
Identification		Flags	Fragment Offset	
Time to Live Protocol Header Checksum		ecksum		
Source Address				
Destination Address				
Options Pad			Padding	

IPv6 Header:

Version	Prio		Flow Label	
Payload Length		Next Header	Hop Limit	
Source Address				
Destination Address				

IPng includes an improved options mechanism over IPv4. IPv6 options are placed in separate extension headers that are located between the IPv6 header and the transport–layer header in a packet. Most extension headers are not examined or processed by any router along a packet's delivery path until it arrives at its final destination. This

mechanism facilitates a major improvement in router performance for packets containing options. In IPv4 the presence of any options requires the router to examine all options.

Another improvement is that, unlike IPv4 options, IPv6 extension headers can be of arbitrary length and the total amount of options carried in a packet is not limited to 40 bytes. This feature, plus the manner in which it is processed, permits IPv6 options to be used for functions that were not practical in IPv4, such as the IPv6 Authentication and Security Encapsulation options.

To improve the performance when handling subsequent option headers and the transport protocol which follows, IPv6 options are always an integer multiple of eight octets long to retain this alignment for subsequent headers.

By using extension headers instead of a protocol specifier and options fields, newly defined extensions can be integrated more easily.

Current specifications define extension headers in the following ways:

- Hop-by-hop options that apply to each hop (router) along the path
- Routing header for loose/strict source routing (used infrequently)
- A fragment defines the packet as a fragment and contains information about the fragment (IPv6 routers do not fragment)
- Authentication IP Security
- Encryption IP Security
- Destination options for the destination node (ignored by routers)

Improved Quality-of-Service/Traffic Control

While quality of service can be controlled by use of a control protocol such as RSVP, IPv6 provides for explicit priority definition for packets by using the priority field in the IP header. A node can set this value to indicate the relative priority of a particular packet or set of packets, which can then be used by the node, one or more routers, or the destination to make choices concerning the packet (that is, dropping it or not).

IPv6 specifies two types of priorities, those for congestion—controlled traffic, and those for non—congestion—controlled traffic. No relative ordering is implied between the two types.

Congestion—controlled traffic is defined as traffic that responds to congestion through some sort of "back—off" or other limiting algorithm. Priorities for congestion—controlled traffic are:

- 0 uncharacterized traffic
- 1 "filler" traffic (for example, netnews)
- 2 unattended data transfer (for example, electronic mail)
- 3 (reserved)
- 4 attended bulk transfer (for example, FTP)
- 5 (reserved)
- 6 interactive traffic (for example, Telnet)
- 7 control traffic (for example, routing protocols)

Non-congestion-controlled traffic is defined as traffic that responds to congestion by dropping (or simply not resending) packets, such as video, audio, or other real-time traffic. Explicit levels are not defined with examples, but the ordering is similar to that for congestion-controlled traffic:

- The lowest value should be used for traffic that the source is most willing to have discarded.
- The highest value should be used for traffic that the source is least willing to have discarded.

This priority control is only applicable to traffic from a particular source address. Control traffic from one address is not an explicitly higher priority than attended bulk transfer from another address.

Flow Labeling

Outside of basic prioritization of traffic, IPv6 defines a mechanism for specifying a particular flow of packets. In IPv6 terms, a *flow* is defined as "a sequence of packets sent from a particular source to a particular (unicast or multicast) destination for which the source desires special handling by the intervening routers."

This flow identification may be used for priority control, but may also be used for any number of other controls.

The flow label is chosen randomly, and should not be construed as identifying any characteristic of the traffic other than the flow to which it belongs. This means that a router cannot determine that a packet is a particular type (for example, FTP) by examining the flow label. It will, however, be able to determine that it is part of the same sequence of packets as the last packet containing that label.

Note: In AIX 4.3 and until IPv6 is in general use, the flow label is mostly experimental. Uses and controls involving flow labels have not yet been defined nor standardized.

Jumbograms

Host-to-Host

Router-to-Host

An IPv4 packet size is limited to 64K. Using the jumbo payload extension header, an IPv6 packet can be up to 2³² octets (slightly over 4 gigabytes).

Tunneling

The key to a successful IPv6 transition is compatibility with the existing installed base of IPv4 hosts and routers. Maintaining compatibility with IPv4 while deploying IPv6 streamlines the task of transitioning the Internet to IPv6.

In most cases, the IPv6 routing infrastructure will evolve over time. While the IPv6 infrastructure is being deployed, the existing IPv4 routing infrastructure can remain functional, and can be used to carry IPv6 traffic. Tunneling provides a way to use an existing IPv4 routing infrastructure to carry IPv6 traffic.

IPv6/IPv4 hosts and routers can tunnel IPv6 datagrams over regions of IPv4 routing topology by encapsulating them within IPv4 packets. Tunneling can be used in a variety of ways:

Router-to-Router	IPv6/IPv4 routers interconnected by an IPv4 infrastructure can tunnel IPv6 packets between themselves. In this case, the tunnel spans one segment of the end–to–end path that the IPv6 packet takes.
Host-to-Router	IPv6/IPv4 hosts can tunnel IPv6 packets to an intermediary IPv6/IPv4 router that is reachable through an IPv4 infrastructure. This type of tunnel spans the first segment of the packet's end-to-end path.

IPv6/IPv4 hosts that are interconnected by an IPv4 infrastructure can tunnel IPv6 packets between themselves. In this case, the tunnel spans the entire end—to—end path that the packet takes.

IPv6/IPv4 routers can tunnel IPv6 packets to their final destination IPv6/IPv4 host. This tunnel spans only the last segment of the end-to-end path.

Tunneling techniques are usually classified according to the mechanism by which the encapsulating node determines the address of the node at the end of the tunnel. In router—to—router or host—to—router methods, the IPv6 packet is being tunneled to a router. In

host-to-host or router-to-host methods, the IPv6 packet is tunneled all the way to its final destination.

The entry node of the tunnel (the encapsulating node) creates an encapsulating IPv4 header and transmits the encapsulated packet. The exit node of the tunnel (the decapsulating node) receives the encapsulated packet, removes the IPv4 header, updates the IPv6 header, and processes the received IPv6 packet. However, the encapsulating node needs to maintain soft state information for each tunnel, such as the maximum transmission unit (MTU) of the tunnel, to process IPv6 packets forwarded into the tunnel.

IPv6 Security

For details about IP Security, versions 4 and 6, see Chapter 4. Internet Protocol (IP) Security, on page 4-1.

IPv6 Multihomed Link Local and Site Local Support

A host can have more than one interface defined. A host with two or more active interfaces is called multihomed. Each interface has a link local address associated with it. Link local addresses are sufficient for allowing communication among nodes attached to the same link.

A multihomed host will have two or more associated link local addresses. In AIX's IPv6 implementation, there are 4 options to handling how link–layer address resolution is resolved on multihomed hosts. Option 1 is turned on by default.

Option 0

No multihomed actions are taken. Transmits will go out on the first link local interface. When the Neighbor Discovery Protocol must perform address resolution, it multicasts a Neighbor Solicitation message out on each interface with a link local address defined. NDP queues the data packet until the first Neighbor Advertisement message is received. The data packet is then sent out on this link.

Option 1

When the Neighbor Discovery Protocol must perform address resolution (i.e., when sending a data packet to a destination and the the link–layer information for the next–hop is not in the Neighbor Cache), it multicasts a Neighbor Solicitation message out on each interface with a link local address defined. NDP then queues the data packet until it gets the link–layer information. NDP then waits until a response is received for each interface. This guarantees that the data packets are sent on the appropriate outgoing interfaces. If NDP did not wait, but responded to the first Neighbor Advertisement received, it would be possible for a data packet to be sent out on a link not associated with the packet's source address. Because NDP must wait, a delay in the first packet being sent will occur. However, the delay will have occurred anyway in waiting for the first response.

Option 2

Multihomed operation is allowed, but dispatching of a data packet is limited to the interface specified by main_if6. When the Neighbor Discovery Protocol must perform address resolution, it multicasts a Neighbor Solicitation message out on each interface with a link local address defined. It then waits for a Neighbor Advertisement message from the interface specified by main_if6 (see the **no** command). Upon receiving a response from this interface, the data packet is sent out on this link.

Option 3

Multihomed operation is allowed, but dispatching of a data packet is limited to the interface specified by main_if6 and site local addresses will only be routed for the interface specified by main_site6 (see the **no** command). The Neighbor Discovery Protocol will operate just as it does for Option 2. For applications that route data packets using site local addresses, on a multihomed host only the site local address specified by main_site6 will be used.

IPv6 in AIX: Supplementary Information

Address Format Definitions

IPv6 Address

An IPv6 address is 128 bits long, and is written as 8 hexadecimal 16-bit fields separated by colons:

```
XXXX:XXXX:XXXX:XXXX:XXXX:XXXX:XXXX
```

An IPv6 address may be shortened by compressing zeros. This can occur only once per address. Zeros are compressed by "::" (for example 4e80::54f:1234).

The prefix corresponds to the first bits of the address. It is represented by the notation "<IPv6 address>/<length>". For example:

- In 4e80::123:456:789/10, the prefix corresponds to the 10 first bits of the 4e80::123:456:789 address.
- The prefix of multicast addresses is FF00::/8, meaning that a multicast address starts with FF.

The prefix defines either:

- · the address type,
- or, roughly, the equivalent of the IPv4 "network" and "subnet" address parts,
- or the aggregate that defines the route.

The interface ID corresponds to the last bits of the address, and represents a unique host identifier. It is roughly equivalent to the IPv4 "host" part.

Aggregatable Global Unicast Address

The IPv6 aggregatable global unicast address format will be used in the future for the addressing plan in the Internet Protocol.

See draft-ietf-ipngwg-unicast-aggr-04.txt and draft-ietf-ipngwg-tla-assignment-03.txt.

EUI-64

EUI-64 is a 64-bit interface ID automatically built from the MAC IEEE 802 address.

The implementation is described in the addressing architecture draft draft—ietf—ipngwg—addr—arch—v2—02.txt.

Special Addresses

Mapped Address

A mapped address is an IPv4-based address and is used for compatibility with IPv4 hosts:

- An IPv6 socket can communicate with an IPv4—only host using a mapped address as a destination address.
- An IPv6 socket can bind to a mapped address to receive packets from IPv4—only host.

A mapped address is represented either as:

• ::FFFF:a.b.c.d where a.b.c.d is an IPv4 address.

or

 ::FFFF:XXXX:YYYY where XXXXYYYY is the hexadecimal representation of the a.b.c.d IPv4 address.

Compat Address

A compat address is an IPv4-based address and is used for compatibility with IPv4 routing in SIT tunneling. When an IPv6 packet is sent to a compat address, it is encapsulated in an IPv4 packet. The destination must decapsulate the packet to process it.

A compat address is represented either as:

• ::a.b.c.d where a.b.c.d is an IPv4 address.

or

• :: XXXX: YYYY where XXXXYYYY is the hexadecimal representation of the a.b.c.d IPv4 address.

Unspecified Address

The unspecified address is composed of 16 null bytes, represented as "0:0:0:0:0:0:0:0:0", or simply "::".

Loopback Address

The loopback address, represented as "0:0:0:0:0:0:0:1", (or simply "::1"), may be used by a node to send an IPv6 datagram to itself.

Link-local Address

The stations that are not yet configured with an IPv6 address may use the link-local address.

The link-local addresses are only defined within a link and can only be used by stations connected to the same link or to the same local network. They are automatically configured when the interface is initialised and allow to communicate between neighbor nodes.

These addresses are composed of a link-local prefix (fe80::/64) appended to a 64 bits interface ID which is generally at the EUI format. For example:

fe80::4260:8cff:fe2c:9c38

Configuration Methods

An IPv6 network is configured in the following manner:

- Perform all IPv4—related configuration
- Run the autoconf6 command to generate link-local addresses, the sit0 interface (on the existing IPv4), and to install basic unicast and multicast routes.
- Run ndpd-host daemon to listen for Router Advertisements (RA). As RAs are received, addresses are added to the appropriate interfaces, and routes are created as necessary.

To perform these operations using SMIT, proceed as follows:

- 1. smit tcpip
- 2. Select the **IPv6 configuration** item and follow the menu.
- 3. Select the **IPv6 daemon/Process Configuration** item to start the **autoconf6** process and the **ndpd–host** subsystem.

Tunneling

Tunneling allows IPv6 communication in a primarily IPv4 world, and facilitates gradual transition to IPv6.

A route directs a packet through a tunnel interface, which encapsulates the packet (by prepending an appropriate IPv4 header), and then calls the IPv4 output routine.

AIX 4.3 supports two kinds of IPv6 tunnels:

- Simple Internet Transition (SIT) interfaces
- · Configured Tunnel Interface (CTI) interfaces.

SIT Tunneling

SIT tunnels are well suited for host-to-host communication through an IPv4 routing hierarchy.

The interface is configured as a compat address, using IPv4 address from the associated IPv4 interface. All traffic destined for a compat address is routed through this interface.

The **sit0** interface is typically added by **autoconf6** when the system is started.

A SIT tunnel is open-ended, in that only the local side is configured.

CTI Tunneling

CTI tunnels are well suited for router-to-router communication.

The interface is configured using IPv6 and IPv4 source and destination addresses specified by the user. Traffic is routed through an appropriate cti interface by a particular routing table entry.

The cti0 (and any other cti interfaces) must be added by a system administrator.

A cti tunnel is closed-ended, in that both ends must be configured.

Name Servers

The name servers process is very similar to IPv4.

- The forward zone uses IPv6 AAAA records instead of A records (see the example).
- The reverse zone uses **IP6.INT** domain instead of **in–addr.arpa** domain, but still using **PTR** records (see the example).
- **IP6.INT** domain **PTR** records are the reversed IPv6 address, with each nibble separated by a dot.
- The IPv4 connections are used to transmit requests.

Example of forward zone file

```
$ORIGIN merl.century.com

ipv6 9999999 IN SOA gtw.merl.century.com.

root.gtw.merl.century.com.

( 1072 3600 300 3600000 3600)

9999999 IN NS elvis.merl.century.com.

$ORIGIN ipv6.merl.century.com.

nobody-v6-11 IN AAAA fe80::99:9999:1997

danny-v6-co IN AAAA ::9.3.114.56

danny-v6-11 IN AAAA fe80::260:97ff:fea3:9234
```

Example of reverse zone file

Name Service Clients

For binary compatibility, the resolvers only return IPv4 addresses unless the **RES_USE_INET6** resolver option is set. This option can be set using different methods:

 Setting the environment variable RES_OPTIONS to the value inet6 will turn on the option for all processes running in the current shell. • Calling **res_init()** and setting the **RES_USE_INET6** flag in **_res.options** (in application code) will turn on the option for the calling process.

```
(void) res_init();
_res.options |= RES_USE_INET6;
```

Porting Applications

For most (but not all) tokens that contain "in" or "inet", use the IPv6 name instead. Some examples:

v4 token	v6 token
AF_INET	AF_INET6
sockaddr_in	sockaddr_in6
in_addr	in6_addr
<netinet in.h=""></netinet>	<netinet in.h=""></netinet>
<netinet ip.h=""></netinet>	<netinet ip6.h=""></netinet>
INADDRSZ	IN6ADDRSZ
sin_addr	sin6_addr
sin_len	sin6_len

Use **inet_ntop()** and **inet_pton()** instead of **inet_ntoa()** and **inet_aton()** for conversions between ASCII and binary formats.

Use gethostbyname2() or getaddrinfo() instead of gethostbyname().

Also read RFC 2133.

Path MTU Discovery

The path MTU discovery algorithm is always enabled in IPv6.

AIX Implementations Limitations

Network Interfaces

IPv6 supports the following network interfaces:

- Token-ring
- Ethernet (Standard Ethernet and IEEE 802.3)
- FDDI

Applications

IPv6 supports the following applications:

- ping
- telnet/telnetd
- ftp/ftpd
- tftp/tftpd
- crash/ndb
- iptrace/ipreport/tcpdump
- traceroute
- resolver routines/named
- inetd
- rsh/rcp/rshd
- rexec/rexecd
- rlogin/rlogind
- ifconfig/netstat/route/nslookup
- mail/sendmail
- autoconf6/ndpd-host/ndp
- nslookup
- ndpd-router from AIX 4.3.2 and later
- gated with RIPng, BGP4+ from AIX 4.3.2 and later

Routing in AIX 4.3.0 and 4.3.1

- AIX IPv6 can only be used as an IPv6 host or as an IPv6/IPv4 router by using tunneling. IPv6 cannot be used as an IPv6/IPv6 router.
- It is not possible to forward and to route packets between two IPv6 interfaces.
- Also, the routing applications are not updated to understand IPv6 routing protocols.
- The router advertisements are not generated and the ndpd-router application is not supported.

Routing from AIX 4.3.2

- AIX IPv6 can be a host and a router by allowing the station to be multi-homed. IPv6 is able to forward and route packets between two IPv6 intefaces.
- The ndpd-router command is supported.
- RIPng routing protocol is supported by npd-router and gated daemons.
- BGP4+ is supported by the gated daemon.

Anycast address

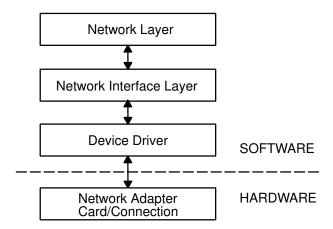
This function is not supported.

Packet Tracing

Packet tracing is the process by which you can verify the path of a packet through the layers to its destination. The **iptrace** command performs network interface level packet tracing. The **ipreport** command issues output on the packet trace in both hexadecimal and ASCII format. The **trpt** command performs transport protocol level packet tracking for the TCP. The **trpt** command output is more detailed, including information on time, TCP state, and packet sequencing.

Network Interface Packet Headers

At the Network Interface layer, packet headers are attached to outgoing data (see figure). Packets are then sent through the network adapter to the appropriate network. Packets may pass through many gateways before reaching their destinations. At the destination network, the headers are stripped from the packets and the data is sent to the appropriate host.



Packet Flow through Network interface structure

Packet header information for several of the more common network interfaces follows.

Ethernet Adapter Frame Headers

The following table represents an Internet Protocol (IP) or Address Resolution Protocol (ARP) frame header for the Ethernet adapter.

Ethernet Adapter Frame Header			
Field	Length	Definition	
DA	6 bytes	Destination address.	
SA	6 bytes	Source address. If bit 0 of this field is set to 1, it indicates that routing information (RI) is present.	
Туре	2 bytes	Specifies whether the packet is IP or ARP. The type number values are listed below.	

Type field numbers:

IP 0800 ARP 0806

Token-Ring Frame Headers

The medium access control (MAC) header for the token–ring adapter is composed of five fields, as shown in the following table.

Token-Ring MAC Header			
Field	Length	Definition	
AC	1 byte	Access control. The value in this field x'00' gives the header priority 0.	
FC	1 byte	Field control. The value in this field x'40' specifies the Logical Link Control frame.	
DA	6 bytes	Destination address.	
SA	6 bytes	Source address. If bit 0 of this field is set to 1, it indicates that routing information (RI) is present.	
RI	18 bytes	Routing information. The valid fields are discussed below.	

The MAC header consists of two routing information fields of 2 bytes each: routing control (RC) and segment numbers. A maximum of eight segment numbers may be used to specify recipients of a limited broadcast. RC information is contained in bytes 0 and 1 of the RI field. The settings of the first two bits of the RC field have the following meanings:

bit $(0) = 0$	Use the nonbroadcast route specified in the RI field.
bit (0) = 1	Create the RI field and broadcast to all rings.
bit $(1) = 0$	Broadcast through all bridges.
bit (1) = 1	Broadcast through limited bridges.

The logical link control (LLC) header is composed of five fields, as shown in the following LLC header table.

802.3 LLC Header			
Field	Length	Definition	
DSAP	1 byte	Destination service access point. The value in this field is x'aa'.	
SSAP	1 byte	Source service access point. The value in this field is x'aa'.	
CONTROL	1 byte	Determines the LLC commands and responses. The three possible values for this field are discussed below.	
PROT_ID	3 bytes	Protocol ID. This field is reserved. It has a value of x'0'.	
TYPE	2 bytes	Specifies whether the packet is IP or ARP.	

Control Field Values

x'03'	Unnumbered Information (UI) frame. This is the normal, or unsequenced, way in which token-ring adapter data is transmitted through the network. TCP/IP sequences the data.
x'AF'	Exchange identification (XID) frame. This frame conveys the characteristics of the sending host.
x'E3'	Test frame. This frame supports testing of the transmission path, echoing

back the data that is received.

802.3 Frame Headers

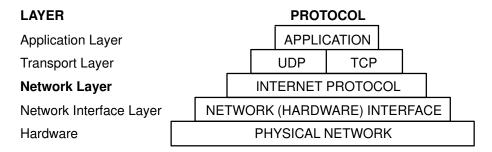
The MAC header for the 802.3 adapter is composed of three fields, as shown in the following MAC header table.

802.3 MAC Header			
Field	Length	Definition	
DA	6 bytes	Destination address.	
SA	6 bytes	Source address. If bit 0 of this field is set to 1, it indicates that routing information (RI) is present.	

The LLC header for 802.3 is the same as for Token–Ring MAC header.

Internet Network-Level Protocols

The Internet network–level protocols handle machine–to–machine communication. In other words, this layer implements TCP/IP routing. These protocols accept requests to send packets (along with the network address of the destination machine) from the Transport layer, convert the packets to datagram format, and send them down to the Network Interface layer for further processing (see figure).



Network Layer of the TCP/IP Suite of Protocols

TCP/IP provides the protocols that are required to comply with RFC 1100, *Official Internet Protocols*, as well as other protocols commonly used by hosts in the Internet community.

Note: The use of Internet network, version, socket, service, and protocol numbers in TCP/IP also complies with RFC 1010, *Assigned Numbers*.

Address Resolution Protocol

The first network–level protocol is the Address Resolution Protocol (ARP). ARP dynamically translates Internet addresses into the unique hardware addresses on local area networks.

To illustrate how ARP works, consider two nodes, jim and fred. If node jim wishes to communicate with fred, and jim and fred are on different local area networks (LANs), jim and fred communicate through *bridges*, *routers*, or *gateways*, using IP addresses. If, on the other hand, jim and fred are on the same LAN, an IP address alone is insufficient for direct communication. Within a LAN, nodes communicate using low–level hardware addresses.

Nodes on the same segment of the same LAN use ARP to determine the hardware address of other nodes. First, node <code>jim</code> broadcasts an ARP request for node <code>fred</code>'s hardware address. The ARP request contains <code>jim</code>'s IP and hardware addresses, and <code>fred</code>'s IP address. When <code>fred</code> receives the ARP request, it places an entry for <code>jim</code> in its ARP cache (which is used to map quickly from IP address to hardware address), then responds directly to <code>jim</code> with an ARP response containing <code>fred</code>'s IP and hardware addresses. When node <code>jim</code> receives <code>fred</code>'s ARP response, it places an entry for <code>fred</code> in its ARP cache.

Once an ARP cache entry exists at jim for fred, node jim is able to send packets directly to fred without resorting again to ARP (unless the ARP-cache entry for fred is deleted, in which case ARP is reused to contact fred).

Unlike most protocols, ARP packets do not have fixed–format headers. Instead, the message is designed to be useful with a variety of network technologies, such as:

- Ethernet LAN adapter (supports both Ethernet and 802.3 protocols)
- Token-ring network adapter
- Fiber Distributed Data Interface (FDDI) network adapter

However, ARP does not translate addresses for Serial Line Interface Protocol (SLIP) or Serial Optical Channel Converter, since these are point-to-point connections.

The kernel maintains the translation tables, and the ARP is not directly available to users or applications. When an application sends an Internet packet to one of the interface drivers, the driver requests the appropriate address mapping. If the mapping is not in the table, an

ARP broadcast packet is sent through the requesting interface driver to the hosts on the local area network.

Entries in the ARP mapping table are deleted after 20 minutes; incomplete entries are deleted after 3 minutes. To make a permanent entry in the ARP mapping tables, use the **arp** command with the *pub* parameter:

```
arp -s 802.3 host2 0:dd:0:a:8s:0 pub
```

When any host that supports ARP receives an ARP request packet, the host notes the IP and hardware addresses of the requesting system and updates its mapping table, if necessary. If the receiving host IP address does not match the requested address, the host discards the request packet. If the IP address does match, the receiving host sends a response packet to the requesting system. The requesting system stores the new mapping and uses it to transmit any similar pending Internet packets.

Internet Control Message Protocol

The second network—level protocol is the Internet Control Message Protocol (ICMP). ICMP is a required part of every Internet Protocol implementation. ICMP handles error and control messages for IP. This protocol allows gateways and hosts to send problem reports to the machine sending a packet. ICMP does the following:

- Tests whether a destination is alive and reachable.
- Reports parameter problems with a datagram header.
- Performs clock synchronization and transit time estimations.
- Obtains Internet addresses and subnet masks.

Note: ICMP uses the basic support of IP as if it were a higher–level protocol. However, ICMP is actually an integral part of IP and must be implemented by every IP module.

ICMP provides feedback about problems in the communications environment, but does not make IP reliable. That is, ICMP does not guarantee that an IP packet is delivered reliably or that an ICMP message is returned to the source host when an IP packet is not delivered or is incorrectly delivered.

ICMP messages may be sent in any of the following situations:

- When a packet cannot reach its destination.
- When a gateway host does not have the buffering capacity to forward a packet.
- When a gateway can direct a host to send traffic on a shorter route.

TCP/IP sends and receives several ICMP message types. ICMP is embedded in the kernel, and no application programming interface (API) is provided to this protocol.

Internet Control Message Protocol Message Types

ICMP sends and receives the following message types:

ecno request	Sent by nosts and gateways to test whether a destination is
	alive and reachable.

information request	Sent by hosts and gateways to obtain an Internet address
	fan a maturaul, ta culai da thace ana atta da ad. Thia maga agus

for a network to which they are attached. This message type is sent with the network portion of IP destination

address set to a value of 0.

timestamp requestSent to request that the destination machine return its

current value for time of day.

address mask request Sent by host to learn its subnet mask. The host can either

send to a gateway, if it knows the gateway address, or send

a broadcast message.

destination unreachable Sent when a gateway cannot deliver an IP datagram.

source quench Sent by discarding machine when datagrams arrive too

quickly for a gateway or host to process, in order to request that the original source slow down its rate of sending

datagrams.

redirect message Sent when a gateway detects that some host is using a

nonoptimum route.

echo reply Sent by any machine that receives an echo request in reply

to the machine which sent the request.

information reply Sent by gateways in response to requests for network

addresses, with both the source and destination fields of

the IP datagram specified.

timestamp reply Sent with current value of time of day.

address mask reply Sent to machines requesting subnet masks.

parameter problem Sent when a host or gateway finds a problem with a

datagram header.

time exceeded Sent when the following are true:

 Each IP datagram contains a time—to—live counter (hop count), which is decremented by each gateway.

A gateway discards a datagram because its hop count

has reached a value of 0.

Internet Timestamp Used to record the time stamps through the route.

Internet Protocol

The third network–level protocol is the Internet Protocol (IP), which provides unreliable, connectionless packet delivery for the Internet. IP is connectionless because it treats each packet of information independently. It is unreliable because it does not guarantee delivery (that is, it does not require acknowledgments from the sending host, the receiving host, or intermediate hosts).

IP provides the interface to the network interface level protocols. The physical connections of a network transfer information in a frame with a header and data. The header contains the source address and the destination address. IP uses an Internet datagram, which contains information similar to the physical frame. The datagram also has a header containing Internet addresses of both source and destination of the data.

IP defines the format of all the data sent over the Internet (see figure).

1 8	3 1	6 1	9 3		
Length	Type of Service		Total Length		
Identification			Fragment Offset		
o Live	Protocol		Header Checksum		
Source Address					
Destination Address					
Options					
Data					
	Length Identif	Length Type of Service Identification o Live Protocol Source A Destination Optio	Length Type of Service Identification Flags o Live Protocol Source Address Destination Address Options		

Internet Protocol Packet Header

IP Header Field Definitions

Version Specifies the version of the IP used. The current version of the

IP protocol is 4.

Length Specifies the datagram header length, measured in 32–bit

words.

Type of Service Specifies the Type of Service field contains five subfields that

specify the type of precedence, delay, throughput, and reliability desired for that packet. (The Internet does not guarantee this request.) The default settings for these five subfields are routine precedence, normal delay, normal throughput, and normal reliability. This field is not generally used by the Internet at this time. This implementation of IP complies with the requirements

of the IP specification, RFC 791, Internet Protocol.

Total Length Specifies the length of the datagram including both the header

and the data measured in octets. Packet fragmentation at gateways, with reassembly at destinations, is provided. The total length of the IP packet can be configured on an interface—by—interface basis with the Web-based System Manager fast path, **wsm network**, the **ifconfig** command, or the System Management Interface Tool (SMIT) fast path, **smit chinet**. Use Web-based System Manager or SMIT to set the values permanently in the configuration database; use the **ifconfig** command to set or change the values in the running

system.

Identification Contains a unique integer that identifies the datagram.

Fragment Flags Controls datagram fragmentation, along with the Identification

field. The Fragment Flags specify whether the datagram may be fragmented and whether the current fragment is the last one.

Fragment Offset Specifies the offset of this fragment in the original datagram

measured in units of 8 octets.

Time to Live Specifies how long the datagram can remain on the Internet.

This keeps misrouted datagrams from remaining on the Internet

forever. The default time to live is 255 seconds.

Protocol Specifies the high–level protocol type.

Header Checksum Indicates a number computed to ensure the integrity of header

values.

Source Address Specifies the Internet address of the sending host.

Destination Address Options

Specifies the Internet address of the receiving host.

Provides network testing and debugging. This field is not required for every datagram.

End of Option List

Indicates the end of the option list. It is used at the end of all options, not at the end of each option individually, and need be used only if the end of the options would not otherwise coincide with the end of the IP header. End of Option List is used if options exceed the length of the datagram.

No Operation

Provides alignment between other options; for example, to align the beginning of a subsequent option on a 32-bit boundary.

Loose Source and Record Route

Provides a means for the source of an Internet datagram to supply routing information to be used by the gateways in forwarding the datagram to a destination and for recording the route information. This is a *loose* source route: the gateway or host IP is allowed to use any route of any number of other intermediate gateways in order to reach the next address in the route.

Strict Source and Record Route

Provides a means for the source of an Internet datagram to supply routing information to be used by the gateways in forwarding the datagram to a destination and for recording the route information. This is a *strict* source route: the gateway or host IP must send the datagram directly to the next address in the source route through only the directly connected network indicated in the next address in order to reach the next gateway or host specified in the route.

Record Route Provides trace routes. This option provides a means to record the route of an Internet datagram.

Stream Identifier

Indicates a stream identifier. This option provides a way for a stream identifier to be carried through networks that do not support the stream concept.

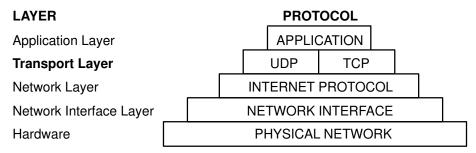
Internet Timestamp

Provides a record of the time stamps through the route.

Outgoing packets automatically have an IP header prefixed to them, and incoming packets have their IP header removed before being sent to the higher-level protocols. The IP protocol provides for the universal addressing of hosts in the Internet network.

Internet Transport-Level Protocols

The TCP/IP transport–level protocols (see figure) allow application programs to communicate with other application programs. User Datagram Protocol (UDP) and the TCP are the basic transport–level protocols for making connections between Internet hosts. Both TCP and UDP allow programs to send messages to and receive messages from applications on other hosts. When an application sends a request to the Transport layer to send a message, UDP and TCP break the information into packets, add a packet header including the destination address, and send the information to the Network layer for further processing. Both TCP and UDP use protocol ports on the host to identify the specific destination of the message.



Transport Layer of the TCP/IP Suite of Protocols

Higher–level protocols and applications use UDP to make datagram connections and TCP to make stream connections. The operating system sockets interface implements these protocols.

User Datagram Protocol

Sometimes an application on a network needs to send messages to a specific application or process on another network. The UDP provides a datagram means of communication between applications on Internet hosts. Because senders do not know which processes are active at any given moment, UDP uses destination protocol ports (or abstract destination points within a machine), identified by positive integers, to send messages to one of multiple destinations on a host. The protocol ports receive and hold messages in queues until applications on the receiving network can retrieve them.

Since UDP relies on the underlying IP to send its datagrams, UDP provides the same connectionless message delivery as IP. It offers no assurance of datagram delivery or duplication protection. However, UDP does allow the sender to specify source and destination port numbers for the message and calculates a checksum of both the data and header. These two features allow the sending and receiving applications to ensure the correct delivery of a message (see figure). Applications that require reliable delivery of datagrams must implement their own reliability checks when using UDP. Applications that require reliable delivery of streams of data should use TCP.

Bits 0	1	6	31
SOURCE	PORT NUMBER	DESTINATION PORT NUMBER	
	LENGTH	CHECKSUM	

User Datagram Protocol (UDP) Packet Header

UDP Header Field Definitions

Source Port Number Address of the protocol port sending the information.

Address of the protocol port receiving the information.

Length Length in octets of the UDP datagram.

Checksum Provides a check on the UDP datagram using the same

algorithm as the Internet Protocol.

The applications programming interface (API) to UDP is a set of library subroutines provided by the sockets interface.

Transmission Control Protocol

TCP provides reliable stream delivery of data between Internet hosts. Like UPD, TCP uses Internet Protocol, the underlying protocol, to transport datagrams, and supports the block transmission of a continuous stream of datagrams between process ports. Unlike UDP, TCP provides reliable message delivery. TCP ensures that data is not damaged, lost, duplicated, or delivered out of order to a receiving process. This assurance of transport reliability keeps applications programmers from having to build communications safeguards into their software.

The following are operational characteristics of TCP:

each direction between its users by packaging some number of bytes into segments for transmission through the Internet system. TCP implementation allows a segment size of at least 1024 bytes. In general, TCP decides when to block and forward packets at its own convenience.

Reliability TCP must recover data that is damaged, lost, duplicated, or

delivered out of order by the Internet. TCP achieves this reliability by assigning a sequence number to each octet it transmits and requiring a positive acknowledgment (ACK) from the receiving TCP. If the ACK is not received within the time—out interval, the data is retransmitted. The TCP

retransmission time—out value is dynamically determined for each connection, based on round—trip time. At the receiver,

the sequence numbers are used to correctly order segments that may be received out of order and to eliminate duplicates. Damage is handled by adding a checksum to each segment transmitted, checking it at the

receiver, and discarding damaged segments.

Flow Control TCP governs the amount of data sent by returning a

window with every ACK to indicate a range of acceptable sequence numbers beyond the last segment successfully received. The window indicates an allowed number of octets that the sender may transmit before receiving further

permission.

Multiplexing TCP allows many processes within a single host to use

TCP communications facilities simultaneously. TCP receives a set of addresses of ports within each host. TCP combines the port number with the network address and the host address to uniquely identify each socket. A pair of

sockets uniquely identifies each connection.

ConnectionsTCP must initialize and maintain certain status information

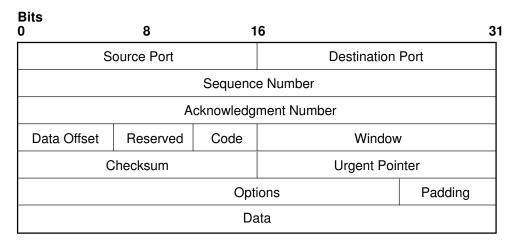
for each data stream. The combination of this information, including sockets, sequence numbers, and window sizes, is called a connection. Each connection is uniquely specified

by a pair of sockets identifying its two sides.

their communications. Default values are used when these

features are not needed.

The **TCP Packet Header** figure illustrates these characteristics.



Transmission Control Protocol (TCP) Packet Header

TCP Header Field Definitions

Source Port Identifies the port number of a source application program.

Destination Port Identifies the port number of a destination application

program.

Sequence Number Specifies the sequence number of the first byte of data in

this segment.

Acknowledgment Number Identifies the position of the highest byte received.

Data Offset Specifies the offset of data portion of the segment.

Reserved Reserved for future use.

Code Control bits to identify the purpose of the segment:

URG Urgent pointer field is valid.

ACK Acknowledgement field is valid.

PSH Segment requests a PUSH.

RTS Resets the connection.

SYN Synchronizes the sequence numbers.

FIN Sender has reached the end of its byte

stream.

Window Specifies the amount of data the destination is willing to

accept.

Checksum Verifies the integrity of the segment header and data.

Urgent Pointer

Indicates data that should be delivered as quickly as possible. This pointer specifies the position where urgent data ends.

Options

End of Option List

Indicates the end of the option list. This message is used at the end of all options, not at the end of each option individually. This option need only be used if the end of the options would not otherwise coincide with the end of the TCP header.

No Operation Indicates boundaries between options. May be used between other options; for example, to align the beginning of a subsequent option on a word boundary. There is no guarantee that senders will use this option, so receivers must be prepared to process options even if they do not begin on a word boundary.

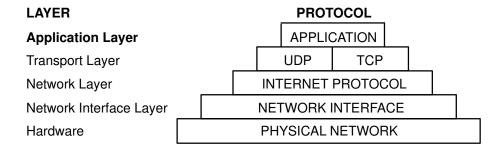
Maximum Segment Size

Indicates the maximum segment size TCP can receive. This should only be sent in the initial connection request.

The applications programming interface to TCP consists of a set of library subroutines provided by the sockets interface.

Internet Application–Level Protocols

TCP/IP implements higher-level Internet protocols at the application program level (see figure). When an application needs to send data to another application on another host, the applications send the information down to the transport level protocols to prepare the information for transmission.



Application Layer of the TCP/IP Suite of Protocols

The official Internet application—level protocols include:

- Domain Name Protocol (DOMAIN)
- Exterior Gateway Protocol (EGP)
- File Transfer Protocol (FTP)
- Name/Finger Protocol (FINGER)
- Telnet Protocol (TELNET)
- Trivial File Transfer Protocol (TFTP)

TCP/IP implements other higher–level protocols that are not official Internet protocols but are commonly used in the Internet community at the application program level. These protocols include:

- Distributed Computer Network (DCN) Local—Network Protocol (HELLO)
- Remote Command Execution Protocol (EXEC)
- Remote Login Protocol (LOGIN)
- Remote Shell Protocol (SHELL)
- Routing Information Protocol (RIP)
- Time Server Protocol (TIMED)

TCP/IP does not provide APIs to any of these application-level protocols.

Domain Name Protocol

The Domain Name Protocol (DOMAIN) allows a host in a domain to act as a *name server* for other hosts within the domain. DOMAIN uses UDP or TCP as its underlying protocol and allows a local network to assign host names within its domain independently from other domains. Normally, the DOMAIN protocol uses UDP. However, if the UDP response is truncated, TCP may be used. The DOMAIN protocol in TCP/IP supports both.

In the Domain hierarchical naming system, local resolver routines may resolve Internet names and addresses using a local name resolution database maintained by the **named** daemon. If the name requested by the host is not in the local database, the resolver routine queries a remote DOMAIN name server. In either case, if the name resolution information is unavailable, the resolver routines attempt to use the /etc/hosts file for name resolution.

Note: TCP/IP configures local resolver routines for the DOMAIN protocol if the local file /etc/resolv.conf exists. If this file does not exist, the TCP/IP configures the local resolver routines to use the /etc/hosts database.

TCP/IP implements the DOMAIN protocol in the **named** daemon and in the resolver routines and does not provide an API to this protocol.

Exterior Gateway Protocol

Exterior Gateway Protocol (EGP) is the mechanism that allows the exterior gateway of an *autonomous system* to share routing information with exterior gateways on other autonomous systems.

Autonomous Systems

An autonomous system is a group of networks and gateways for which one administrative authority has responsibility. Gateways are *interior neighbors* if they reside on the same autonomous system and *exterior neighbors* if they reside on different autonomous systems. Gateways that exchange routing information using EGP are said to be EGP *peers* or *neighbors*. Autonomous system gateways use EGP to provide access information to their EGP neighbors.

EGP allows an exterior gateway to ask another exterior gateway to agree to exchange access information, continually checks to ensure that its EGP neighbors are responding, and helps EGP neighbors to exchange access information by passing routing update messages.

EGP restricts exterior gateways by allowing them to advertise only those destination networks reachable entirely within that gateway's autonomous system. Thus, an exterior gateway using EGP passes along information to its EGP neighbors but does not advertise access information about its EGP neighbors outside its autonomous system.

EGP does not interpret any of the distance metrics that appear in routing update messages from other protocols. EGP uses the distance field to specify whether a path exists (a value of 255 means that the network is unreachable). The value cannot be used to compute the shorter of two routes unless those routes are both contained within a single autonomous

system. Therefore, EGP cannot be used as a routing algorithm. As a result, there will be only one path from the exterior gateway to any network.

In contrast to the Routing Information Protocol (RIP), which can be used within an autonomous system of Internet networks that dynamically reconfigure routes, EGP routes are predetermined in the /etc/gated.conf file. EGP assumes that IP is the underlying protocol.

EGP Message Types

Neighbor Acquisition Used by exterior gateways to request to become neighbors

Request of each other.

Neighbor Acquisition Used by exterior gateways to accept the request to become

Reply neighbors.

Neighbor Acquisition Used by exterior gateways to deny the request to become

Refusal neighbors. The refusal message includes reasons for

refusal, such as out of table space.

Neighbor Cease Used by exterior gateways to cease the neighbor

relationship. The cease message includes reasons for

ceasing, such as going down.

Neighbor Cease Used by exterior gateways to acknowledge the request to

Acknowledgment cease the neighbor relationship.

Neighbor HelloUsed by exterior gateways to determine connectivity. A

gateway issues a Hello message and another gateway

issues an I Heard You message.

I Heard You Used by exterior gateways to reply to a Hello message.

The I Heard You message includes the access of the answering gateway and, if the gateway is unreachable, a

reason for lack of access, such as You are

unreachable due to problems with $\ensuremath{\mathsf{my}}$ network

interface.

NR Poll Used by exterior gateways to query neighbor gateways

about their ability to reach other gateways.

Network Reachability Used by exterior gateways to answer the NR Poll

message. For each gateway in the message, the <code>Network</code> <code>Reachability</code> message contains information on the addresses that gateway can reach through its neighbors.

EGP Error Used by exterior gateways to respond to EGP messages

that contain bad checksums or have fields containing

incorrect values.

TCP/IP implements the EGP protocol in the **gated** server command and does not provide an API to this protocol.

File Transfer Protocol

File Transfer Protocol (FTP) allows hosts to transfer data among dissimilar hosts, as well as files between two foreign hosts indirectly. FTP provides for such tasks as listing remote directories, changing the current remote directory, creating and removing remote directories, and transferring multiple files in a single request. FTP keeps the transport secure by passing user and account passwords to the foreign host. Although FTP is designed primarily to be used by applications, it also allows interactive user—oriented sessions.

FTP uses reliable stream delivery (TCP/IP) to send the files and uses a Telnet connection to transfer commands and replies. FTP also understands several basic file formats including NETASCII, IMAGE, and Local 8.

TCP/IP implements FTP in the **ftp** user command and the **ftpd** server command and does not provide an applications programming interface (API) to this protocol.

When creating anonymous ftp users and directories please be sure that the home directory for users ftp and anonymous (for example, /u/ftp) is owned by root and does not allow write permissions (for example, dr-xr-xr-x). The script /usr/samples/tcpip/anon.ftp can be used to create these accounts, files and directories.

Telnet Protocol

The Telnet Protocol (TELNET) provides a standard method for terminal devices and terminal—oriented processes to interface. TELNET is commonly used by terminal emulation programs that allow you to log into a remote host. However, TELNET can also be used for terminal—to—terminal communication and interprocess communication. TELNET is also used by other protocols (for example, FTP) for establishing a protocol control channel.

TCP/IP implements TELNET in the **tn**, **telnet**, or **tn3270** user commands. The **telnetd** daemon does not provide an API to TELNET.

TCP/IP supports the following Telnet options which are negotiated between the client and server:

BINARY TRANSMISSION (Used in tn3270 sessions)	Transmits characters as binary data.
SUPPRESS GO_AHEAD (AIX suppresses GO-AHEAD options.)	Indicates that when in effect on a connection between a sender of data and the receiver of the data, the sender need not transmit a GO_AHEAD option. If the GO_AHEAD option is not desired, the parties in the connection will probably suppress it in both directions. This action must take place in both directions independently.
TIMING MARK (Recognized, but has a negative response)	Makes sure that previously transmitted data has been completely processed.
EXTENDED OPTIONS LIST	Extends the TELNET option list for another 256 options. Without this option, the TELNET option allows only 256 options.
ECHO (User-changeable command)	Transmits echo data characters already received back to the original sender.
TERM TYPE	Enables the server to determine the type of terminal connected to a user TELNET program.
SAK (Secure Attention Key)	Establishes the environment necessary for secure communication between you and the system.
NAWS (Negotiate About Window Size)	Enables client and server to negotiate dynamically for the window size. This is used by applications that support

Note: Telnet must allow transmission of eight bit characters when not in binary mode in order to implement ISO 8859 Latin code page. This is necessary for internationalization of the TCP/IP commands.

changing the window size.

Trivial File Transfer Protocol

The Trivial File Transfer Protocol (TFTP) can read and write files to and from a foreign host. Since TFTP uses the unreliable User Datagram Protocol to transport files, it is generally quicker than FTP. Like FTP, TFTP can transfer files as either NETASCII characters or as 8-bit binary data. Unlike FTP, TFTP cannot be used to list or change directories at a foreign host and it has no provisions for security like password protection. Also, data can be written or retrieved only in public directories.

The TCP/IP implements TFTP in the **tftp** and **utftp** user commands and in the **tftpd** server command. The **utftp** command is a form of the **tftp** command for use in a pipe. TCP/IP does not provide an API to this protocol.

Name/Finger Protocol

The Name/Finger Protocol (FINGER) is an application—level Internet protocol that provides an interface between the **finger** command and the **fingerd** daemon. The **fingerd** daemon returns information about the users currently logged in to a specified remote host. If you execute the finger command specifying a user at a particular host, you will obtain specific information about that user. The FINGER Protocol must be present at the remote host and at the requesting host. FINGER uses Transmission Control Protocol (TCP) as its underlying protocol.

Note: TCP/IP does not provide an API to this protocol.

Distributed Computer Network Local-Network Protocol

Local—Network Protocol (HELLO) is an interior gateway protocol designed for use within autonomous systems. (For more information, see "Autonomous Systems", on page 3-32.) HELLO maintains connectivity, routing, and time—keeping information. It allows each machine in the network to determine the shortest path to a destination based on time delay and then dynamically updates the routing information to that destination.

The **gated** daemon provides the Distributed Computer Network (DCN) local network protocol.

Remote Command Execution Protocol

The **rexec** user command and the **rexecd** daemon provide the remote command execution protocol, allowing users to run commands on a compatible remote host.

Remote Login Protocol

The **rlogin** user command and the **rlogind** daemon provide the remote login protocol, allowing users to log in to a remote host and use their terminals as if they were directly connected to the remote host.

Remote Shell Protocol

The **rsh** user command and the **rshd** daemon provide the remote command shell protocol, allowing users to open a shell on a compatible foreign host for running commands.

Routing Information Protocol

Routing Information Protocol (RIP) and the **routed** and **gated** daemons that implement it keep track of routing information based on gateway hops and maintain kernel–routing table entries.

Time Server Protocol

The **timed** daemon is used to synchronize a host's time with the time of other hosts. It is based on the client/server concept.

Assigned Numbers

For compatibility with the general network environment, well–known numbers are assigned for the Internet versions, networks, ports, protocols, and protocol options. Additionally, well–known names are also assigned to machines, networks, operating systems, protocols, services, and terminals. TCP/IP complies with the assigned numbers and names defined in RFC 1010, *Assigned Numbers*.

The Internet Protocol defines a 4-bit field in the IP header that identifies the version of the general Internetwork protocol in use. For IP, this version number in decimal is 4. For details on the assigned numbers and names used by TCP/IP, refer to the /etc/protocols and /etc/services files included with TCP/IP. For further details on the assigned numbers and names, refer to RFC 1010 and the /etc/services file.

TCP/IP Network Adapter Cards

The topics discussed in this section are:

- Installing a Network Adapter, on page 3-36
- Configuring a Token–Ring or Ethernet Adapter, on page 3-37
- Using Turboways 100 and Turboways 155 ATM Adapters, on page 3-38

The network adapter card is the hardware that is physically attached to the network cabling. It is responsible for receiving and transmitting data at the physical level. The network adapter card is controlled by the network adapter device driver.

A machine must have one network adapter card (or connection) for each network (not network type) to which it connects. For instance, if a host attaches to two token–ring networks, it must have two network adapter cards.

TCP/IP uses the following network adapter cards and connections:

- Standard Ethernet Version 2
- IEEE 802.3
- Token–ring
- · Asynchronous adapters and native serial ports
- Fiber Distributed Data Interface (FDDI)
- Serial Optical Channel Converter
- Turboways 100 and Turboways 155 ATM

The Ethernet and 802.3 network technologies use the same type of adapter.

Each machine provides a limited number of expansion slots, some or all of which you may wish to use for communications adapters. Additionally, each machine supports a limited number of communications adapters of a given type. Each machine supports up to eight Ethernet/802.3 adapters, up to eight token—ring adapters, and one asynchronous adapter card with up to 64 connections. Within these limits (software limitations), you can install any combination of these adapters up to the total number of expansion slots available in your machine (hardware limitations).

Only one Transmission Control Protocol/Internet Protocol (TCP/IP) interface is configurable regardless of the number of Serial Optical Channel Converters supported by the system. The **Serial Optical** device driver makes use of both channel converters even though only one logical TCP/IP interface is configured.

Installing a Network Adapter

To install a network adapter:

- 1. Shut down the computer. See the **shutdown** command for information on how to shut down a system.
- 2. Turn off the computer's power.
- 3. Remove the computer's cover.
- 4. Find a free slot on the Micro Channel bus and insert the network adapter into it, being careful to seat the adapter properly.
- 5. Replace the computer's cover.
- 6. Restart the computer.

Configuring and Managing Adapters

To configure and manage token–ring or Ethernet adapters, use the tasks in the following table.

Configuring and Managing Ad	Configuring and Managing Adapters Tasks		
Web-based System Manager: (wsm network application)	Web-based System Manager: wsm network fast path		
-OR- Task	SMIT Fast Path	Command or File	
Configure an Adapter	smit chgtok (token ring) smit chgenet (Ethernet)	1. Determine adapter name:1 lsdev -C -c adapte r -t tokenring -H or lsdev -C -c adapte r -t ethernet -H	
		2. Reset ring speed (token ring) or connector type (Ethernet), if necessary. For example: chdev -l tok0 -a r ing_speed=16 -P or chdev -l ent0 -a b nc_select=dix -P	
Determining a Network Adapter's Hardware Address	smit chgtok (token ring) smit chgenet (Ethernet)	lscfg -l tok0 -v (token ring) ² lscfg -l ent0 -v (Ethernet) ²	
Setting an Alternate Hardware Address	smit chgtok (token ring) smit chgenet (Ethernet)	1. Define the alternate hardware address. For example, for token ring:2,3 chdev -l tok0 -a alt_addr=0X10005A4 F1B7F For Ethernet:2,3 chdev -l ent0 -a alt_addr=0X10005A4 F1B7F -p 2. Begin using alternate address, for token ring:4 chdev -l tok0 -a use_alt_addr=yes For Ethernet:4 chdev -l ent0 -a use_alt_addr=yes	

Note:

- 1. The name of a network adapter can change if you move it from one slot to another or remove it from the system. If you ever do so, issue the **diag** –a command to update the configuration database.
- 2. Substitute your adapter name for tok0 and ent0.
- 3. Substitute your hardware address for 0x10005A4F1B7F.
- 4. After performing this procedure, you may experience a disruption of communication with other hosts until they flush their Address Resolution Protocol (ARP) cache and obtain this host's new hardware address.

Turboways 100 and 155 Asynchronous Transfer Mode (ATM) Adapters

The Turboways 100 and 155 adapters provide connectivity to ATM networks. These adapters implement multi-mode fiber-physical interfaces. ATM is used in networking environments requiring greater bandwidth than available using today's local area networks (LANs).

The Turboways 100 adapter provides 100 Mbps, full duplex connectivity for Micro Channel bus—based servers or clients using permanent virtual circuits (PVCs) and switched virtual circuits (SVCs). The PVC and SVC implementation is compliant with "ATM Forum UNI 3.0" specification. Similarly, Turboways 155 adapters provide 155 Mbps full—duplex connectivity to ATM networks. The adapters support a maximum of 1024 virtual circuits.

ATM Technology

Asynchronous Transfer Mode (ATM) is a cell–switching, connection–oriented technology. In ATM networks, end stations attach to the network using dedicated full duplex connections. The ATM networks are constructed using switches, and switches are inter–connected using dedicated physical connections. Before any data transfers can begin, end–to–end connections need to be established. Multiple connections can and do exist on a single physical interface. Sending stations transmit data by segmenting Protocol Data Units (PDUs) into 53–byte cells. Payload stays in the form of cells during network transport. Receiving stations reassemble cells into PDUs. The connections are identified using a virtual path identifier (VPI) and a virtual channel identifier (VCI). The VPI field occupies one byte in the ATM cell's five–byte header; whereas, the VCI field occupies two bytes in the ATM cell's five–byte header. Basically, a VPI:VCI pair identifies the source of the ATM cell. The function of the ATM switch is to recognize the source of the cell, determine the next hop, and output the cell to a port. The VPI:VCI changes on a hop–by–hop basis. Thus, VPI:VCI values are not universal. Each virtual circuit is described as a concatenation of VPI:VCI values across the network.

ATM Connections

In ATM architecture, there are two kinds of virtual circuits: permanent (PVCs) and switched (SVCs).

Permanent Virtual Circuits PVCs are statically and manually configured. The switches forming the ATM network must first be set up to recognize each endpoint's VPI:VCI combination and to route that endpoint's ATM cells to the destination endpoint through the ATM network. Once a link connection through the network has been established from one endpoint to another, ATM cells can be transmitted through the ATM network and ATM switches. The network switches translate the VPI:VCI values in the appropriate way so as to route the cell to its destination.

Switched Virtual Circuits

SVCs are dynamically set up on an as needed basis. The ATM end stations are assigned 20-byte addresses. There is a concept of control plane and data plane. The control plane uses a signalling channel's VPI:VCI 0:5. SVCs involve on demand call setup, whereby an ATM station sends information elements specifying destination ATM address (and optionally, source ATM address). There are many other information elements for specifying ATM adaptation layer (AAL) parameters, Bandwidth and quality of service (QOS) parameters, etc. In general, calling station, network, and called station participate in a negotiation. Finally, a call is either accepted or rejected. If a call is accepted, network assigns VPI:VCI values for the data plane to the calling station and called station. In the control plane, the ATM network routes (or switches) signalling packets on the basis of the ATM addresses. While these packets are being routed, the switches set up data plane cell routing tables. In the data plane, ATM networks switch cells on the basis of VPI:VCI much like in the case of PVCs. When data transfer is over, connection is terminated.

The ATM address is constructed by registering with the ATM network and acquiring the most significant 13 bytes. The next 6 bytes are the hardware address "burnt" in the adapter. The least significant byte is the selector; the use of this byte is left to the discretion of the end station. ATM networks do not interpret this byte.

TCP/IP over ATM

The Internet Engineering Task Force RFC1577: Classical IP and ARP over ATM standard specifies the mechanism for implementing Internet Protocol (IP) over ATM. Since ATM is connection-oriented technology and IP is a datagram-oriented technology, mapping the IP over ATM is not trivial.

In general, the ATM network is divided into logical IP subnetworks (LISs). Each LIS is comprised of some number of ATM stations. LISs are analogous to traditional LAN segments. LISs are interconnected using routers. A particular adapter (on an ATM station) can be part of multiple LISs. This feature may be very useful for implementing routers.

RFC1577 specifies RFC1483, which specifies logical link control/Sub-Network Access Protocol (LLC/SNAP) encapsulation as the default. In PVC networks for each IP station, all PVCs must be manually defined by configuring VPI:VCI values. If LLC/SNAP encapsulation is not being used, the destination IP address associated with each VPI:VCI must be defined. If LLC/SNAP encapsulation is being used, the IP station can learn the remote IP address by an InARP mechanism. For SVC networks, RFC1577 specifies an ARP server per LIS. The purpose of the ARP server is to resolve IP addresses into ATM addresses without using

broadcasts. Each IP station is configured with the ATM address of the ARP server. IP stations set up SVCs with the ARP server, which in turn, sends InARP requests to the IP stations. Based on InARP reply, an ARP server sets up IP to ATM address maps. IP stations send ARP packets to the ARP server to resolve addresses, which returns ATM addresses. IP stations then set up a SVC to the destination station and data transfer begins. The ARP entries in IP stations and the ARP server age based on a well defined mechanism. For both the PVC and SVC environments, each IP station has at least one virtual circuit per destination address.

The Internet Engineering Task Force RFC2225 replaces the RFC1577 mainly by adding the support of ATM ARP Request Address list. The ATM ARP Request Address list is a list containing one or more ATM addresses of individual ATMARP servers located within the LIS. The RFC2225 client eliminates the single point of failure associated with the 1577 clients' ATMARP services. The 2225 clients have the ability to switch to backup ARP servers when the current ATM ARP server fails.

ESCALA sets the first entry in the ATM ARP Request Address list as the Primary ATM ARP server and the rest of the entries as Secondary ATM ARP servers.

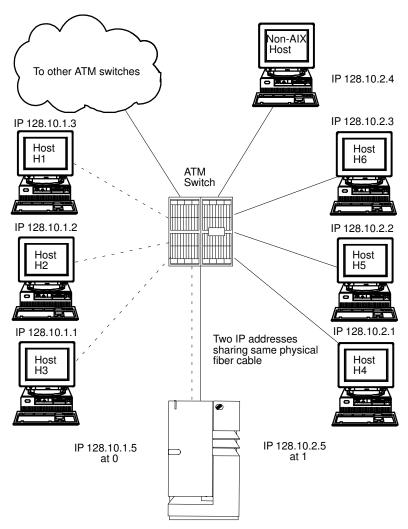
The client will always try to use the Primary ATM ARP server. If the effort to connect to the Primary ATM ARP server falis, the client will try to connect to the first Secondary server (the position in the ATM ARP Request Address list determines the order of the Secondary ATM ARP server). If the connection to the first Secondary ATM ARP server fails, the client will try to contact the next Secondary ATM ARP server in the list. This will go on in a round robin fashion.

If the connection to the Primary ATM ARP server fails, regardless of which Secondary ATM ARP server it is connected to or attempting to connect to, the client will continue to retry the Primary ATMA ARP server every 15 minutes. If it finally connects to the Primary ATM ARP server, then the connection to the current Secondary ATM ARP server will be dropped.

The ATM ARP Request Address list is entered manually either through SMIT or by using the **ifconfig** command. The ATM ARP Request Address list cannot be configured with the Management Information Base (MIB).

PVC Network

Use the Representative ATM Network figure as an example to configure your network.



ATM Gateway box with one TURBOWAYS 155 ATM Adapter

Representative ATM Network

Within the Representative ATM Network figure, one logical IP subnet is represented by dashed lines from each host to the switch. The other IP subnet is represented by solid lines from each host to the switch.

The following Representative Host Configuration table indicates how AIX hosts, H3 and H4, should be configured to communicate with a gateway and with each host on its own logical IP subnet.

Representative Host Configuration		
Network Interface Driver	VPI:VCI	Comment
Host H3		
at0	0:40	Connection to 128.10.1.5 (gateway)
at0	0:42	Connection to 128.10.1.2
at0	0:43	Connection to 128.10.1.3

Host H4		
at0	0:50	Connection to 128.10.2.5 (gateway)
at0	0:52	Connection to 128.10.2.2
at0	0:53	Connection to 128.10.2.3
at0	0:54	Connection to 128.10.2.4

To reach hosts on another logical IP subnet, only a VPI:VCI connection to the gateway needs to be created. (The VPI:VCIs are for illustration purposes only.)

The ATM gateway box has one ATM with two IP addresses sharing the same physical cable.

SVC Network

Using the Representative ATM Network figure as an example, imagine that AIX host H3 wishes to call H4. H1 is the ARP server for subnet 1 and H6 is the ARP server for subnet 2. Assuming a subnet mask of 255.255.255.0, stations with addresses 128.10.1.X are members of one subnet, whereas stations with addresses of 128.10.2.X are members of a second subnet. See the following list of representative host configurations using SVCs.

List of Representative Host Configurations				
Network Interface Driver	IP Address	ARP Server	ARP Server Address	Gateway Address
Host H1				
at0	128.10.1.3	Yes		128.10.1.5
Host H3				
at0	128.10.1.1	No	ATM address of H1	128.10.1.5
Gateway	Gateway			
at0	128.10.1.5	No	ATM address of H1	
at1	128.10.2.5	No	ATM address of H6	
Host H4				
at0	128.10.2.1	No	ATM address of H6	128.10.2.5
Host H6	Host H6			
at0	128.10.2.3	Yes		128.10.2.5

Note: Each subnet requires one and only one ARP server.

Because H3 recognizes that address 128.10.2.1 is not on its subnet, H3 consults H1 to resolve the default gateway's IP address to an ATM address. H3 then places a call to the gateway. The gateway recognizes that the data is bound for the second subnet and consults H6 to successfully resolve H4's IP address to an ATM address. Connections are then established between H3 and the gateway and between the gateway and H4.

Configuring an ATM Adapter

To configure your ATM 100, ATM 155, or PCI ATM adapter, use the Web-based System Manager fast path **wsm network** or the SMIT fast path **smit chg_atm**. Select an adapter name, then use the online help and multiple—choice lists to decide which changes to make for your configuration.

ATM Adapter Statistics

The **atmstat** command can be used for getting ATM adapter statistics. Using the **atmstat** command with the **-r** flag resets the statistics. The format of the command is **atmstat** *DeviceName*. This command returns the following sets of statistics:

Transmit Statistics

Packets: his field contains the number of packets (or PDUs) transmitted.

Bytes: This field contains the count of bytes transmitted. These are the user bytes.

The ATM overhead, for example, ATM cell header, AAL 5 PDU trailer, etc.,

are excluded.

Interrupts:

This field is not used.

Transmit Errors:

This field contains the number of Transmit Errors for this device.

Packets Dropped:

This field contains the number of Transmit Packets that were dropped, for instance, due to out of buffers condition.

Max Packets on S/W Transmit Queue:

This field does not apply to ATM.

S/W Transmit Queue Overflow:

This field does not apply to ATM.

Current S/W + H/W Transmit Queue Length:

This field contains the current transmit queue length.

Cells Transmitted:

This field contains the number of cells transmitted by this device.

Out of Xmit Buffers:

This field contains the number of packets dropped because of Out of Xmit Buffers condition.

Current HW Transmit Queue Length:

This field contains the current number of Transmit Packets on the hardware queue.

Current SW Transmit Queue Length:

This field does not apply to ATM.

Receive Statistics

Packets: This field contains the number of packets (or PDUs) received.

Bytes: This field contains the count of Bytes received. These are the user bytes.

The ATM overhead, for example, ATM cell header, AAL 5 PDU trailer, etc.,

are excluded.

Interrupts:

This field contains the number of Interrupts taken by the system for the Adapter–to–System indications. Some of the events that cause these interrupts are packet received, transmit done indication, and so on.

Receive Errors:

This field contains the number of Receive Errors for this device.

Packets Dropped:

This field contains the number of Received Packets dropped, for instance, due to out of buffers condition.

Bad Packets:

This field does not apply to ATM.

Cells Received:

This field contains the number of Cells Received by this device.

Out of Rcv Buffers:

This field contains the number of Packets Dropped because of out of receive buffers condition.

CRC Errors:

This field contains the number of Received Packets that encountered CRC errors.

Packets Too Long:

This field contains the number of Received Packets that exceeded the maximum PDU size.

Incomplete Packets:

This field contains the number of Incomplete received packets.

Cells Dropped:

This field contains the number of dropped cells. Cells could be dropped for a number of reasons, such as bad header error control (HEC), out of buffer condition, and so on.

General Statistics

No mbuf Errors:

This field contains the number of mbuf requests that were denied.

Adapter Loss of Signals:

This field contains the number of times the adapter encountered Loss of Signal.

Adapter Reset Count:

This field contains the number of times the adapter has been reset.

Driver Flags: Up Running Simplex

This field contains the NDD flags.

Virtual Connections in use:

This field contains the number of VCs that are currently allocated or in use.

Max Virtual Connections in use:

This field contains the maximum number of VCs allocated since the last reset of the statistics.

Virtual Connections Overflow:

This field contains the number of allocate VC requests that have been denied.

SVC UNI Version:

This field contains the current UNI version of the signaling protocol in use.

Additional ATM Statistics

Using the **atmstat** command with the **-d** flag provides detailed statistics.

Turboways ATM Adapter Specific Statistics:

Packets Dropped - No small DMA buffer:

This field contains the number of received packets dropped because the adapter did not have small system buffers for DMA.

Packets Dropped - No medium DMA buffer:

This field contains the number of received packets dropped because the adapter did not have medium system buffers for DMA.

Packets Dropped - No large DMA buffer:

This field contains the number of received packets dropped because the adapter did not have large system buffers for DMA.

Receive Aborted - No Adapter Receive buffer:

This field contains the number of received packets aborted because the adapter did not have on–card receive buffers.

Transmit Aborted - No small DMA buffer:

This field contains the number of transmit packets dropped because of the lack of small system buffers for DMA.

Transmit Aborted - No medium DMA buffer:

This field contains the number of transmit packets dropped because of the lack of medium system buffers for DMA.

Transmit Aborted - No large DMA buffer:

This field contains the number of transmit packets dropped because of the lack of large system buffers for DMA.

Transmit Aborted - No MTB DMA buffer:

This field contains the number of transmit packets dropped because of the lack of large system buffers for DMA.

Transmit Aborted - No Adapter Transmit buffer:

This field contains the number of transmit packets dropped because of the lack of adapter on–card transmit buffers.

Max Hardware Transmit Queue Length:

This field contains the maximum number of transmit packets queued in the hardware queue.

Small Mbufs in Use:

This field contains the number of small mbufs currently in use. The adapter device driver allocates these buffers according to the configuration information provided by system administrators. This information can be used to tune the configuration information.

Medium Mbufs in Use:

This field contains the number of medium mbufs currently in use. The adapter device driver allocates these buffers according to the configuration information provided by system administrators. This information can be used to tune the configuration information.

Large Mbufs in Use:

This field contains the number of large mbufs currently in use. The adapter device driver allocates these buffers according to the configuration information provided by system administrators. This information can be used to tune the configuration information.

Huge Mbufs in Use:

This field contains the number of huge mbufs currently in use. The adapter device driver allocates these buffers according to the configuration information provided by the system administrators. This information can be used to tune the configuration information.

MTB Mbufs in Use:

This field contains the number of MTB mbufs currently in use. The adapter device driver allocates these buffers according to the configuration information provided by the system administrators. This information can be used to tune the configuration information.

Max Small Mbufs in Use:

This field contains the maximum number of small mbufs that have been used. The adapter device driver allocates these buffers according to the

configuration information provided by the system administrators. This information can be used to tune the configuration information.

Max Medium Mbufs in Use:

This field contains the maximum number of medium mbufs that have been used. The adapter device driver allocates these buffers according to the configuration information provided by system administrators. This information can be used to tune the configuration information.

Max Large Mbufs in Use:

This field contains the maximum number of large mbufs that have been used. The adapter device driver allocates these buffers according to the configuration information provided by system administrators. This information can be used to tune the configuration information.

Max Huge Mbufs in Use:

This field contains the maximum number of huge mbufs that have been used. The adapter device driver allocates these buffers according to the configuration information provided by system administrators. This information can be used to tune the configuration information.

MTB Mbufs in Use:

This field contains the maximum number of MTB mbufs that have been used. The adapter device driver allocates these buffers according to the configuration information provided by system administrators. This information can be used to tune the configuration information.

Small Mbufs overflow:

This field contains the number of times that a small mbuf could not be allocated. This information can be used to tune the configuration information.

Medium Mbufs overflow:

This field contains the number of times that a medium mbuf could not be allocated. This information can be used to tune the configuration information.

Large Mbufs overflow:

This field contains the number of times that a large mbuf could not be allocated. This information can be used to tune the configuration information.

Huge Mbufs overflow:

This field contains the number of times that a huge mbuf could not be allocated. This information can be used to tune the configuration information.

MTB Mbufs overflow:

This field contains the number of times that an MTB mbuf could not be allocated. This information can be used to tune the configuration information.

PCI ATM Adapter Specific Statistics:

Total Receive Buffers: 48 Using: 32

This message contains the number of receive buffers allocated as well as the number that are currently in use.

TCP/IP Network Interfaces

The TCP/IP Network Interface layer formats IP datagrams at the Network layer into packets that specific network technologies can understand and transmit. A network interface is the network—specific software that communicates with the network—specific device driver and the IP layer in order to provide the IP layer with a consistent interface to all network adapters that may be present.

The IP layer selects the appropriate network interface based on the destination address of the packet to be transmitted. Each network interface has a network address. The Network Interface layer is responsible for adding or removing any link layer protocol header required to deliver a message to its destination. The **network adapter** device driver controls the network adapter card.

A network interface is usually associated with a network adapter, but there is no requirement for one. For instance, the loopback interface has no network adapter associated with it. A machine must have one network adapter card for each network (not network type) to which it connects. However, a machine requires only one copy of the network interface software for each network adapter it uses. For instance, if a host attaches to two token—ring networks, it must have two network adapter cards. However, only one copy of the token—ring network interface software and one copy of the token—ring device driver is required.

TCP/IP supports types of network interfaces:

- Standard Ethernet Version 2 (en)
- IEEE 802.3 (et)
- Token-ring (tr)
- Serial Line Internet Protocol, or SLIP (sl)
- Loopback (lo)
- FDDI
- · Serial Optical (so)
- ATM (at)
- Point-to-Point Protocol (PPP)

The Ethernet, 802.3, and token–ring interfaces are for use with local area networks (LANs). The SLIP interface is for use with serial connections. The loopback interface is used by a host to send messages back to itself. The Serial Optical interface is for use with optical point–to–point networks using the Serial Optical Link device handler. The ATM interface is for use with 100 Mbits/sec and 155 Mbits/sec ATM connections.

Automatic Configuration of Network Interfaces

When a new network adapter is physically installed in the system, the operating system automatically adds the appropriate network interface for that adapter. For example, if you install a token–ring adapter in your system, the operating system will assign it the name tok0 and add a token–ring network interface named tr0. If you install an Ethernet adapter in your system, the operating system will assign it the name ent0 and add both an Ethernet Version 2 and an IEEE 802.3 interface, named en0 and et0 respectively.

In most cases, there is a one—to—one correspondence between adapter names and network interface names. For example, token—ring adapter tok0 corresponds to interface tr0, adapter tok1 corresponds to interface tr1, and so on. Similarly, Ethernet adapter ent0 corresponds to interface en0 (for Ethernet Version 2) and et0 (for IEEE 802.3), and adapter ent1 corresponds to interface en1 (for Ethernet Version 2) and et1 (for IEEE 802.3).

In the case of ATM, according to RFC1577, it is possible for an ATM station to be part of multiple Logical IP Subnetworks. When this is the case, there will be multiple interfaces associated with a device. This will require that an interface be specifically added and a device name be assigned to it.

Note: Under normal circumstances, you will never need to delete or add a network interface manually. However, some problem determination procedures may require you to do so. In this case, use the Web-based System Manager fast path, **wsm network**, or the SMIT fast path, **smit inet**, to delete and re—add the appropriate interface.

At each system startup, the operating system automatically configures the network interface software based upon the information in the ODM database. Initially, the network interface is configured with default values. In order to communicate through a given network interface, the Internet address must be set. This is the only attribute that you need to set. All other necessary attributes can use the default values. The default values for each network interface follow.

Ethernet Default Configuration Values

The following is a list of valid Ethernet network adapter attributes along with their default values, which may be changed using the Web-based System Manager fast path **wsm network** or the Network Interface Selection menu in SMIT.

Attribute	Default Value	Possible Values
netaddr		
state	down	up, down, detach
arp	yes	yes, no
netmask		
broadcast		

The following valid Ethernet network device driver attribute is shown along with its default values, which may be changed using the Web-based System Manager fast path **wsm network** or the Network Interface Drivers menu in SMIT.

Attribute	Default Value	Possible Values
mtu	1500	60 through 1500

802.3 Default Configuration Values

The following is a list of valid 802.3 network adapter attributes along with their default values, which may be changed using the Web-based System Manager fast path **wsm network** or the Network Interface Selection menu in SMIT.

Attribute	Default Value	Possible Values
netaddr		
state	down	up, down, detach
arp	yes	yes, no
netmask		
broadcast		

The following valid 802.3 network device driver attribute is shown along with its default values, which may be changed using the Web-based System Manager fast path **wsm network** or the Network Interface Drivers menu in SMIT.

Attribute	Default Value	Possible Values
mtu	1492	60 through 1492

Token–Ring Default Configuration Values

The following is a list of valid token–ring network adapter attributes along with their default values, which may be changed using the Web-based System Manager fast path **wsm network** or the Network Interface Selection menu in SMIT.

Attribute	Default Value	Possible Values
netaddr		
netmask		
state	down	up, down, detach
arp	yes	yes, no
hwloop	no	yes, no
netmask		
broadcast		
allcast	no	yes, no

The following valid token–ring network device driver attributes are shown along with its default values, which may be changed using the Web-based System Manager fast path **wsm network** or the Network Interface Drivers menu in SMIT.

Attribute	Default Value	Possible Values
mtu (4Mbps)	1500	60 through 4056
mtu (16Mbps)	1500	60 through 17960

Note: When operating through a bridge, the default value of 1500 for the maximum transmission unit (MTU) should be changed to a value that is 8 less than the maximum information field (maximum I–frame) advertised by the bridge in the routing control field. For example, if the maximum I–frame value is 1500 in the routing control field, the MTU size should be set to 1492. This is for token–ring network interfaces only. For more information, refer to "Problems with a Token–Ring/Token–Ring Bridge", on page 3-158.

When using the IBM 16/4 PowerPC token-ring adapter (ISA) or the IBM PCMCIA 16/4 token-ring card, the MTU is restricted to 2000.

SLIP Default Configuration Values

The following is a list of valid SLIP network adapter attributes along with their default values as shown under the Web-based System Manager fast path **wsm network** or the Network Interface Selection menu in SMIT.

Attribute	Default Value	Possible Values
netaddr		
dest		
state	up	up, down, detach
netmask		

The following valid SLIP network device driver attribute is shown along with its default values as displayed under the Web-based System Manager fast path **wsm network** or the Network Interface Drivers menu in SMIT.

Attribute	Default Value	Possible Values
mtu	1006	60 through 4096

Serial Optical Default Configuration Values

The following is a list of valid Serial Optical network channel converter attributes along with their default values as shown under the Web-based System Manager fast path **wsm network** or the Network Interface Selection menu in SMIT.

Attribute	Default Value	Possible Values
netaddr		
state	down	up, down, detach
netmask		

The following valid serial optical network device handler attribute is shown along with its default values as displayed under the Web-based System Manager fast path **wsm network** or the Network Interface Drivers menu in SMIT.

Attribute	Default Value	Possible Values
mtu	61428	1 through 61428

ATM Default Configuration Values

The following is a list of valid ATM network adapter attributes along with their default values as shown under the Web-based System Manager fast path **wsm network** or the Network Interface Selection menu in SMIT.

Attribute	Default Value	Possible Values
netaddr		
netmask		
state	ир	up, down, detach
Connection Type	svc_s	svc_c, svc_s, pvc
ATM Server Address		
Alternate Device		
idle timer	60	1 through 60
Best Effort Bit Rate (UBR) in kbits/sec	0	1 through 155,000

The following valid ATM network device driver attribute is shown along with its default values as displayed under the Web-based System Manager fast path **wsm network** or the Network Interface Drivers menu in SMIT.

Attribute	Default Value	Possible Values
mtu	9180	1 through 64K

Note: Network Administrators need to exercise caution while changing MTU size from default to some other value. The MTU size needs to be coordinated with other stations on the network.

If PVCs are to be used on an interface, the VPI:VCIs needs to be defined. This is part of the Network Interface Selection Menu. The last option on this menu is PVCs for IP over ATM Network, which helps you list, add, change, or remove PVCs.

Implications of Multiple Network Interfaces on the Same Network

Sometimes network managers feel the need to provide greater availability and performance by adding a second network adapter to a particular machine. For example, they may want to have two token—ring adapters attached to the same physical network. While it is possible to have more than one interface on the same network, in general this is not recommended for two reasons:

- 1. Having two interfaces on the same network is a violation of TCP/IP architecture.
 - In TCP/IP architecture, a host machine with two network adapters is defined as an IP router. Different network adapters must be attached to different physical networks. In the case of token–ring, TCP/IP addresses multiple rings bridged together as a single logical ring (as if it were a single physical ring).
- 2. Having two interfaces on the same network can cause broadcast storms.

Whenever an IP host sees traffic for a network whose IP address is different from its own network, it generates an Internet Control Message Protocol (ICMP) packet announcing this conflict. Since every host on the network sees the traffic that is misaddressed, every host generates ICMP packets. If the amount of misaddressed traffic is significant, the ICMP traffic can grow to the point that network performance degrades.

It is possible to avoid the broadcast storms created when multiple interfaces are connected to the same network. However, doing so is still a violation of TCP/IP architecture. The solution is to give the different interfaces different IP addresses on the same network. For example, you might have two token—ring adapters on the same network named tr0 and tr1. You must assign distinct IP addresses and names to tr0 and tr1. (TCP/IP architecture requires that each interface have a unique IP address and name; otherwise, unpredictable behavior will result.) For instance, you might give interface tr0 the IP address 10.10.10.1 and the name laurel.foo.bar.com, and interface tr1 the IP address 10.10.10.2 and the name hardy.foo.bar.com.

Managing Network Interfaces

To manage network interfaces, use the tasks in the following table.

Managing Network Interfaces	Tasks	
Web-based System Manager: (wsm network application) OR-	wsm network fast path	
Task	SMIT Fast Path	Command or File
List all network devices	smit Isinet	Isdev -C -c if
Configure a network device	smit chinet	See the ifconfig command and the rc.net file
Changing network interface info with remotely mounted /usr	smit chdev ^{1,2}	chgif ^{1,2}
Obtaining statistics for a network interface		netstat -v

Note:

- 1. Changes from a remotely mounted /usr affect only the Information Database (ODM) until the network is rebooted or until the ifconfig command is used to make the changes take affect right away.
- 2. When using a remotely mounted /**usr**, be careful not to modify the interface being used, since that is the location of the libraries, commands, and kernel.

TCP/IP Addressing

TCP/IP includes an Internet addressing scheme that allows users and applications to identify a specific network or host to communicate with. An Internet address works like a postal address, allowing data to be routed to the chosen destination. TCP/IP provides standards for assigning addresses to networks, subnetworks, hosts, and sockets, and for using special addresses for broadcasts and local loopback.

Internet addresses are made up of a network address and a host (or local) address. This two-part address allows a sender to specify the network as well as a specific host on the network. A unique, official network address is assigned to each network when it connects to other Internet networks. However, if a local network is not going to connect to other Internet networks, it can be assigned any network address that is convenient for local use.

The Internet addressing scheme consists of Internet Protocol (IP) addresses and two special cases of IP addresses: broadcast addresses and loopback addresses.

Internet Addresses

The Internet Protocol (IP) uses a 32-bit, two-part address field. The 32 bits are divided into four *octets* as in the following:

01111101 00001101 01001001 00001111

These binary numbers translate into:

125 13 73 15

The two parts of an Internet address are the network address portion and the host address portion. This allows a remote host to specify both the remote network and the host on the remote network when sending information. By convention, a host number of 0 (zero) is used to refer to the network itself.

TCP/IP supports three classes of Internet addresses: Class A, Class B, and Class C. The different classes of Internet addresses are designated by how the 32 bits of the address are allocated. The particular address class a network is assigned depends on the size of the network.

Class A Addresses

A Class A address consists of an 8-bit network address and a 24-bit local or host address. The first bit in the network address is dedicated to indicating the network class, leaving 7 bits for the actual network address. Since the highest number that 7 bits can represent in binary is 128, there are 128 possible Class A network addresses. Of the 128 possible network addresses, two are reserved for special cases: the network address 127 is reserved for local loopback addresses, and a network address of all ones indicates a broadcast address.

Therefore, there are 126 possible Class A network addresses and 16,777,216 possible local host addresses. In a Class A address (see figure), the highest order bit is set to 0.

Network Address	Local Host Address		
(8 bits)	(24 bits)		
01111101	00001101	01001001	00001111

Note: The high-order bit (or first bit) will always be 0 in a Class A address.

Class A Address

In other words, the first octet of a Class A address is in the range 1 to 126.

Class B Addresses

A Class B address consists of a 16-bit network address and a 16-bit local or host address. The first two bits in the network address are dedicated to indicating the network class, leaving 14 bits for the actual network address. Therefore, there are 16,384 possible network addresses and 65,536 local host addresses. In a Class B address (see figure), the highest order bits are set to 1 and 0.

Network Address (16 bits)		ost Address 6 bits)
10011101 00001101	01001001	00001111
10011101 00001101	01001001	

Note: The two highest order bits (or first two bits) will always be 1 and 0 in a Class B address.

Class B Address

In other words, the first octet of a Class B address is in the range 128 to 191.

Class C Addresses

A Class C address consists of a 24—bit network address and an 8—bit local host address. The first two bits in the network address are dedicated to indicating the network class, leaving 22 bits for the actual network address. Therefore, there are 2,097,152 possible network addresses and 256 possible local host addresses. In a Class C address (see figure), the highest order bits are set to 1 and 1.

	Network Address (24 bits)		Local Host Address (8 bits)
11011101	00001101	01001001	00001111
<u></u>			

Note: The two highest order bits (or first two bits) will always be 1 and 1 in a Class C address.

Class C Address

In other words, the first octet of a Class C address is in the range 192 to 223.

When deciding which network address class to use, you need to consider how many local hosts there will be on the network and how many subnetworks there will be in the organization. If the organization is small and the network will have fewer than 256 hosts, a Class C address is probably sufficient. If the organization is large, then a Class B or Class A address may be more appropriate.

Note: Class D (1–1–1–0 in the highest order bits) addresses provide for multicast addresses and are supported by UDP/IP under AIX.

Machines read addresses in binary code. The conventional notation for Internet host addresses is the *dotted decimal*, which divides the 32–bit address into four 8–bit fields. The following binary value:

0001010 00000010 00000000 00110100

can be expressed as:

010.002.000.052 or 10.2.0.52

where the value of each field is specified as a decimal number and the fields are separated by periods.

Note: The **hostent** command does recognize the following addresses: .08, .008, .09, and .009. Addresses with leading zeros are interpreted as octal, and numerals in octal cannot contain 8s or 9s.

TCP/IP requires a unique Internet address for each network interface (adapter) on a network. These addresses are determined by entries in the configuration database, which must agree with entries in the /etc/hosts file or the named database if the network is using a name server.

Internet Addresses Using Zeros

When a C class Internet address contains a 0 as the host address portion, (for example, 192.9.200.0), TCP/IP sends a wildcard address on the network. This means that all machines with a Class C address of 192.9.200.*X* (where *X* represents a value between 0 and 254) should respond to the request. This results in a network flooded with requests to nonexistent machines.

Similarly, problems occur for Class B addresses such as 129.5.0.0. All machines with a Class B address of 129.5. X. X. (where X represents a value between 0 and 254) are obliged to respond to the request. In this case, since Class B addresses account for bigger networks than Class C addresses, the network is flooded with significantly more requests to nonexistent machines than for a Class C network.

Subnet Addresses

Subnet addressing allows an autonomous system made up of multiple networks to share the same Internet address. The subnetwork capability of TCP/IP also makes it possible to divide a single network into multiple logical networks (subnets). For example, an organization can have a single Internet network address that is known to users outside the organization, yet configure its network internally into departmental subnets. In either case, fewer Internet network addresses are required while local routing capabilities are enhanced.

A standard Internet Protocol address field has two parts: a network address and a local address. To make subnets possible, the local address part of an Internet address is divided into a subnet number and a host number. The subnet is identified so that the local autonomous system can route messages reliably.

In the basic Class A Internet address (see figure), which consists of an 8-bit network address and 24-bit local address, the local address identifies the specific host machine on the network.

Network Address (8 bits)	Lo	ocal Host Address (24 bits)	
01111101	00001101	01001001	00001111

Class A Address

To create a subnet address for this Class A Internet address, the local address can be divided into a number identifying the physical network (or subnet) and a number identifying the host on the subnet. Senders route messages to the advertised network address, and the local system takes responsibility for routing messages to its subnets and their hosts. When deciding how to partition the local address into subnet address and host address, you should consider the number of subnets and the number of hosts on those subnets.

In the following figure, the local address is partitioned into a 12–bit subnet address and a 12–bit host address.

Network Address (8 bits)	Local Host Address (24 bits)		
Network Address	Subnet Address	Host Address	
01111101	00001101 0100	1001 00001111	

Note: The high-order bit (or first bit) will always be 0 in a Class A address.

Class A Address with Corresponding Subnet Address

You have great flexibility when assigning subnet addresses and host addresses. The bits of the local address can be divided according to the needs and potential growth of the organization and its network structure. The only restrictions are:

- network_address is the Internet address for the network.
- subnet_address is a field of a constant width for a given network.
- host_address is a field that is at least 1-bit wide.

If the width of the <code>subnet_address</code> field is 0, the network is not organized into subnets, and addressing to the network is performed using the Internet network address.

The bits that identify the subnet are specified by a bit mask and, therefore, are not required to be adjacent in the address. However, it is generally desirable for the subnet bits to be contiguous and located as the most significant bits of the local address.

Subnet Masks

When a host sends a message to a destination, the system must determine whether the destination is on the same network as the source or if the destination can be reached directly through one of the local interfaces. The system compares the destination address to the host address using the *subnet mask*. If the destination is not local, the system sends the message on to a gateway. The gateway performs the same comparison to see if the destination address is on a network it can reach locally.

The subnet mask tells the system what the subnet partitioning scheme is. This bit mask consists of the network address portion and subnet address portion of the Internet address (see figure). For example, the subnet mask of the Class A address with the partitioning scheme defined above is shown in this figure.

Network Address (8 bits)	Local Host Address (24 bits)		
Network Address	Subnet Address	Host Address	
01111101	00001101 0100	1001 00001111	

Class A Address with Corresponding Subnet Address

Network Address (8 bits)	Local Host Address (24 bits)				
Network Address	Subnet Address		Host Address		
Subnet Mask				Host Address	
01111101	00001101	0100	1001	00001111	

Class A Address with Corresponding Subnet Mask

The subnet mask is a set of 4 bytes, just like the Internetwork address. The subnet mask consists of high bits (1's) corresponding to the bit positions of the network and subnetwork address, and low bits (0's) corresponding to the bit positions of the host address. A subnet mask for the above address looks like the following figure.

Network Address (8 bits)	Local Host Address (24 bits)		
Network Address	Subnet Address	Host Address	
11111111	11111111 1111	0000 00000000	

Example Subnet Mask

Address Comparison

The destination address and the local network address are compared by performing the logical AND and exclusive OR on the subnet mask of the source host.

The comparison process is outlined below:

- 1. Perform a logical AND of the destination address and the mask of the local subnet address.
- 2. Perform an exclusive OR on the result of the previous operation and the local net address of the local interface.

If the result is all 0's, the destination is assumed to be reachable directly through one of the local interfaces.

3. If an autonomous system has more than one interface (and therefore more than one Internet address), the comparison process is repeated for each local interface.

For example, assume there are two local interfaces defined for a host network, T125. Their Internet addresses and the binary representations of those addresses are shown in the following example:

Local Network Interface Addresses

```
CLASS A 73.1.5.2 = 01001001 00000001 00000101 00000010
CLASS B 145.21.6.3 = 10010001 00010101 00000110 00000011
```

The corresponding subnet masks for the local network interfaces are shown in the following example:

Local Network Interface Addresses

```
CLASS A 73.1.5.2 = 11111111 11111111 11100000 00000000 CLASS B 145.21.6.3 = 11111111 11111111 11111111 11000000
```

If the source network, T125, is requested to send a message to a destination network with the host address 114.16.23.8 (represented in binary as: 01110010 00010000 00010111 00001000), the system checks whether the destination can be reached through a local interface.

Note: The **subnetmask** keyword must be set in the configuration database of each host that is to support subnets. Before the subnetwork capability can be used, all hosts on the network must support it. Set the subnet mask permanently in the configuration database using the Web-based System Manager wsm network application or the Network Interface Selection menu in SMIT. The subnet mask can also be set in the running system using the **ifconfig** command. Using **ifconfig** to set the subnet mask is not a permanent change.

Broadcast Addresses

The TCP/IP can send data to all hosts on a local network or to all hosts on all directly connected networks. Such transmissions are called *broadcast messages*. For example, the **routed** routing daemon uses broadcast messages to guery and respond to routing gueries.

For data to be broadcast to all hosts on all directly connected networks, User Datagram Protocol (UDP) and Internet Protocol (IP) are used to send the data, and the host destination address in the IP header has all bits set to 1. For data to be broadcast to all hosts on a specific network, all the bits in the local address part of the IP address are set to 0. There are no user commands that use the broadcast capability, although such commands, or programs, can be developed.

The broadcast address can be changed temporarily by changing the *broadcast* parameter in the **ifconfig** command. Change the broadcast address permanently by using the Web-based System Manager fast path, **wsm network**, or the SMIT fast path **smit chinet**. Changing the broadcast address may be useful if you need to be compatible with older versions of software that use a different broadcast address; for example, the host IDs are all set to 0.

Local Loopback Addresses

The Internet Protocol defines the special network address, 127.0.0.1, as a local loopback address. Hosts use local loopback addresses to send messages to themselves. The local loopback address is set by the configuration manager during the system startup process. Local loopback is implemented in the kernel and can also be set with the **ifconfig** command. Loopback is invoked when the system is started.

Getting an Official Internet Address

The Internet Registry (IR) is the organization that is responsible for assigning identifiers, such as IP network numbers and autonomous system numbers, to networks. Currently, the Defense Data Network, Network Information Center (DDN NIC) serves as the IR. To obtain an official Internet address, contact the NIC at INFO@INTERNIC.NET or at 1–800–444–4345 (USA).

TCP/IP Address and Parameter Assignment – Dynamic Host Configuration Protocol (DHCP)

Transmission Control Protocol/Internet Protocol (TCP/IP) enables communication between machines with configured addresses. Part of the burden a network administrator must face is address assignment and parameter distribution for all machines on the network. Commonly, this is a process in which the administrator dictates the configuration to each user, allowing the user to configure his own machine. However, misconfigurations and misunderstandings can generate service calls that the administrator must deal with individually. The Dynamic Host Configuration Protocol (DHCP) gives the network administrator a method to remove the end user from this configuration problem and maintain the network configuration in a centralized location.

DHCP is an application—layer protocol that allows a machine on the network, called a *client*, to get an IP address and other configuration parameters from the server. It gets information by exchanging packets between a daemon on the client and another on the server. Most operating systems now provide a DHCP client in their base package.

To obtain an address, the DHCP client daemon (**dhcpcd** on AIX) broadcasts a DHCP discover message, which is received by the server and processed. (Multiple servers can be configured on the network for redundancy.) If there is an address available for that client, a DHCP offer message is created, which contains an IP address and other options that are appropriate for that client. The client receives the DHCP offer and stores it while waiting for other offers. When the client chooses the best offer, it broadcasts a DHCP request that specifies which server's offer it wants.

All configured DHCP servers receive the request. Each checks to see if it is the requested server. If not, the server frees the address assigned to that client. The requested server marks the address as assigned and returns a DHCP acknowledgement, at which time, the transaction is complete and the client has an address for a period of time (lease) described by the server.

When half of the lease time is used, the client sends the server a *renew* packet to extend the lease time. If the server is willing to renew, it sends a DHCP acknowledgement. If the client does not get a response from the server that owns its current address, it broadcasts a DHCP rebind packet to reach the server if, for example, the server has been moved from one network to another. If the client has not renewed its address after the full lease time, the interface is brought down and the process starts over. This cycle prevents multiple clients on a network from being assigned the same address.

The DHCP server assigns addresses based on *keys*. Four common keys are network, class, vendor, and client ID. The server uses these keys to get an address and a set of configuration options to return to the client.

network Identifies which network segment the packet came from. The network key

allows the server to check its address database and assign an address by

network segment.

class Is completely client configurable. It can specify an address and options.

This key could be used to denote machine function in the network or describe how machines are grouped for administrative purposes. For example, the network administrator might want to create a netbios class that contains options for NetBIOS clients or an accounting class that represents Accounting department machines that need access to a specific

printer.

vendor Helps identify the client by its hardware/software platform (for example, a

Windows 95 client or an OS/2 Warp client).

client ID

Identifies the client either through the machine's host name or its medium access control (MAC) layer address. The client ID is specified in the configuration file of the **dhcpcd** daemon. Also, the client ID can be used by the server to pass options to a specific client or prohibit a particular client from receiving any parameters.

These keys can be used by the configuration either singularly or in combinations. If multiple keys are provided by the client and multiple addresses can be assigned, only one is chosen, and the option set is derived from the chosen key first. For more detailed information about the selection of keys and addresses, see Configuring DHCP, on page 3-62.

The other piece of the picture is a BOOTP relay agent. A relay agent is needed so initial broadcasts from the client can leave the local network. The relay agents act as forwarding agents for DHCP and BOOTP packets.

The DHCP Server

For AIX version 4.3.1, the DHCP server has been segmented into three main pieces, a database, a protocol engine, and a set of service threads, each with its own configuration information.

The DHCP Database

The **db_file.dhcpo** database is used to track clients and addresses and for access control (for example, allowing certain clients on some networks but not others or disabling BOOTP clients on a particular network). Options are also stored in the database for retrieval and delivery to clients. The database is implemented as a dynamically loadable object, which allows for easy server upgrade and maintenance.

Using the information in the configuration file, the database is primed and verified for consistency. A set of checkpoint files handles updates to the database and reduces the overhead of writes to the main storage file. The database also contains the address and option pools, but these are static and are discussed in Configuring DHCP, on page 3-62.

The main storage file and its back up are flat ASCII files that can be edited, if necessary. The format for the database main storage files are:

```
DF01
"CLIENT ID" "0.0.0.0" State LeaseTimeStart LeaseTimeDuration LeaseTimeEnd
"Server IP Address" "Class ID" "Vendor ID" "Hostname" "Domain Name"
"CLIENT ID" "0.0.0.0" State LeaseTimeStart LeaseTimeDuration LeaseTimeEnd
"Server IP Address" "Class ID" "Vendor ID" "Host Name" "Domain Name"
...
```

The first line is a version identifier for the file: DF01c. The lines that follow are client record definition lines. The server reads from the second line to the end of the file. (The parameters in quotes must be enclosed in quotes.)

"CLIENT ID" The ID the client uses to represent itself to the server.

"0.0.0.0" is the IP address currently assigned to the DHCP server. If no address has been assigned, it will be "0.0.0.0".

State

The current state of the client. The DHCP protocol engine contains the allowable set, and the states are maintained in the DHCP database. The number next to *State* represents its value. The states can be:

- (1) FREE Represents addresses that are available for use. In general, clients do not have this state unless they have no address assigned. **dadmin** and the output from **Issrc** report this state as "Free".
- (2) RESERVED Indicates client and address are tied, but loosely. The client has issued a DHCP discover message and the DHCP server has responded, but the client has not yet responded with a DHCP request for that address. **dadmin** and the output from **Issrc** report this state as "Reserved".

- (3) BOUND Indicates client and address are tied and that the client has been assigned this address for some amount of time.

 dadmin and the output from Issrc report this state as "Leased".
- (4) RELEASED Indicates the client and address are tied for informational purposes only. The DHCP protocol suggests that DHCP servers maintain information about the clients it has served for future reference (mainly to try giving the same address to that client that has been assigned that address in the past). This state indicates that the client has released the address. The address is available for use by other clients, if no other addresses are available. **dadmin** and the output from **Issrc** report this as "Released".
- (5) EXPIRED Indicates the client and address are tied together, but only for informational purposes, in a similar manner to released addresses. The expired state, however, represents clients that let their leases expire. An expired address is available for use and is reassigned after all free addresses are unavailable and before released addresses are reassigned. dadmin and the output from Issrc report this state as "Expired".
- (6) BAD Represents an address that is in use in the network but has not been handed out by the DHCP server. This state also represents addresses that clients have rejected. Clients should not have this state. **dadmin** and the output from **Issrc** report this state as "Used" and "Bad", respectively.

The one state not listed above is UNKNOWN (value 0), which represents clients that have no address assigned. This state should never apply to addresses. **dadmin** reports "Unknown" and **Issrc** reports "Corrupt" for this state.

LeaseTimeStart Is the start of the current lease time (in the number of seconds since January 1, 1970).

LeaseTimeDuration

Represents the duration of the lease (in seconds).

LeaseTimeEnd Uses the same format as LeaseTimeStart, but it represents the end of the lease. Some configuration options use different values for the start and end of a lease and these values can be overridden by configuration file options. See DHCP Server File Syntax for db file Database, on page 3-75.

"Server IP Address"

Is the IP address of the DHCP server that owns this record.

[&]quot;Class ID"

[&]quot;Vendor ID"

[&]quot;Host Name"

[&]quot;Domain Name" Are values that the server uses to determine which options are sent to the server (stored as quoted strings). These parameters increase performance, since option lists can be pregenerated for these clients when the DHCP server starts up.

Checkpoint Files

The syntax for the checkpoint files is not specified. If the server crashes or you have to kill it and cannot do a normal closing of the database, the server can process the checkpoint and backup files to reconstruct a valid database. The worst case could be to lose one client (the client that was being written to the checkpoint file when it crashed). The default files are:

/etc/db_file.cr normal database operation
/etc/db_file.crbk backups for the database
/etc/db_file.chkpt and /etc/db_file.chkpt2
rotating checkpoint files

The DHCP server for AIX Versions 4.3.1 and later is threaded. To maintain high throughput, database operations (including save operations) are thread efficient. When a save is in requested, the existing checkpoint file is rotated to the next checkpoint file, the existing database file is copied to the backup file, and the new save file is created. Each client record is then logged and a bit is toggled to indicate that the client should use the new checkpoint file for logging. When all client records are recorded, the save is closed, and the backup and old checkpoint files are deleted. This way, clients can still be processed and, depending on whether the client's record has been saved, database changes will go into a new save file or to a new checkpoint file.

The DHCP Protocol Engine

For AIX Versions 4.3.1 and later, the DHCP protocol engine has been updated to RFC 2131, but is still compatible with RFC 1541. (The server can also process options as defined in RFC 2132.) The protocol engine uses the database to determine what information should be returned to the client.

The configuration of the address pools have some configuration options that affect the state machine. For example, the DHCP server pings addresses before it hands them out. The amount of time the server waits for a response is now configurable for each address pool.

DHCP Threaded Operations

The last piece of the DHCP server is actually a set of operations that are used to keep things running. Since the DHCP server is threaded, these operations are actually set up as threads that occasionally do things to make sure everything is together.

The first thread, the **main** thread, handles the SRC requests (such as **startsrc**, **stopsrc**, **Issrc**, **traceson**, and **refresh**). This thread also coordinates all operations that affect all threads and handles signals. For example,

- A SIGHUP (-1) causes a refresh of all databases in the configuration file.
- A SIGTERM (-15) will cause the server to gracefully stop.

The next thread, the **dadmin** thread, interfaces with **dadmin** client program and the DHCP server. The dadmin tool can be used to get status as well as modify the database to avoid editing the database files manually. Previous versions of the DHCP server prevented any clients from getting addresses if a status request was running. With the addition of the **dadmin** and **src** threads, the server can handle service requests and still handle client requests.

The next thread is the **garbage** thread, which runs timers that periodically clean the database, save the database, purge clients that do not have addresses, and remove reserved addresses that have been in reserve state for too long. All these timers are configurable (see Configuring DHCP, on page 3-62). The other threads are packet processors. The number of these is configurable; the default is 10. Each of these can handle a request from a DHCP client. The number of packet processors required is somewhat load— and machine—dependent. If the machine is used for other services than DHCP, it may not be wise to start up 500 threads.

Planning DHCP

To use this protocol, the network administrator needs to set up a DHCP server and configure BOOTP relay agents on links that do not have a DHCP server. Advance planning can reduce DHCP load on the network. For example, one server can be configured to handle all your clients, but all packets must be passed through it. If you have a single router between two large networks, it may be wiser to place two servers in your network, one on each link.

Another aspect to consider is that DHCP implies a pattern of traffic. For example, if you set your default lease time to fewer than two days and your machines are powered off for the weekend, Monday morning will be a period of high DHCP traffic. Although DHCP traffic does not cause huge overhead for the network, it needs to be considered when deciding where to place DHCP servers on a network and how many to use.

The goal of DHCP is that, after enabling DHCP to get the client on the network, a client has no requirement to enter anything. The DHCP client, <code>dhcpcd</code>, reads the <code>dhcpcd.ini</code> file, which contains information on logging and other parameters needed to start running. After installation, you need to decide which method to use for TCP/IP configuration: minimum configuration or DHCP. If DHCP is selected, you must choose an interface and you can specify some optional parameters. To choose the interface, you can select the keyword <code>any</code>, which tells <code>dhcpcd</code> to find the first interface that works and use it. This method minimizes the amount of input on the client side.

Configuring DHCP

By default, the DHCP server is configured by reading the /etc/dhcpsd.cnf file, which specifies the server's initial database of options and addresses. The server is started in the /etc/rc.tcpip file, or it can be started from Web-based System Manager, from SMIT, or through SRC commands. The DHCP client can be configured by running Web-based System Manager, the System Management Interface Tool (SMIT), or editing a flat ASCII file.

Configuring the DHCP server is usually the hardest part of using DHCP in your network. First, figure out what networks you need to have DHCP clients on. Each subnet in your network represents a pool of addresses that the DHCP server must add to its database. For example:

The example above shows a subnet, 9.3.149.0, with a subnet mask 255.255.255.0. All addresses in this subnet, 9.3.149.1 through 9.3.149.254, are in the pool. Optionally, a range can be specified on the end of the line or a range or exclude statement can be included in the subnet container. See DHCP Server File Known Options, on page 3-70 for common configuration methods and definitions.

The database clause with <code>db_file</code> indicates which database method to use for processing this part of the configuration file. Comments begin with a # (pound sign). From the # to the end of the line are ignored by the DHCP server. Each <code>option</code> line is used by the server to tell the client what to do. DHCP Server File Known Options, on page 3-70 describes the currently supported and known options. See DHCP Server File Syntax for General Server Operation, on page 3-73 for ways to specify options that the server does not know about.

If the server does not understand how to parse an option, it uses default methods to send the option to the client. This also allows the DHCP server to send site—specific options that are not RFC defined, but may be used be certain clients or client configurations.

The Configuration File

The configuration file has an address section and an option definition section, which are based on the concept of containers that hold options, modifiers, and, potentially, other containers.

A *container* (basically, a method to group options) uses an identifier to classify clients into groups. The container types are subnet, class, vendor, and client. Currently, there is not a generic user–definable container. The identifier uniquely defines the client so that the client can be tracked if, for example, it moves between subnets. More than one container type may be used to define client access.

Options are identifiers that are returned to the client, such as default gateway and DNS address.

Modifiers are single statements that modify some aspect of a container, such as lease time default.

Containers

When the DHCP server receives a request, the packet is parsed and identifying keys determine which containers, options, and addresses are extracted.

The previous example shows a subnet container. Its identifying key is the client's position in the network. If the client is from that network, then it falls into that container.

Each type of container uses a different option to identify a client:

- The subnet container uses the giaddr field or the interface address of the receiving interface to determine which subnet the client came from.
- The class container uses the value in option 77 (User Site Class Identifier).
- The vendor uses the value in option 60 (Vendor Class Identifier).
- The client container uses the option 61 (Client Identifier) for DHCP clients and the chaddr field in the BOOTP packet for BOOTP clients.

Except for subnets, each container allows the specification of the value that it will match including regular expression matching.

There is also an implicit container, the *global* container. Options and modifiers placed in the global container apply to all containers unless overridden or denied. Most containers can be placed inside other containers implying a scope of visibility. Containers may or may not have address ranges associated with them. Subnets, by their nature, have ranges associated with them.

The basic rules for containers and subcontainers are as follows:

- · All containers are valid at the global level.
- Subnets can never be placed inside other containers.
- Restricted containers cannot have regular containers of the same type within them. (For
 example, a container with an option that only allows a class of Accounting cannot
 include a container with an option that allows all classes that start with the letter "a". This
 is illegal.)
- Restricted client containers cannot have subcontainers.

Given the above rules, you can generate a hierarchy of containers that segment your options into groups for specific clients or sets of clients.

If a client matches multiple containers, how are options and addresses handed out? The DHCP server receives messages, it passes the request to the database (**db_file** in this case), and a container list is generated. The list is presented in order of depth and priority. Priority is defined as an implicit hierarchy in the containers. Strict containers are higher priority than regular containers. Clients, classes, vendors, and finally subnets are sorted, in

that order, and within container type by depth. This generates a list ordered by most specific to least specific. For example:

```
Subnet 1
--Class 1
--Client 1
Subnet 2
--Class 1
----Vendor 1
----Client 1
---Client 1
```

The example shows two subnets, <code>Subnet 1</code> and <code>Subnet 2</code>. There is one class name, <code>Class 1</code>, one vendor name, <code>Vendor 1</code>, and one client name, <code>Client 1</code>. <code>Class 1</code> and <code>Client 1</code> are defined in multiple places. Because they are in different containers, their names can be the same but values inside them may be different. If <code>Client 1</code> sends a message to the <code>DHCP</code> server from <code>Subnet 1</code> with <code>Class 1</code> specified in its option list, the <code>DHCP</code> server would generate the following container path:

```
Subnet 1, Class 1, Client 1
```

The most specific container is listed last. To get an address, the list is examined in reverse hierarchy to find the first available address. Then, the list is examined in forward hierarchy to get the options. Options override previous values unless an option *deny* is present in the container. Also, since Class 1 and Client 1 are in Subnet 1, they are ordered according to the container priority. If the same client is in Subnet 2 and sends the same message, the container list generated is:

```
Subnet 2, Class 1, Client 1 (at the Subnet 2 level), Client 1 (at the Class 1 level)
```

Subnet 2 is listed first, then Class 1, then the Client 1 at the Subnet 2 level (because this client statement is only one level down in the hierarchy). The hierarchy implies that a client matching the first client statement is less specific than the client matching Client 1 of Class 1 within Subnet 2.

Priority selected by depth within the hierarchy is not superseded by the priority of the containers themselves. For example, if the same client issues the same message and specifies a vendor identifier, the container list is:

```
Subnet 2, Class 1, Vendor 1, Client 1 (at Subnet 2 level), Client 1 (at Class 1 level)
```

Container priority improves search performance because it follows a general concept that client containers are the most specific way to define one or more clients. The class container holds less specific addresses than a client container; vendor is even less specific; and subnet is the least specific.

Addresses and Address Ranges

Any container type may have associated addresses ranges; subnets must have. Each range within a container must be a subset of the parent container's range and must not overlap with other containers' ranges. For example, if a class is defined within a subnet and the class has a range, the range must be a subset of the subnet's range. Also, the range within that class container cannot overlap with any other ranges at its level.

Ranges can be expressed on the container line and modified by range and exclude statements to allow for disjoint address sets associated with a container. So, if you have the top ten addresses and the second ten addresses of a subnet available, the subnet could specify these addresses by range in the subnet clause to reduce both memory use and the chance of address collision with other clients not in the specified ranges.

Once an address has been selected, any subsequent container in the list that contains address ranges is removed from the list along with its children. The reason for this is that network–specific options in removed containers are not valid if an address is not used from within that container.

Options

After the list has been culled to determine addresses, a set of options is generated for the client. In this selection process, options overwrite previously selected options unless a *deny* is encountered, in which case, the denied option is removed from the list being sent to the client. This method allows inheritance from parent containers to reduce the amount of data that must be specified.

Modifiers

Modifiers are items that modify some aspect of a particular container, such as access or lease time. After defining the address and option pools, start thinking about the container modifiers. The most common are **leasetimedefault**, **supportBootp**, and **supportUnlistedclients**.

leasetimedefault Defines the amount of time an address is to be leased to a client.

supportBootp Defines whether the server should respond to BOOTP clients.

supportUnlistedclients

Indicates whether clients should be explicitly defined by a client statement to receive addresses. The value for supportUnlistedClients can be **none**, **dhcp**, **bootp**, or **both**. This allows for you to restrict access to bootp client and allow all DHCP clients to get addresses.

Other modifiers are listed in DHCP Server File Syntax for db file Database, on page 3-75.

Logging

After selecting modifiers, the next item to set up is logging. Logging parameters are specified in a container like the database, but the container keyword is <code>logging_info</code>. When learning to configure DHCP, it is advisable to turn logging to its highest level. Also, it is best to specify the logging configuration prior to any other configuration file data to ensure that configuration errors are logged after the logging subsystem is initialized. Use the <code>logitem</code> keyword to turn on a logging level or remove the <code>logitem</code> keyword to disable a logging level. Other keywords for logging allow the specification of the log filename, file size, and the number of rotating log files.

Server-specific Options

The last set of parameters to tweak are server—specific options, which allow the user to control the number of packet processors, how often the garbage collection threads are run, etc.

For example, two server-specific options are:

reservedTime Indicates how long an address should stay in reserved state after sending an OFFER to the DHCP client

reservedTimeInterval

Indicates how often the DHCP server should scan through the addresses to see if there are any that have been in the reserved state long than **reservedTime**.

These options are useful if you have several clients that broadcast DISCOVER messages, but either they do not broadcast their REQUEST message or their REQUEST message gets lost in the network. Using these parameters keeps addresses from being reserved indefinitely for a noncompliant client.

Another particularly useful option is **SaveInterval**, which indicates how often saves should occur. All server–specific options are listed in DHCP Server File Syntax for General Server Operation, on page 3-73 with the logging keywords.

Performance Considerations

It is important to understand that certain configuration keywords and the structure of the configuration file have an effect on the memory use and performance of the DHCP server.

First, excessive memory use can be avoided by understanding the inheritance model of options from parent to child containers. In an environment that supports no unlisted clients, the administrator must explicitly list each client in the file. When options are listed for any specific client, the server uses more memory storing that configuration tree than when options are inherited from a parent container (for example, the subnet, network, or global containers). Therefore, the administrator should verify whether any options are repeated at the client level within the configuration file and, if so, determine whether these options can be specified in the parent container and shared by the set of clients as a whole.

Also, when using the **logItem** entries INFO and TRACE, numerous messages are logged during the processing of every DHCP client's message. Appending a line to the log file can be an expensive operation; therefore, limiting the amount of logging improves the performance of the DHCP server. When an error with the DHCP server is suspected, logging can be dynamically re—enabled using either the SRC traceson or dadmin commands.

Finally, selecting a **numprocessors** value should depend on the size of the DHCP–supported network, the **pingTime db_file** configuration parameter, and the typical propagation delay on the network. Because each packet processor thread issues an ICMP Echo Request to verify the status of a server–owned address before offering it to a client, the amount of time that any Echo Response is waited on directly affects the amount of processing time for a DISCOVER message. Essentially, the packet processor thread is able to do nothing more than wait for any response or for the **pingTime** timeout. Therefore, lowering the **numprocessors** value improves the response time of the server, lowering the number of retransmissions by clients, yet still maintaining the ping benefit of the server's design.

For best performance, select a **pingTime** based on the propagation delay of any remote networks supported by the DHCP server. Also, select the **numprocessors** value based on this **pingTime** value and the size of the network. A selected value that is too small can cause all packet processing threads to be stopped waiting for any Echo Responses while incoming DHCP client messages queue on the server port, which would result in the server handling client messages in batches rather than in a constant stream.

To prevent this situation, the value for **numprocessors** should be set to a number higher than the estimated number of DISCOVER messages that can be received within one **pingTime** interval during a period of high DHCP client activity. However, the **numprocessors** value should not be set so high that it could burden the kernel with thread management.

For example, the values **numprocessors 5** and **pingTime 300** cause poor performance in an environment with a potential 10 DISCOVER messages per second because at peak demand, only 5 messages are handled every 3 seconds. This environment should be configured with values similar to **numprocessors 20** and **pingTime 80**.

Customizing a Configuration File

Many system administrators deal with networks that include multiple client types; for example, a single network may include Windows 95 computers, AIX computers, OS/2 computers, Java OS clients, and IBM Thin Client machines. Each of these require unique vendor identifiers (the field used to identify the type of machine to the DHCP server). Java OS clients and Thin Client machines can require unique parameters such as bootfiles, and configuration options may need to be tailored specifically for them. But the Windows 95 computers would not handle Java—specific options well. So, machine—specific options can be encapsulated within vendor containers.

To take this example further, imagine a primary use for certain machines based on the type of user for those machines. For instance, the development staff may use the AIX clients for programming, the marketing staff may use the OS/2 clients, sales may use Java OS clients

and IBM Thin Client machines, and accounting may use Windows 95 machines. Each of these user families could need different configuration options (different printers, nameservers, or default web servers, etc.). In this case, such options could be included in the vendor container, since each group uses a different machine type. But if the same machine type is used by multiple groups, placing the options within a subordinate class identifier would allow your marketing managers, for example, to use a specific set of printers that other employees could not access.

Note: The following fictional example represents part of a configuration file. Comments are preceded by a pound sign (#) and describe how each line defines the installation.

```
vendor "AIX CLIENT"
# No specific options, handles things based on class
vendor "OS/2 Client"
# No specific options, handles things based on class
vendor "Windows 95"
{ option 44 9.3.150.3  # Default NetBIOS Nameserver
vendor "Java OS"
{ bootstrapserver 9.3.150.4 # Default TFTP server for the Java OS boxes
 option 67 "javaos.bin" # The bootfile of the Java OS box
vendor "IBM Thin Client"
{ bootstrapserver 9.3.150.5  # Default TFTP server for Thin Client boxes option 67 "thinos.bin"  # Default bootfile for the Thin Client boxes
subnet 9.3.149.0 255.255.255.0
{ option 3 9.3.149.1  # The default gateway for the subnet option 6 9.3.150.2  # This is the nameserver for the subnet
  class accounting 9.3.149.5-9.3.149.20
                                  # The accounting class is limited to address
                                    range 9.3.149.5-9.3.149.20
                                  # The printer for this group is also in this
                                   range, so it is excluded.
     exclude 9.3.149.15
     option 9 9.3.149.15
                                # The LPR server (print server)
     vendor "Windows 95"
     option 9 deny
                                 # This installation of Windows 95 does not
                                   support
                                  # this printer, so the option is denied.
     }
  }
```

DHCP and the Dynamic Domain Name System (DDNS)

The DHCP server provides options to enable operation in a DDNS environment. To use DHCP in a DDNS environment, you must set and use a Dynamic Zone on a DNS server.

After the DDNS server is configured, you must decide if the DHCP server is going to do A record updates, PTR record updates, updates for both record types, or none at all. This decision depends on whether a client machine can do part or all of this work.

- If the client can share update responsibility, you probably want the server to do the PTR record updates and the client to do theA record updates.
- If the client can do both updates, configure the server to do none.
- If the client cannot do updates, configure the server to do both.

The DHCP server has a set of configuration keywords that allow you to specify a command to run when an update is required. These are:

updatedns (Deprecated.) Represents the command to issue to do any type of update.

It will be called for both the PTR and the A record update.

updatednsA Specifies the command to update the A record.

updatednsP Specifies the command to update the PTR record.

These keywords specify executable strings that the DHCP server executes when an update is required. The keyword strings must contain four <code>%s</code> (percent symbol, letter s). The first <code>%s</code> is the hostname; the second, the domain name; the third, the IP address; and the fourth, the lease time. These are used as the first four parameters for the AIX command, **dhcpaction**. The remaining two parameters for **dhcpaction** indicate the record to update (A, PTR, NONE, or BOTH) and whether NIM should be updated (NIM or NONIM). See DHCP and Network Installation Management (NIM) Interactions and Suggestions, on page 3-89 for more information about NIM and DHCP interaction. For example:

The DHCP server also has a set of keywords to remove the DNS entries when a lease is released or expires. The keywords are:

releasednsA Removes the A record.

releasednsP Removes the PTR record.

removedns Removes both record types.

These keywords specify executable strings that the DHCP server executes when an address is released or expired. AIX provides a command, **dhcpremove**, that works similarly to **dhcpaction**, but only takes three parameters:

- 1. The IP address, specified as a %s in the command string
- 2. Which record to remove (A, PTR, NONE, or BOTH).
- 3. Whether NIM should be updated (NIM or NONIM).

```
releasednsA "/usr/sbin/dhcpremove '%s' '%s' '%s' A NONIM"

# This does the dhcpremove command only the A record releasednsP "/usr/sbin/dhcpremove '%s' '%s' '%s' PTR NONIM"

# This does the command only on the PTR record removedns "/usr/sbin/dhcpremove '%s' '%s' BOTH NIM"

# This does the command on both records and updates NIM
```

The **dhcpaction** and **dhcpremove** scripts do some parameter checking, then set up a call to **nsupdate**, which has been updated to work with the AIX and OS/2 DDNS servers. See the **nsupdate** command description for more information.

If NIM interaction is **NOT** required by the name update, the DHCP server can be configured to use a socket transfer between the DHCP daemon and the **nsupdate** command to improve performance and enable DNS updates to be retried upon failure. To configure this option, the **updateDNSA**, **updateDNSP**, **releaseDNSA**, or the **releaseDNSP** keyword must specify "nsupdate_daemon" as the first quoted word. The parameters and flags for this update are identical to those that are accepted by the **nsupdate** command. Additionally, the following variable names can be used for substitution:

\$hostname Replaced by the host name of the client on DNS update or

the host name previously associated with the client for DNS

removal.

\$domain Replaced by the DNS domain for the update or the

previously used domain of the client's host name for a DNS

removal.

\$ipadress Replaced by the IP address to be associated or

disassociated from the DHCP client's name.

\$leasetime Replaced by the lease time (in seconds).

\$clientid Replaced by the string representation of the DHCP client's

identifier or the combination hardware type and hardware

address for BOOTP clients.

For example:

```
updateDNSA "nsupdate_daemon -p 9.3.149.2 -h $hostname -d $domain -s"d;a;*;a;a;$ipaddress;s;$leasetime;3110400""

updateDNSP "nsupdate_daemon -p 9.3.149.2 -r $ipaddress -s"d;ptr;*;a;ptr;$hostname.$domain.;s;$leasetime;3110400""

releaseDNSA "nsupdate_daemon -p 9.3.149.2 -h $hostname -d $domain -s"d;a;*;s;1;3110400""

releaseDNSP "nsupdate_daemon -p 9.3.149.2 -r $ipaddress -s"d;ptr;*;s;1;3110400""
```

See the **nsupdate** command description for more information.

Also, administrator—defined policies have been added for hostname exchanges between the server and the clients. By default, the hostname that is returned to the client and used for a DDNS update is option 12 (defined in the server configuration file). Alternatively, the default hostname can be a the client suggested hostname, either through option 81 (the DHCPDDNS option) or through option 12 (the HOSTNAME option). However, the administrator can override the default hostname by using the **hostnamepolicy**, **proxyarec**, and **appenddomain** configuration keywords. These options and their parameters are defined in DHCP Server File Syntax for db file Database, on page 3-75.

DHCP Compatibility with Older Versions

The DHCP server for AIX Versions 4.3.1 and later recognizes previous versions' configuration and database files, **dhcps.ar** and **dhcps.cr**. It parses the old configuration files and generates new database files in the old locations. The old databases are converted automatically to the new file. The configuration file itself is not converted.

The DHCP server database module, **db_file**, can read the old format. The DHCP server can recognize when a database container is not in the configuration file and treat the whole file as configuring the server parameters, logging parameters, and the **db_file** database parameters.

Note: Some old configuration file syntax is deprecated, but is still supported. Other deprecations are:

- The network container is completely deprecated. To specify correctly, either convert the network clause with a range into a valid subnet container with a subnet address, subnet netmask, and the range. If the network container has subnet containers, remove the network container keyword and its braces and then place the subnet mask in the appropriate place on the line. To start using the database container, group everything that pertains to networks and client accessinto one database container of type **db file**.
- The updatedns and removedns keywords are deprecated and replaced in favor of specifying the action for A and PTR records separately.

- The clientrecorddb and addressrecorddb keywords have been deprecated to clientrecorddb and backupfile, respectively.
- The option sa and option ga keywords have been replaced by bootstrapserver and giaddrfield keywords, respectively. See DHCP Server File Syntax for General Server Operation, on page 3-73 and DHCP Server File Syntax for db_file Database, on page 3-75 for more information.

DHCP Server File Known Options

Note: Note: The options that are shown in the following table as not allowed to be specified (No in the Can Specify? column) can be specified in the configuration file, but are overwritten by the proper value. For a better definition of each option, see RFC 2132.

Option Number	Default Data Type	Can Specify?	Description/Use	
0	None	No	Pads out the option field, if necessary. The server adds padding when necessary.	
1	Dotted quad	No	Net mask of the subnet from which the address was drawn.	
2	32-bit integer	Yes	Specifies the offset of the client's subnet, in seconds from Coordinated Universal Time (UTC).	
3	One or more dotted quads	Yes	A list of the default gateways' IP addresses.	
4	One or more dotted quads	Yes	A list of time servers' IP addresses.	
5	One or more dotted quads	Yes	A list of name servers' IP addresses.	
6	One or more dotted quads	Yes	A list of DNS IP addresses.	
7	One or more dotted quads	Yes	A list of log servers' IP addresses.	
8	One or more dotted quads	Yes	A list of cookie servers' IP addresses.	
9	One or more dotted quads	Yes	A list of LPR servers' IP addresses.	
10	One or more dotted quads	Yes	A list of Impress servers' IP addresses.	
11	One or more dotted quads	Yes	A list of Resource Location servers' IP addresses.	
12	An ASCII string	Yes	A hostname for the client to use.	
13	16-bit unsigned integer	Yes	Size of the bootfile.	
14	An ASCII string	Yes	Path for Merit Dump file.	
15	An ASCII string	Yes	The default DNS domain name.	
16	An IP address	Yes	The address of the Swap server.	
17	An ASCII string	Yes	The default root path.	
18	An ASCII string	Yes	The path to extensions for the client.	
19	Yes, No, True, False, 1, 0	Yes	Whether IP Forwarding should be turned on.	
20	Yes, No, True, False, 1, 0	Yes	Whether non-local source routing should be used.	
21	One or more pairs of dotted quads, in the form DottedQuad: DottedQuad	Yes	Filter policies for IP addresses.	

Option Number	Default Data Type	Can Specify?	Description/Use	
22	16-bit unsigned integer	Yes	Maximum size to allow for datagram fragments.	
23	8-bit unsigned integer	Yes	IP time-to-live (TTL).	
24	32-bit unsigned integer	Yes	Number of seconds to use in the Path MTU aging timeout.	
25	List of one or more 16–bit unsigned integers	Yes	Path MTU Plateau table. Specifies a set of values that represent the MTU sizes to use when using Path MTU discovery.	
26	16-bit unsigned integer	Yes	MTU size for the receiving interface.	
27	Yes, No, True, False, 1, 0	Yes	Whether all subnets are local.	
28	An IP address (dotted quad)	Yes	Broadcast address for the interface.	
29	Yes, No, True, False, 1, 0	Yes	Whether ICMP netmask discovery should be used.	
30	Yes, No, True, False, 1, 0	Yes	Whether client should become an ICMP netmask supplier.	
31	Yes, No, True, False, 1, 0	Yes	Whether ICMP Router Discovery messages should be used.	
32	IP address (dotted quad)	Yes	Address to use for router solicitation.	
33	One or more IP address pairs, in the form DottedQuad:DottedQuad	Yes	Each address pair represents a static route.	
34	Yes/No, True/False, 1/0	Yes	Whether trailer encapsulation should be used.	
35	32-bit unsigned integer	Yes	ARP cache timeout value.	
36	Yes/No, True/False, 1/0	Yes	Whether Ethernet encapsulation should be used.	
37	8-bit unsigned integer	Yes	TCP time-to-live (TTL).	
38	32-bit unsigned integer	Yes	TCP keep alive interval.	
39	Yes/No, True/False, 1/0	Yes	Whether TCP keep alive should be used.	
40	An ASCII string	Yes	The NIS default domain.	
41	One or more dotted quads	Yes	The IP addresses of the NIS servers.	
42	One or more dotted quads	Yes	The IP addresses of the NTP servers.	
43	hex string of digits, in the form of hex "digits", hex "digits", or 0xdigits	Yes, but really only specified with vendor container	Encapsulated option container for the vendor container.	
44	One or more dotted quads	Yes	NetBIOS name servers' IP addresses.	
45	One or more dotted quads	Yes	NetBIOS datagram distribution servers' IP addresses.	
46	8-bit unsigned integer	Yes	NetBIOS Node Type.	
47	hex string of digits, in form of hex "digits", hex "digits", or 0xdigits	Yes	NetBIOS Scope.	
48	One or more dotted quads	Yes	X Windows font servers' IP addresses.	

Option Number	Default Data Type	Can Specify?	Description/Use
49	One or more dotted quads	Yes	X Windows Display Manager.
50	None	No	Requested IP Address, used by client to indicate the address it wants.
51	32-bit unsigned integer	Yes	Lease time for the returned address. By default, the DHCP server uses the leasetimedefault keyword, but direct specification of option 51 overrides it.
52	None	No	Option overload. Client uses this to indicate the sname and file fields of the BOOTP packet may have options.
53	None	No	DHCP server or client uses this option to indicate the type of DHCP message.
54	None	No	DHCP server or client uses this option to indicate the server's address or the server to which the message is directed.
55	None	No	DHCP client uses this to indicate desired options.
56	An ASCII string	Yes	A string the DHCP server sends to the client. In general, this can be used by the DHCP server and client to indicate problems.
57	No	No	DHCP client uses this option to tell the DHCP server the maximum DHCP packet size the client can receive.
58	32-bit unsigned integer	Yes	The number of seconds until the client should send a renew packet.
59	32-bit unsigned integer	Yes	The number of seconds until the client should send a rebind packet.
60	None	No	DHCP client uses this option to indicate its vendor type. The DHCP server uses this field to match vendor containers.
61	None	No	DHCP client uses this to uniquely identify itself. The DHCP server uses this field to match client containers.
64	An ASCII string	Yes	Specifies the NIS+ domain.
65	One or more dotted quads	Yes	IP Addresses of NIS+ servers.
66	An ASCII string	Yes	Specifies the TFTP server name. This is a hostname and is used instead of the siaddr field if the client understands this option.
67	An ASCII string	Yes	Specifies the bootfile name. This can be used instead of the bootfile keyword, which places the file in the filename field of the packet.
68	One or more dotted quads, or NONE	Yes	Addresses of home agents.
69	One or more dotted quads	Yes	Default SMTP servers to use.

Option Number	Default Data Type	Can Specify?	Description/Use
70	One or more dotted quads	Yes	Default POP3 servers to use.
71	One or more dotted quads	Yes	Default NNTP servers to use.
72	One or more dotted quads	Yes	Default WWW servers to use.
73	One or more dotted quads	Yes	Default Finger servers to use.
74	One or more dotted quads	Yes	Default IRC servers to use.
75	One or more dotted quads	Yes	Default Street Talk servers to use.
76	One or more dotted quads	Yes	Default Street Talk directory assistance servers to use.
77	An ASCII string	Yes	The user site class identifier. The DHCP server uses this field to match class containers.
81	An ASCII string plus other items	No	The DHCP client uses this option to define the policy the DHCP server should use with respect to DDNS.
255	None	No	DHCP server and client use this option to indicate the end of an option list.

DHCP Server File Syntax for General Server Operation

Note: Time Units (*time_units*) shown in the following table are optional and represent a modifier to the actual time. The default time unit is minutes. Valid values are seconds (1), minutes (60), hours (3600), days (86400), weeks (604800), months (2392000), and years (31536000). The number shown in parentheses is a multiplier applied to the specified value *n* to express the value in seconds.

Keyword	Form	Subcont ainers?	Default Value	Meaning
database	database <i>db type</i>	Yes	None	The primary container that holds the definitions for the address pools, options, and client access statements. db type is the name of a module that is loaded to process this part of the file. The only value available to current AIX versions is db_file.
logging_info	logging_info	Yes	None	The primary logging container that defines the logging parameters.

Keyword	Form	Subcont ainers?	Default Value	Meaning
logitem	logitem NONE	No	All default	Enables the logging level. Multiple
	logitem SYSERR	1	to not enabled.	lines are allowed.
	logitem OBJERR		oriabioa.	
	logitem PROTOCOL	1		
	logitem PROTERR	1		
	logitem WARN	1		
	logitem WARNING	1		
	logitem CONFIG	1		
	logitem EVENT	1		
	logitem PARSEERR	1		
	logitem ACTION	1		
	logitem ACNTING	1		
	logitem STAT	1		
	logitem TRACE	1		
	logitem RTRACE	1		
	logitem START	1		
numLogFiles	numLogFiles n	No	0	Specifies the number of log files to create. The log rotates when the first one fills. n is the number of files to create.
logFileSize	logFileSize n	No	0	Specifies the size of each log file in 1024–byte units.
logFileName	logFileName path	No	None	Specifies the path to the first log file. The original log file is named filename or filename.extension. The filename must be eight or fewer characters. When a file is rotated, it is renamed beginning with the base filename, then either appending a number or replacing the extension with a number. For example, if the original file name is file, the rotated file name becomes file01. If the original file name is file.log, it becomes file.01.
CharFlag	charflag yes	No	true	Not applicable to the AIX DHCP
	charflag true			server, but the OS/2 DHCP server uses it to produce debug windows.
	charflag false			l
	charflag no			
StatisticSnapS hot	StatisticSnapShot n	No	–1, never	Specifies how often statistics are written to the log file in seconds.
UsedIpAddress ExpireInterval	UsedlpAddressExpireInt erval <i>n time_units</i>	No	–1, never	Specifies how often addresses placed in the BAD state are recouped and retested for validity.

Keyword	Form	Subcont ainers?	Default Value	Meaning
leaseExpireInte rval	leaseExpireInterval <i>n</i> time_units	No	900 seconds	Specifies how often addresses in the BOUND state are checked to see if they have expired. If the address has expired, the state is moved to EXPIRED.
reservedTime	reservedTime n time_units	No	–1, never	Specifies how long addresses should sit in RESERVED state before being recouped into the FREE state.
reservedTimeIn terval	reservedTimeInterval n time_units	No	900 seconds	Specifies how often addresses in the RESERVE state are checked to see if they should be recouped into the FREE state.
saveInterval	saveInterval <i>n</i> time_units	No	3600 seconds	Specifies how often the DHCP server should force a save of the open databases. For heavily loaded servers, this should be 60 or 120 seconds.
clientpruneintv	clientpruneintv n time_units	No	3600 seconds	Specifies how often the DHCP server has the databases remove clients are not associated with any address (in the UNKNOWN state). This reduces the memory use of the DHCP server.
numprocessors	numprocessors n	No	10	Specifies the number of packet processors to create. Minimum of one.
userObject	userObject <i>obj_name</i>	Yes	None	Indicates that the server should load a user-defined shared object and call routines within this object through each interaction with DHCP clients. The object to be loaded is located in the /usr/sbin directory by the name obj_name.dhcpo. See the DHCP Server User-Defined Extension API for more information.

DHCP Server File Syntax for db_file Database

Note: Time Units (*time_units*) shown in the following table are optional and represent a modifier to the actual time. The default time unit is minutes. Valid values are seconds (1), minutes (60), hours (3600), days (86400), weeks (604800), months (2392000), and years (31536000). The number shown in parentheses is a multiplier applied to the specified value *n* to express the value in seconds.

Also, items that are specified in one container can be overridden inside a subcontainer. For example, you could globally define BOOTP clients, but within a certain subnet allow BOOTP clients by specifying the supportBootp keywork in both containers.

The client, class, and vendor containers allow for regular expression support. For class and vendor, a quoted string with the first character after the quote being an exclamation point (!) indicates that the rest of the string should be treated as a regular expression. The client container allows for regular expressions on both the hwtype and the hwaddr fields. A single string is used to represent both fields with the following format:

decimal_number-data

If decimal_number is zero, then data is an ASCII string. If any other number, data is hex digits.

Keyword	Form	Subcont ainers?	Default Value	Meaning
subnet	subnet subnet id netmask	Yes	None	Specifies a subnet and a pool of addresses. All addresses are assumed to be in the pool unless a range is specified on
	subnet subnet id netmask range			the line or addresses are modified later in the container by a range or exclude statement. The optional range is a pair of IP addresses in dotted guad format separated
	subnet subnet id netmask label:priority			by a dash. An optional label and priority can be specified. These are used by virtual subnets to identify and order the subnets in the virtual
	subnet subnet id netmask range label:priority			subnet. The label and priority are separated by a colon. These containers are only allowed at the global or database container level.
subnet	subnet subnet id range	Yes	None	Specifies a subnet that goes within a network container. It defines a range of addresses that is the whole subnet unless the optional range part is specified. The netmask associated with the subnet is taken from the surrounding network container.
				Note: This method is deprecated in favor of the other subnet forms.

Keyword	Form	Subcont ainers?	Default Value	Meaning
option	option number data option numberdeny option * deny	No	None	Specifies an option to send to a client or, in the case of deny, an option to prevent from being sent to the client. The option * deny clause means all options not specified in the current container are not to be returned to the client. option numberdeny only denies the specified option. number is an unsigned 8—bit integer. data is specific to the option (see above) or can be specified as a quoted string (indicating ASCII text) or 0xhexdigits or hex"hexdigits" or hex "hexdigits". If the option is in a vendor container, the option will be encapsulated with other options in an option 43.
exclude	exclude an IP address exclude dotted_quad-dotted_quad	No	None	Modifies the range on the container in which the exclude statement is in. The exclude statement is not valid in the global or database container levels. The exclude statement removes the specified address or range from the current range on the container. The exclude statement allows you to create non—contiguous ranges for subnets or other containers.

Keyword	Form	Subcont ainers?	Default Value	Meaning
range	range IP_address range dotted_quad-dotted_quad	No	None	Modifies the range on the container in which the range statement is in. The range statement is not valid in the global or database container levels. If the range is the first in the container that does not specify a range on the container definition line, then the range for the container becomes the range specified by the range statement. Any range statement after the first range or all range statements for a containers that specifies ranges in its definition are added to the current range. With the range statement, a single address or set of addresses can be added to the range. The range must fit inside the subnet container definition.
client	client hwtype hwaddr NONE	Yes	None	Specifies a client container that denies the client specified by the <i>hwaddr</i> and <i>hwtype</i> from getting an address. If <i>hwtype</i> is 0, then <i>hwaddr</i> is an
	client hwtype hwaddr ANY client hwtype hwaddr dotted_quad			ASCII string. Otherwise, hwtype is the hardware type for the client and hwaddr is the hardware address of the client. If the hwaddr is a string, then quotes are accepted around the string. If the hwaddr is a hexstring, then the address may be specified by 0xhexdigits or hex digits. range allows the client
	client hwtype hwaddr range			specified by the hwaddr and hwtype to get an address in the range. Must be regular expressions to match multiple clients.

Keyword	Form	Subcont ainers?	Default Value	Meaning
class	class string class string range	Yes	None	Specifies a class container with name <i>string</i> . String can be quoted or not. If quoted, the quotes are removed before comparison. Quotes are required for strings with spaces or tabs. This container is valid at any level. A range can be supplied to indicate a set of addresses to hand out to a client with this class. The range is either a single dotted quad IP addresses separated by a dash.
network	network network id netmask	Yes	None	Specifies a network ID using class information (for example, 9.3.149.0 with a netmask of 255.255.255.0 would be network 9.0.0.0 255.255.255.0). This version of the network container is
	network network id network network id range			used to hold subnets with the same network ID and netmask. When a range is provided, all the addresses in the range are in the pool. The range must be in the network ID's network. This uses class full addressing. This is only valid in the global or database container level.
				Note: The network keyword is deprecated in favor of the subnet container.

Keyword	Form	Subcont ainers?	Default Value	Meaning
vendor	vendor vendor_id hex"" vendor vendor_id hex "" vendor vendor_id 0xdata vendor vendor_id "" vendor vendor_id range vendor vendor_id range hex "" vendor vendor_id range hex "" vendor vendor_id range in range vendor vendor_id range in range vendor vendor_id range	Yes	None	Specifies a vendor container. Vendor containers are used to return option 43 to the client. The vendor id may be specified in a quoted string or a binary string in the form $0xhexdigits$ or $hex"digits"$. An optional range may be placed after the vendor id. The range is specified as two dotted quads separated by a dash. After the optional range, an optional hexstring or ASCII string can be specified as the first part of the option 43. If options are in the container, they are appended to the option 43 data. After all options are processed an End Of Option List Option is appended to the data. To return options outside of an option 43, use a regular expression client that matches all clients to specify normal options to return based on the vendor ID.

Keyword	Form	Subcont ainers?	Default Value	Meaning
inoption	inoption number option_data	Yes	None	Specifies a container to be matched against any arbitrary incoming option specified by the client. <i>number</i> specifies the option number. <i>option_data</i> specifies the key to match for this container to be selected during address and option selectoin for the client. <i>option_data</i> is specified in expected form — quoted string, IP address, integer value — for well known options, or it can be optionally speicifed as a hexadecimal string of bytes if preceded by the characters 0x. For options
	inoption number option_data range			that are not well known to the server, a hexadecimal string of bytes can be specified in the same fashion. Additionally, the option_data can indicate a regular expression to be compared against the string representation of the client's option data. Regular expressions are specified in a quoted string beginning "! (double quote followed by an exclamation mark). The string form of options not well known to the server will be a hexadecimal string of bytes NOT preceded with the characters 0x.

Keyword	Form	Subcont ainers?	Default Value	Meaning
virtual	virtual fill <i>id id</i> virtual sfill <i>id id</i> virtual rotate <i>id id</i> virtual srotate <i>id id</i>	No	None	Specifies a virtual subnet with a policy. fill means use all addresses in the container before going to the next container. rotate means select an address from the next pool in the list on each request. sfill and srotate are the same as fill and rotate, but a search is done to see if the client matches containers, vendors, or classes in the subnet. If a match is found that can supply an address, the address is taken from that container instead of following the policy. There can be as many IDs as needed. id is either the subnet ID from the subnet definition or the label from the subnet definition. The label is required if there are multiple subnets with the same subnet id.
inorder:	inorder: id id	No	None	Specifies a virtual subnet with a policy of fill, which means use all addresses in the container before going to the next container. There can be as many IDs as needed. <i>id</i> is either the subnet ID from the subnet definition or the label from the subnet definition. The label is required if there are multiple subnets with the same subnet ID.
balance:	balance: id id	No	None	Specifies a virtual subnet with a policy of rotate, which means use the next address in the next container. There can be as many IDs as needed. <i>id</i> is either the subnet ID from the subnet definition or the label from the subnet definition. The label is required if there are multiple subnets with the same subnet ID.

Keyword	Form	Subcont ainers?	Default Value	Meaning
supportBootp	supportBootp true	No	Yes	Specifies whether the current
	supportBootp 1			container and all below it (until overridden) should support
	supportBootp yes]		BOOTP clients.
	supportBootp false	1		
	supportBootp 0			
	supportBootp no			
supportUnlisted clients	supportUnlistedclients BOTH	No	Both	Specifies whether the current container and all below it (until
	supportUnlistedclients DHCP			overridden) should support unlisted clients. The value
	supportUnlistedclients BOOTP			indicates whether all clients should be allowed access
	supportUnlistedclients NONE			without specific client statements, DHCP clients
	supportUnlistedclients true			only, BOOTP clients only, or no one.
	supportUnlistedclients yes			Note: The true and false
	supportUnlistedclients 1			values are supported for
	supportUnlistedclients false			compatability with previous versions and are deprecated. The true value coresponds to BOTH and the false value cooresponds to NONE.
	supportUnlistedclients no			
	supportUnlistedclients 0			
addressrecordd b	addressrecrddb <i>path</i>	No	None	If specified, it works like the backupfile keyword. Only valid in the global or database container level.
				Note: This method is deprecated.
backupfile	backupfile <i>path</i>	No	/etc/db_file .crbk	Specifies the file to use for database backups. Only valid in the global or database container level.
checkpointfile	checkpointfile <i>path</i>	No	/etc/db_file .chkpt	Specifies the database checkpoint files. The first checkpoint file is the <i>path</i> . The second checkpoint file is <i>path</i> with the last character replaced with a 2. So, the checkpoint file should not end in 2. Only valid in the global or database container level.
clientrecorddb	clientrecorddb <i>path</i>	No	/etc/db_file .cr	Specifies the database save file. The file contains all the client records the DHCP server has serviced. Only valid in the global or database container level.

Keyword	Form	Subcont ainers?	Default Value	Meaning
bootstrapserver	bootstrapserver IP address	No	None	Specifies the server clients should use from which to TFTP files after receiving BOOTP or DHCP packets. This value fills in the siaddr field in the packet. This is valid at any container level.
giaddrfield	giaddrfield <i>IP address</i>	No	None	Specifies the giaddrfield for response packets. Note: This specification is illegal in the BOOTP and DHCP protocols, but some clients require the giaddr field to be the default gateway for the network. Because of this potential conflict, giaddrfield should only be used within a client container, although it can work at any level.
pingTime	pingTime <i>n time_unit</i>	No	3 seconds	Specifies the amount of time to wait for a ping response before handing out an address. The default time unit is hundredths of a second. The time unit value is defined in the note preceding this table. This is valid at any container level. The time_unit parameter is optional.
bootptime	bootptime <i>n time_unit</i>	No	−1, infinite	Specifies the amount of time to lease an address to a BOOTP client. The default is –1, which means infinite. The normal time unit values are available. The time unit parameter is optional. This is valid at any container level.
AllRoutesBroad cast	allroutesbroadcast no allroutesbroadcast false allroutesbroadcast 0 allroutesbroadcast yes allroutesbroadcast true allroutesbroadcast 1	No	0	Specifies whether responses should be broadcast to all routes, if a broadcast response is required. This is valid at any container level. This is ignored by AIX DHCP servers, because the actual MAC address of the client, including RIFs, are stored for the return packet. This is valid at any container level.

Keyword	Form	Subcont ainers?	Default Value	Meaning
addressassigne d	addressassigned "string"	No	None	Specifies a quoted string to execute when an address is assigned to a client. The string should have two %s. The first %s is the client id in the form type—string. The second %s is an IP address in dotted quad format. This is valid at any container level.
addressrelease d	addressreleased "string"	No	None	Specifies a quoted string to execute when an address is released by a client. The string should have one %s. The %s is the IP address being released in dotted quad format. This is valid at any container level.
appenddomain	appenddomain 0	No	No	Specifies whether to append the defined option 15 domain name to the client–suggested hostname in the event that the client does not suggest a domain name as well. This is valid at any container level.
	appenddomain no			
	appenddomain false			
	appenddomain 1			
	appenddomain yes			
	appenddomain true			
canonical	canonical 0	No	0	Specifies that the client id is in
	canonical no			canonical format. This is valid only in the client container.
	canonical false			,
	canonical 1			
	canonical yes			
	canonical true	1		
leaseTimeDefau It	leaseTimeDefault n time_unit	No	86400 seconds	Specifies the default lease time for clients. This is valid at any container level. The time_unit parameter is optional.

Keyword	Form	Subcont ainers?	Default Value	Meaning
proxyarec	proxyarec never proxyarec usedhcpddns proxyarec usedhcpddnsplus proxyarec always proxyarec usedhcpddnsprotected proxyarec usedhcpddnsprotected proxyarec alwaysprotected proxyarec alwaysprotected proxyarec standard proxyarec protected	No	usedhcpdd nsplus	Specifies what options and methods should be used for A record updates in the DNS. never means never update the A record. usedhcpddns means use option 81 if the client specifies it. usedhcpddnsplus means use option 81 or option 12 and 15, if specified. always means do the A record update for all clients. XXXXprotected modifies the nsupdate command to make sure the client is allowed. standard is a synonym for always. protected is a synonym for alwaysprotected. This is valid at any container level.
releasednsA	releasednsA "string"	No	None	Specifies the execution string to use when an address is released. The string is used to remove the A record associated with the address released. This is valid at any container level.
releasednsP	releasednsP "string"	No	None	Specifies the execution string to use when an address is released. The string is used to remove the PTR record associated with the address released. This is valid at any container level.
removedns	removedns "string"	No	None	Specifies the execution string to use when an address is released. The string is used to remove the PTR and A record associated with the address released. This is valid at any container level. Note: This is deprecated in favor of the releasednsA and releasednsP keywords.

Keyword	Form	Subcont ainers?	Default Value	Meaning
updatedns	updatedns " <i>string</i> "	No	None	Specifies the execution string to use when an address is bound. The string is used to update both the A and the PTR record associated with the address. This is valid at any container level.
				Note: This is deprecated in favor of the updatednsA and updatednsP keywords.
updatednsA	updatednsA "string"	No	None	Specifies the execution string to use when an address is bound. The string is used to update the A record associated with the address. This is valid at any container level.
updatednsP	updatednsP "string"	No	None	Specifies the execution string to use when an address is bound. The string is used to update the PTR record associated with the address. This is valid at any container level.
hostnamepolicy	hostnamepolicy suggested hostnamepolicy resolved	No	default	Specifies which hostname to return to the client. Default policy is to prefer the defined hostname and domain name over suggested names. Other policies imply strict adherence
	hostnamepolicy always_resolved		return th none if r the conf policies	(for example: defined will return the defined name or none if no name is defined in the configuration). Also, policies using the always
	hostnamepolicy defined			modifier will dictate the server to return the hostname option regardless of whether the client requested it through the
	hostnamepolicy always_defined			parameter list option. Note that suggesting a hostname also implies requesting it, and hostnames can be suggested through option 81 or through
	hostnamepolicy default			options 12 and 15. This keyword is valid at any container level.

Keyword	Form	Subcont ainers?	Default Value	Meaning
bootfilepolicy	bootfilepolicy suggested bootfilepolicy merge bootfilepolicy defined bootfilepolicy always	No	suggested	Specifies a preference for returning the bootfile name to a client. suggested prefers the client—suggested bootfile name to any server—configured name. merge appends the client suggested name to the server—configured home directory. defined prefers the defined name over any suggested bootfile name. always returns the defined name regardless of whether the client requests the bootfile option through the parameter list option.
stealfromchildre n	stealfromchildren true stealfromchildren 1 stealfromchildren yes stealfromchildren false stealfromchildren 0 stealfromchildren no	No	No	Specifies whether the parent container should "steal" from children containers when the parent container runs out of addresses. This means that if you have a subnet with class defined with a range of addresses, those addresses are reserved for those clients that specify that class. If stealfromchildren is true, then addresses will be pulled from a child to try and satisfy the request. The default is to not steal an address.
homedirectory	homedirectory path	No	None	Specifies the home directory to use in the file section of the response packet. This can be specified at any container level. The bootfile policy defines how items specified in the file section of the incoming packet interact with the bootfile and the home directory statements.
bootfile	bootfile <i>path</i>	No	None	Specifies the bootfile to use in the file section of the response packet. This can be specified at any container level. The bootfile policy defines how items specified in the file section of the incoming packet interact with the bootfile and the home directory statements.

DHCP and Network Installation Management (NIM) Interactions and Suggestions

The concept of dynamically assigning Internet Protocol (IP) addresses is fairly new. The following suggestions are provided to help with DHCP and NIM interaction.

- 1. When configuring objects in the NIM environment, use host names whenever possible. This allows you to use a dynamic nameserver that updates the IP addresses when the host name is converted to an IP address in the NIM environment.
- 2. You should place the NIM master and the DHCP server on the same system. The DHCP server has an option in the update DNS string that, when set to NIM, attempts to keep the NIM objects out of those states that need static IP addresses when those addresses change.
- 3. For NIM clients, you should set the default lease time to twice the time it takes to install a client. This allows a leased IP address to be valid during the installation. After the installation, the client will reboot, and DHCP will be started or will need to be configured, depending on the type of installation.
- 4. The dhcpsd server should be responsible for both the PTR and the A DNS records. When NIM reinstalls the machine, the file containing the RSA is deleted, and the client cannot update its records. Therefore, the server should update the system records. To do this, change the updatedns line in /etc/dhcpcd.ini to:

```
updatedns "/usr/sbin/dhcpaction '%s' '%s' '%s' NONE NONIM" In the /etc/dhcpsd.cnf file, change the updatedns line to:

updatedns "/usr/sbin/dhcpaction '%s' '%s' '%s' BOTH NIM"
```

Note: When a NIM object is placed into the BOS installation pending state, the <code>dhcpsd</code> server may pass arguments that are different from those originally intended. To avoid this, try to minimize the time the client is in this pending state.

Following these suggestions should allow the NIM environment to work with dynamic clients.

For more information on Network Installation Management, see AIX Network Installation Management Guide and Reference.

Configuring TCP/IP

If you installed the Transmission Control Protocol/Internet Protocol (TCP/IP) and Network File System (NFS) software, you can configure your system to communicate over a network.

After you install TCP/IP and NFS software, use the Web-based System Manager fast path, **wsm networks**, or the System Management Interface Tool (SMIT) fast path, **smit tcpip**, to configure your system. Online help will guide you through the process.

Prerequisites

The TCP/IP software must be installed. If you need to install TCP/IP software, you will have to install the TCP/IP Optional Support optional software product.

You must have root authority to configure TCP/IP.

Updating the Hosts List

A nameserver is a machine on your network that stores the names and addresses of all the network machines. The names are stored in a Hosts List. When one machine wants to communicate with another, it sends that machine name to the nameserver. The nameserver refers to the Hosts List and responds with the address of the machine name requested. Having a nameserver is an advantage because the Hosts List is stored and updated at one location, but is accessible to all machines on the network. This saves time and storage space.

- If you are using a nameserver for network communications, you do not need to perform this procedure. You have finished configuring TCP/IP.
- If you are not using a nameserver for network communications, you must update the
 hosts list to include the names of the systems on the network. With root authority, use the
 Web-based System Manager fast path, wsm network, or the System Management
 Interface Tool (SMIT) fast path, smit hostent.

TCP/IP Daemons

Daemons (also known as *servers*) are processes that run continuously in the background and perform functions required by other processes. Transmission Control Protocol/Internet Protocol (TCP/IP) provides daemons for implementing certain functions in the operating system. These daemons are background processes and thus run without interrupting other processes (unless that is part of the daemon's function).

Daemons are invoked by commands at the system management level, by other daemons, or by shell scripts. You can also control daemons with the **inetd** daemon, the **rc.tcpip** shell script, and the System Resource Controller (SRC).

Subsystems and Subservers

A *subsystem* is a daemon, or server, that is controlled by the SRC. A *subserver* is a daemon that is controlled by a subsystem. (Daemon commands and daemon names are usually denoted by a **d** at the end of the name.) The categories of subsystem and subserver are mutually exclusive. That is, daemons are not listed as both a subsystem and as a subserver. The only TCP/IP subsystem that controls other daemons is the **inetd** daemon. Thus, all TCP/IP subservers are also **inetd** subservers.

TCP/IP daemons controlled by the SRC are the following:

Subsystems

timed

gated	Provides gateway routing functions and supports the Routing Information Protocol (RIP), the Routing Information Protocol Next Generation (RIPng), Exterior Gateway Protocol (EGP), the Border Gateway Protocol (BGP) and BGP4+, the Defense Communications Network Local–Network Protocol (HELLO), Open Shortest Path First (OSPF), Intermediate System to Intermediate System (IS–IS), and Internet Control Message Protocol (ICMP and ICMPv6)/Router Discovery routing protocols. In addition, the gated daemon supports the Simple Network Management Protocol (SNMP). The gated daemon is one of two routing daemons available for routing to network addresses and is the preferred routing daemon. The gated daemon is preferred over the routed daemon because the gated daemon supports more gateway protocols.
inetd	Invokes and schedules other daemons when requests for the daemons' services are received. This daemon can also start other daemons. The inetd daemon is also known as the super daemon.
iptrace	Provides interface—level packet—tracing function for Internet protocols.
named	Provides the naming function for the Domain Name Server Protocol (DOMAIN).
routed	Manages the network routing tables and supports the Routing Information Protocol (RIP). The gated daemon is preferred over the routed daemon because the gated daemon supports more gateway protocols.
rwhod	Sends broadcasts to all other hosts every three minutes and stores information about logged—in users and network status. Use

amounts of a machine's resources.

Provides the timeserver function.

the **rwhod** daemon with extreme care, as it can steal significant

Note: Both the **routed** and **gated** daemons are listed as TCP/IP subsystems. Do not run the **startsrc –g tcpip** command, which initiates both of these routing daemons, along with all the other TCP/IP subsystems. Running both daemons simultaneously on one machine can produce unpredictable results.

TCP/IP daemons controlled by the **inetd** subsystem are the following:

inetd Subservers

comsat Notifies users of incoming mail.

fingerd Provides a status report on all logged–in users and network status

at the specified remote host. This daemon uses the Finger

protocol.

ftpd Provides the file transfer function for a client process using the

File Transfer Protocol (FTP).

rexecd Provides the foreign host server function for the rexec command.

rlogind Provides the remote login facility function for the rlogin command.

rshd Provides the remote command execution server function for the

rcp and rsh commands.

talkd Provides the conversation function for the **talk** command.

syslogd Reads and logs system messages. This daemon is in the RAS

group of subsystems.

telnetd Provides the server function for the TELNET protocol.

tftpd Provides the server function for the Trivial File Transfer Protocol

(TFTP).

uucpd Handles communications between the Basic Network Utilities

(BNU) and TCP/IP.

System Resource Control (SRC)

Among other functions, SRC allows you to start daemons, stop them, and trace their activity. In addition, SRC provides the ability to group daemons into subsystems and subservers.

System Resource Control is a tool designed to aid the person who manages your system in controlling daemons. SRC allows control beyond the flags and parameters available with each daemon command.

Refer to the System Resource Controller Overview in *AIX 4.3 System Management Concepts: Operating System and Devices* for more information concerning the System Resource Controller.

SRC Commands

SRC commands can affect one daemon, a group of daemons, or a daemon and those daemons it controls (subsystem with subservers). In addition, some TCP/IP daemons do not respond to all SRC commands. The following is a list of SRC commands that can be used to control TCP/IP daemons and their exceptions.

startsrc Starts all TCP/IP subsystems and **inetd** subservers. The **startsrc**

command works for all TCP/IP subsystems and inetd subservers.

stopsrc Stops all TCP/IP subsystems and **inetd** subservers. This

command is also called the **stop normal**. The **stop normal** command allows subsystems to process all outstanding work and

terminate gracefully. For inetd subservers, all pending

connections are allowed to start and all existing connections are allowed to complete. The **stop normal** command works for all

TCP/IP subsystems and **inetd** subservers.

stopsrc –f Stops all TCP/IP subsystems and **inetd** subservers. This

command is also called the **stop force**. The **stop force** command immediately terminates all subsystems. For **inetd** subservers, all pending connections and existing connections are terminated

immediately.

refresh Refreshes the following subsystems and subservers: the **inetd**,

syslogd, named, dhcpsd, and gated subsystems.

Issrc Provides short status for subsystems, which is the state of the

specified subsystem (active or inoperative). Also provides short status for **inetd** subservers. The short status for **inetd** subservers includes: subserver name, state, subserver description, command

name, and the arguments with which it was invoked.

Issrc –I Provides the short status plus additional information (long status)

for the following subsystems:

gated State of debug or trace, routing protocols

activated, routing tables, signals accepted and

their function.

inetd State of debug, list of active subservers and their

short status; signals accepted and their function.

named State of debug, **named.conf** file information.

dhcpsd State of debug, all controlled IP addresses and

their current state.

routed State of debug and trace, state of supplying

routing information, routing tables.

syslogd syslogd configuration information.

The **Issrc** –**I** command also provides long status for **inetd** subservers. The long status includes short status information and active connection information. Some subservers will provide additional information. The additional information by subserver includes:

ftpd State of debug and logging telnetd Type of terminal emulating

rlogind State of debug

fingerd State of debug and logging

The **rwhod** and **timed** subservers do not provide long status.

traceson Turns on socket–level debugging. Use the trpt command to

format the output. The timed and iptraced subsystems do not

support the **traceson** command.

tracesoff Turns off socket–level debugging. Use the trpt command to

format the output. The timed and iptraced subsystems do not

support the **tracesoff** command.

For examples of how to use these commands, see the articles on the individual commands. For more information on the System Resource Controller, refer to System Resource Controller Overview in AIX 4.3 System Management Concepts: Operating System and Devices.

Configuring the inetd Daemon

To configure the **inetd** daemon:

- 1. Specify which subservers it will invoke by adding an **inetd** subserver.
- 2. Specify what its restart characteristics are by changing restart characteristics of the **inetd** daemon.

Configuring the inetd Daemor	n Tasks					
Web-based System Manager: (wsm network application)	Web-based System Manager: wsm network fast path (wsm network application)					
-OR-						
Task	SMIT Fast Path	Command or File				
Starting the inetd Daemon	smit mkinetd	startsrc –s inetd				
Changing Restart Characteristics of the inetd Daemon	smit chinetd or smit Isinetd					
Stopping the inetd Daemon	smit rminetd	stopsrc –s inetd				
Listing All inetd Subservers	smit inetdconf					
Adding an inetd Subserver ¹	smit mkinetdconf	edit /etc/inetd.conf then run refresh –s inetd or kill –1 inetdPID ²				
Change/Show Characteristics of an inetd Subserver	smit inetdconf	edit /etc/inetd.conf then run refresh –s inetd or kill –1 inetdPID ²				
Removing an inetd Subserver	smit rminetd	edit /etc/inetd.conf then run refresh –s inetd or kill –1 inetdPID ²				

Note:

- 1. Adding an **inetd** subserver configures the **inetd** daemon so that it will invoke the subserver when it is needed.
- 2. Both **refresh** and **kill** commands inform the **inetd** daemon of changes to its configuration file.

Client Network Services

Client Network Services (accessible using the Web-based System Manager fast path, wsm network, or the SMIT fast path, smit clientnet) refers to the TCP/IP protocols available for use by an AIX system. Each protocol (or service) is known by the port number it uses on the network, hence the term well-known port. As a convenience to programmers, the port numbers can be referred to by names as well as numbers. For example, the TCP/IP mail protocol uses port 25 and is known by the name smtp. If a protocol is listed (uncommented) in the /etc/services file, then a host can use that protocol.

By default, all the TCP/IP protocols are defined in the /etc/services file. So, you do not have to do any configuration of this file. However, suppose you wrote your own client/server programs. You may want to add your service to the /etc/services file, and reserve a specific port number and name for your service. If you do decide to add your service to /etc/services, note that port numbers 0 through 1024 are reserved for system use.

Client Network Services Tasks	Client Network Services Tasks				
Web-based System Manager: (wsm network application) -OR-					
Task	SMIT Fast Path	Command or File			
Listing All Services	smit Isservices	view /etc/services			
Adding a Service	smit mkservices	edit /etc/services			
Change/Show Characteristics of a Service	smit chservices	edit /etc/services			
Removing a Service	smit rmservices	edit /etc/services			

Server Network Services

Server Network Services includes controlling remote access, starting or stopping TCP/IP, and managing the **pty** device driver, as shown in the following table.

The **pty** device driver is installed automatically with the system. By default, it is configured to support 16 BSD–style symbolic links, and it is available for use by the system at boot time.

Server Network Services Task	Server Network Services Tasks					
Web-based System Manager: (wsm network application)	Web-based System Manager: wsm network fast path (wsm network application)					
-OR-						
Task	SMIT Fast Path	Command or File				
Controlling Remote Access	See "Remote Command Execution Access", on page 3-148 and "Restricted File Transfer Program Users", on page 3-148.					
Start, Restart, or Stop TCP/IP Subsystems	smit otherserv	See "System Resource Control", on page 3-92.				
Change/Show Characteristics of the pty Device Driver	smit chgpty	chdev –I pty0 –P –a num=X where X ranges from 0 to 64				
Make the pty Device Driver Unavailable for Use	smit pty then select Remove the PTY; Keep Definition					
Make the pty Device Driver Available for Use	smit pty then select Configure the Defined PTY					
Generate an Error Report	smit errpt					
Trace the pty	smit trace					

TCP/IP Name Resolution

Although 32—bit Internet addresses provide machines an efficient means of identifying the source and destination of datagrams sent across an internetwork, users prefer meaningful, easily remembered names. Transmission Control Protocol/Internet Protocol (TCP/IP) provides a naming system that supports both flat and hierarchical network organizations.

The topics discussed in this section are:

- Naming, on page 3-97
- Performing Local Name Resolution (/etc/hosts), on page 3-103
- Planning for DOMAIN Name Resolution, on page 3-104
- Configuring Name Servers, on page 3-105
- Configuring a Forwarder, on page 3-115
- Configuring a Forward Only Server, on page 3-116
- Configuring a Host to Use a Name Server, on page 3-118
- Configuring Dynamic Zones on the DNS Name Server, on page 3-119

Naming

Naming in flat networks is very simple: Host names consist of a single set of characters and generally are administered locally. In flat TCP/IP networks, each machine on the network has a file (/etc/hosts) containing the name—to—Internet—address mapping information for every host on the network. As a TCP/IP network grows, the administrative burden of keeping each machine's naming file current grows. When TCP/IP networks become very large, as on the Internet, naming is divided hierarchically. Typically, the divisions follow the network's organization. In TCP/IP, hierarchical naming is known as the *domain name system* (DNS) and uses the DOMAIN protocol. The DOMAIN protocol is implemented by the named daemon in TCP/IP.

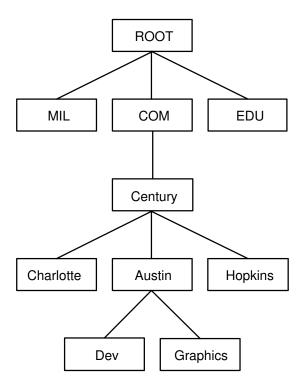
As in naming for flat networks, the domain name hierarchy provides for the assignment of symbolic names to networks and hosts that are meaningful and easy for users to remember. However, instead of each machine on the network keeping a file containing the name—to—address mapping for all other hosts on the network, one or more hosts are selected to function as *name servers*. Name servers translate (resolve) symbolic names assigned to networks and hosts into the efficient Internet addresses used by machines. A name server has complete information about some part of the domain, referred to as a *zone*, and it has *authority* for its zone.

Naming Authority

In a flat network, all hosts in the network are administered by one central authority. This form of network requires that all hosts in the network have unique host names. In a large network, this requirement creates a large administrative burden on the central authority.

In a domain network, groups of hosts are administered separately within a tree—structured hierarchy of domains and subdomains. In this case, host names need to be unique only within the local domain, and only the *root domain* is administered by a central authority. This structure allows subdomains to be administered locally and reduces the burden on the central authority. For example, the root domain of the Internet consists of such domains as com (commercial organizations), edu (educational organizations), gov (governmental organizations), and mil (military groups). New top—level domains can only be added by the central authority. Naming at the second level is delegated to designated agents within the respective domains. For example, in the following figure, com has naming authority for all commercial organization subdomains beneath it. Likewise, naming at the third level (and so on) is delegated to agents within that level. For example, in the Domain Structure of the

Internet figure, Century has naming authority for its subdomains Austin, Hopkins, and Charlotte.



Domain Structure of the Internet

Century's Austin subdomain might also be divided into zones, for example, Dev and Graphics. In this case, the zone <code>austin.century.com</code> would have all the data contained in the domain <code>austin.century.com</code>, except that which was delegated to Dev and Graphics. The zone <code>dev.century.com</code> would contain only the data delegated to Dev; it would know nothing about Graphics, for example. The zone <code>austin.century.com</code> (as opposed to the domain of the same name) would contain only that data not delegated to other zones.

Naming Conventions

In the hierarchical domain name system, names consist of a sequence of case—insensitive subnames separated by periods with no embedded blanks. The DOMAIN protocol specifies that a local domain name must be less than 64 characters and that a host name must be less than 32 characters in length. The host name is given first, followed by a . (period), a series of local domain names separated by periods, and finally the root domain. A fully specified domain name for a host, including periods, must be less than 255 characters in length and in the following form:

```
host.subdomain1.[subdomain2 . . . subdomain].rootdomain
```

Since host names must be unique within a domain, you can use an abbreviated name when sending messages to a host within the same domain. For example, instead of sending a message to smith.eng.lsu.edu, a host in the eng domain could send a message to smith. Additionally, each host can have several aliases that other hosts can use when sending messages.

Choosing Names for the Hosts on Your Network

The purpose of using names for hosts is to provide a quick, easy, and unambiguous way to refer to the computers in your network. Internet system administrators have discovered that there are good, as well as poor, choices for host names. These suggestions are intended to help you avoid common pitfalls in choosing host names.

Here are some suggestions for choosing unambiguous, easy to remember host names:

- Terms that are rarely used, for example, sphinx or eclipse.
- Theme names, such as colors, elements (for example, helium, argon, or zinc), flowers, fish, and others.
- Real words (as opposed to random strings of characters).

Here are some examples of poor choices. In general, these are poor choices because they are difficult to remember or are confusing (either to humans or computers):

- Terms that are already in common use, for example, up, down, or crash.
- · Names containing only numbers.
- Names that contain punctuation marks.
- Names that rely on case distinction, for example, Orange and orange.
- The name or initials of the primary user of the system.
- Names having more than 8 characters.
- Unusual or purposefully incorrect spellings, for example, czek, which could be confused with "check" or "czech."
- Names that are, or resemble, domain names, for example, yale.edu.

Name Servers

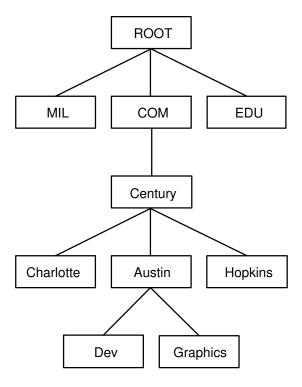
In a flat name space, all names must be kept in the /etc/hosts file on each host on the network. If the network is very large, this can become a burden on the resources of each machine.

In a hierarchical network, certain hosts designated as *name servers* resolve names into Internet addresses for other hosts. This has two advantages over the flat name space: It keeps the resources of each host on the network from being tied up in resolving names, and it keeps the person who manages the system from having to maintain name resolution files on each machine on the network. The set of names managed by a single name server is known as its *zone of authority*.

Note: Although the host machine that performs the name resolution function for a zone of authority is commonly referred to as a *name server* host, the process controlling the function, the **named** daemon, is the actual name server process.

To further reduce unnecessary network activity, all name servers *cache* (store for a period of time) name—to—address mappings. When a client asks a server to resolve a name, the server checks its cache first to see if the name has been resolved recently. Since domain and host names do change, each item remains in the cache for a limited length of time specified by the record's time to live. In this way, authorities can specify how long they expect the name resolution to be accurate.

Within any autonomous system there can be multiple name servers. Typically, name servers are organized hierarchically and correspond to the network's organization. Referring to the figure Domain Structure of the Internet, each domain might have a name server responsible for all subdomains within the domain. Each subdomain name server communicates with the name server of the domain above it (called the *parent* name server), as well as with the name servers of other subdomains.



Domain Structure of the Internet

For example, in the figure Domain Structure of the Internet, Austin, Hopkins, and Charlotte are all subdomains of the domain Century. If the tree hierarchy is followed in the network design, the Austin name server communicates with the name servers of Charlotte and Hopkins as well as with the parent Century name server. The Austin name server will also communicate with the name servers responsible for its subdomains.

There are several types of name servers:

Master Name ServerLoads its data from a file or disk and may delegate authority to other servers in its domain.

Slave Name Server Receives its information at system startup time for the given

zone of authority from a master name server, and then periodically asks the master server to update its information. On expiration of the refresh value in the start of authority (SOA) Resource Record on a slave name server, or upon receipt of a Notify message from the master name server, the slave will reload the database from the master if the serial number of the database on the master is greater than the serial number in the current database on the slave. If it becomes necessary to force a new zone transfer from

refresh the **named** daemon on the slave name server.

Although its method of database replication is similar to that of the slave name server, the stub name server only replicates the Name Server records of the master database

the master, simply remove the existing slave databases and

rather than the whole database.

Hint Server Indicates a name server that relies only on the hints that it

has built from previous queries to other name servers. The hint name server responds to queries by asking other servers that have the authority to provide the information

needed if a hint name server does not have a name—to—address mapping in its cache.

Stub Name Server

Forwarder or Client Server Forwards gueries it cannot satisfy locally to a fixed list of forwarding servers. Forwarding-only servers (a forwarder that obtains information and passes it on to other clients, but that is not actually a server) will not interact with the master name servers for the root domain and other domains. The gueries to the forwarding servers are recursive. There may be one or more forwarding servers, which are tried in turn until the list is exhausted. A client and forwarder configuration is typically used when you do not wish all the servers at a given site to interact with the rest of the Internet servers, or when you may wish to build a large cache on a select number of name servers.

Remote Server

Runs all the network programs that use the name server without the name server process running on the local host. All queries are serviced by a name server that is running on another machine on the network.

One name server host can perform in different capacities for different zones of authority. For example, a single name server host can be a master name server for one zone and a slave name server for another zone.

Name Resolution

The process of obtaining an Internet address from a host name is known as name resolution and is done by the **gethostbyname** subroutine. The process of translating an Internet address into a host name is known as reverse name resolution and is done by the gethostbyaddr subroutine. These routines are essentially accessors into a library of name translation routines known as resolvers.

Resolver routines on hosts running TCP/IP normally attempt to resolve names using the following sources:

- 1. BIND/DNS (named)
- 2. Network Information Service (NIS)
- 3. Local /etc/hosts file

When NIS+ is installed, lookup preferences are set using the irs.conf file. For more information, see AIX 4.3 NIS/NIS+ Guide.

To resolve a name in a domain network, the resolver routine first queries the domain name server database, which may be local if the host is a domain name server or may be on a foreign host. Name servers translate domain names into Internet addresses. The group of names for which a name server is responsible is its zone of authority. If the resolver routine is using a remote name server, the routine uses the Domain Name Protocol (DOMAIN) to query for the mapping. To resolve a name in a flat network, the resolver routine checks for an entry in the local /etc/hosts file. When NIS or NIS+ is used, the /etc/hosts file on the master server is checked.

By default, resolver routines attempt to resolve names using the above resources. BIND/DNS will be tried first, if the /etc/resolv.conf file does not exist or if BIND/DNS could not find the name, NIS is queried if it is running. NIS is authoritative over the local /etc/hosts, so the search will end here if it is running. If NIS is not running, then the local /etc/hosts file is searched. If none of these services could find the name, then the resolver routines return with HOST NOT FOUND. If all of the services are unavailable, then the resolver routines return with SERVICE UNAVAILABLE.

The default order described above can be overwritten by creating the /etc/irs.conf configuration file and specifying the desired order. Also, both the default and /etc/irs.conf orderings can be overwritten with the environment variable, NSORDER. If either the /etc/irs.conf file or environment variable, NSORDER, are defined, then at least one value must be specified along with the option.

To specify host ordering with the /etc/irs.conf file:

```
hosts value [ continue ]
```

The order is specified with each method indicated on a line by itself. The value is one of the listed methods and the continue keyword indicates that another resolver method follows on the next line.

To specify host ordering with the **NSORDER** environment variable:

```
NSORDER=value, value, value
```

The order is specified on one line with values separated by commas. White spaces are permitted between the commas and the equal sign.

For example, if the local network is organized as a flat network, then only the /etc/hosts file is needed.

In this example, the /etc/irs.conf file would contain the following line:

```
hosts local
```

Alternatively, the **NSORDER** environment variable could be set as:

```
NSORDER=local
```

If the local network is a domain network using a name server for name resolution and an /etc/hosts file for backup, then both services should be specified.

The /etc/irs.conf file would contain the following lines:

```
hosts dns continue
hosts local
```

The **NSORDER** environment variable would be set as:

```
NSORDER=bind, local
```

Note: The values listed must be in lowercase.

When following any defined or default resolver ordering, the search algorithm will continue from one resolver to the next only if the:

- current service is not running, therefore, it is unavailable.
- current service could not find the name and is not authoritative.

If the /etc/resolv.conf file does not exist, then BIND/DNS is considered to be not set up or running and therefore not available. If the **getdomainname** and **yp bind** subroutines fail, then it is assumed that the NIS service is not set up or running and therefore not available. If the /etc/hosts file could not be opened, then a local search is impossible and therefore the file and service are unavailable.

When a service is listed as authoritative, it means that this service is the expert of its successors and should have all pertinent names and addresses. Resolver routines do not try successor services since they may contain only a subset of the information in the authoritative service. Name resolution ends at service listed as authoritative, even if it does not find the name (in which case, the resolver routine returns HOST_NOT_FOUND). If an authoritative service is not available, then the next service specified will be queried.

An authoritative source is specified with the string "=auth" directly behind a value. The entire word, "authoritative" can be typed in, but only the "auth" will be used. For example, if the **NSORDER** environment variable contains the following:

```
hosts = nis=auth,dns,local
```

The search is ended after the NIS query, if NIS is running, regardless of whether the name was found. If NIS is not running, then the next source is queried, which is DNS.

TCP/IP name servers use caching to reduce the cost of searching for names of hosts on remote networks. Instead of searching for a host name each time a request is made, a name server first looks at its cache to see if the host name has been resolved recently.

Since domain and host names do change, each item remains in the cache for a limited length of time specified by the record's time-to-live (TTL). In this way, name servers can specify how long they expect their responses to be considered authoritative.

Potential Host Name Conflict Between nameserver and sendmail

In a DNS environment, a host name that is set using the **hostname** command from the command line or in the rc.net file format must be the official name of the host as returned by the name server. Generally, this name is the full domain name of the host in the form:

host.subdomain.subdomain.rootdomain

Note: Resolver routines require the default domain to be set. If the default domain is not set in the hostname command, then it must be set in the /etc/resolv.conf file.

If the host name is not set up as a fully qualified domain name, and if the system is set up to use a domain name server in conjunction with the sendmail program, the sendmail configuration file (/etc/sendmail.cf) must be edited to reflect this official host name. In addition, the domain name macros in this configuration file must be set for the sendmail program to operate correctly.

Note: The domain specified in the /etc/sendmail.cf file takes precedence over the domain set by the hostname command for all sendmail functions.

Potential Domain Name Conflict Between nameserver and sendmail

For a host that is in a DOMAIN network but is not a name server, the local domain name and domain name server are specified in the /etc/resolv.conf file. In a DOMAIN name server host, the local domain and other name servers are defined in files read by the named daemon when it starts.

Reverse Address Resolution Protocol

The Reverse Address Resolution Protocol (RARP) translates unique hardware addresses into Internet addresses on the Ethernet local area network (LAN) adapter (Ethernet protocol only). Standard Ethernet protocol is supported with the following restrictions. The server will:

- Only reply to RARP requests
- Only use permanent ARP table entries
- · Not use dynamic ARP table entries
- Not automatically reply for itself

The system administrator must manually build and maintain a table of permanent ARP entries using the arp command. A specific ARP table entry must be added on the AIX server for each host that requires RARP replies from an authoritative source.

Performing Local Name Resolution (/etc/hosts)

You should configure the /etc/hosts file if your network is small, and you are using a flat naming scheme. Even if you are using a hierarchical (or domain) naming scheme with name servers, you may want to configure the /etc/hosts file to identify hosts that are not known by the name servers.

You can configure your system for local host resolution using Web-based System Manager. System Management Interface Tool (SMIT), or commands. If you choose the command method, be sure to preserve the format of the /etc/hosts file, as described in "Hosts File Format for TCP/IP" in the AIX Files Reference.

Local Name Resolution Tasks				
Web-based System Manager: wsm network fast path (wsm network application) -OR-				
Task	SMIT Fast Path	Command or File		
List All the Hosts	smit Ishostent	view /etc/hosts		
Add a Host	smit mkhostent	edit /etc/hosts		
Change/Show Characteristics of a Host	smit chhostent	edit /etc/hosts		
Remove a Host	smit rmhostent	edit /etc/hosts		

Planning for DOMAIN Name Resolution

If you are part of a larger internetwork, you need to coordinate setting up your domain and name servers with their central authority.

Some hints in planning your own DOMAIN name resolution system:

- Obtain a copy of the book DNS and BIND in a Nutshell. Configuring domain name resolution is one of the more esoteric tasks required of a TCP/IP network administrator. Because of the vast possibilities in domain name architecture and name server configurations, you will find this book invaluable. TCP/IP Network Administration is another book with useful information about domain name resolution. You may also want to obtain a copy of the book Managing NFS and NIS in order to understand the interaction between DNS and NIS. For more information about NIS+, see AIX 4.3 NIS/NIS+ Guide.
- · Plan ahead.

Changing a name is *much* more difficult than setting up the initial one. Obtain consensus from your organization on network, gateway, name server, and host names before you set up your files.

Set up redundant name servers.

If you cannot set up redundant name servers, be sure to set up slave and hint name servers so you have some type of backup.

- In selecting the name servers, keep the following in mind:
 - Choose machines that are physically closest to exterior systems.
 - The name servers should be as independent as possible. Try for different power supplies and independent cabling.
 - Find another network to back up your name resolution service, and do the same for other networks.
- Test the servers.
 - Test both regular and reverse name resolution.
 - Test zone transfer from master to slave name servers.
 - Test each name server after a system crash and reboot.
- Send name resolution requests to forwarder servers before they go to exterior name servers. This allows your name servers to share caches and improve performance by reducing the load on your master name servers.

Configuring Name Servers

In a hierarchical network, certain hosts are designated as *name servers*. These hosts resolve names into Internet addresses for other hosts. The **named** daemon controls the name server function and, therefore, must be run on a name server host.

Before you configure a name server, decide which type or types best fit the network it will serve. There are several types of name servers.

A *master name server* actually stores the database containing name—to—address mapping information. It loads its data from a file or disk and may delegate authority to other servers in its domain. A *slave name server* or *stub name server* receives its information at system startup time for a given zone of authority from a master name server, and then periodically asks the master server to update its information. A *hint name server* responds to requests to resolve names by querying other servers that have the authority to provide the information needed.

Note: Previous generations of the **named** name server specified the master name server as the primary name server, the slave name server as the secondary name server, and the hint name server as the caching—only name server. Any reference to the **named.conf** file in this documentation is specific to AIX version 4.3.2 and above.

Keep in mind that a name server may function in different capacities for different zones of authority. For example, one name server host can be a master name server for one zone and a slave name server for another zone. If your system has NIS or NIS+ installed, these services can also provide name resolution. For more information, see *AIX 4.3 NIS/NIS+ Guide*.

There are several files involved in configuring name servers.

conf

This file is read when the **named** daemon starts. The records in the conf file tell the **named** daemon which type of server it is, which domains it has authority over (its zones of authority), and where to get the data for initially setting up its database. The default name of this file is /etc/named.conf. However, you can change the name of this file by specifying the name and path of the file on the command line when the **named** daemon is started. When using the /etc/named.conf as the conf file, if it does not exist, a message is generated in **syslog** and **named** terminates. However, if an alternative conf file is specified, and the alternative file does not exist, no error message will be generated and **named** will continue.

cache

Contains information about the local cache. The local cache file contains the names and addresses of the highest authority name servers in the network. The cache file uses the Standard Resource Record Format. The name of the cache file is set in the conf file.

domain data

There are three typical domain data files, also referred to as the **named** data files. The **named** *local* file contains the address resolution information for local loopback. The **named** data file contains the address resolution data for all machines in the name server's zone of authority. The named reverse data file contains the reverse address resolution information for all machines in the name server's zone of authority. The domain data files use the Standard Resource Record Format. Their file names are user definable and are set in the conf file. By convention, the names of these files generally include the name of the daemon (named), and in the extension the type of file and name of the domain. For example, the name server for the domain abc might have the following files:

```
named.abc.data
named.abc.rev
named.abc.local
```

When modifying the **named** data files the serial number in the SOA Resource Record must be incremented for slave name servers to properly realize the new zone changes.

resolv.conf

The presence of this file indicates to a host that it should go to a name server to resolve a name first. If the resolv.conf file does not exist, the host looks in the /etc/hosts file for name resolution. On a name server, the **resolv.conf** file must exist and can contain the local host's address, the loopback address (127.0.0.1), or be empty.

Note: The resolver routines require the default domain be set. If the default domain is not set in the /etc/resolv.conf file, then it must be set in the hostname.

Time-to-live (TTL) is specified in resource records. If TTL is not specified in a record, the length of this time period defaults to the minimum field as defined in the start of authority (SOA) record for that zone. TTL is used when data is stored outside a zone (in a cache) to ensure that the data does not stay around indefinitely.

Configuring a Master Name Server

To configure a master name server, use the Web-based System Manager fast path wsm network or use the following procedure, which edits a series of files and then uses the System Management Interface Tool (SMIT) or the command line to start the **named** daemon.

1. Edit the /etc/named.conf file. If there is no named.conf file in the /etc directory, copy the /usr/samples/tcpip/named.conf sample file into the /etc directory and edit it. See "named.conf File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a configuration file.

This file is read each time the **named** daemon starts. It tells the server which type of server it is, the zone or zones for which it is responsible, and where to get its initial information.

a. Using the options configuration clause, specify the directory in which the named data files can be located (optional). Use this if you want the **named** data files to use paths relative to this directory. For example:

```
options {
    directory "/usr/local/domain";
```

b. Optionally specify the name of the hint zone file. This allows for the caching of record data outside of the defined zones. For example:

```
zone "." IN {
    type hint;
    file "/etc/named.ca";
};
```

c. Specify a number of zones. To configure a server as the master for a zone, specify each zone and its domain data file. For example, a master server for both forward and reverse zones might resemble:

```
zone "abc.aus.century.com" in {
    type master;
    file "/etc/named.abcdata";
};
zone "201.9.192.in-addr.arpa" in {
    type master;
    file "/etc/named.abcrev";
};
```

d. Define the name of the **named** local file. For example:

```
zone "0.0.127.in-addr.arpa" in {
    type master;
    file "/etc/named.local";
};
```

2. Edit the /etc/named.ca file. Refer to the "DOMAIN Cache File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a cache file.

This file contains the addresses of the servers that are authoritative, or root name servers for the domain. For example:

```
; root name servers.
        1
                   ΙN
                                relay.century.com.
relay.century.com.
                    3600000
                                            129.114.1.2
```

Note: All lines in this file must be in Standard Resource Record Format.

- 3. Edit the /etc/named.local file. See the "DOMAIN Local Data File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a local data file.
 - a. Specify the start of authority of the zone and the default time-to-live information. For example:

```
@ IN SOA venus.abc.aus.cntry.com. gail.zeus.abc.aus.cntry.com.
(
                           1.1
                                  ;serial
                           3600
                                    ;refresh
                            600
                                    ; retry
                            3600000 ; expire
                           86400)
                                   ; minimum
```

b. Specify the name server (NS) record. For example:

```
NS
       venus.abc.aus.century.com.
```

c. Specify the pointer (PTR) record.

```
1
                      localhost.
       TN
              PTR
```

Note: All lines in this file must be in Standard Resource Record Format.

4. Edit the /etc/named.data file. The /usr/samples/tcpip/hosts.awk file contains directions for creating the /etc/named.data file. Use the /usr/samples/tcpip/named.data sample file as an example when creating the /etc/named.data file. See the "DOMAIN Data File Format for TCP/IP" in the AIX Files Reference for more information and an example of a hosts data file.

a. Specify the start of authority of the zone and the default time-to-live information for the zone. This record designates the start of a zone. There should only be one start of authority record per zone. For example:

```
bob.robert.abc.aus.century.com.
@
   ΙN
         SOA
              venus
(
                                 ;serial
                           1.1
                           3600 ; refresh
                           600
                                   ;retry
                           3600000 ; expire
                           86400) ; minimum
```

b. Include name-to-address resolution information on all hosts in the name server's zone of authority. For example:

```
venus
             ΙN
                             192.9.201.1
                             128.114.100.1
venus
             ΤN
```

c. Include name server records for all master name servers in the zone. For example:

```
ΙN
               venus.abc.century.com
ΙN
      NS
               kronos.xyz.century.com
```

d. Include other types of entries, such as canonical name records and mail exchanger records as needed.

Note: All lines in this file must be in Standard Resource Record Format.

- 5. Edit the /etc/named.rev file. The /usr/samples/tcpip/addrs.awk file contains directions for creating the /etc/named.rev file. See the "DOMAIN Reverse Data File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a reverse hosts data file.
 - a. Specify the start of authority of the zone and the default time-to-live information. This record designates the start of a zone. There should only be one start of authority record per zone. For example:

```
@ IN SOA venus.abc.aus.centurv.com.
bob.robert.abc.aus.century.com.
                           1.1
                                  ;serial
                           3600
                                  ;refresh
                           600
                                  ;retry
                           3600000 ; expire
                           86400) ; minimum
```

b. Include address-to-name resolution information on all hosts to be in the name server's zone of authority.

```
; ABC.AUS.CENTURY.COM Hosts
                           PTR
                     ΙN
                                 venus.abc.aus.century.com.
                     IN
                           PTR
                                 kronos.abc.aus.century.com.
```

c. Include other types of entries, such as name server records and canonical name records (optional).

Note: All lines in this file must be in Standard Resource Record Format.

6. Create an /etc/resolv.conf file by issuing the following command:

```
touch /etc/resolv.conf
```

The presence of this file indicates that the host should use a name server, not the /etc/hosts file, for name resolution. This file must exist on a name server host and may contain either the local host's address, the loopback address (127.0.0.1), be empty.

Alternatively, the /etc/resolv.conf file may contain the following entry:

```
nameserver 127.0.0.1
```

The 127.0.0.1 address is the loopback address, which causes the host to access itself as the name server. The /etc/resolv.conf file may also contain an entry like the following:

```
domain domainname
```

In the previous example, domainname would be aus.century.com.

- 7. Perform one of the following steps:
 - Enable the **named** daemon using the following SMIT fast path:

```
smit stnamed
```

This initializes the daemon with each system startup. Indicate whether you want to start the **named** daemon now, at the next system restart, or both.

 Edit the /etc/rc.tcpip file. Uncomment the line for the named daemon by removing the comment (#) symbol from the following line:

```
#start /etc/named "$src running"
```

This initializes the daemon with each system startup.

8. If you chose not to initialize the **named** daemon through SMIT, start the daemon for this session by running the following command:

```
startsrc -s named
```

Configuring a Slave Name Server

To configure a slave name server, use the Web-based System Manager fast path wsm network or use the following procedure, which edits a series of files and then uses SMIT or the command line to start the named daemon.

1. Edit the /etc/named.conf file. If there is no named.conf file in the /etc directory, copy the /usr/samples/tcpip/named.conf sample file into the /etc directory and edit it. Refer to the "named.conf File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a conf file.

This file is read each time the named daemon starts. It tells the server which type of server it is, the zone for which it is responsible, and where to get its initial information.

a. Using the options configuration clause, specify the directory where the named data files can be located (optional). Use this if you want the **named** data files to use paths relative to this directory. For example:

```
options {
    directory "/usr/local/domain";
};
```

b. Optionally specify the name of the hint zone file for the name server. This allows for the caching of record data outside of the defined zones. For example:

```
zone "." IN {
    type hint;
    file "/etc/named.ca";
};
```

c. Specify the slave zone clauses. Each stanza includes the zone type, an optional file name to which the name server can back up its data, and the list of master server Internet addresses. This list of addresses define the hosts from which the zone will be replicated. For example:

```
zone "abc.aus.century.com" IN {
    type slave;
    file "/etc/named.abc.bak";
   masters { 192.9.201.1; 192.9.201.2; };
};
zone "xyz.aus.century.com" IN {
    type slave;
    file "/etc/named.xyz.bak";
    masters { 192.9.201.1; 192.9.201.2; };
};
```

d. Include slave zone clauses to define the reverse name resolution information for the name server. For example:

```
zone "201.9.192.in-addr.arpa" IN {
    type slave;
    file "named.rev.bak";
    masters { 192.9.201.1; 192.9.201.2; };
};
zone "100.114.128.in-addr.arpa" IN {
   type slave;
    file "named.rev.bak";
    masters { 192.9.201.1; 192.9.201.2; };
};
```

e. To support resolving the loopback network address, specify a zone of type master with a source of /etc/named.local as well as the domain for which the name server will be responsible.

```
zone "0.0.127.in-addr.arpa" IN {
    type master;
    file "/etc/named.local";
};
```

2. Edit the /etc/named.ca file. Refer to the "DOMAIN Cache File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a cache file.

This file contains the addresses of the servers that are authoritative name servers for the root domain of the network. For example:

```
; root name servers.
                  IN
                         NS
                               relay.century.com.
        1
relay.century.com. 3600000
                                           129.114.1.2
                               ΤN
```

Note: All lines in this file must be in Standard Resource Record Format.

- 3. Edit the /etc/named.local file. See the "DOMAIN Local Data File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a local data file.
 - a. Specify the start of authority of the zone and the default time-to-live information. For example:

```
@ IN SOA venus.abc.aus.cntry.com. gail.zeus.abc.aus.cntry.com.
(
                           1.1
                                   ;serial
                           3600
                                   ;refresh
                           600
                                   ;retry
                           3600000 ;expire
                           86400)
                                   ; minimum
```

b. Specify the name server (NS) record. For example:

```
ΙN
             venus.abc.aus.century.com.
```

c. Specify the pointer (PTR) record.

```
1 IN PTR localhost.
```

Note: All lines in this file must be in Standard Resource Record Format.

4. Create an /etc/resolv.conf file by issuing the following command:

```
touch /etc/resolv.conf
```

The presence of this file indicates that the host should use a name server, not the /etc/hosts file, for name resolution. You may want to enter records to specify the name, domain, and address of the name server.

- 5. Perform one of the following steps:
 - Enable the **named** daemon using the following SMIT fast path:

```
smit stnamed
```

This initializes the daemon with each system startup. Indicate whether you want to start the **named** daemon now, at the next system restart, or both.

 Edit the /etc/rc.tcpip file. Uncomment the line for the named daemon by removing the comment (#) symbol from the following line:

```
#start /etc/named "$src_running"
```

This initializes the daemon with each system startup.

6. If you chose not to initialize the **named** daemon through SMIT, start the daemon for this session by running the following command:

```
startsrc -s named
```

Configuring a Hint Name Server

To configure a hint, or cache—only, name server, use the Web-based System Manager fast path **wsm network** or use the following procedure, which edits a series of files and then uses SMIT or the command line to start the **named** daemon.

Procedure

Configure a hint name server according to the following steps:

- Edit the /etc/named.conf file. If there is no named.conf file in the /etc directory, copy
 the /usr/samples/tcpip/named.conf sample file into the /etc directory and edit it. Refer
 to the "named.conf File Format for TCP/IP" in the AIX Files Reference for more
 information and a detailed example of a conf file.
 - To support resolving the loopback network address, specify a zone of type master with a source of /etc/named.local as well as the domain for which the name server will be responsible. For example:

```
zone "." IN {
    type hint;
    file "/etc/named.ca";
};
```

2. Edit the /etc/named.ca file. Refer to the "DOMAIN Cache File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a cache file.

This file contains the addresses of the servers that are authoritative name servers for the root domain of the network. For example:

```
; root name servers.  1 \hspace{1cm} \text{IN} \hspace{1cm} \text{NS} \hspace{1cm} \text{relay.century.com.}  relay.century.com.  3600000 \hspace{1cm} \text{IN} \hspace{1cm} \text{A} \hspace{1cm} 129.114.1.2
```

Note: All lines in this file must be in Standard Resource Record Format.

3. Edit the /etc/named.local file. Refer to the "DOMAIN Local Data File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a local data file.

a. Specify the start of authority of the zone and the default time-to-live information. For example:

```
@ IN SOA venus.abc.aus.cntry.com. gail.zeus.abc.aus.cntry.com.
(
                            1.1 ; serial ; refresh
                                    ;retry
                            600
                            3600000 ;expire
                            86400) ; minimum
```

b. Specify the name server (NS) record. For example:

```
TN
      NS
             venus.abc.aus.century.com.
```

c. Specify the pointer (PTR) record.

```
ΙN
      PTR
              localhost.
```

Note: All lines in this file must be in Standard Resource Record Format.

4. Create an /etc/resolv.conf file by issuing the following command:

```
touch /etc/resolv.conf
```

The presence of this file indicates that the host should use a name server, not the /etc/hosts file, for name resolution. You may want to enter records to specify the name, domain, and address of the name server.

- 5. Perform one of the following steps:
 - Enable the **named** daemon using the following SMIT fast path:

```
smit stnamed
```

This initializes the daemon with each system startup. Indicate whether you want to start the **named** daemon now, at the next system restart, or both.

- Edit the /etc/rc.tcpip file. Uncomment the line for the named daemon by removing the comment (#) symbol from the following line:

```
#start /etc/named "$src_running"
```

This initializes the daemon with each system startup.

6. If you chose not to initialize the **named** daemon through SMIT, start the daemon for this session by issuing the following command:

```
startsrc -s named
```

Configuring a Domain Mail Server

Configuring a domain mail server provides users external to your organization a simple method for addressing mail to your organization's users. That is, without a domain mail server, the mail address must specify a particular host in your organization. For example sam@orange.widget.com, where widget.com is your organization's domain name, and orange is the host that sam uses. But with a domain mail server, users outside your organization can simply specify the user name and domain name, without having to know which host the user uses, for example, sam@widget.com.

To configure a domain mail server, use the Web-based System Manager fast path wsm **network** or use one of the following procedures.

To Configure a Domain Mail Server

1. Create a mail exchanger (MX) record and an address (A) record for the mail server (black.widget.com):

```
widget.com IN MX 10 black.widwidget.com IN A 192.10.143.9 black.widget.com IN A 192.10.143.9
                                                         10 black.widget.com
```

2. Edit sendmail.cf on the mail server (black.widget.com) to add the domain alias (the w class):

```
Cw $w $?D$w.$D$. widget.com
```

3. Mail clients must know where to send their non-local mail, so edit sendmail.cf on each client to point to the mail server (the S macro):

```
DRblack.widget.com
```

- 4. Use the NameServOpt option to configure the sendmail daemon so everyone can use the MX records defined in the name server brown.widget.com.
- 5. Add aliases for users in the domain that do not have accounts on the mail server using the aliases file, for example:

```
sam:sam@orange.widget.com
david:david@green.widget.com
judy:judy@red.widget.com
```

Note: Mailbox (MB) records can serve the same function.

- 6. The serial number in the SOA Resource Record should be incremented because the database has been modified.
- 7. Refresh the name server database by issuing the **refresh** –**s named** command.
- 8. Perform the following steps:
 - a. Enter the command **sendmail** -bz to recompile the **sendmail.cf** file on the mail server.
 - b. Enter the command **sendmail** -bi to recompile the aliases database on the mail server.
 - c. Run the **refresh –s sendmail** command to make the changes take effect.
- 9. On the clients, recompile (sendmail) and run the refresh -s sendmail command to make the changes take effect.

There are other methods to configure a domain mail server. The following procedures use experimental mailbox (MB), mail rename (MR), and mail group (MG) records.

To Configure a Domain Mail Server Using mailbox (MB) Records

1. Define a mailbox (MB) record for each user in the domain. Add entries such as:

```
sam IN MB orange.widget.com.
```

to the /etc/named.data file on host brown.widget.com. These entries identify to the mail server (black.widget.com) where to send mail for each user in the domain.

- 2. Configure the sendmail daemon on the mail server (black.widget.com) to use the MB records defined in the name server (brown.widget.com). Use the NameServOpt option.
- 3. Increment the serial number in the SOA Resource Record, because the database has been modified.
- 4. Refresh the name server database by running the refresh -s named command.
- 5. Enter the command sendmail -bz to recompile the sendmail.cf file on the mail server, and then run the refresh -s sendmail command to make the changes take effect.

Defining a Mail Rename (MR) Record for a User

- 1. Edit the /etc/named.data file on your domain name server.
- 2. Add a Mail Rename record for each alias. For example, if a user sam has an alias sammy, the Mail Rename record would be:

```
sammy IN MR sam
```

This record will cause all mail addressed to sammy to be delivered to sam. Each MR record should be entered on a line by itself.

- 3. Since the database has been modified, the serial number in the SOA Resource Record should be incremented.
- 4. Refresh the name server database by issuing the **refresh –s named** command.
- 5. Enter the command sendmail -bz to recompile the sendmail.cf file on the mail server, and then run the refresh -s sendmail command to make the changes take effect.

Defining Mail Group (MG) Member Records

- 1. Edit the /etc/named.data file on your domain name server.
- 2. Add MG records for each mail group. MG records function like the /etc/aliases file, with the aliases maintained on the name server. For example:

```
users IN HINFO users-request widget.com
users IN MG sam
users IN MG david
users IN MG judy
```

This example would cause all mail addressed to users@widget.com to be delivered to sam, david, and judy. Each MG record should be on a line by itself.

Note: Users sam, david, and judy must have MB records defined.

- 3. The serial number in the SOA Resource Record should be incremented, because the database has been modified.
- 4. Refresh the name server database by issuing the **refresh** –**s named** command.
- 5. Enter the command **sendmail -bz** to recompile the **sendmail.cf** file on the mail server, and then run the refresh -s sendmail command to make the changes take effect.

Defining Mail Exchanger (MX) Records

- 1. Edit the /etc/named.data file on your domain name server.
- 2. Add MX records for each machine not directly connected to your network to which you wish to forward mail. For example, if mail addressed to users on purple.widget.com should be forwarded to post.office.widget, the MX record would look like the following:

```
purple.widget.com IN MX 0 post.office.widget.
```

You must specify both host and machine names when using MX records. Each MX record should be on a line by itself. You may use wildcards, for example:

```
*.widget.com IN MX 0 post.office.widget.
```

This example would cause mail to an unknown host (a host without an explicit MX record) in the widget.com domain to be forwarded to post.office.widget.

Note: Wildcard MX records are not appropriate for use on the Internet.

- 3. The serial number in the SOA Resource Record should be incremented because the database has been modified
- 4. Refresh the name server database using the **refresh –s named** command.
- 5. Enter the command **sendmail -bz** to recompile the **sendmail.cf** file on the mail server, and then run the refresh -s sendmail command to make the changes take effect.

Configuring a Forwarder

To configure a forwarder, use the Web-based System Manager fast path **wsm network** or use the following procedure, which edits a series of files and then uses SMIT or the command line to start the **named** daemon.

- Edit the /etc/named.conf file. If there is no named.conf file in the /etc directory, copy
 the /usr/samples/tcpip/named.conf sample file into the /etc directory and edit it. Refer
 to the "named.conf File Format for TCP/IP" in the AIX Files Reference for more
 information and a detailed example of a conf file.
 - Specify a forwarders line in the options stanza of the /etc/named.conf file that lists the IP addresses of the name servers that should receive the forwarded requests. For example:

```
options {
    ...
    forwarders { 192.100.61.1; 129.35.128.222; };
    ...
};
```

Specify the hint zone. For example:

```
zone "." IN {
   type hint;
   file "/etc/named.ca";
};
```

2. Edit the /etc/named.ca file. Refer to the "DOMAIN Cache File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a cache file.

This file contains the addresses of the servers that are authoritative name servers for the root domain of the network. For example:

```
; root name servers.

1 IN NS relay.century.com.
relay.century.com. 3600000 IN A 129.114.1.2
```

Note: All lines in this file must be in Standard Resource Record Format.

- 3. Edit the /etc/named.local file. Refer to the "DOMAIN Local Data File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a local data file.
 - a. Specify the start of authority of the zone and the default time—to—live information. For example:

```
@ IN SOA venus.abc.aus.cntry.com. gail.zeus.abc.aus.cntry.com.

1.1 ;serial
3600 ;refresh
600 ;retry
3600000 ;expire
86400) ;minimum
```

b. Specify the name server (NS) record. For example:

```
IN NS venus.abc.aus.century.com.
```

c. Specify the pointer (PTR) record.

```
1 IN PTR localhost.
```

Note: All lines in this file must be in Standard Resource Record Format.

4. Create an /etc/resolv.conf file by issuing the following command:

```
touch /etc/resolv.conf
```

The presence of this file indicates that the host should use a name server, not the /etc/hosts file, for name resolution.

Alternatively, the /etc/resolv.conf file may contain the following entry:

```
nameserver 127.0.0.1
```

The 127.0.0.1 address is the loopback address, which causes the host to access itself as the name server. The /etc/resolv.conf file may also contain an entry like the following:

```
domain domainname
```

In the previous example, domainname would be austin.century.com.

- Perform one of the following steps:
 - Enable the **named** daemon using the following SMIT fast path:

```
smit stnamed
```

This initializes the daemon with each system startup. Indicate whether you want to start the **named** daemon now, at the next system restart, or both.

- Edit the /etc/rc.tcpip file. Uncomment the line for the named daemon by removing the comment (#) symbol from the following line:

```
#start /etc/named "$src_running"
```

This initializes the daemon with each system startup.

6. If you chose not to initialize the named daemon through SMIT, start the daemon for this session by issuing the following command:

```
startsrc -s named
```

Configuring a Forward Only Server

To configure a forward only name server, use the Web-based System Manager fast path wsm network or use the following procedure, which edits a series of files and then uses SMIT or the command line to start the **named** daemon.

Note: You can achieve a similar configuration without running a forward only name server. Instead, create an /etc/resolv.conf file that contains name server lines that point to the forwarders you wish to use.

- 1. Edit the /etc/named.conf file. If there is no named.conf file in the /etc directory, copy the /usr/samples/tcpip/named.conf sample file into the /etc directory and edit it. Refer to the "named.conf File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a conf file.
 - Specify the forwarders and forward only lines in the options stanza of the /etc/named.conf file listing the IP addresses of the name servers receiving the forwarded requests. For example:

```
options {
    forwarders { 192.100.61.1; 129.35.128.222; };
    forward only;
};
```

2. Edit the /etc/named.ca file. Refer to "DOMAIN Cache File Format for TCP/IP" in AIX Files Reference for more information and a detailed example of a cache file. This file contains the addresses of the servers that are authoritative name servers for root domain of the network. For example:

```
; root name servers.
1 IN NS relay.century.com. relay.century.com. 3600000 IN A 129.114.1.2
```

Note: All lines in this file must be in Standard Resource Record Format.

- 3. Edit the /etc/named.local file. Refer to the "DOMAIN Local Data File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a local data file.
 - a. Specify the start of authority of the zone and the default time-to-live information. For example:

```
@ IN SOA venus.abc.aus.cntry.com. gail.zeus.abc.aus.cntry.com.

(

1.1 ;serial
3600 ;refresh
600 ;retry
3600000 ;expire
86400) ;minimum
```

b. Specify the name server (NS) record. For example:

```
IN NS venus.abc.aus.century.com.
```

c. Specify the pointer (PTR) record.

```
1 IN PTR localhost.
```

Note: All lines in this file must be in Standard Resource Record Format.

4. Create an /etc/resolv.conf file by issuing the following command:

```
touch /etc/resolv.conf
```

The presence of this file indicates that the host should use a name server, not the /etc/hosts file, for name resolution.

Alternatively, the /etc/resolv.conf file may contain the following entry:

```
nameserver 127.0.0.1
```

The 127.0.0.1 address is the loopback address, which causes the host to access itself as the name server. The /etc/resolv.conf file may also contain an entry such as:

```
domain domainname
```

In the previous example, domainname would be austin.century.com.

- 5. Perform one of the following steps:
 - Enable the named daemon using the following SMIT fast path:

```
smit stnamed
```

This initializes the daemon with each system startup. Indicate whether you want to start the **named** daemon now, at the next system restart, or both.

 Edit the /etc/rc.tcpip file. Uncomment the line for the named daemon by removing the comment (#) symbol from the following line:

```
#start /etc/named "$src_running"
```

This initializes the daemon with each system startup.

6. If you chose not to initialize the **named** daemon through SMIT, start the daemon for this session by issuing the following command:

```
startsrc -s named
```

Configuring a Host to Use a Name Server

To configure a host to use a name server, use the Web-based System Manager fast path wsm network or use the following procedure.

- 1. Create an /etc/resolv.conf file.
- 2. If this host will use more than one name server, add the other name servers that this host will use.
- 3. Assuming the name server is operational, you can test the communication between the host and the name server using the following command:

```
host hostname
```

Use the name of a host that should be resolved by the name server to see if the process is working. The output you should receive will appear similar to the following:

```
brown.abc.aus.century.com is 129.35.145.95
```

Other configuration tasks are shown in the following table.

Configuring a Host to Use a Name Server Tasks				
Web-based System Manager (wsm network application)	: wsm network fast path			
–OR–				
Task	SMIT Fast Path	Command or File		
Create an /etc/resolv.conf File	smit stnamerslv2	create and edit /etc/resolv.conf ¹		
List All the Name Servers Used by a Host	smit Isnamerslv	view /etc/resolv.conf		
Add a Name Server	smit mknamerslv	edit /etc/resolv.conf ²		
Remove a Name Server	smit rmnamerslv	edit /etc/resolv.conf		
Start/Restart Using Domain Name Resolution	smit stnamersly			
Stop Using Domain Name Resolution	smit spnamerslv			
Change/Show the Domain	smit mkdomain	edit /etc/resolv.conf		
Remove the Domain	smit rmdomain	edit /etc/resolv.conf		

Note:

1. On the first line of the /etc/resolv.conf file, enter the word domain followed by the full name of the domain that this host is in. For example:

```
domain abc.aus.century.com
```

2. On any blank line below the line which starts with the word domain, enter the word nameserver, followed by at least one space, followed by the dotted decimal Internet address of the name server that this host will use (the name server must serve the domain indicated by the domain statement). You may have up to 16 name server entries. For example, your /etc/resolv.conf file might contain the entries:

```
nameserver 192.9.201.1
nameserver 192.9.201.2
```

The system will query the name servers in the order listed.

Configuring Dynamic Zones on the DNS Name Server

The **named** command allows for dynamic updates. The named database and configuration files need to be configured to allow for client machines to issue updates. A zone can be set to dynamic or static. The default zone is static.

To make a zone dynamic, you need to add the keyword allow-update to that zone's stanza in the /etc/named.conf file. The allow-update keyword specifies an Internet address match list that define hosts allowed to submit updates. Refer to the "named.conf File Format for TCP/IP" in the AIX Files Reference for more information and a detailed example of a conf file. In the following example, all hosts are allowed to update the dynamic zone:

```
zone "aoot.austin.ibm.com" IN {
    type master;
    file "named.data";
    allow-update { any; };
};
```

After a zone is marked dynamic, three modes of security can be initiated:

Unsecured Allows anyone at anytime to update any information in the

zone.

Warning: Use of this mode is not recommended. It can lead to data loss, data interception, and user frustration. At the least, an unsecured zone should be limited to updates

only from specific internet addresses.

Controlled Allows for the creation of new information and the

> replacement of existing information. This is probably the easiest mode to use for a secure transition environment. This mode also requires that all incoming updates be

timestamped and have keyed signatures.

Presecured Requires all updates to existing information be replaced

with similar information. This mode also requires that all incoming updates be timestamped and have keyed

signatures.

A dynamic zone defaults to unsecured mode. To use one of the other modes, add "controlled" or "presecured" after the keyword update-security in the zone stanza of the /etc/named.conf file. This tells the named server the level of security to use with that zone. For example:

```
zone "aoot.austin.ibm.com" IN {
    type master;
    file "named.boot";
    allow-update { any; };
    update-security controlled;
};
```

After a mode is selected, the actual data files need to be modified for your level of security. In unsecured mode, the data files are used "as is." For controlled or presecured mode, you need to generate a set of master server/hostname key pairs per name in the zone. This should be done with the **nsupdate** command using the **-g** option. This command generates the key pair, a private and a public key. These keys are needed to authentically sign for updates. After generating all the keys for your list of zone names, you need to add them to the data file. The KEY format is as follows:

Index ttl Class Type KeyFlags Protocol Algorithm KeyData where:

Specifies the name used to reference the data in the zone. Index ttl Specifies the time to live for this data. This is an optional field. Class Specifies the class of the data. This is dependent on the zone, but

usually it is IN.

Type Indicates the type of the record. In this case, it is KEY.

KeyFlags Gives named information about the key. 0x0000 defines the typical

key record used for a host. 0x0100 defines the key record

associated with the zone name.

Protocol Specifies the protocol to use. Currently, there is only one, 0.

Algorithm Specifies the algorithm of the key. Currently, there is only one, 1.

This is the MD5 Private/Public authentication method.

KeyData Indicates the key in base64 representation. The **nsupdate**

> command generates both the public and private keys in base64 representation. The public key is listed last in the output file.

Example

To ensure security over a host name in a dynamic zone, a line similar to the following needs to be added to the zone file for the zone containing the hostname.

4660 KEY 0x0000 bears ΙN AQOtg....

The above example indicates that bears has a KEY record defined. Someone wanting to update bears would need to sign his update with the private key matching the this public key in the database. For the nsupdate command to succeed, the private key needs to be placed on the client in a keyfile (defaults to /etc/keyfile). It should be in the format:

hostname mastername base64 key

A similar KEY entry is required in the zone definition section. A zone key is required for both presecured and controlled modes or the mode is considered to be unsecured. This can be done as shown in the previous bears example, but the private key is left for the administrator to use with the **nsupdate** command's administrative mode.

To generate a key pair using the **nsupdate** command, enter:

nsupdate -g -h ZoneName -p ServerName -k AdminKeyFile

This generates a key for the zone. Place the last key of the pair in the beginning section for the zone as follows:

0x0100 0 1 Key IN

The zone is ready to be loaded. The administrator should use the zone key to apply updates and maintenance operations on the zone.

TCP/IP Routing

The topics discussed in this section are:

- Static and Dynamic Routing, on page 3-121
- Gateways, on page 3-122
- Planning for Gateways, on page 3-123
- · Configuring a Gateway, on page 3-124
- Restricting Route Use, on page 3-126
- Manually Removing Dynamic Routes, on page 3-126
- Configuring the routed Daemon, on page 3-126

A *route* defines a path for sending packets through the Internet network to an address on another network. A route does not define the complete path, only the path segment from one host to a gateway that can forward packets to a destination (or from one gateway to another). There are three types of routes:

host route Defines a gateway that can forward packets to a specific

host or gateway on another network.

network route Defines a gateway that can forward packets to any of the

hosts on a specific network.

default route Defines a gateway to use when a host or network route to a

destination is not otherwise defined.

Routes are defined in the kernel *routing table*, which can hold up to 32 route definitions. These route definitions include information on networks reachable from the local host, gateways that can be used to reach remote networks, and the hop count (or distance metric) to those networks. When a gateway receives a datagram, it checks the routing tables to find out where next to send the datagram along the path to its destination.

Static and Dynamic Routing

In TCP/IP, routing can be one of two types: *static* or *dynamic*. With static routing, you maintain the routing table manually using the **route** command. Static routing is practical for a single network communicating with one or two other networks. However, as your network begins to communicate with more networks, the number of gateways increases, and so does the amount of time and effort required to maintain the routing table manually.

With dynamic routing, daemons update the routing table automatically. Routing daemons continuously receive information broadcast by other routing daemons, and so continuously update the routing table.

TCP/IP provides two daemons for use in dynamic routing, the **routed** and **gated** daemons. The **gated** daemon supports Routing Information Protocol (RIP), Routing Information Protocol Next Generation (RIPng), Exterior Gateway Protocol (EGP), Border Gateway Protocol (BGP) and BGP4+, Defense Communications Network Local—Network Protocol (HELLO), Open Shortest Path First (OSPF), Intermediate System to Intermediate System (IS—IS), and Internet Control Message Protocol (ICMP and ICMPv6)/Router Discovery routing protocols simultaneously. In addition, the **gated** daemon supports the Simple Network Management Protocol (SNMP). The **routed** daemon only supports Routing Information Protocol.

Routing daemons can operate in one of two modes, *passive* or *active*, depending upon the options you use when starting the daemons. In active mode, routing daemons both broadcast routing information periodically about their local network to gateways and hosts, and receive routing information from hosts and gateways. In passive mode, routing

daemons receive routing information from hosts and gateways, but do not attempt to keep remote gateways updated (they do not advertise their own routing information).

These two types of routing can be used not only for gateways, but for other hosts on a network as well. Static routing works the same for gateways as for other hosts. Dynamic routing daemons, however, must be run in the passive (quiet) mode when run on a host that is not a gateway.

Gateways

Gateways are a type of router. Routers connect two or more networks and provide the routing function. Some routers, for example, route at the network interface level or at the physical level.

Gateways, however, route at the network level. Gateways receive IP datagrams from other gateways for delivery to hosts on the local network, and route IP datagrams from one network to another. For example, a gateway connecting two Token-Ring networks has two Token-Ring adapter cards, each with its own Token-Ring network interface. To pass on information, the gateway receives datagrams through one network interface and sends them out through the other network interface. Gateways periodically verify their network connections through interface status messages.

Gateways route packets according to the destination network, not according to the destination host. That is, a gateway machine is not required to keep track of every possible host destination for a packet. Instead, a gateway routes packets according to the network of the destination host. The destination network then takes care of sending the packet to the destination host. Thus, a typical gateway machine requires only limited disk storage capacity (if any) and limited main memory capacity.

The distance a message must travel from originating host to destination host depends upon the number of gateway hops it must make. A gateway is zero hops from a network to which it is directly attached, one hop from a network that is reachable through one gateway, and so on. Message distance is usually expressed in the number of gateway hops required, or hop counts (also called the metric).

Interior and Exterior Gateways

Interior gateways are gateways that belong to the same autonomous system. They communicate with each other using the Routing Information Protocol (RIP), Routing Information Protocol Next Generation (RIPng), Intermediate System to Intermediate System protocol, Open Shortest Path First protocol (OSPF), or the HELLO Protocol (HELLO). Exterior gateways belong to different autonomous systems. They use the Exterior Gateway Protocol (EGP), the Border Gateway Protocol (BGP), or BGP4+.

For example, consider two autonomous systems. The first is all the networks administered by the Widget Company. The second is all the networks administered by the Gadget Company. The Widget Company has one machine, called apple, which is Widget's gateway to the Internet. The Gadget Company has one machine, called orange, which is Gadget's gateway to the Internet. Both companies have several different networks internal to the companies. The gateways connecting the internal networks are interior gateways. But apple and orange are exterior gateways.

Each exterior gateway does not communicate with every other exterior gateway. Instead, the exterior gateway acquires a set of neighbors (other exterior gateways) with which it communicates. These neighbors are not defined by geographic proximity, but rather by their established communications with each other. The neighboring gateways, in turn, have other exterior gateway neighbors. In this way, the exterior gateways' routing tables are updated and routing information is propagated among the exterior gateways.

The routing information is sent in a pair, (N.D), where N is a network and D is a distance reflecting the cost of reaching the specified network. Each gateway advertises the networks it can reach and the costs of reaching them. The receiving gateway calculates the shortest paths to other networks and passes this information along to its neighbors. Thus, each

exterior gateway is continually receiving routing information, updating its routing table and then passing that information to its exterior neighbors.

Gateway Protocols

All gateways, whether interior or exterior, use protocols to communicate with each other. Here are brief descriptions of the more commonly used TCP/IP gateway protocols:

HELLO Protocol (HELLO)

HELLO is one protocol that the interior gateways use to communicate among themselves. HELLO calculates the shortest path to other networks by determining the path that has the least delay time.

Routing Information Protocol (RIP)

Routing Information Protocol is a protocol that the interior gateways use to communicate among themselves. Like the HELLO Protocol, RIP calculates the shortest path to other networks. Unlike HELLO, RIP estimates distance not by delay time, but by hop counts. Because the gated daemon stores all metrics internally as time delays, it converts RIP hop counts into time delays.

Routing Information Protocol Next Generation

RIPng is the RIP protocol that is enhanced to support IPv6.

Open Shortest Path First (OSPF)

OPSF is a protocol that the interior gateways use to communicate among themselves. It is a link-state protocol that is bettter suited than RIP for complex networks with many routers. It provides equal cost multipath routing.

Exterior Gateway Protocol (EGP)

The exterior gateways can use the Exterior Gateway Protocol to communicate among themselves. The EGP does not calculate the shortest path to other networks. Instead, it merely indicates whether a particular network is reachable or not.

Border Gateway Protocol (BGP)

The exterior gateways can use this protocol to communicate among themselves. It exchanges reachability information between automomous systems, but provides more capabilities than EGP. BGP uses path attributes to provide more information about each route as an aid in selecting the best route.

Border Gateway Protocol 4+

BGP4+ is the BGP protocol version 4, which supports IPv6 and has other enhancements over past versions of the protocol.

Intermediate System to Intermediate System (IS-IS)

Interior gateways use IS-IS protocol to communicate among themselves. It is a link-state protocol that can route IP and ISO/CLNP packets and, like OSPF, uses a "shorter path first" algorithm to determine routes.

Planning for Gateways

Before you configure the gateways for your network, you must first:

- 1. Consider the number of gateways to use.
- 2. Decide on the type of routing to use.

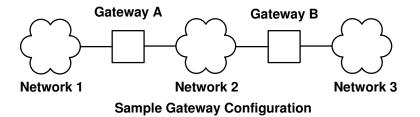
Consider the Number of Gateways to Use

The number of gateways you need to configure will depend upon:

- The number of networks you want to connect.
- How you want to connect the networks.

The level of activity on the connected networks.

For example, suppose users on Network 1, Network 2, and Network 3 all need to communicate with each other (as illustrated in the Sample Gateway Configuration figure). To connect Network 1 directly to Network 2, you would use a single gateway (Gateway A). To connect Network 2 directly to Network 3, you would use another gateway (Gateway B). Now, assuming the proper routes are defined, all the users on all three networks can communicate.



However, if Network 2 is very busy, communication between Network 1 and Network 3 may suffer unacceptable delays. Furthermore, if most of the inter-network communication occurs between Network 1 and Network 3, you may want to connect Network 1 directly to Network 3. To do this, you could use an additional pair of gateways, Gateway C (on Network 1) and Gateway D (on Network 3), with a direct connection between these two additional gateways. This may be an inefficient solution, however, because one gateway can connect more than two networks.

A more efficient solution may be to connect Gateway A to Gateway B directly, as well as to Network 2. This would require a second network adapter in both Gateway A and Gateway B. In general, the number of networks you connect through a single gateway is limited by the number of network adapter cards the gateway machine can support.

Decide on the Type of Routing to Use

If your network is small, and its configuration rarely changes, you will probably want to use static routing. But if you have a large network whose configuration changes frequently, you will probably want to use dynamic routing. You may decide to use a combination of static and dynamic routing. That is, you might want to give static definitions to a few specific routes, while allowing other routes to be updated by the daemons. The static routes you create are not advertised to other gateways and are not updated by the routing daemons.

If Using Dynamic Routing

Choose the routing daemon according to the type of gateway you need and the protocols your gateway must support. If the gateway is an interior gateway, and only needs to support RIP, choose the **routed** daemon. If the gateway must support any other protocol, or is an exterior gateway, choose the gated daemon.

Note: Unpredictable results can occur if the gated and routed daemons run on the same host at the same time.

Configuring a Gateway

To configure a machine to act as a gateway, use the following the instructions below. For clarity, this procedure assumes that the gateway machine will connect two networks, and that the gateway machine has already been minimally configured (see "Configuring TCP/IP", on page 3-90) on one of the networks.

1. Install and configure the second network adapter, if you have not done so already. (See "Installing a Network Adapter", on page 3-36, "Configuring a High-Performance Token-Ring Adapter", on page 3-37, and "Configuring a High-Performance Ethernet Adapter", on page 3-37.)

- 2. Choose an IP address for the second network interface, and then configure the network interface by following the instructions in "Managing Network Interfaces", on page 3-51.
- 3. Add a route to the second network.
- 4. To use an AIX machine as an internetwork router over TCP/IP networks,enter:
 - no -o ipforwarding=1
- 5. The gateway machine can now access both of the networks to which it is directly attached.
 - a. If you want to use static routing to communicate with hosts or networks beyond these two networks, add any other routes you want.
 - b. If you want to use dynamic routing, follow the instructions in either "Configuring the routed Daemon", on page 3-126 or "Configuring the gated Daemon", on page 3-126. If your internetwork will join the Internet, you should also follow the instructions in "Getting an Autonomous System Number", on page 3-129.

Configuring a Gateway Tasks				
Web-based System Manager: wsm network fast path (wsm network application)				
-OR-				
Task	SMIT Fast Path	Command or File		
Displaying the Routing Table	smit Isroute	netstat -rn1		
Adding a Static Route	smit mkroute	route add destination gateway ²		
Removing a Static Route	smit rmroute	route delete destination gateway ²		
Flushing the Routing Table	smit fshrttbl	route flush		

Note:

- 1. The table is divided into columns for destination address, gateway address, flags, reference count (hop count), and network interface. (For a detailed discussion of each of these columns, see the netstat command in the AIX Commands Reference.) If frames are not reaching their destination and the routing tables indicate the correct route, there is a good chance that one or more of the following conditions exist:
 - Network is failing.
 - Remote host or gateway is failing.
 - Remote host or gateway is down or not ready to receive frames.
 - Remote host does not have a route back to the source network.
- 2. The destination value is the dotted decimal address or symbolic name of the destination host or network, and the gateway value is the dotted decimal address or symbolic name of the gateway. (A default route specifies 0 as the destination.)

Restricting Route Use

Routes can be restricted so they can be used only by some users. The restrictions are based on the primary and auxiliary group IDs of users. Using the route command, you can specify a list of up to 32 group IDs that are allowed or not allowed to use a route. If the list is of allowed groups, any user that belongs to any group on the list can use the route. If the list is of disallowed groups, only users that do not belong to any of the groups on the list can use the route. The root user can use any route.

Groups may also be associated with an interface using the ifconfig command. In this case, a forwardable packet may use any route allowed for the groups associated with its incoming interface.

If there are two or more routes to the same destination, any ICMP redirects that are received for that destination will be ignored and path MTU discovery will not be done on those routes.

Manually Removing Dynamic Routes

If you are using the **routed** daemon, a manually deleted route is *not* replaced by incoming RIP information (since ioctl's are used). If you are using the **gated** daemon, and the -n flag is not used, the manually deleted route is replaced by the route as discovered in incoming RIP information.

Configuring the routed Daemon

To configure the **routed** daemon:

- 1. Remove the comment symbol (#) and modify the routed clause in the /etc/rc.tcpip shell script. This automatically starts the **routed** daemon with each system startup.
 - Specify whether you want the gateway to run in active (-s flag) or passive (-q flag) mode.
 - Specify whether you want packet tracing on or off (-t flag). Packet tracing can also be turned on after the routed daemon is already started by using the kill command to send a SIGUSR1 signal to the daemon. This signal can also be used to increment the level of tracing through four levels. In addition, packet tracing can be turned off while the routed daemon is running by using the kill command to send a SIGUSR2 signal to the daemon. For more information, see the routed daemon and the kill command.
 - Specify whether you want debugging turned on or off (-d flag). If you use this flag, specify which log file you want debugging information stored in, or choose for it to be directed to the console display.
 - Specify whether you are running the routed daemon on a gateway (-q flag).

Note: A host that is not a gateway can run the **routed** daemon, but it must be run in passive mode.

- 2. Identify any known networks by listing them in the /etc/networks file. Refer to "Networks File Format for TCP/IP" in the AIX Files Reference for more information. A sample networks file is located in the /usr/samples/tcpip directory.
- 3. Set up routes in the /etc/gateways file to any known gateways that are not directly connected to your network. Refer to "Gateways File Format for TCP/IP" in AIX Files Reference for detailed examples of entries in the /etc/gateways file. A sample gateways file is located in the /usr/samples/tcpip directory.

Attention: Do not run the routed daemon and the gated daemon on the same machine. Unpredictable results may occur.

Configuring the gated Daemon

To configure the gated daemon:

1. Decide which gateway protocols are most appropriate for your system. The choices for routing protocols are EGP, BGP, RIP, RIPng, HELLO, OSPF, ICMP/Router Discovery, and IS-IS. You can also use SNMP, a protocol allowing you to change or show management information for a network element from a remote host.

Note: Use EGP, BGP, or BGP4+ to advertise addresses of networks in an autonomous system to gateways in other autonomous systems. If you are on the Internet, EGP, BGP, or BGP4+ must be used to advertise network reachability to the core gateway system. Use the interior routing protocols to advertise reachability information within an autonomous system.

- 2. Identify any known networks by listing them in the /etc/networks file. Refer to "Networks File Format for TCP/IP" in AIX Files Reference for more information. A sample networks file is located in the /usr/samples/tcpip directory.
- 3. Edit the /etc/gated.conf file to reflect the desired gated daemon configuration.
 - a. Specify the level of trace output you want. If tracing is needed before the gated.conf file is parsed, use the -t flag to turn tracing on when the daemon starts. Refer to "gated Daemon" in AIX Commands Reference for more information.
 - b. Specify the routing protocols you want to use. Each protocol has its own protocol statement. Remove the comment symbols (#) and modify the statements corresponding to the protocols you want to use.
 - If using EGP:
 - Set up the EGP autonomoussystem clause. Obtain an autonomous system number from the Internet authority if you are on the Internet, or if not, assign an autonomous system number considering the autonomous system numbers of other systems on your network.
 - Set the EGP statement to yes.
 - Set up a group clause for each autonomous system.
 - Set up a neighbor clause for each neighbor in that autonomous system. For example:

```
autonomoussystem 283;
egp yes {
        group maxup 1 {
             neighbor nogendefault 192.9.201.1;
             neighbor nogendefault 192.9.201.2;
       } ;
       group {
             neighbor 192.10.201.1;
             neighbor 192.10.201.2;
       } ;
} ;
```

- If using RIP or HELLO:
 - Set the RIP or HELLO statement to yes.
 - Specify quiet in the RIP or HELLO statement if you want the gateway only to accept routing information, not broadcast information. Or specify supplier in the RIP or HELLO statement if you want the gateway to broadcast routing information as well as accept routing information.

 Specify pointopoint if you want the gateway to broadcast only to source gateways specified in the sourcegateways statement. If not, omit this value. If you include pointopoint, specify a gateway name or Internet address in dotted decimal in the sourcegateways clause. For example:

```
# Broadcast to specific gateways
rip/hello pointopoint {
         sourcegateways
               101.25.32.1
               101.25.32.2;
} ;
# Broadcast to all
rip/hello supplier {
         interface en0 noripout ;
        trustedgateways
                101.25.33.1
               101.25.33.2;
} ;
```

These first two examples could both be active in the **gated.conf** file.

```
# Broadcast to no one
rip/hello quiet {
         interface tr0 noripin ;
} ;
```

- If using BGP:
 - Set up the BGP autonomoussystem clause. Obtain an autonomous system number from the Internet authority if you are on the Internet, or if not, assign an autonomous system number considering the autonomous system numbers of other systems on your network.
 - Set the BGP statement to yes.
 - Set up a peer clause for each neighbor in that autonomous system. For example:

```
# Perform all BGP operations
bgp yes {
       peer 192.9.201.1 ;
} ;
```

- If using SNMP:
 - Set the SNMP statement to yes.

```
snmp yes ;
```

Configuring the gated Daemon To Run IPv6

To configure the gated daemon to run under Internet Protocol version 6 (IPv6), first ensure that your system has been configured for IPv6 and IPv6 routing:

- 1. Run **autoconf6** to automatically configure your interfaces for IPv6.
- 2. Configure site local addresses for each IPv6 interface on which you want to use IPv6 routing using the following command:

```
ifconfig interface inet6 fec0:n::address/64 alias
where
interface
             Is the name of the interface, such as tr0 or en0.
```

Is any decimal number; for example, 11 n

address Is the portion of the IPv6 interface address that follows the double

colons; for example, given the IPv6 address

fe80::204:acff:fe86:298d, the address entry would be

204:acff:fe86:298d.

Note: You can use the command netstat -i to see what your IPv6 address is for each configured interface.

So, if token ring tr0 has an IPv6 address of fe80::204:acff:fe86:298d, you would issue the following command:

```
ifconfig tr0 inet6 fec0:13::204:acff:fe86:298d/64 alias
```

3. Turn on IPv6 forwarding with the following command:

```
no -o ip6forwarding=1
```

4. Start **ndpd-router** with the following command:

```
ndpd-router -q
```

See **ndpd-router** to determine which flags to use for your network configuration.

Starting **ndpd-router** will allow your system to act as a router for the Neighbor Discovery Protocol. Neighbor Discovery Protocol routers inform Neighbor Discovery hosts with routing information so hosts can route IPv6 packets.

Any hosts on the network that you want to be part of the IPv6 network must run **ndpd-host**. Hosts on the network that run **ndpd-host** will recognize themselves as part of an IPv6 network and use Neighbor Discovery Protocol, which allows them to determine and monitor link-layer addresses both to allow neighbor routing and to find neighboring routers for forwarding packets.

See ndpd-router, ndpd-host, or read RFC 1970, Neighbor Discovery, for more information.

Next, configure the gated daemon:

- 1. Decide which IPv6 gateway protocols are most appropriate for your system. The choices for IPv6 routing protocols are Border Gateway Protocol enhanced for IPv6 (BGP4+) and Routing Information Protocol Next Generation (RIPng).
- 2. Edit the etc/gated.conf file to reflect the desired gated daemon configuration.

Note: AIX versions 4.3.2 and later run gated version 3.5.9. The syntax of the gated.conf file has changed slightly from earlier versions. Read the gated.conf documentation or use the sample file that is shipped in the /usr/sample/tcpip directory for correct syntax.

When configuring BGP4+ or RIPng, use IPv6 addresses in which the syntax specifies an IP address.

Note: By default, RIPng multicasts its packets.

Once the /etc/gated.conf file has been modified, the gated daemon can be started.

Getting an Autonomous System Number

If you use EGP or BGP, you should obtain an official autonomous system number for your gateway. To obtain an official autonomous system number, contact the NIC at INFO@INTERNIC.NET, or at 1-800-444-4345 (USA).

Path MTU Discovery

For two hosts communicating across a path of multiple networks, a transmitted packet will become fragmented if its size is greater than the smallest MTU of any network in the path. Since packet fragmentation can result in reduced network performance, it is desirable to avoid fragmentation by transmitting packets with a size is no greater than the smallest MTU in the network path. This size is called the path MTU.

AIX supports a path MTU discovery algorithm as described in RFC 1191. Path MTU discovery can be enabled for TCP and UDP applications by modifying the tcp_pmtu_discover and udp_pmtu_discover options of the no command. When enabled for TCP, path MTU discovery will automatically force the size of all packets transmitted by TCP applications to not exceed the path MTU. Since UDP applications themselves determine the size of their transmitted packets, UDP applications must be specifically written to utilize path MTU information by using the IP FINDPMTU socket option, even if the udp_pmtu_discover no option is enabled. By default, the tcp_pmtu_discover and udp_pmtu_discover options are disabled.

When the path MTU has been discovered for a network route, a separate host route is cloned for the path. These cloned host routes, as well as the path MTU value for the route, can be displayed using the netstat -r command. Accumulation of cloned routes can be avoided by allowing unused routes to expire and be deleted. Route expiration is controlled by the route_expire option of the no command. Route expiration is disabled by default.

Since routes can change dynamically, the path MTU value for a path may also change over time. Decreases in the path MTU value will result in packet fragmentation, so discovered path MTU values are periodically checked for decreases. By default, decreases are checked for every 10 minutes, and this value can be changed by modifying the value of the pmtu default age option of the no command.

Increases in the path MTU value can result in a potential increase in network performance, so discovered path MTU values are periodically checked for increases. By default, increases are checked for every 30 minutes, and this value can be changed by modifying the value of the pmtu rediscover interval option of the no command.

If not all of the routers in the network path support RFC 1191, then it may not be possible to determine an exact path MTU value. In these cases, the mmtu command can be used to add or delete path MTU values that are attempted.

Note:

- 1. Path MTU discovery cannot be used on duplicate routes, including those configured for group routing.
- 2. Enabling path MTU discovery will set the value of the arpgsize option of the no command to a minimum value of 5. This value will not be decreased if path MTU discovery is subsequently disabled.

SLIP

Configuring SLIP over a Modem

To configure Serial Line Interface Protocol (SLIP) between two systems that communicate through a modem, you can use the Web-based System Manager fast path wsm network or use the following procedure, which alternates between the System Management Interface Tool (SMIT) interface and the command line to complete the configuration. For clarity, the following instructions use the names bronze and gold for the two hosts.

- 1. Physically connect the modems to bronze and gold.
- 2. To create a tty on bronze using SMIT:
 - a. Enter:

```
smit maktty
```

- b. Select **rs232** as the type of tty you wish to create.
- c. Select an available serial port, for example **sa0** (system serial port 1).
- d. Select a port number for this tty from the list.
- e. Set the BAUD rate to the baud rate of your modem.
- Set Enable LOGIN to disable.
- g. Exit SMIT.
- Create a tty on gold.

Follow the same procedure as you did for bronze (in step 2 above), except set Enable LOGIN to enable.

The rest of these instructions assume that the tty number on both bronze and gold is tty1.

- 4. Test the physical connection with ATE.
 - a. On bronze, enter:

```
ate
```

- b. At the Unconnected Main Menu, select the Alter subcommand. Set the Rate to the baud rate of your modem and the Device to tty1.
- c. At the Unconnected Main Menu, select the Connect subcommand. When ATE prompts you for a phone number, enter the phone number of gold and press Enter.
- d. At this point, you should receive a login prompt for gold. Login.
- e. Finally, return to the connected screen, logout from gold, press Ctrl-v (to get to the ATE CONNECTED MAIN MENU), press t to terminate the connection, and press q to exit ATE.

Note: If you do not receive a login prompt, return to step 1 and verify that your configuration is correct. Do not proceed until you can login to gold.

Because the tty configuration for use with ATE is slightly different from the configuration for use with SLIP, you must make the following changes:

a. On bronze, enter:

```
smit chgtty
```

b. On gold, enter:

```
smit chatty-pdisable tty1
```

Select tty1, then select Change/Show TTY Program. Set Enable LOGIN to disable, then exit SMIT.

5. Add the following line to the /usr/lib/uucp/Devices file on both bronze and gold:

```
Direct tty1 - 9600 direct
```

or replace 9600 with whatever your modem speed is.

- 6. Create a SLIP network interface on bronze.
 - a. Enter:

```
smit mkinet1sl
```

- b. For TTY PORT for SLIP Network Interface, select tty1.
- c. Specify an INTERNET ADDRESS, for example, 130.130.130.1.
- d. Specify the DESTINATION address (of gold), for example, 130.130.130.2.
- e. Specify the BAUD RATE of your modem.
- f. Specify the DIAL STRING, for example:
 - "" AT OK ATDT555–1234 CONNECT ""
 - The meaning of this command is: Use tty1 at 9600 baud. Send AT to the modem. The modem should respond with OK. Dial the phone number 555-1234. The modern should respond with CONNECT. The spaces before and after the "" characters are necessary.
- g. Exit SMIT.
- 7. Create a SLIP network interface on gold.

Follow the same procedure as you did for bronze (in step 5 above), except exchange the INTERNET ADDRESS and the DESTINATION address.

8. Add the following two entries to the /etc/hosts file on both bronze and gold:

```
130.130.130.1
                bronze
130.130.130.2
                gold
```

The name you assign should be unique. In other words, if the Token-Ring interface on bronze is already assigned the name bronze, assign the SLIP interface a name such as bronze_slip.

Note: For a simplified interface to the **slattach** command, you may use the script /usr/sbin/slipcall.

- 9. Test the SLIP connection.
 - a. On bronze, enter:

```
ping gold
```

b. On gold, enter:

```
ping bronze
```

If both tests succeed, the SLIP connection is ready for use. If not, return to step 5 and verify that the configuration on both bronze and gold is correct.

Configuring SLIP over a Null Modem Cable

To configure SLIP between two systems that are attached using a null modem cable, you can use the Web-based System Manager fast path wsm network or use the following procedure, which alternates between the System Management Interface Tool (SMIT) interface and the command line to complete the configuration. For clarity, these instructions use the names bronze and gold for the two hosts.

- 1. Physically connect bronze and gold by the null modem cable. You will need the following cables. (The cables are listed in the order they will be connected from bronze to gold.)
 - a. Cable B (part number 00G0943 for example). Serial Port Jumper Cable; two are provided with each system, except some models that do not require them.
 - b. Cable D (part number 6323741, feature code 2936 for example). Asynchronous Cable EIA-232/V.24.
 - c. Cable E (part number 59F2861, feature code 2937 for example). Printer/Terminal Interposer EIA-232 (null modem cable).
 - d. Changer Adapter (both sides of the adapter are sockets).
- 2. Create a tty on bronze.
 - a. Enter:

```
smit maktty
```

- b. Select **rs232** as the type of tty you wish to create.
- c. Select an available serial port, for example **sa0** (system serial port 1).
- d. Select a port number for this tty from the list.
- e. Set the BAUD rate to 19200. (Later, you will change this to 38400. But for now, use 19200.)
- f. Set Enable LOGIN to disable, and then exit SMIT.
- 3. Create a tty on gold.

Follow the same steps as you did for bronze (in step 2 above), except set Enable LOGIN to enable.

Note: The rest of these instructions assume that the tty number on both bronze and gold is tty1.

- 4. Test the physical connection with ATE.
 - a. On bronze, enter:

ate

- b. At the Unconnected Main Menu, select the Alter subcommand. Set the Rate to 19200, and the Device to tty1.
- c. At the Unconnected Main Menu, select the Connect subcommand. When ATE prompts you for a phone number, press Enter. You should receive the message:

```
ate: 0828-010 The Connect command has made a connection
through port ttyl
```

- d. Press Enter. You should receive a login prompt for gold. Login to gold.
- e. Finally, return to the connected screen, logout from gold, press Ctrl-v (to get to the ATE CONNECTED MAIN MENU), press t to terminate the connection, and press q to exit ATE.

Note: If you do not receive a login prompt, return to step 1 and verify that your configuration is correct. Do not proceed until you can login to gold.

Since the tty configuration for use with ATE is slightly different from the configuration for use with SLIP, you must make the following changes:

a. On bronze, enter:

```
smit chatty
```

Select **tty1**. Set the BAUD rate to 38400, and then exit SMIT.

b. On gold, enter:

```
pdisable tty1
```

c. On gold, enter:

```
smit chatty
```

Select tty1. Set Enable LOGIN to disable, set the BAUD rate to 38400, and then exit SMIT.

5. Add the following line to the /usr/lib/uucp/Devices file on both bronze and gold:

```
Direct tty1 - 38400 direct
```

- 6. Create a SLIP network interface on bronze.
 - a. Enter:

```
smit mkinet1sl
```

- b. For TTY PORT for SLIP Network Interface, select tty1.
- c. Specify an INTERNET ADDRESS, for example 130.130.130.1.
- d. Specify the DESTINATION address (of gold), for example, 130.130.130.2, and then select OK or Enter.
- 7. Create a SLIP network interface on gold.

Follow the same procedure as you did for bronze (in step 5 above), except exchange the INTERNET ADDRESS and the DESTINATION address.

8. Add the following two entries to the /etc/hosts file on both bronze and gold:

```
130.130.130.1
                bronze
130.130.130.2
                gold
```

The name you assign should be unique. In other words, if the Token-Ring interface on bronze is already assigned the name bronze, assign the SLIP interface a name such as bronze_slip.

9. Start SLIP on both bronze and gold.

Enter:

```
slattach tty1
```

- 10. Test the SLIP connection.
 - a. On bronze, enter:

```
ping gold
```

b. On gold, enter:

```
ping bronze
```

If both tests succeed, the SLIP connection is ready for use. If not, return to step 5 and verify that the configuration on both bronze and gold is correct.

Deactivating a SLIP Connection

To deactivate a SLIP connection:

1. Enter:

```
ps -ef | grep slatt
```

Note the process numbers of processes associated with the **slattach** command.

2. For each process number, enter:

```
kill process_number
```

Do not use the -9 flag of the kill command.

If slattach is accidentally killed with a -9 flag, a slip lock may remain in /etc/locks. Delete this lock file to clean up after slattach.

Removing a TTY

To remove a tty, you can use the Web-based System Manager fast path, wsm network, or the System Management Interface Tool (SMIT) fast path, smit rminet.

Asynchronous Point-to-Point Protocol (PPP) Subsystem

The Asynchronous Point-to-Point Protocol (PPP) subsystem provides an alternative to SLIP. PPP provides a standard method for transporting multiprotocol datagrams over point-to-point media. PPP is comprised of three main layers:

- 1. A method for encapsulating multiprotocol datagrams. AIX PPP supports the TCP/IP network layer protocols.
- 2. A Link Control Protocol (LCP) for establishing, configuring, and testing the data-link connection. AIX PPP implements this through streams kernel extensions.
- 3. A family of Network Control Protocols (NCPs) for establishing and configuring different network layer protocols. AIX PPP supports Internet Protocol Control Protocol (IPCP) for negotiating a TCP/IP connection.

AIX PPP supports the following Request for Comments (RFCs):

- RFC 1661, "The Point-to-Point Protocol, LCP"
- RFC 1332, "The PPP Internet Protocol Control Protocol (IPCP)"
- RFC 1662, "PPP in HDLC-like Framing"
- RFC 1334, "PPP Authentication Protocols"
- RFC 1990, "PPP Multilink"

AIX PPP differentiates between client and server. An AIX system can act as both a client and a server. The distinction is made to simplify configuration. PPP servers tend to allocate a pool of IP addresses among the connections that are being made. There is some correlation between the media devices. AIX PPP breaks this correlation. All server PPP connections are allocated on a first-available basis. This facilitates the separation of PPP from the media. The attachment process must request to be linked to the proper type of link.

User-Level Processes

The Asynchronous Point-to-Point Protocol on AIX utilizes three user-level processes:

- 1. A control daemon (pppcontrold) executed by root under the System Resource Controller (startsrc -s pppcontrold). The control daemon's function encompasses loading and configuring all kernel extensions associated with the subsystem. It remains running as long as PPP function is required by the operating system.
- 2. An attachment process (pppattachd) that binds a TTY stream to an instance of the Link Control Protocol, Network Control Protocol, and a datagram protocol. An instance of pppattachd exists for each active PPP connection in the system. Any user of the attachment process must belong to the uucp group and contain /usr/sbin within their **PATH** environment variable.
- 3. A dialer process (pppdial) that establishes an outgoing connection. The dialer is intended to be executed by **pppattachd** as the connector program. Its purpose is to interact over the asynchronous device prior to PPP negotiation. This interaction is defined similarly to the UUCP chat dialog format. The dialer capability is provided to assist in establishing a connection with a remote system. The actual session establishment is out of the scope of PPP.

Configuring the Asynchronous Point-to-Point Protocol

You can use Web-based System Manager or SMIT to configure the Asynchronous Point-to-Point Protocol. The following table shows all tasks that you may need when configuring your system. You must have root privileges to perform the tasks in this table. At a minimum, when you initially configure your system, you will choose the following tasks from the table:

- · Add a Link Configuration
- Add a Server Interface (if you are setting up the machine as a PPP server)
- Add a Demand Interface (if you want the machine to support demand connections)
- Manipulate PAP or CHAP Users/Passwords (if you want the machine to support PPP authentication)
- Start PPP to effect your changes (or Stop then Start PPP, if PPP is currently running)

Configuring the Asynchronous PPP Tasks			
Web-based System Manager: wsm network fast path (wsm network application)			
-OR-			
Task	SMIT Fast Path		
Create Link Control Configuration	smit ppplcp		
Add a Link Configuration	smit addlcp		
Change/Show a Link Configuration	smit chglcp		
Remove a Link Configuration ¹	smit rmlcp		
Create PPP IP Interfaces	smit pppip		
Add a Server Interface	smit addpppserver		
Change/Show a Server Interface	smit listserver		
Remove a Server Interface ¹	smit rmlistserver		
Add a Demand Interface	smit addpppdemand		
Change/Show a Demand Interface	smit listdemand		
Remove a Demand Interface ¹	smit rmlistdemand		
Manipulate PAP users/passwords	smit ppppap		
Add a PAP User	smit addpapuser		
Change/Show a PAP User	smit listpapuser		
Remove a PAP User	smit rmpapuser		
Manipulate CHAP users/passwords	smit pppchap		
Add a CHAP User	smit addchapuser		
Change/Show a CHAP User	smit listchapuser		
Remove a CHAP User	smit rmchapuser		
Start PPP ²	smit startppp		
Stop PPP ³	smit stopppp		

Note:

- 1. Selecting this task destroys the existing information.
- 2. An alternative way to start PPP is to issue the **startsrc –s pppcontrold** command. However, the SMIT interface also allows you to set PPP to start at boot time.
- 3. An alternative way to stop PPP is to issue the **stopsrc -s pppcontrold** command. However, the SMIT interface also allows you to have PPP not start at boot time.

PPP and SNMP

PPP can interact with the TCP/IP SNMP daemon to report PPP link layer configuration information as well as information about active Link Control Protocol (LCP) interfaces. Providing that both the TCP/IP SNMP and the SNMP management software are configured correctly, PPP SNMP enables:

- retrieval of PPP Link Configuration information (Maximum Receive Unit size, Asynchronous Character Mapping, etc.)
- setting of PPP Link Configuration information
- retrieval of LCP interface information for active LCP links
- changing of the state of active LCP links can be changed to "down" by setting the appropriate ifAdminStatus Management Information Base (MIB) object

Not all objects defined by RFC1471 for the PPP MIB are supported. Only the pppLink table applies to the AIX PPP subsystem, thus the pppLqr and pppTests portions are not supported. The **pppLink** portion is supported with the following exceptions:

- The pppLinkConfigMagicNumber object is read only. In AIX PPP, magic number negotiation is always performed and cannot be disabled.
- The pppLinkConfigFcsSize object is read only. AIX PPP will only support FCS sizes of 16.

Enabling PPP SNMP

By default, SNMP for PPP is disabled. To enable PPP SNMP, you can use the Web-based System Manager fast path wsm network or use the following procedure. You must have root privileges to perform this procedure.

Note: The following procedure assumes that PPP Link Configuration is already set. If not, perform the procedure in "Configuring the Asynchronous Point-to-Point Protocol", on page 3-136 before enabling PPP SNMP.

1. Start the SMIT Interface and display the Change/Show a Link Configuration screen by entering:

```
smit chglcp
```

- 2. Toggle the Enable PPP SNMP subagent field to yes.
- 3. Accept your changes and exit SMIT.

PPP SNMP will not be enabled, however, until PPP is restarted.

- If PPP is currently running,
 - a. Stop PPP using the **smit stopppp** fast path (see the table in Configuring the Asynchronous Point-to-Point Protocol, on page 3-136).
 - b. Periodically check to see if the subsystem has completed shutdown by entering:

```
lssrc -s pppcontrold
```

The amount of time it takes to completely stop the subsystem is dependent on the number of links defined in the PPP configuration. The subsystem is completely shut down when the output of this command shows a status of inoperative.

- c. Start PPP using the **smit startppp** fast path (see the table in Configuring the Asynchronous Point-to-Point Protocol, on page 3-136).
- If PPP is not currently running, start PPP using the smit startppp fast path (see the table in Configuring the Asynchronous Point-to-Point Protocol, on page 3-136).

TCP/IP Quality of Service (QoS)

Quality of Service (QoS) is a family of evolving Internet standards that provides ways to give preferential treatment to certain types of IP traffic. With the proper support for QoS along a route, this can ameliorate the effects of variable queueing delays and congestion that contribute to poor network performance. AIX provides host support for QoS to classify outbound traffic into distinct classes of service and to announce and establish resource reservations as requested by client applications.

QoS can be used by an institution to deploy and enforce network policies governing the use of network bandwidth. With QoS, an AIX host can:

- Regulate the amount of traffic of a certain type injected into the network:
- Mark selected packets according to some policy so that subsequent routers can deliver the indicated service:
- Support services such as the virtual leased line service with proper QoS support along the route; and
- Participate in the resource reservation requests from receivers and announce sender sessions available for resource reservation requests.

The QoS support in AIX provides the following functions:

- Differentiated services as defined in RFC 2474
- Traffic policing
- In–profile and out–of–profile packet marking
- Traffic shaping
- Metering
- Integrated services for client and server applications as defined in RFC 1633
- RSVP signaling (RFC 2205)
- Guaranteed service (RFC 2212)
- Controlled-Load service (RFC 2211)
- Policy-based networking
- RAPI shared library for application

The QoS subsystem for AIX consists of four components:

QoS kernel extension (/usr/lib/drivers/gos)

The QoS kernel extension resides in /usr/lib/drivers/gos and is loaded and unloaded using the **cfgqos** and **ucfgqos** configuration methods. This kernel extension enables QoS support on the AIX host.

Policy agent (/usr/sbin/policyd)

The policy agent is a user-level daemon that resides in /usr/sbin/policyd. It provides support for policy management and interfaces with the QoS kernel extension to install, modify, and delete policy rules. Policy rules may be defined in the local configuration file (/etc/policyd.conf) or retrieved from a central network policy server using LDAP or both.

RSVP agent (/usr/sbin/rsvpd)

The RSVP agent is a user-level daemon that resides in /usr/sbin/rsvpd. It implements the RSVP signaling protocol semantics.

RAPI shared library (/usr/lib/librapi.a)

Applications may use the RSVP API (RAPI) in order to request enhanced quality of service as defined by the Integrated Services Internet QoS model. This library interacts with the local RSVP agent to propagate the QoS request along the path of the data flow using the RSVP protocol. This API is an open standard.

Note: The QoS support for AIX is based on a set of evolving Internet standards and draft standards currently under development by the Internet Engineering Task Force (IETF) and its various working groups. This technology will become more consistent and well defined as these standardization efforts progress within the IETF. It is also important to note that QoS is an emerging Internet technology that is just beginning to be deployed within the Internet. There are many benefits of QoS at all stages of deployment. However, true end-to-end services can only be realized when QoS support exists all along a particular route.

QoS Models

The QoS models for the Internet are open standards defined by the IETF. There are two Internet QoS models currently being standardized within the IETF: integrated services and differentiated services. These two Internet QoS models augment the traditional best-effort service model described in RFC 1812.

Integrated Services

Integrated Services (IS) is a dynamic resource reservation model for the Internet described in RFC 1633. Hosts use a signaling protocol called Resource ReSerVation Protocol (RSVP) to dynamically request a specific quality of service from the network. QoS parameters are carried in these RSVP messages and each network node along the path installs the parameters to obtain the requested quality of service. These QoS parameters describe one of two currently defined services, guaranteed service and controlled-load service. An important characteristic of IS is that this signaling is done for each traffic flow and reservations are installed at each hop along the route. Although this model is well-suited for meeting the dynamically changing needs of applications, there exist some significant scaling issues that imply it cannot be deployed in the network in which single routers handle many simultaneous flows.

Differentiated Services

Differentiated Services (DS) removes the per-flow and per-hop scalability issues, replacing them with a simplified mechanism of classifying packets. Rather than a dynamic signaling approach, DS uses bits in the IP type of service (TOS) byte to separate packets into classes. The particular bit pattern in the IP TOS byte is called the DS codepoint and is used by routers to define the quality of service delivered at that particular hop, in much the same way routers do IP forwarding using routing table lookups. The treatment given to a packet with a particular DS codepoint is called a per-hop behavior (PHB) and is administered independently at each network node. When the effects of these individual, independent PHBs are concatenated, this results in an end-to-end service.

Differentiated services is being standardized by an IETF working group, which has defined three PHBs: the Expedited Forwarding (EF) PHB, the Assured Forwarding (AF) PHB group, and the Default (DE) PHB. The EF PHB can be used to implement a low latency, low jitter, low loss, end-to-end service such as a virtual leased line (VLL). AF is a family of PHBs, called a PHB group, that is used to classify packets into various drop precedence levels. The drop precedence assigned to a packet determines the relative importance of a packet within the AF class. It can be used to implement the so-called Olympic service, which consists of three classes: bronze, silver, and gold. The DE PHB is the traditional best-effort service model as standardized in RFC 1812.

Supported Standards and Draft Standards

The following RFCs and Internet drafts describe the standards on which the AIX QoS implementation is based.

RFC 2474	Definition of the Differentiated Services Field (DS Field) in the IPv4 and IPv6 Headers
RFC 2475	An Architecture for Differentiated Services
RFC 1633	Integrated Services in the Internet Architecture: an Overview
RFC 2205	Resource ReSerVation Protocol (RSVP)
RFC 2210	The Use of RSVP with IETF Integrated Services
RFC 2211	Specification of the Controlled–Load Network Element Service
RFC 2212	Specification of Guaranteed Quality of Service
RFC 2215	General Characterization Parameters for Integrated Service Network Elements
draft-ietf-diffserv-framewor k-01.txt, October 1998	A Framework for Differentiated Services
draft-ietf-diffserv-rsvp-01.t xt, November 1998	A Framework for Use of RSVP with DIFF-serv Networks
draft-ietf-diffserv-phb-ef-0 1.txt	An Expedited Forwarding PHB
draft-ietf-diffserv-af-04.txt	Assured Forwarding PHB Group
draft-rajan-policy-qossche ma-00.txt, October 1998	Schema for Differentiated Services and Integrated Services in Networks
draft-ietf-rap-framework-0 1.txt, November 1998	A Framework for Policy-based Admission Control[25]
draft-ietf-rap-rsvp-ext-01.t xt, November 1998	RSVP Extensions for Policy Control

Note: QoS is an emerging Internet technology that is just beginning to be deployed within the Internet. There are many benefits of QoS at all stages of deployment. However, true end-to-end services can only be realized when QoS support exists all along a particular route.

QoS Installation

QoS for AIX is packaged with bos.net.tcp.server. This fileset must be installed in order to use QoS. To use the RAPI shared library, **bos.adt.include** must also be installed.

QoS Configuration

Stopping and Starting the AIX QoS Subsystem

QoS can be started or stopped through SMIT with the **smit qos** fast path or with the **mkqos** and rmgos commands.

To disable the QoS subsystem now and on the next system restart:

```
/usr/sbin/rmqos -B
```

To enable the QoS subsystem now only:

```
/usr/sbin/mkqos -N
```

See the command descriptions for mkqos and rmqos for the startup and removal command flags.

The policyd and rsvpd daemons are configured through the /etc/policyd.conf and /etc/rsvpd.conf configuration files, respectively. These configuration files must be edited to customize the QoS subsystem to the local environment. QoS will not work properly with the supplied example configurations.

Configuring the RSVP agent

The RSVP agent is required if the host is to support the RSVP protocol. The /etc/rsvpd.conf configuration file is used to configure the RSVP agent. The syntax of the configuration file is described in the sample configuration file installed in /etc/rsvpd.conf.

Example Configuration

```
interface 1.2.3.1
interface 1.2.3.2 disabled
interface 1.2.3.3 disabled
interface 1.2.3.4
     trafficControl
   }
rsvp 1.2.3.1
   {
    maxFlows 64
   }
rsvp 1.2.3.4
   {
     maxFlows 100
```

The above example illustrates a possible RSVP configuration in which the AIX hosts has 4 interfaces (virtual or physical) given by the 4 IP addresses, 1.2.3.1, 1.2.3.2, 1.2.3.3, and 1.2.3.4.

Interface 1.2.3.1 has been enabled for RSVP. However, traffic control has not been specified and incoming RSVP RESV messages do not cause resource reservation within the AIX TCP subsystem. This interface can support a maximum of 64 simultaneous RSVP sessions.

Interfaces 1.2.3.2 and 1.2.3.3 have been disabled. The RSVP agent cannot use this interface to transmit or receive RSVP messages.

Interface 1.2.3.4 has been enabled for RSVP. In addition, it can install resource reservations into the AIX TCP subsystem in response to an RSVP RESV message. This interface may support up to 100 RSVP sessions.

Any other interfaces present on the host but not mentioned explictly in /etc/rsvpd.conf will be disabled.

Configuring the Policy Agent

The policy agent is a required component of the AIX QoS subsystem. The /etc/policyd.conf configuration file is used to configure the policy agent. The syntax of this configuration file is described in the sample configuration file installed in /etc/policyd.conf.

Example Configurations

In the following example, a premium service category is created and used in the tcptraffic policy rule. This service category has a maximum rate of 110000 Kbps, a token bucket depth of 10000 bits, and an outgoing IP TOS value of 11100000 in binary. The tcptraffic policy rule gives this premimum service to all traffic with source IP address given by 1.2.3.6, destination address 1.2.3.3, and destination port in the range 0 to 1024. This rule is only in effect from 8:00am to 11:00pm in the local time zone.

```
ServiceCategories premium
    {
       PolicyScope DataTraffic MaxRate 110000
       MaxTokenBucket 10000
       OutgoingTOS 11100000
    }
ServicePolicyRules tcptraffic
    {
       PolicyScope DataTraffic
Direction Outgoing
Permission Allowed
       ProtocolNumber 6 # tcp
       TimeOfDayRange 8:00-23:00 SourceAddressRange 1.2.3.6-1.2.3.6
       DestinationAddressRange 1.2.3.3-1.2.3.3
       DestinationPortRange 0-1024
        ServiceReference premium
    }
```

The following statements set up a default service category and use it to restrict the UDP traffic flowing from interfaces 1.2.3.1 through 1.2.3.4 to IP addresses 1.2.3.6 through 1.2.3.10, port 8000.

```
ServiceCategories default
       PolicyScope DataTraffic MaxRate 110000
       MaxTokenBucket 10000
        OutgoingTOS 00000000
    }
ServicePolicyRules udptraffic
       PolicyScope DataTraffic
Direction Outgoing
Permission Allowed
       ProtocolNumber 17 # udp
        SourceAddressRange 1.2.3.1-1.2.3.4
        DestinationAddressRange 1.2.3.6-1.2.3.10
        DestinationPortRange 8000-8000
ServiceReference default
        ServiceReference
```

The example configuration below can be used to download rules from an LDAP server using the distinguished subtree name,

```
ou=NetworkPolicies, o=myhost.mydomain.com, c=us
```

to lookup the policies on the LDAP server host.

```
ReadFromDirectory
 LDAP_Server 1.2.3.27
 Base ou=NetworkPolicies,o=myhost.mydomain.com,c=us
}
```

QoS Problem Determination

The **qosstat** command may be used to display status information about the installed and active policies in the QoS subsystem. This information may be useful to you in determining where a problem exists if you are troubleshooting your QoS configuration. qosstat can be used to generate the following report.

```
Action:
  Token bucket rate (B/sec): 10240
  Token bucket depth (B): 1024
  Peak rate (B/sec): 10240
 Min policied unit (B): 20
 Max packet size (B): 1452
  Type: IS-CL
 Flags: 0x00001001 (POLICE, SHAPE)
  Statistics:
    Compliant packets: 1423 (440538 bytes)
  Conditions:
    Source address Dest address
                                           Protocol
    192.168.127.39:8000 192.168.256.29:35049 tcp (1connection)
Action:
 Token bucket rate (B/sec): 10240
 Token bucket depth (B): 1024
 Peak rate (B/sec): 10240
 Outgoing TOS (compliant): 0xc0
  Outgoing TOS (non-compliant): 0x00
  Flags: 0x00001011 (POLICE, MARK)
  Type: DS
  Statistics:
    Compliant packets: 335172 (20721355 bytes)
    Non-compliant packets: 5629 (187719 bytes)
  Conditions:
   Source address Dest address Protocol
192.168.127.39:80 *:* tcp (1 co
192.168.127.40:80 *:*
                                              tcp (1 connection)
```

tcp (5 connections)

QoS Reference

Commands

- qosstat
- mkqos
- rmqos

Methods

- cfgqos
- ucfgqos

TCP/IP Security

For any number of reasons, the person who administers your system may have to meet a certain level of security. For instance, the security level might be a matter of corporate policy. Or a system might need access to U.S. government systems and thus be required to communicate at a certain security level. These security standards might be applied to the network, the operating system, application software, even programs written by the person who administers your system.

This section describes the security features provided with Transmission Control Protocol/Internet Protocol (TCP/IP), both in standard mode and as a secure system, and discusses some security considerations that are appropriate in a network environment.

The topics discussed in this section are:

- Operating System–Specific Security, on page 3-145
- TCP/IP—Specific Security, on page 3-146
- TCP/IP Command Security, on page 3-147
- Trusted Processes, on page 3-149
- The Network Trusted Computing Base (NTCB), on page 3-150
- Data Security and Information Protection, on page 3-151

Operating System-Specific Security

Many of the security features available for TCP/IP are based on those available through the operating system. The following sections outline TCP/IP security.

Access Control

The security policy for networking is an extension of the security policy for the operating system, and it consists of the following major components:

- User authentication
- Connection authentication
- Data import and export security

User authentication is provided at the remote host by a user name and password, the same as when a user logs in to the local system. Trusted TCP/IP commands, such as ftp, rexec, and telnet, have the same requirements and go through the same verification process as trusted commands in the operating system.

Connection authentication is provided to ensure that the remote host has the expected Internet Protocol (IP) address and name. This prevents a remote host from masquerading as another remote host.

Data import and export security permits data at a specified security level to flow to and from network interface adapters at the same security and authority levels. For example, top secret data can flow only between adapters that are set to the top secret security level.

Auditing

Network auditing is provided by TCP/IP, using the audit subsystem to audit both kernel network routines and application programs. The purpose of auditing is to record those actions that affect the security of the system and the user responsible for those actions.

The following types of events are audited:

Kernel Events

- Change configuration
- Change host ID

- Change route
- Connection
- Create socket
- Export object
- Import object

Application Events

- Access the network
- Change configuration
- Change host ID
- Change static route
- Configure mail
- Connection
- Export data
- Import data
- · Write mail to a file

Creation and deletion of objects are audited by the operating system. Application audit records suspend and resume auditing to avoid redundant auditing by the kernel.

Network Trusted Computing Base (NTCB)

The Network Trusted Computing Base consists of hardware and software for ensuring network security. The hardware security features are provided by the network adapters used with TCP/IP. The software portion of the NTCB contains only trusted processes and their associated files.

Trusted Path, Trusted Shell, and Secure Attention Key (SAK)

The operating system provides the trusted path to prevent unauthorized programs from reading data from a user terminal. This path is used when a secure communication path with the system is required, such as when you are changing passwords or logging in to the system. The operating system also provides the trusted shell feature (tsh), which executes only trusted programs that have been tested and verified as secure. TCP/IP supports both of these features, along with the secure attention key (SAK), which establishes the environment necessary for secure communication between you and the system. The local SAK is available whenever you are using TCP/IP. A remote SAK is available through the telnet command.

The local SAK has the same function in **telnet** that it has in other operating system application programs: it terminates the telnet process and all other processes associated with the terminal in which telnet was running. Inside the telnet program, however, you can send a request for a trusted path to the remote system using the telnet send sak command (while in **telnet** command mode). You can also define a single key to initiate the SAK request using the telnet set sak command.

TCP/IP-Specific Security

Some portions of security are specific to TCP/IP. These features (TCP/IP commands and TCP/IP trusted processes) work together with the operating system security features discussed to provide the security for TCP/IP.

TCP/IP Command Security

Some commands in TCP/IP provide a secure environment during operation. These commands are **ftp**, **rexec**, and **telnet**. The **ftp** function provides security during file transfer. The **rexec** command provides a secure environment for executing commands on a foreign host. The **telnet** (TELNET) function provides security for login to a foreign host.

These commands provide security during their operation only. That is, they do not set up a secure environment for use with other commands. For securing your system for other operations, use the **securetcpip** command. This command gives you the ability to secure your system by disabling the nontrusted daemons and applications, and by giving you the option of securing your IP layer network protocol as well.

The **ftp**, **rexec**, **securetcpip**, and **telnet** commands provide the following forms of system and data security:

securetcpip

The **securetcpip** command enables TCP/IP security features. Access to commands that are not trusted is removed from the system when this command is issued. Each of the following commands are removed by running the **securetcpip** command:

- · rlogin and rlogind
- rcp, rsh, and rshd
- · tftp and tftpd
- trpt

The **securetcpip** command is used to convert a system from the standard level of security to a higher security level. Once your system has been converted, you do not need to issue the **securetcpip** command again unless you reinstall TCP/IP.

ftp

The **ftp** command provides a secure environment for transferring files. When a user invokes the **ftp** command to a foreign host, the user is prompted for a login ID. A default login ID is shown: the user's current login ID on the local host. The user is prompted for a password for the remote host.

The automatic login process searches the local user's **\$HOME**/.netrc file for the user's ID and password to use at the foreign host. For security, the permissions on the **\$HOME**/.netrc file must be set to 600 (read and write by owner only). Otherwise, automatic login fails.

Note: Since use of the **.netrc** file requires storage of passwords in a nonencrypted file, the automatic login feature of the **ftp** command is not available when your system has been configured with the **securetcpip** command. This feature can be reenabled by removing the **ftp** command from the tcpip: stanza in the **/etc/security/config** file.

To use the file transfer function, the **ftp** command requires two TCP/IP connections, one for the File Transfer Protocol (FTP) and one for data transfer. The protocol connection is primary and is secure because it is established on reliable communicating ports. The secondary connection is needed for the actual transfer of data, and both the local and remote host verify that the other end of this connection is established with the same host as the primary connection. If the primary and secondary connections are not established with the same host, the **ftp** command first displays an error message stating that the data connection was not authenticated, and then it exits. This verification of the secondary connection prevents a third host from intercepting data intended for another host.

rexec

The **rexec** command provides a secure environment for executing commands on a foreign host. The user is prompted for both a login ID and a password.

An automatic login feature causes the **rexec** command to search the local user's \$HOME/.netrc file for the user's ID and password on a foreign host. For security, the permissions on the **\$HOME**/.netrc file must be set to 600 (read and write by owner only). Otherwise, automatic login fails.

Note: Because use of the .netrc file requires storage of passwords in a nonencrypted file, the automatic login feature of rexec is not available when your system is operating in secure. This feature can be reenabled by removing the **rexec** entry form the tcpip: stanza in the /etc/security/config file.

telnet or tn

The **telnet** (TELNET) command provides a secure environment for login to a foreign host. The user is prompted for both a login ID and a password. The user's terminal is treated just like a terminal connected directly to the host. That is, access to the terminal is controlled by permission bits. Other users (group and other) do not have read access to the terminal, but they can write messages to it if the owner gives them write permission. The telnet command also provides access to a trusted shell on the remote system through the secure attention key (SAK). This key sequence differs from the sequence that invokes the local trusted path and can be defined within the telnet command.

Remote Command Execution Access (/etc/hosts.equiv)

Users on the hosts listed in the /etc/hosts.equiv file can run certain commands on your system without supplying a password.

Remote Command Execution Access Tasks				
Web-based System Manager: wsm network fast path (wsm network application) -OR-				
Task	SMIT Fast Path	Command or File		
List Remote Hosts That Have Command Execution Access	smit Ishostsequiv	view /etc/hosts.equiv		
Add a Remote Host for Command Execution Access **edit /etc/hosts.equiv				
Remove a Remote Host from Command Execution Access	smit rmhostsequiv	*edit /etc/hosts.equiv		

For more information about file procedures preceded by an asterisk (*), refer to the "hosts.equiv File Format for TCP/IP" in the AIX Files Reference.

Restricted File Transfer Program Users (/etc/ftpusers)

Users listed in the /etc/ftpusers file are protected from remote FTP access. For example, suppose user ross is logged into a remote system, and he knows the password of user carl on your system. If carl is listed in /etc/ftpusers, ross will not be able to FTP files to or from carl's account, even though ross knows carl's password.

Remote FTP Users Tasks				
Web-based System Manager: wsm network fast path (wsm network application)				
-OR-				
Task	SMIT Fast Path	Command or File		
List Restricted FTP Users smit Isftpusers view /etc/ftpusers				
Add a Restricted User smit mkftpusers *edit /etc/ftpusers				
Remove a Restricted User	smit rmftpusers	*edit /etc/ftpusers		

For more information about file procedures preceded by an asterisk (*), refer to the "ftpusers File Format for TCP/IP" in the AIX Files Reference.

Trusted Processes

A trusted program, or trusted process, is a shell script, a daemon, or a program that meets a particular standard of security. These security standards are set and maintained by the U.S. Department of Defense, which also certifies some trusted programs.

Trusted programs are trusted at different levels. Security levels include A1, B1, B2, B3, C1, C2, and D, with level A1 providing the highest security level. Each security level must meet certain requirements. For example, the C2 level of security incorporates the following standards:

program integrity Er	nsures that the process will do	what it is supposed to do.
----------------------	---------------------------------	----------------------------

no more and no less.

modularity Means that the process source code is broken down into

modules that cannot be directly affected or accessed by

other modules.

principle of least privilege States that at all times a user is operating at the lowest

level of privilege authorized. That is, if a user has access only to view a certain file, then the user does not

inadvertently also have access to alter that file.

Keeps a user from, for example, accidentally stumbling limitation of object reuse

> across a section of memory that has been flagged for overwriting but not yet cleared, and may contain sensitive

material.

TCP/IP contains several trusted daemons and many nontrusted daemons. The trusted daemons have been tested to ensure that they operate within particular security standards.

Examples of trusted daemons are:

- ftpd
- rexecd
- telnetd

Examples of nontrusted daemons are:

- rshd
- rlogind
- tftpd

For a system to be trusted, it must operate with a trusted computing base. This means, for a single host, that the machine must be secure. For a network, this means that all file servers, gateways, and other hosts must be secure.

Network Trusted Computing Base (NTCB)

The network contains both hardware and software mechanisms to implement the networking security features. This section defines the components of the Network Trusted Computing Base as they relate to TCP/IP.

The hardware security features for the network are provided by the network adapters used with TCP/IP. These adapters are programmed to control incoming data by receiving only data destined for the local system and to broadcast data receivable by all systems.

The software component of the NTCB consists of only those programs that are considered trusted. The programs and associated files that are part of a secure system are listed in the following tables on a directory-by-directory basis.

/etc Directory				
Name	Owner	Group	Mode	Permissions
gated.conf	root	system	0664	rw-rw-r
gateways	root	system	0664	rw-rw-r
hosts	root	system	0664	rw-rw-r
hosts.equiv	root	system	0664	rw-rw-r
inetd.conf	root	system	0644	rw-rr
named.conf	root	system	0644	rw-rr
named.data	root	system	0664	rw-rw-r
networks	root	system	0664	rw-rw-r
protocols	root	system	0644	rw-rr
rc.tcpip	root	system	0774	rwxrwxr—
resolv.conf	root	system	0644	rw-rw-r
services	root	system	0644	rw-rr
3270.keys	root	system	0664	rw-rw-r
3270keys.rt	root	system	0664	rw-rw-r

/usr/bin Directory				
Name	Owner	Group	Mode	Permissions
host	root	system	4555	r-sr-xr-x
hostid	bin	bin	0555	r–xr–xr–x
hostname	bin	bin	0555	r-xr-xr-x
finger	root	system	0755	rwxr-xr-x
ftp	root	system	4555	r-sr-xr-x
netstat	root	bin	4555	r-sr-xr-x
rexec	root	bin	4555	r-sr-xr-x
ruptime	root	system	4555	r-sr-xr-x
rwho	root	system	4555	r-sr-xr-x
talk	bin	bin	0555	r-xr-xr-x
telnet	root	system	4555	r-sr-xr-x

/usr/sbin Directory				
Name	Owner	Group	Mode	Permissions
arp	root	system	4555	r-sr-xr-x
fingerd	root	system	0554	r-xr-xr-
ftpd	root	system	4554	r-sr-xr-
gated	root	system	4554	r-sr-xr-
ifconfig	bin	bin	0555	r-xr-xr-x
inetd	root	system	4554	r-sr-xr-
named	root	system	4554	r-sr-x-
ping	root	system	4555	r-sr-xr-x
rexecd	root	system	4554	r-sr-xr-
route	root	system	4554	r-sr-xr-
routed	root	system	0554	r-xr-x
rwhod	root	system	4554	r-sr-xr-
securetcpip	root	system	0554	r-xr-xr-
setclock	root	system	4555	r-sr-xr-x
syslogd	root	system	0554	r-xr-xr-
talkd	root	system	4554	r-sr-xr-
telnetd	root	system	4554	r-sr-xr-

/usr/ucb Directory				
Name Owner Group Mode Permissions				
tn	root	system	4555	r-sr-xr-x

/var/spool/rwho Directory				
Name	Owner	Group	Mode	Permissions
rwho (directory)	root	system	0755	drwxr–xr–x

Data Security and Information Protection

The security feature for TCP/IP does not encrypt user data transmitted through the network. Therefore, it is suggested that users identify any risk in communication that could result in the disclosure of passwords and other sensitive information, and based on that risk, apply appropriate countermeasures.

The use of this product in a Department of Defense (DOD) environment may require adherence to DOD 5200.5 and NCSD-11 for communications security.

TCP/IP Problem Determination

This section contains information about diagnosing common problems in a Transmission Control Protocol/Internet Protocol (TCP/IP) network environment.

The **netstat** command is a good tool to use to help determine in which area the problem is. Once you have isolated the problem to an area, you can use more sophisticated tools to proceed. For example, you might use the **netstat -i** and **netstat -v** to determine if you have a problem with a particular hardware interface, and then run diagnostics to further isolate the problem. Or, if the **netstat -s** command shows that there are protocol errors, you could then use the trpt or iptrace commands.

The topics discussed in this section are:

- Communication Problems, on page 3-152
- Name Resolution Problems, on page 3-152
- Routing Problems, on page 3-154
- Problems with System Resource Controller (SRC) Support, on page 3-155
- telnet or rlogin Problems, on page 3-155
- Configuration Problems, on page 3-156
- Common Problems with Network Interfaces, on page 3-157
- Problems with Packet Delivery, on page 3-159
- Problems with Dynamic Host Configuration Protocol (DHCP), on page 3-160

Communication Problems

If you cannot communicate with a host on your network:

- Try to contact the host, using the ping command. Run the ping command on the local host to verify that the local interface to the network is up and running.
- Try to resolve the name of the host, using the host command. If the name does not resolve, you have a name resolution problem. See "Name Resolution Problems", on page 3-152 for more information.

If the name resolves and you are trying to contact a host on another network, you may have a routing problem. See "Routing Problems", on page 3-154 for more information.

- If your network is a token-ring network, check to see if the target host is on another ring. If so, the allcast field is probably set incorrectly. Use the Web-based System Manager fast path, wsm network, or the System Management Interface Tool (SMIT) fast path smit chinet to access the Network Interfaces menu. Then, set the Confine Broadcast to Local Ring field to **no** in the token-ring dialog.
- If there are a large number of Address Resolution Protocol (ARP) packets on your network, verify that your subnet mask is set correctly. This condition is known as a broadcast storm and may affect your system's performance.

Name Resolution Problems

Resolver routines on hosts running TCP/IP attempt to resolve names, using the following sources in the order listed:

- DOMAIN name server (named)
- 2. Network Information Service (NIS)
- 3. Local /etc/hosts file

When NIS+ is installed, lookup preferences are set using the irs.conf file. For more information, see AIX 4.3 NIS/NIS+ Guide.

Client Host

If you cannot get a host name resolved, and you are using flat name resolution (using the /etc/hosts file), verify that the host name and correct Internet Protocol (IP) address information is in the /etc/hosts file.

If you cannot get a host name resolved, and you are using a name server:

- 1. Verify that you have a **resolv.conf** file specifying the domain name and Internet address of a name server.
- 2. Verify that the local name server is up by issuing the **ping** command with the IP address of the name server (found in the local resolv.conf file).
- 3. If the local name server is up, verify that the named daemon on your local name server is active by issuing the **Issrc -s named** command on the name server.
- 4. If you are running the **syslogd**, there could be error messages logged. The output for these messages is defined in the /etc/syslog.conf file.

If these steps do not identify the problem, start looking at the name server host.

Name Server Host

If you cannot get a host name resolved:

1. Verify that the **named** daemon is active by issuing the following command:

```
lssrc -s named
```

2. Verify that the address of the target host exists and is correct in the name server's database. Send a SIGINT signal to the **named** daemon to dump the database and cache to the file /var/tmp/named dump.db. Verify that the address you are trying to resolve is there and is correct.

Add or correct name-to-address resolution information in the named hosts data file for the master name server of the domain. Then issue the following SRC command to reread the data files:

```
refresh -s named
```

3. Verify that the name resolution requests are being processed. To do this, enter the named daemon from the command line and specify a debugging level. Valid debug levels are 1 through 9. The higher the level, the more information the debug mechanism logs.

```
startsrc -s named -a "-d DebugLevel"
```

4. Check for configuration problems in the **named** data files. For more information, see "Configuring Name Servers", on page 3-105, the "DOMAIN Data File Format," "DOMAIN Reverse Data File Format," "DOMAIN Cache File Format," and the "DOMAIN Local Data File Format" in the AIX Files Reference.

Note: A common error is the incorrect use of the . (period) and the @ (at sign) in the DOMAIN data files.

If external users cannot reach your domains:

 Make sure that all your non-master name servers (slave, hint) have equal time-to-live (TTL) information in the DOMAIN data files.

If external resolvers query your servers constantly:

Make sure your servers are distributing DOMAIN data files with reasonable TTL values. If the TTL is zero or another small value, the data you transfer times out very quickly. Set the minimum value in your start of authority (SOA) records to a week or more to solve this problem.

Routing Problems

If you cannot reach a destination host, consider the following situations:

- If you receive a Network Unreachable error message, make sure that a route to the gateway host has been defined and is correct. Check this by using the netstat -r command to list kernel routing tables.
- If you receive a No route to host error message, verify that the local network interface is up by issuing the ifconfig interface name command. The output should indicate that the interface is up. Use the ping command to try and reach another host on your network.
- If you receive a Connection timed out error message:
 - Verify that the local gateway is up using the **ping** command with the name or Internet address of the gateway.
 - Make sure that a route to the gateway host has been defined and is correct. Check this by using the **netstat -r** command to list kernel routing tables.
 - Make sure the host you want to communicate with has a routing table entry back to your machine.
- If you are using static routing, make sure that a route to the target host and gateway host has been defined. Check this by using the **netstat** -r command to list kernel routing tables.

Note: Make sure the host you want to communicate with has a routing table entry to your machine.

- · If you are using dynamic routing, verify that the gateway is listed and correct in the kernel routing tables by issuing the **netstat –r** command.
- If the gateway host is using the Routing Information Protocol (RIP) with the routed daemon, make sure that a static route to the target host is set up in the /etc/gateways file.

Note: You need to do this only if the routing daemon cannot identify the route to a distant host through queries to other gateways.

- If the gateway host is using the RIP with the gated daemon, make sure that a static route to the target host is set up in the gated.conf file.
- If you are using dynamic routing with the **routed** daemon:
 - If **routed** cannot identify the route through queries (for example, if the target host is not running the RIP, check the /etc/gateways file to verify that a route to the target host is defined.
 - Make sure that gateways responsible for forwarding packets to the host are up and running the RIP. Otherwise, you'll need to define a static route.
 - Run the **routed** daemon using the debug option to log such information as bad packets received. Invoke the daemon from the command line using the following command:

```
startsrc -s routed -a "-d"
```

- Run the **routed** daemon using the -t flag, which causes all packets sent or received to be written to standard output. When routed is run in this mode, it remains under the control of the terminal that started it. Therefore, an interrupt from the controlling terminal kills the daemon.
- If you are using dynamic routing with the gated daemon:
 - Verify that the /etc/gated.conf file is configured correctly and that you are running the correct protocols.

- Make sure the gateway on the source network is using the same protocol as the gateway on the destination network.
- Make sure that the machine with which you are trying to communicate has a route back to your host machine.
- Verify that the gateway names in the gated.conf file correspond to the gateway names listed in the /etc/networks file.
- If you are using the RIP or HELLO protocols, and routes to the destination cannot be identified through routing queries, check the gated.conf file to verify that a route to the target host is defined. You should set static routes under the following conditions:
 - The destination host is not running the same protocol as the source host so cannot exchange routing information.
 - The host must be reached by a distant gateway (a gateway that is on a different autonomous system than the source host). The RIP can be used only among hosts on the same autonomous system.

Other Possibilities

If all else fails, you may want to turn on tracing for your routing daemon (either routed or gated). Use the SRC traceson command from the command line, or send a signal to the daemon to specify different levels of tracing. See the gated daemon or the routed daemon for specifics on sending signals to these daemons.

Problems with SRC Support

• If changes to the /etc/inetd.conf file do not take effect:

Update the inetd daemon by issuing the refresh -s inetd command or the kill -1 InetdPID command.

• If the **startsrc –s** [*subsystem name*] returns the following error message:

```
0513-00 The System Resource Controller is not active.
```

The System Resource Controller subsystem has not been activated. Issue the srcmstr & command to start SRC, then reissue the **startsrc** command.

You may also want to try starting the daemon from the command line without SRC support.

• If the refresh -s [subsystem name] or Issrc -Is [subsystem name] returns the following error message:

```
[subsystem name] does not support this option.
```

The subsystem does not support the SRC option issued. Check the subsystem documentation to verify options the subsystem supports.

• If the following message is displayed:

```
SRC was not found, continuing without SRC support.
```

A daemon was invoked directly from the command line instead of using the **startsrc** command. This is not a problem. However, SRC commands, such as stopsrc and refresh, will not manipulate a subsystem that is invoked directly.

telnet or rlogin Problems

The following explanations may be useful in solving problems with the telnet or rlogin command.

Screen Distortion

If you are having trouble with screen distortion in full-screen applications:

1. Check the **TERM** environment variable by issuing one of the following commands:

OR

echo \$TERM

2. Verify that the **TERM** variable is set to a value that matches the type of terminal display you are using.

telnet Debugging

telnet subcommands that may help in debugging problems include:

display Displays set and toggle values.

toggle Toggles the display of all network data in hex.

Toggles the display of internal telnet process options. toggle options

Programs Using Extended Curses

Problems with function and arrow keys may arise when using the rlogin and telnet commands with programs using extended curses. Function and arrow keys generate escape sequences, which are split if too little time is allotted for the entire key sequence. Curses waits a specific amount of time to decide whether an Esc indicates the escape key only or the start of a multibyte escape sequence generated by other keys, such as cursor keys, the action key, and function keys.

If no data, or data that is not valid, follows the Esc in the allotted amount of time, curses decides that the Esc is the escape key, and the key sequence is split. The delay resulting from the **rlogin** or **telnet** command is network dependent. Sometimes arrow and function keys work and sometimes they do not, depending on the speed of the network to which you are connecting. Setting the ESCDELAY environment variable to a large value (1000 to 1500) effectively solves this problem.

Configuration Problems

Network interfaces are automatically configured during the first system startup after the adapter card is installed. However, you still need to set some initial values for TCP/IP including the host name, the Internet address, and the subnet mask. To do this, you can use the Web-based System Manager fast path, wsm network, or you can use the SMIT interface in the following ways:

- Use the smit mktcpip fast path to set the initial values for the host name, the Internet address, and the subnet mask.
- Use the **smit mktcpip** fast path to specify a name server to provide name resolution service. (Note that **smit mktcpip** configures one network interface only.)
- Use the smit chinet fast path to set other network attributes.

You may also want to set up any static routes the host needs for sending transmitting information, such as a route to the local gateway. Use the Web-based System Manager fast path, wsm network, or the SMIT fast path, smit mkroute, to set these up permanently in the configuration database.

If you are having other problems with your configuration, see the "Configuring a TCP/IP Network Checklist", on page 3-4 for more information.

Common Problems with Network Interfaces

Network interfaces are configured automatically during the first system startup after the adapter card is installed. However, there are certain values that must be set in order for TCP/IP to start. These include the host name and Internet address and can be set using the Web-based System Manager fast path, wsm network, or the SMIT fast path. smit mktcpip.

If you choose the SMIT method, use the **smit mktcpip** fast path to set these values permanently in the configuration database. Use the smit chinet and smit hostname fast paths to change them in a running system. The **smit mktcpip** fast path minimally configures TCP/IP. To add adapters, use the Further Configuration menu, which can be reached with the **smit tcpip** fast path.

If you have already checked these to verify accuracy and you are still having trouble sending and receiving information, check the following:

- Verify that your network adapter has a network interface by executing the netstat -i command. The output should list an interface, such as tr0, in the Name column. If it does not, create a network interface through Web-based System Manager or by entering the SMIT fast path **smit mkinet**.
- Verify that IP address for the interface is correct by executing the netstat -i command. The output should list the IP address in the Network column. If it is incorrect, set the IP address through Web-based System Manager or by entering the SMIT fast path smit chinet.
- Check the error log by running the **errpt** command to ensure that there are no adapter problems.
- · Verify that the adapter card is good by running diagnostics.

If these steps do not identify the problem, refer to "Problems with a SLIP Network Interface", on page 3-157, "Problems with an Ethernet Network Interface", on page 3-158, or "Problems with a Token-Ring Network Interface", on page 3-158.

Problems with a SLIP Network Interface

In general, the most effective method for debugging problems with a Serial Line Interface Protocol (SLIP) interface is to retrace your configuration, verifying each step. However, you can also:

- Verify that the slattach process is running and using the correct tty port by issuing the ps -ef command. If it is not, run the slattach command. (See "Configuring SLIP over a Modem", on page 3-131, or "Configuring SLIP over a Null Modem Cable", on page 3-133 for the exact syntax you should use.)
- Verify that the point-to-point addresses are specified correctly by entering the smit chinet fast path.

Select the SLIP interface. Make sure that the INTERNET ADDRESS and DESTINATION Address fields are correct.

If the modem is not functioning correctly:

- Make sure that the modem was installed properly. See the modem's operator manual.
- Verify that any flow control the modem does is turned off.

If the tty is not functioning properly, verify that the tty's baud rate and modem characteristics are set correctly in the configuration database by entering the **smit tty** fast path.

Problems with an Ethernet Network Interface

If the network interface has been initialized, the addresses correctly specified, and you have verified that the adapter card is good:

- Verify that you are using a T-connector plugged directly into the inboard/outboard transceiver.
- Make sure you are using Ethernet cable. (Ethernet cable is 50 OHM.)
- Make sure you are using Ethernet terminators. (Ethernet terminators are 50 OHM.)
- Ethernet adapters may be used with either the transceiver that is on the card or with an external transceiver. There is a jumper on the adapter to specify which you are using. Verify that your jumper is set correctly (see your adapter manual for instructions).
- Verify that you are using the proper Ethernet connector type (thin is BNC; thick is DIX). If you change this connector type, use the Web-based System Manager fast path, wsm devices, or the SMIT fast path, smit chgenet, to set the Apply Change to Database Only field. (The field should be checked in Web-based System Manager or set to yes in SMIT.) Reboot the machine to apply the configuration change. (See "Configuring and Managing Adapters", on page 3-37.)

Problems with a Token-Ring Network Interface

If you cannot communicate with some of the machines on your network although the network interface has been initialized, the addresses correctly specified, and you have verified that the adapter card is good:

- Check to see if the hosts with whom you cannot communicate are on a different ring. If they are, use the Web-based System Manager fast path, wsm devices, or the SMIT fast path smit chinet to check the Confine BROADCAST to Local Token-Ring field. The field should not be checked in Web-based System Manager or set to no in SMIT.
- Check to see whether the token-ring adapter is configured to run at the correct ring speed. If it is configured incorrectly, use the Web-based System Manager wsm network application or SMIT to change the adapter ring speed attribute (see "Configuring a High-Performance Token-Ring Adapter", on page 3-37). When TCP/IP is restarted, the token-ring adapter will have the same ring speed as the rest of the network.

Problems with a Token–Ring/Ethernet Bridge

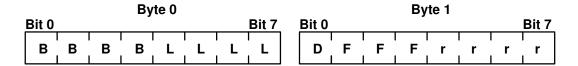
If you cannot communicate between a token-ring and an Ethernet network, using a bridge, and you have verified that the bridge is functioning properly, the Ethernet adapter may be dropping packets. A machine drops packets if the incoming packet (including headers) is greater than the network adapter's maximum transmission unit (MTU) value. For instance, a 1500-byte packet sent by a token-ring adapter over the bridge collects an 8-byte logical link control (LLC) header, making the total packet size 1508. If the receiving Ethernet adapter's MTU is set to 1500, the packet is dropped.

Check the MTU values of both network adapters. To allow for the eight-byte LLL header, the token-ring adapter attaches to outgoing packets, the MTU value for the token-ring adapter should be set at least eight bytes lower than the MTU value for the Ethernet adapter. For example, the MTU for a token-ring adapter should be set to 1492 to communicate with an Ethernet adapter with an MTU of 1500.

Problems with a Token-Ring/Token-Ring Bridge

When operating through a bridge, the default value of 1500 for the maximum transmission unit (MTU) should be changed to a value that is eight less than the maximum information field (maximum I-frame) advertised by the bridge in the routing control field.

To find the routing control field value, use the **iptrace** daemon to look at incoming packets. Bits 1, 2, and 3 of Byte 1 are the Largest Frame Bits, which specify the maximum information field that can be transmitted between two communicating stations on a specific route. Refer to the following figure for the format of the routing control field:



Routing Control Fields

Values for the Largest Frame Bits are as follows:

000	Specifies a maximum of 516 bytes in the information field.
001	Specifies a maximum of 1500 bytes in the information field.
010	Specifies a maximum of 2052 bytes in the information field.
011	Specifies a maximum of 4472 bytes in the information field.
100	Specifies a maximum of 8144 bytes in the information field.
101	Reserved.
110	Reserved.
111	Used in all-routes broadcast frames.

For example, if the maximum I-frame value is 2052 in the routing control field, the MTU size should be set to 2044. This is for token-ring network interfaces only.

Note: When using **iptrace**, the output file must *not* be on a Network File System (NFS).

Problems with Packet Delivery

Communicating with a Remote Host

If you cannot communicate with a remote host, try the following:

- Run the ping command on the local host to verify that the local interface to the network is up and running.
- Use the **ping** command for hosts and gateways that are progressively more hops from the local host to determine the point at which communication fails.

If you are having trouble with packet loss or are experiencing delays in packet delivery, try the following:

- Use the trpt command to trace packets at the socket level.
- Use the **iptrace** command to trace all protocol layers.

If you cannot communicate between a token-ring and an Ethernet network using a bridge, and you have verified that the bridge is good:

 Check the MTU values of both adapters. The MTU values should be compatible to allow communication. A machine drops packets if the incoming packet (including headers) is greater than the adapter's MTU values. For instance, a 1500-byte packet sent over the bridge collects an 8-byte LLC header, making the total packet size 1508. If the receiving machine's MTU is set to 1500, a packet of 1508 bytes is dropped.

snmpd Response to Queries

If **snmpd** is not responding to queries and there are no log messages received, the packet may be to large for the kernel User Datagram Protocol (UDP) packet handler. If this is the case, increase the kernel variables, **udp_sendspace** and **udp_recvspace** by issuing the following commands:

```
no -o udp_sendspace=64000
no -o udp_recvspace=64000
```

The maximum size for a UPD packet is 64K. If your query is larger than 64K, it will be rejected. The packet should be split into smaller packets to avoid this problem.

Problems with Dynamic Host Configuration Protocol (DHCP)

If you cannot get an IP address or other configuration parameters:

- Check to see that you have specified an interface to be configured. This can be done through the Web-based System Manager wsm network application, by editing the /etc/dhcpcd.ini file, or by using the SMIT fast path smit dhcp.
- Check to see that there is a server on the local network or a relay agent configured to get your requests off the local network.
- Check to see that the **dhcpcd** program is running. If it is not, use the **startsrc –s dhcpcd** command.

TCP/IP Reference

Transmission Control Protocol/Internet Protocol (TCP/IP) topics discussed in this section are:

- List of TCP/IP Commands, on page 3-161
- List of TCP/IP Daemons, on page 3-162
- List of Methods, on page 3-162
- List of TCP/IP Files, on page 3-163
- List of RFCs, on page 3-163
- Getting RFCs, on page 3-163

List of TCP/IP Commands

chnamsv Changes Transmission Control Protocol/Internet Protocol

(TCP/IP) based name service configuration on a host.

chprtsv Changes a print service configuration on a client or server

machine.

hostent Directly manipulates address—mapping entries in the

system configuration database.

ifconfig Configures or displays network interface parameters for a

network, using TCP/IP.

mknamsv Configures TCP/IP—based name service on a host for a

client.

mkprtsv Configures TCP/IP—based print service on a host.

mktcpip Sets the required values for starting TCP/IP on a host.

no Configures network options.

rmnamsvUnconfigures TCP/IP-based name service on a host.rmprtsvUnconfigures a print service on a client or server machine.

slattach Attaches serial lines as network interfaces.

arp Displays or changes the Internet address to hardware

address translation tables used by the Address Resolution

Protocol (ARP).

gettable Gets Network Information Center (NIC) format host tables

from a host.

hostidSets or displays the identifier of the current local host.hostnameSets or displays the name of the current host system.htableConverts host files to the format used by network library

routines.

ipreport Generates a packet trace report from the specified packet

trace file.

iptrace Provides interface—level packet tracing for Internet

protocols.

Isnamsv Shows name service information stored in the database. **Isprtsv** Shows print service information stored in the database.

mkhosts Generates the host table file.

namerslv Directly manipulates domain name server entries for local

resolver routines in the system configuration database.

netstat Shows network status.

route Manually manipulates the routing tables.

Directly manipulates entries in three separate system ruser

databases that control foreign host access to programs.

ruptime Displays the status of each host on a network.

Enables the network security feature. securetcpip

setclock Sets the time and date for a host on a network. timedc Returns information about the **timed** daemon.

Performs protocol tracing on Transmission Control Protocol trpt

(TCP) sockets.

List of TCP/IP Daemons

fingerd Provides remote user information.

ftpd Provides the server function for the Internet File Transfer

Protocol (FTP) protocol.

gatedn Provides gateway routing functions for the Routing

Information Protocol (RIP), Hello Protocol (HELLO),

Exterior Gateway Protocol (EGP), Border Gateway Protocol (BGP), and Simple Network Management Protocol (SNMP).

inetd Provides Internet service management for a network.

named Provides the server function for the Domain Name Protocol

(DOMAIN).

rexecd Provides the server function for the **rexec** command. rlogind Provides the server function for the **rlogin** command.

routed Manages network routing tables.

rshd Provides the server function for remote command

execution.

rwhod Provides the server function for the rwho and ruptime

commands.

syslogd Reads and logs system messages.

talkd Provides the server function for the talk command. telnetd Provides the server function for the TELNET protocol. tftpd Provides the server function for the Trivial File Transfer

Protocol (TFTP).

timed_n Invokes the **timeserver** daemon at system startup time.

SMIT Fast Paths for TCP/IP

The Summary of Fast Paths in AIX 4.3 System Management Guide: Operating System and Devices provides the fast paths for TCP/IP system management tasks.

List of Methods

Device methods are programs associated with a device that perform basic device configuration operations. See "List of TCP/IP Programming References" in AIX Communications Programming Concepts for information about TCP/IP methods.

List of TCP/IP Files

/etc/rc.bsdnet

See "List of TCP/IP Programming References" in AIX Communications Programming Concepts for information about TCP/IP files and file formats.

List of RFCs

For a list of the RFCs (Request for Comments) supported by AIX, see the "List of TCP/IP Programming References" in AIX Communications Programming Concepts.

- RFC 1359 "Connecting to the Internet: What connecting institutions should anticipate"
- RFC 1325 "FYI on questions and answers: Answers to commonly asked 'new Internet user' questions"
- RFC 1244 "Site Security Handbook"
- RFC 1178 "Choosing a Name for Your Computer"
- RFC 1173 "Responsibilities of host and network managers: A summary of the 'oral tradition' of the Internet"

Getting RFCs

Many RFCs are available online. Paper copies of all RFCs are available from SRI, either individually or on a subscription basis. For more information, contact nisc@nisc.sri.com, or call 1-415-859-6387.

Online copies are available using FTP from ftp.nisc.sri.com as rfc/rfcnnnn.txt or rfc/rfcnnnn.ps (where nnnn is the RFC number without leading zeros). Additionally, RFCs may be requested through electronic mail from SRI's automated mail server by sending a message to mail-server@nisc.sri.com. In the body of the message, indicate the RFC to be sent, for example, send rfcnnnn (where nnnn is the number of the RFC). For PostScript RFCs, specify the extension, for example, send rfcnnnn.ps. Multiple requests can be sent in a single message by specifying each request on a separate line. The RFC index can be requested by typing send rfc-index.

Dynamic Load API

AIX supports name resolution from five different maps:

- Domain Name Server (DNS),
- Network Information Server (NIS),
- NIS+,
- local methods of name resolution, and
- · dynamically loaded, user-defined APIs.

With the Dynamic Load Application Programming Interface (API) you can load your own modules to provide routines that supplement the maps provided by AIX. The Dynamic Load API enables you to create dynamically loading APIs in any of five map classes: hosts, services, protocols, networks, and netgroup. You can build your own user modules containing APIs for any or all of the map classes. The following sections define an API's function names and prototypes for each of the five classes. See Configuring a Dynamic API for information about configuring a dynamically loading API with AIX.

Function Names and Prototypes

AIX supports five map classes for user—loadable modules: services, protocols, hosts, networks and netgroup. To instantiate each map accesssor, AIX requires that a user provided module must use the following AIX specified function names and function prototypes for each map class.

Services Map Type

The following is the required prototype for a user-defined services map class:

```
void *sv_pvtinit();
void sv_close(void *private);
struct servent * sv_byname(void *private, const char *name, const char
*proto);
struct servent * sv_byport(void *private, int port, const char *proto);
struct servent * sv_next(void *private);
void sv_rewind(void *private);
void sv_minimize(void *private);
```

Function *sv_pvtinit* is always required, though it is not required to return anything more than NULL (that is, the function can do nothing more than return NULL if the calling routine does not need private data), but the routine must exist.

Functions other than *sv_pvtinit* in a user–defined module for services map class are optional. It is permissible for a user to provide none or only part of the optional functions definition.

Protocols Map Type

The following is the required prototype for a user-defined protocols map class:

```
void * pr_pvtinit();
void pr_close(void *private);
struct protoent * pr_byname(void *private, const char *name);
struct protoent * pr_bynumber(void *private, int num);
struct protoent * pr_next(void *private);
void pr_rewind(void *private);
void pr_minimize(void *private);
```

Function *pr_pvtinit* is always required. The user–defined modules need not return anything more than NULL (that is, the function can do nothing more than return NULL if the calling routine does not need private data), but the routine must exist.

Functions other than $pr_pvtinit$ in the user module for the protocols map class are optional. It is permissible for a user to provide none or only part of the optional functions definition.

Hosts Map Type

The following is the required prototype for a user-defined hosts map class:

```
void * ho_pvtinit();
void ho_close(void *private);
struct hostent * ho_byname(void *private, const char *name);
struct hostent * ho_byname2(void *private, const char *name, int af);
struct hostent * ho_byaddr(void *private, const void *addr, size_t len, int
struct hostent * ho_next(void *private);
void ho_rewind(void *private);
void ho_minimize(void *private);
```

Function ho pvtinit is always required. The user-defined function is not required to return anything more than NULL (that is, the function can do nothing more than return NULL if the calling routine does not need private data), but the routine must exist.

Functions other than ho pvtinit in the user-defined module for hosts map class are optional. It is permissible for a user to provide none or only part of the optional functions definition.

Networks Map Type

The following is the required prototype for a user-defined networks map class:

```
void * nw_pvtinit();
void nw_close(void *private);
struct nwent * nw_byname(void *private, const char *name, int addrtype);
struct nwent * nw_byaddr(void *private, void *net, int length, int
addrtype);
struct nwent * nw_next(void *private);
void nw_rewind(void *private);
void nw_minimize(void *private);
```

Function nw pvtinit is always required. The function is not required to return anything more than NULL (that is, the function can do nothing more than return NULL if the calling routine does not need private data), but the routine must exist.

Functions other than *nw_pvtinit* in the user module for the networks map class are optional. It is permissible for a user to provide none or part of the optional functions definition.

Data Structure

AIX provides a data structure required to implement the networks map class, which uses this structure to communicate with the operating system.

```
struct nwent {
                              /* official name of net */
        char *name;
        char **n_aliases; /* alias list */
        int n_addrtype;  /* net address type */
void *n_addr;  /* network address */
int n_length;  /* address length, in bits */
 };
```

Netgroup Map Type

The following is the required prototype for a user-defined netgroup map class:

```
void * ng_pvtinit();
void ng_rewind(void *private, const char *group);
void ng_close(void *private);
int ng_next(void *private, char **host, char **user, char **domain);
int ng_test (void *private, const char *name, const char *host, const char
*user, const char *domain);
void ng_minimize(void *private);
```

Function ng_pvtinit is always required. The user is not required to return anything more than NULL (that is, the function can do nothing more than return NULL if the calling routine does not need private data), but the routine must exist.

Functions other than ng pvtinit in the user module for the netgroup map class are optional. It is permissible for a user to provide none or part of the optional functions definition.

Using the Dynamic Load API

You need to name your user-defined module according to a pre-established convention, and you also need to configure it into the operating system before it will work. The following sections explain API module naming and configuration.

Naming the User-Provided Module

The names of modules containing user-defined APIs follow this general form:

NameAddressfamily

where Name is the name of the dynamic loadable module name. The length of the Name can be between one to eight characters.

The following key words are reserved as the user option name and may not be used as the name of the dynamically loadable module:

- local
- bind
- dns
- nis
- Idap

Addressfamily, which represents the address family, can be either 4 or 6. If no number is specified, address family is AF UNSPEC. If the number is 4, address family is AF INET. If the number is 6, address family is AF INET6.

Any other format for user options is not valid.

If user called *gethostbyname2* system call in the application, whatever the address family the user passed to the gethostbyname2 system call will overwrite the address family in the user option. (for example, we have a user option david6, in the application, there is a system call gethostbyname2(name, AF_INET). Then the address family AF_INET will overwrite the user option's address family (6, same as AF INET6).

Configuring a Dynamic API

There are three ways that a user can specify user-provided, dynamically loading resolver routines:

NSORDER Environment variable

You can specify zero or more user options in the environment variable NSORDER. You are not restricted in the number of options that you can enter, nor in the sequence in which they are entered. For example, on the command line, type:

```
export NSORDER=local, bind, bob, nis, david4, jason6
```

In this example, AIX will invoke the listed name resolution modules left to right until the name is resolved. The modules named local, bind, and nis are reserved by AIX:, but bob, david4, and jason6 are user-provided modules.

/etc/netsvc.conf Configuration File

You can specify zero or more user options in the configuration file /etc/netsvc.conf. You are not restricted in the number of options that you can enter, nor in the sequence in which they are entered. For example:

```
hosts=nis, jason4, david, local, bob6, bind
```

/etc/irs.conf Configuration File

You can specify zero or more user options in the configuration file /etc/irs.conf. You are not restricted in the number of options that you can enter, nor in the sequence in which they are entered. For example:

```
hosts dns continue
hosts jason6 merge
hosts david4
```

The maximum number of user modules a user can specify from any one of the above source is 16.

Environment variable NSORDER has the highest priority, next is configuration file /etc/netsve.conf, the configuration file /etc/irs.conf has the lowest priority. If there is a user option specified in a higher priority source (for example, NSORDER), then any other user options specified in the lower priority sources (for example, /etc/netsvc.conf and /etc/irs.conf) will be ignored.

Procedures

To create and install a module containing a dynamically loading API, follow these steps:

- 1. First a user has to create the dynamic loadable module based upon AIX specifications.
- 2. The user must also create an export file (for example, rnd.exp) which exports all the symbols to be used.
- 3. AIX provides the sample Makefile for user to create a dynamic loadable module file. (for example, rnd.so file). The sample Makefile, sample export file and sample user module file are located at /usr/samples/tcpip/dynload.
- 4. After compilation, user must put all the dynamic loadable object files under the file path /usr/lib/netsvc/dynload.
- 5. Next a user needs to configure
 - Environment variable NSORDER or one of the the configuration files,
 - /etc/netsvc.conf or
 - /etc/irs.conf.
- 6. After this, the dynamic loadable API functionality is ready to be used.

Chapter 4. Internet Protocol (IP) Security

IP Security enables secure communications over the Internet and within company networks by securing data traffic at the IP layer. This allows individual users or organizations to secure traffic for all applications, without having to make any modifications to the applications. Therefore the transmission of any data, such as e-mail or application-specific company data, can be made secure.

The mechanism for securing data between two nodes is accomplished by creating a virtual tunnel between two hosts. This is also referred to as creating a Virtual Private Network (VPN). The secure tunnel encapsulates all IP traffic between the two hosts in a manner specified by the user. It provides data integrity, privacy, and authentication depending on how the tunnel is defined.

This chapter discusses the following topics:

- Benefits of a Virtual Private Network (VPN), on page 4-1
- Security, on page 4-2
- IP Security Features, on page 4-3
- Security Associations, on page 4-4
- Tunnels and Key Management, on page 4-4
- · Native Filtering Capability, on page 4-5
- IP Security Installation, on page 4-7
- IP Security Configuration, on page 4-8, which includes sections on:
 - Basic Configuration, on page 4-10
 - Advanced Configuration, on page 4-17
- IP Security Problem Determination, on page 4-29
- IP Security Reference, on page 4-34

Benefits of a Virtual Private Network (VPN)

A virtual private network is an extension of an enterprise's private intranet across a public network such as the Internet. It creates a secure private connection through what is essentially a private tunnel. VPNs securely convey information across the Internet connecting remote users, branch offices, and business partners/suppliers into an extended corporate network. Internet service providers (ISPs) offer cost—effective access to the Internet through direct lines or local telephone numbers. Therefore, companies can eliminate their current networking options that use more expensive leased lines, long—distance calls, and toll—free telephone numbers.

Security

AIX uses IP Security (IP Sec), which is an open, standard security technology developed by the Internet Engineering Task Force (IETF), as an integral element of eNetwork VPN solutions. IP Sec provides cryptography–based protection of all data at the IP layer of the communications stack. It provides secure communications transparently, with no changes required to existing applications. IP Sec is the IETF–chosen industry standard network security framework for use in both the IP Versions 4 and 6 environments.

IP Sec protects your data traffic in three ways using robust cryptographic techniques:

1. Authentication

The process by which the identity of a host or end point is verified

2. Integrity Checking

The process of ensuring that no modifications were made to the data while in–transit across the network

3. Encryption

The process of ensuring privacy by "hiding" data and private IP addresses while in–transit across the network

Authentication algorithms provide proof of the sender's identity and data integrity by using a cryptographic hash function to process a packet of data (with the immutable IP header fields included) with a secret key to produce a unique digest. On the receiver's side, the data is decapsulated using the same function and key. If either the data has been altered or the sender's key was not valid, the datagram is discarded.

Encryption uses a cryptographic algorithm to modify and randomize the data using a certain algorithm and key to produce cyphertext. Encryption makes the data unreadable while in transit. Once received, the data is recovered using the same algorithm and key (with symmetric encryption algorithms). Encryption must occur with authentication to verify the data integrity of the encrypted data.

These basic services are implemented in IP Sec by the use of the Encapsulating Security Payload (ESP) and the Authentication Header (AH). With ESP, confidentiality is provided by encrypting the original IP packet, building an ESP header, and putting the encrypted data (known as *ciphertext*) in the ESP payload.

When it is desired to have authentication and integrity checking without confidentiality, the AH may be used alone. With AH, the static fields of the IP header and the data have a hash algorithm applied to compute a keyed digest. The receiver uses its key to compute and compare the digest to make sure the packet is unaltered and the sender's identity is authenticated.

IP Security Features

The IP Security feature of AIX provides the following functions:

- AH support using RFC 2402, and ESP support using RFC 2406.
- Automatic key refreshment with tunnels using IETF's IKE protocol, and with IBM tunnels.
- Automatic key refreshment with manual tunnels using statically defined keys.
- Tunnel mode and transport mode of encapsulation
- Authentication Algorithms of HMAC MD5 and HMAC SHA1
- Encryption Algorithms include 56 bit DES CBC with 64 bit IV, Triple DES, DES CBC 4 (32 bit IV), CDMF (40 bit DES that is exportable).
- Dual IP Stack Support (IP version 4 and IP version 6). IP Version 6 is configurable with manual tunnels
- Both IP Version 4 and IP Version 6 traffic can be encapsulated and filtered. Because the IP stacks are separate, the IP Security function for each stack can be configured independently.
- Manual tunnels can be configured to provide interoperability with other systems that do not support the automatic IKE key refreshment method, and for use of IP V6 tunnels.
- Filtering of secure and non-secure traffic by a variety of IP characteristics such as source and destination IP addresses, interface, protocol, port numbers, etc.
- · Automatic filter rule creation and deletion with most tunnel types.
- Use of hostnames for the destination address when defining tunnels and filter rules. The hostnames will be converted to IP addresses automatically (as long as DNS is available).
- · Logging of IP Security events to syslog.
- Extensive use of system traces and statistics for problem determination.
- User defined default action allows the user to specify whether traffic that does not match defined tunnels should be allowed or denied.

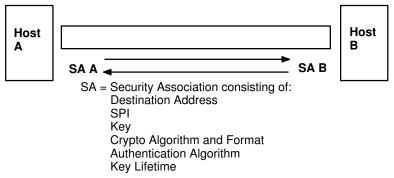
IKE Features

The following features are available with Internet Key Exchange, AIX versions 4.3.2 and later:

- · Authentication with preshared keys
- Use of main mode (identify protect mode) and aggressive mode
- Support for Diffie Hellman groups 1 and 2
- ESP encryption support for DES, Triple DES, Null encryption. ESP authentication support with HMAC MD5 and HMAC SHA1
- AH support for HMAC MD5 and HMAC SHA1
- Separation of security policies with tunnel definition to allow security policy definitions to be reused

Security Associations

The building block on which secure communications is built is a concept known as a security association. Security associations (SAs) relate a specific set of security parameters to a type of traffic. With IP Security–protected data, a separate SA exists for each direction and for each header type, AH or ESP. The information contained in the SA includes the IP addresses of the communicating parties, a unique identifier known as the Security Parameters Index (SPI), the algorithms selected for authentication and/or encryption, the authentication and encryption keys, and the key lifetimes (see figure). The goal of key management is to negotiate and compute the SAs that protect IP traffic.



Establishment of a Secure Tunnel between Hosts A & B

Tunnels and Key Management

To set up a secure communication between two hosts, Security Associations must be negotiated and managed during the use of the tunnel. Three types of tunnels are supported in AIX, and each uses a different key management technique. They are:

- IKE tunnels (dynamically changing keys, IETF standard)
- IBM tunnels (dynamic, proprietary)
- Manual tunnels (static, persistent keys, IETF standard)

IKE Tunnel Support

IKE Tunnels are based on the ISAKMP/Oakley standards developed by the IETF. With this protocol, security parameters are negotiated and refreshed, and keys are exchanged securely. Three types of authentication are described in the standards, preshared key, digital signature and public key. AIX Version 4.3.3 currently implements Preshared Key.

The negotiation uses a two phase approach. The first phase authenticates the communicating parties, and specifies the algorithms to be used for securely communicating in phase 2. During phase 2, IP Security parameters to be used during data transfer are negotiated, security associations and keys are created and exchanged.

Algorithm	AH IP Version 4	ESP IP Version 4
HMAC MD5	Х	Х
HMAC SHA1	Х	Х
DES CBC 8		Х
3DES CBC		Х
ESP Null		Х

Manual Tunnels

Manual tunnels provide backward compatibility and will interoperate with machines that do not support IKE key management protocols. The disadvantage of manual tunnels is that the key values are static. In other words, the encryption and authentication keys are the same for the life of the tunnel and must be manually updated.

Algorithm	AH IP Version 4	AH IP Version 6	ESP IP Version 4	ESP IP Version 6
HMAC MD5	Х	Х	Х	Х
HMAC SHA1				
DES CBC 8			Х	Х
DES CBC 4			Х	Х
CDMF			Х	Х
3 DES CBC			Х	Х

Since IKE tunnels offer more effective security, IKE is the preferred key management method.

Native Filtering Capability

Filtering is a basic function in which incoming and outgoing packets can be accepted or denied based on a variety of characteristics. This allows a user or system administrator to configure the host to control the traffic between this host and other hosts. Filtering is done on a variety of packet properties, such as source and destination addresses, IP version (4 or 6), subnet masks, protocol, port, routing characteristics, fragmentation, interface, and tunnel definition.

Rules, known as *filter rules*, are used to associate certain kinds of traffic with a particular tunnel. In a basic AIX configuration for manual tunnels, when a user defines a host—to—host tunnel, filter rules are autogenerated to direct all traffic from that host through the secure tunnel. If more specific types of traffic are desired (for instance subnet to subnet), the filter rules can be edited or replaced to allow precise control of the traffic using a particular tunnel.

For IKE tunnels, the filter rules are also automatically generated and inserted in the filter table once the tunnel is activated.

Similarly, when the tunnel is modified or deleted, the filter rules for that tunnel are automatically deleted. This greatly simplifies IP Security configuration and helps reduce human error. Tunnel definitions can be propagated and shared among AIX machines and AIX firewalls using import and export utilities. This is especially helpful in the administration of a large number of machines.

Filter rules are necessary to associate particular types of traffic with a tunnel, but data being filtered does not necessarily need to travel in a tunnel. This allows AIX to provide base firewall function for users who want to restrict the flow of certain types of traffic to or from their machine. This is especially useful for the administration of machines in an intranet or machines that do not have the protection of a firewall. In concept, it is somewhat like setting up a demilitarized zone (DMZ); filter rules provide a second barrier around a group of machines in case of a compromise.

Once the filter rules are generated, they are stored in a table and loaded into the kernel. When packets are ready to be sent or received from the network, the filter rules are checked in the list from top to bottom to determine whether the packet should be permitted, denied or sent through a tunnel. The criteria of the rule is compared to the packet characteristics until a match is found or the default rule is reached.

The IP Security function also implements filtering of non–secure packets based on very granular user–defined criteria. This is a useful function to allow the control of IP traffic between networks and machines that do not require the authentication or encryption properties of IP Security.

IP Security Installation

The IP Security feature in AIX is separately installable and loadable. The core filesets that need to be installed are:

- **bos.net.ipsec.rte** The run time environment for the kernel IP Security environment and commands
- bos.msg.LANG.net.ipsec (where LANG is the desired language, such as en US)

Also, the following filesets need to be installed for IKE tunnel support:

- bos.net.ipsec.keymgt
- bos.net.ipsec.websm
- The bos.crypto fileset that is appropriate for your country

Once installed, IP Security can be separately loaded for IP Version 4 and IP Version 6. This is accomplished by issuing **mkdev** commands or through the IP Security SMIT menus.

Loading IP Security

Attention: Loading IP Security will enable the filtering function. Therefore, before loading, it is important to ensure the correct filter rules are created, or all outside communication may be blocked.

If using SMIT or Web-based System Manager, (**wsm network** fast path) the IP security modules will be automatically loaded when IP Security is started. This is the prefrerred method to ensure that the kernel extensions and IKE daemons are loaded in the proper order.

If the loading completed successfully, the **Isdev** command will show the IP Security devices as Available.

```
lsdev -C -c ipsec
  ipsec_v4 Available IP Version 4 Security Extension
  ipsec_v6 Available IP Version 6 Security Extension
```

Once the IP Security kernel extension has been loaded, tunnels and filters are ready to be configured.

IP Security Configuration

To configure IP Sec, tunnels and filters must be configured. When a simple tunnel is defined for all traffic to use, the filter rules can be automatically generated. If more complex filtering is desired, filter rules can be configured separately.

You can configure IP Sec using the Web-based System Manager application wsm network or the System Management Interface Tool (SMIT). If using SMIT, the following fastpaths will take you directly to the configuration panels you need:

ips4_basicips6 basicBasic configuration for IP version 4ips6 basicBasic configuration for IP version 6

This section on IP Security Configuration discusses the following topics:

- Tunnels versus Filters, on page 4-8
- Tunnels and Security Associations, on page 4-9
- Choosing a Tunnel Type, on page 4-9
- Basic Configuration, on page 4-10
- Static Filter Rules and Examples, on page 4-13
- Advanced Manual Tunnel Configuration, on page 4-17
- Configuring IKE Tunnels, on page 4-18
- Predefined Filter Rules, on page 4-23
- Logging Facilities, on page 4-23
- Coexistence of IP Security and IBM Secured Network Gateway 2.2/IBM Firewall 3.1 or 3.2, on page 4-28

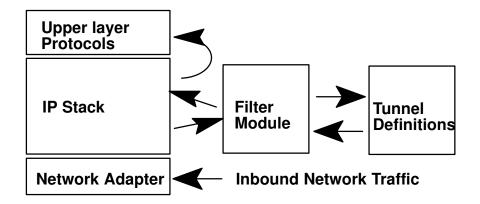
Tunnels versus Filters

There are two related but distinct parts of IP Security: *tunnels* and *filters*. Tunnels require filters, but filters do not require tunnels.

Filtering is a basic function in which incoming and outgoing packets can be accepted or denied based on a variety of characteristics. This allows a system administrator to configure the host to control the traffic between this host and other hosts. Filtering is done on a variety of packet properties, such as source and destination addresses, IP Version (4 or 6), subnet masks, protocol, port, routing characteristics, fragmentation, interface, and tunnel definition. This filtering is done at the IP layer, so no changes are required to the applications.

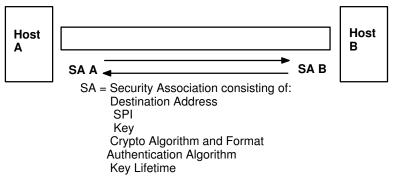
Tunnels define a security association between two hosts. These security associations involve specific security parameters that are shared between end points of the tunnel.

The following illustration indicates how a packet comes in the network adapter to the IP stack. From there, the filter module is called to determine if the packet should be permitted or denied. If a tunnel ID is specified, the packet will be checked against the existing tunnel definitions. If the decapsulation from the tunnel is successful, the packet will be passed to the upper layer protocol. This function will occur in reverse order for outgoing packets. The tunnel relies on a filter rule to associate the packet with a particular tunnel, but the filtering function can occur without passing the packet to the tunnel.



Tunnels and Security Associations

Tunnels are used whenever it is desired to have data authenticated, or authenticated and encrypted. Tunnels are defined by specifying a security association between two hosts (see figure). The security association defines the parameters for the encryption and authentication algorithms and characteristics of the tunnel.



Establishment of a Secure Tunnel between Hosts A & B

The Security Parameter Index (SPI) and the destination address identify a unique security association. Therefore, these two parameters are required for uniquely specifying a tunnel. Other parameters such as cryptographic algorithm, authentication algorithm, keys, and lifetime can be specified or defaults can be used.

Choosing a Tunnel Type

The decision to use IBM tunnels, manual tunnels, or, for AIX versions 4.3.2 and later, IKE tunnels, depends on the tunnel support of the remote end and the type of key management desired. IKE tunnels are preferable (when available) because they offer secure key negotiation and key refreshment in an industry–standard way. They also take advantage of the new IETF ESP and AH header types and support anti–replay protection.

IBM tunnels offer similar security, but their support is limited to a smaller set of encryption and authentication algorithms, but they provide backward compatibility and ease of use with their import/export functions with the IBM Firewall.

If the remote end does not support IBM tunnels, or uses one of the algorithms requiring manual tunnels, manual tunnels should be used. Manual tunnels ensure interoperability with a large number of hosts. Because the keys are static and difficult to change and may be cumbersome to update, they are not as secure.

IBM Tunnels may be used between any two AIX machines running AIX Version 4.3, or between an AIX 4.3 host and a host running IBM Secure Network Gateway 2.2 or IBM Firewall 3.1/3.2. Manual tunnels may be used between a host running AIX Version 4.3 and any other machine running IP Security and having a common set of cryptographic and authentication algorithms. Almost all vendors offer Keyed MD5 with DES, or HMAC MD5 with DES. This is a base subset that works with almost all implementations of IP Security.

When setting up manual or IBM tunnels, the procedure depends on whether you are setting up the first host of the tunnel or setting up the second host, which must have parameters matching the first host's setup. When setting up the first host, the keys may be autogenerated, and the algorithms can be defaulted. When setting up the second host, it is best to import the tunnel information from the remote end, if possible.

Another important consideration is determining whether the remote system is behind a firewall. If it is, the setup must include information about the intervening firewall.

Basic Configuration (Manual or IBM Tunnels)

Setting Up Tunnels and Filters

For the simplest case, setting up a manual tunnel, it is not necessary to separately configure the filter rules. As long as all traffic between two hosts goes through the tunnel, the necessary filter rules are automatically generated. The process of setting up a tunnel is to define the tunnel on one end, import the definition on the other end, and activate the tunnel and filter rules on both ends. Then the tunnel is ready to use.

Information about the tunnel must be made to match on both sides if it is not explicitly supplied (see figure). For instance, the encryption and authentication algorithms specified for the source will be used for the destination if the destination values are not specified. This makes creating the tunnel much simpler.



Manual Tunnel Using HMAC MD5 and DES CBC 8

Creating a Manual Tunnel on Host A

You can configure a tunnel using the Web-based System Manager application wsm network, the SMIT fast path ips4 basic (for IP Version 4) or ips6 basic (for IP version 6), or you can use the following procedure.

The following is a sample of the **gentun** command used to create a manual tunnel:

```
gentun -v 4 -t manual -s 5.5.5.19 -d
5.5.5.8 -a HMAC_MD5 -e DES_CBC_8 -N 23567
```

This will create a tunnel with output (using Istun -v 4) that looks similar to:

```
Tunnel ID : 1
IP Version : IP Version 4
Source : 5.5.5.19
Destination : 5.5.5.8
Policy : auth/encr
Tunnel Mode : Tunnel
Send AH Algo : HMAC_MD5
Send ESP Algo : DES_CBC_8
Receive AH Algo : HMAC_MD5
Receive ESP Algo : DES_CBC_8
Source AH SPI : 300
Source ESP SPI : 300
Dest AH SPI : 23576
Dest ESP SPI : 23576
Tunnel Life Time : 480
Status : Inactive
Target : -
Target Mask : -
Target Mask : -
Replay : No
New Header : Yes
Snd ENC-MAC Algo : -
Che tunnel will be activated when the mktun commands
```

The tunnel will be activated when the **mktun** command is used:

```
mktun -v 4 -t1
```

The filter rules associated with the tunnel are automatically generated and output (using **Isfilt –v 4**) looks similar to:

Rule 4: Rule 4:
Rule action : permit
Source Address : 5.5.5.19
Source Mask : 255.255.255.255 Destination Address : 5.5.5.8 Destination Mask : 255.255.255.255 Source Routing : yes Protocol : all
Source Port : any 0
Destination Port : any 0
Scope : both
Direction : outbound
Logging control : no
Fragment control : all packets
Tunnel ID number : 1
Interface : all Interface : all Auto-Generated : yes Rule 5: Rule action : permit
Source Address : 5.5.5.8
Source Mask : 255.255 : 255.255.255.255 Destination Address : 5.5.5.19 Destination Mask : 255.255.255 Source Routing : yes Source Routing Protocol : all
Source Port : any 0
Destination Port : any 0
Scope : both
Direction : inbound
Logging control : no
Fragment control : all packets
Tunnel ID number : 1
Interface : all Interface : all Auto-Generated : yes

These filter rules in addition to the default filter rules are activated by the **mktun –v 4 –t 1** command.

To set up the other side (when it is another AIX machine), the tunnel definition can be exported on host A then imported to host B.

To export:

```
exptun -v 4 -t 1 -f /tmp
```

This will export the tunnel definition into a file named **ipsec_tun_manu.exp** and any associated filter rules to the file **ipsec_fltr_rule.exp** in the directory indicated by the **-f** flag.

Creating a manual tunnel on Host B

To create the matching end of the tunnel, the export files are copied to the remote side and imported into that remote AIX 4.3 machine by using the command:

```
imptun -v 4 -t 1 -f /tmp
```

where 1 is the tunnel to be imported and /tmp is the directory where the import files reside. This tunnel number is system generated and must be referenced from the output of the **gentun** command, or by using the **Istun** command to list the tunnels and determine the correct tunnel number to import. If there is only one tunnel in the import file, or if all the tunnels are to be imported, then the **–t** option is not needed.

If the remote machine is not AIX 4.3, the export file can be used as a reference for setting up the algorithm, keys, and SPI values for the other end of the tunnel.

Export files from the IBM Secure Network Gateway (SNG) can be imported to create tunnels in AIX 4.3. To do this, use the **–n** option when importing the file:

```
imptun -v 4 -f /tmp -n
```

Creating an IBM tunnel on Host A

Setting up an IBM tunnel is similar to a manual tunnel, but some of the choices are different for the crypto algorithms and the keys are negotiated dynamically, so there is no need to import keys. IBM tunnels are limited to Keyed MD5 for authentication. If the HMAC MD5 or HMAC SHA algorithms are desired, a manual tunnel must be used.

```
gentun -s 9.3.100.1 -d 9.3.100.245 -t IBM -e
DES_CBC_8 -n 35564
```

As with manual tunnels, from this point the tunnel and filter table must be activated to make the tunnel active:

```
mktun -v 4 -t1
```

To set up the other side, if the other host is an AIX 4.3 IP Security machine, the tunnel definition can be exported on host A, then imported to host B.

To export:

```
exptun -v 4 -f /tmp
```

This will export the tunnel definition into a file named **ipsec_tun_ibm.exp** and any associated filter rules to the file **ipsec_fltr_rule.exp** in the directory indicated by the **-f** flag.

Creating an IBM tunnel on Host B

The procedure is the same for creating the second end of the tunnel on host B for an IBM tunnel. The tunnel definition is exported from host A and imported onto host B. The **–n** flag can be used for a file exported by an IBM Secure Network Gateway or an IBM Firewall 3.1/3.2.

Static Filter Rules and Examples

Filtering can be set up to be simple, using mostly autogenerated filter rules, or can be complex by defining very specific filter functions based on the properties of the IP packets. Matches on incoming packets are done by comparing the source address and SPI value to those listed in the filter table. Therefore, this pair must be unique.

Each line in the filter table is known as a *rule*. A collection of rules will determine what packets are accepted in and out of the machine, and how they will be directed. Filter rules can be written based on source and destination addresses and masks, protocol, port number, direction, fragment control, source routing, tunnel, and interface.

Below is a sample set of filter rules. Within each rule, fields are shown in the following order (an example of each field from rule 1 is shown in parentheses): Rule_number (1), Action (permit), Source_addr (0.0.0.0), Source_mask (0.0.0.0), Dest_addr (0.0.0.0), Dest_mask (0.0.0.0), Source_routing (no), Protocol (udp), Src_prt_operator (eq), Src_prt_value (4001), Dst_prt_operator (eq), Dst_prt_value (4001), Scope (both), Direction (both), Logging (no), Fragment (all packets), Tunnel (0), and Interface (all).

```
1 permit 0.0.0.0 0.0.0.0 0.0.0.0 0.0.0.0 no udp eq 4001 eq 4001 both both no all packets 0 all
```

- 2 permit 0.0.0.0 0.0.0.0 0.0.0.0 0.0.0.0 no ah any 0 any 0 both both no all packets 0 all
- 3 permit 0.0.0.0 0.0.0.0 0.0.0.0 0.0.0.0 no esp any 0 any 0 both both no all packets 0 all
- 4 permit 10.0.0.1 255.255.255.255 10.0.0.2 255.255.255.255 no all any 0 any 0 both outbound no all packets 1 all
- 5 permit 10.0.0.2 255.255.255.255 10.0.0.1 255.255.255.255 no all any 0 any 0 both inbound no all packets 1 all
- 6 permit 10.0.0.1 255.255.255.255 10.0.0.3 255.255.255.255 no tcp lt 1024 eq 514 local outbound yes all packets 2 all
- 7 permit 10.0.0.3 255.255.255.255 10.0.0.1 255.255.255.255 no tcp/ack eq 514 lt 1024 local inbound yes all packets 2 all
- 8 permit 10.0.0.1 255.255.255.255 10.0.0.3 255.255.255.255 no tcp/ack lt 1024 lt 1024 local outbound yes all packets 2 all
- 9 permit 10.0.0.3 255.255.255.255 10.0.0.1 255.255.255.255 no tcp lt 1024 lt 1024 local inbound yes all packets 2 all
- 10 permit 10.0.0.1 255.255.255.255 10.0.0.4 255.255.255.255 no icmp any 0 any 0 local outbound yes all packets 3 all
- 11 permit 10.0.0.4 255.255.255.255 10.0.0.1 255.255.255.255 no icmp any 0 any 0 local inbound yes all packets 3 all
- 12 permit 10.0.0.1 255.255.255.255 10.0.0.5 255.255.255.255 no tcp gt 1023 eq 21 local outbound yes all packets 4 all
- 13 permit 10.0.0.5 255.255.255.255 10.0.0.1 255.255.255.255 no tcp/ack eq 21 gt 1023 local inbound yes all packets 4 all
- 14 permit 10.0.0.5 255.255.255.255 10.0.0.1 255.255.255.255 no tcp eq 20 gt 1023 local inbound yes all packets 4 all
- 15 permit 10.0.0.1 255.255.255.255 10.0.0.5 255.255.255.255 no tcp/ack gt 1023 eq 20 local outbound yes all packets 4 all
- 16 permit 10.0.0.1 255.255.255.255 10.0.0.5 255.255.255.255 no tcp gt 1023 gt 1023 local outbound yes all packets 4 all
- 17 permit 10.0.0.5 255.255.255.255 10.0.0.1 255.255.255.255 no tcp/ack gt 1023 gt 1023 local inbound yes all packets 4 all
- 18 permit 0.0.0.0 0.0.0.0 0.0.0.0 0.0.0.0 no all any 0 any 0 both both yes all packets

Rule 1 is for the IBM Session Key daemon and will only appear in IP Version 4 filter tables. It uses port number 4001 to control packets for refreshing the session key. It is an example of how the port number can be used for a specific purpose. This filter rule should not be modified except for logging purposes.

Rules 2 and 3 are used to allow processing of Authentication Headers (AH) and Encapsulating Security Payload (ESP) headers. They should not be modified except for logging purposes.

Rules 4 and 5 are a set of autogenerated rules that filter traffic between addresses 10.0.0.1 and 10.0.0.2 through tunnel #1. Rule 4 is for outbound traffic and rule 5 is for inbound traffic.

Rules 6 through 9 are a set of user–defined rules that filter outbound **rsh**, **rcp**, **rdump**, **rrestore**, and **rdist** services between addresses 10.0.0.1 and 10.0.0.3 through tunnel #2. Note that logging is set to yes so the administrator can monitor this type of traffic.

Rules 10 and 11 are a set of user-defined rules that filter both inbound and outbound **icmp** services of any type between addresses 10.0.0.1 and 10.0.0.4 through tunnel #3.

Rules 12 through 17 are user—defined filter rules that filter outbound FTP service from 10.0.0.1 and 10.0.0.5 through tunnel #4.

Rule 18 is an autogenerated rule always placed at the end of the table. In this case, it permits all packets that do not match the other filter rules. It may be set to deny all traffic not matching the other filter rules.

Each rule may be viewed separately (using Isfilt) to make each field clear. For example:

Rule 1: Rule action : permit
Source Address : 0.0.0.0
Source Mask : 0.0.0.0 Destination Address : 0.0.0.0 Destination Mask : 0.0.0.0 Source Routing : yes Protocol : uup
Source Port : eq 4001
Destination Port : eq 4001
Scope : both
both Direction : both Direction
Logging control : no Fragment control Tunnel ID number : all packets : 0 Interface : all Auto-Generated : yes

b

Below are listed all the parameters that can be specified in a filter rule:

Both

-v IP version: 4 or 6. Action: _a d Deny Permit Source address. Can be an IP address or hostname. -s Source subnet mask. -m Destination address. Can be an IP address or hostname. <u>_d</u> Destination subnet mask. -М Source routing control: y or n. -g Protocol. Values may be udp, icmp, tcp, tcp/ack, ospf, pip, esp, ah -с and all. Source port or ICMP type operation. -0 Source port or ICMP type value. -p Destination port or ICMP code operation. -0 -P Destination port or ICMP code value. –r Routing. Forwarded packets r ı Local destined/originated packets

-1 Log control. Include in log У n Do not include in log. −f Fragmentation. Applies to fragments headers, fragments, and У non-fragments Applies only to fragments and fragment headers 0 n Applies only to non–fragments h Applies only to non-fragments and fragment headers Tunnel ID. -t —i Interface, such as tr0 or en0.

See the AIX Commands Reference for genfilt and chfilt for more information.

Autogenerated Filter Rules And User Specified Filter Rules

Certain rules are autogenerated for the use of the IP Security filter and tunnel code. These include the rules shown above for the session key daemon to refresh the keys in IKE (AIX versions 4.3.2 and later) or IBM tunnels (both IP version 4 only) and the rules for the processing of AH and ESP packets. Filter rules will also be autogenerated when defining tunnels. For manual and IBM tunnels, they will specify the source and destination address and mask values, as well as the tunnel ID. All traffic between those addresses will flow through the tunnel.

IKE filter rules are also autogenerated and will be kept in a separate table. The IKE negotiation has support to specify protocol and port numbers, they will be set accordingly in the filter rules. The IKE filter rules will be searched after the static filter rules and before the autogenerated rules. The position in the static filter table where IKE filter rules will be inserted will be defaulted, but they can be moved by the user.

Since the autogenerated rules permit all traffic over the tunnel, user defined rules may be necessary to place restrictions on certain types of traffic. These user defined rules should be placed before the autogenerated rules because the first rule that applies to the packet will be used. Below is an example of user defined filter rules that filter traffic based on ICMP operation.

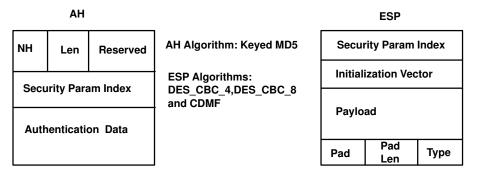
```
1 permit 10.0.0.1 255.255.255.255 10.0.0.4 255.255.255.255 no icmp any 8 any 0 local outbound no all packets 3 all
2 permit 10.0.0.4 255.255.255.255 10.0.0.1 255.255.255.255 no icmp any 0 any 0 local inbound no all packets 3 all
3 permit 10.0.0.4 255.255.255.255 10.0.0.1 255.255.255.255 no icmp any 8 any 0 local inbound no all packets 3 all
4 permit 10.0.0.1 255.255.255.255 10.0.0.4 255.255.255.255 no icmp any 0 any 0 local outbound no all packets 3 all
```

Filter rules will be autogenerated when tunnels are defined. This is done to simplify the configuration of a single tunnel. This function can be suppressed by specifying the **-g** flag in gentun. You can find a sample filter file with **genfilt** commands to generate filter rules for different TCP/IP services in /usr/samples/ipsec/filter.sample.

Advanced Manual Tunnel Configuration

Tunnel configurations allow you to specify a large number of parameters to ensure interoperability with other IP Security implementations. Some examples are given here to explain the functions that can be selected.

The packets are formed using AH headers for authentication, ESP headers for encryption, or the new ESP header that allows both encrypted data and an authentication digest in the same packet (see figure).

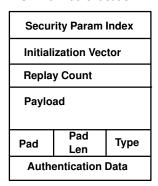


To configure a tunnel using DES CBC MD5, the combined ESP header is used with the ESP encryption algorithm set to DES_CBC_8 and the AH authentication algorithm set to HMAC MD5. To use this authentication algorithm, specify it with the **-b** flag, and its keys with the **-c** flag. For example:

```
gentun -v 4 -t manual -s 5.5.5.19 -d 5.5.5.23 -e DES CBC 8 -b HMAC MD5 -N 2349473
```

This command will generate a tunnel using DES CBC MD5 with autogenerated keys (see figure). The encrypted data and the authentication data are both contained in the same ESP header. It also supports replay preventions, which is not recommended for manual tunnels. It is included in this release for compatibility with other implementations and for testing purposes.

ESP with Authentication



Supported Algorithms:
DES_CBC_MD5,
DES_CBC_8 (w/0 authentication)
CDMF (w/0 authentication)

Configuring IKE Tunnels

Internet Key Exchange (IKE) tunnels have more complex policy parameters; therefore, the Web-based System Manager interface is needed to configure IKE tunnels. Creating IKE tunnels differs from creating IBM or manual tunnels because the setting of security policy is a separate process from defining tunnel endpoints. In IKE, there is a two–step negotiation process, called *phases*, and each phase can have separate security policies.

When the Internet Key negotiation starts, it must set up a secure channel for the negotiations. This is known as the *key management* phase or *phase 1*. During this phase, each party uses preshared keys to authenticate the other and pass ID information. This phase sets up a security association during which the two parties determine how they plan to communicate securely and then which protections are to be used to communicate during the second phase. The result of this phase is an *IKE* or *phase 1* tunnel.

The second phase is known as the *data management* phase or *phase 2* and uses the IKE tunnel to create the IP Sec security associations for AH and ESP that actually protect traffic. The second phase also determines the data that will be using the IP Sec tunnel. For example, it can specify:

- a subnet mask,
- an address range,
- or a protocol and port number combination

(See figure.)

IKE Tunnel Setup Process		
Step 1: Negotiation	Step 2: Key Exchange	
Key Management (Phase 1) IKE SA Parameters Authentication Hash Key Lifetime	Use public key cryptography to establish first shared secret Exchange and authenticate IDs Identify the negotiating parties Result: IKE (phase 1) tunnel	
Data Management (Phase 2) IP Sec Protocols (AH, ESP) Encapsulation Mode Encapsulation Algorithm Authentication Algorithm Key Lifetimes	Generate session keys Exchange and authenticate IDs Identify parties using IP Sec Result: IP Sec (phase 2) tunnel	

In many cases, the endpoints of the key management (IKE) tunnel will be the same as the endpoints of the data management (IP Sec) tunnel. The IKE tunnel endpoints are the IDs of the machines carrying out the negotiation. The IP Sec tunnel endpoints describe the type of traffic that will use the IP Security tunnel. For simple host—to—host tunnels, in which all traffic between two tunnels is protected with the same tunnel, the phase 1 and phase 2 tunnel endpoints are the same. When negotiating parties are two gateways, the IKE tunnel endpoints are the two gateways, and the IP Sec tunnel endpoints are the machines or subnets (behind the gateways) or the range of addresses (behind the gateways) of the tunnel users.

Key Management (Phase 1) Parameters — Tunnel Identification

Key Management (Phase 1) Tunnel

The name of this IKE tunnel. For each tunnel, the endpoints of the negotiation must be specified. These are the two machines that plan to send and validate IKE messages. The name of the tunnel may describe the tunnel endpoints such as VPN Boston or VPN Acme.

Host Identity Type

The ID type that will be used in the IKE exchange. The ID type and value must match the value for the preshared key to ensure that proper key lookup is performed. If a separate ID is used to search a preshared key value, the *host ID* is the key's ID and its *type* is KEY_ID. The KEY_ID type is useful if a single host has more than one preshared key value.

Host Identity

The value of the host ID represented as an IP address, a fully qualified domain name (FQDN), or a user at the fully qualified domain name (user@FQDN).

IP Address

The IP address of the remote host. This value is required when the host ID type is KEY_ID or whenever the host ID type cannot be resolved to an IP address. For example, if the user name cannot be resolved with a local nameserver, the IP address for the remote side must be entered.

Key (Phase 1) Policy

When configuring an IKE tunnel through the Web-based System Manager Network application (**wsm network** fast path), you can choose from several predefined key management policies that will work for most situations. They vary according to the type of encryption and authentication algorithms specified.

Also, you can customize key management policy by specifying the parameters to be used during IKE negotiation. For Phase 1, the user must determine certain key management security properties with which to carry out the exchange. These properties are fully described in the help available with Web-based System Manager and in the documentation for the **ike** command.

Data Management (Phase 2) Policy

The data management policy parameters are the same IP Security parameters used in manual tunnels. They describe the type of protection to be used for protecting data traffic in the tunnel. You can start more than one phase 2 tunnel under the same phase 1 tunnel. Several predefined policies are available to simplify configuration, or they can be used as a template for creating custom security policies.

If a custom security policy is desired, the the specific policy parameters can be specified, in which case the name of the policy must be entered to relate the tunnel endpoints with the appropriate policy.

The following endpoint ID types describe the type of data that is to use the IP Security Data tunnel:

Host, Subnet, or Range

Describes whether the data traffic traveling in the tunnel will be for a particular host, subnet, or address range.

Host/Subnet IDContains the host or subnet identity of the local and remote systems passing traffic over this tunnel. Determines the IDs sent in the phase 2 negotiation and the filter rules that will be built if the negotiation is successful.

Subnet mask Describes all IP addresses within the subnet (for example, host 9.53.250.96 and mask 255.255.255.0)

Starting IP Address Range

Provides the starting IP address for the range of addresses that will be using the tunnel (for example, 9.53.250.96 of 9.53.250.96 to 9.53.250.93)

Ending IP Address Range

Provides the ending IP address for the range of addresses that will be using the tunnel (for example, 9.53.250.93 of 9.53.250.96 to 9.53.250.93)

Port Describes data using a specific port number (for example, 21 or 23)

Describes data being transported with a specific protocol (for example, TCP or UDP). Determines the protocol sent in the phase 2 negotiation and the filter rules that will be built if the negotiation is successful. The protocol for the local endpoint must match the protocol for the remote end point.

The data management parameters are fully described in the help available with Web-based System Manager and in the documentation for the **ike** command.

Example of Setting Up an IKE Tunnel

Protocol

There are typical scenarios that describe the type of situations most customers encounter when trying to set up tunnels. They can be described as the branch office, business partner, and remote access cases.

In the branch office case, the customer has two trusted networks that they want to connect together—the engineering group of one location to the engineering group of another. In this example, there are gateways that connect and all the traffic passing between the gateways use the same tunnel. The traffic at either end of the tunnel is decapsulated and passes in the clear within the company intranet.

In the first phase of the IKE negotiation, the IKE security association is created between the two gateways. The traffic that passes in the IP Security tunnel is the traffic between the two subnets, and the subnet IDs are used in the phase 2 negotiation. After the security policy and tunnel parameters are entered for the tunnel, a tunnel number is created. To start the tunnel, use the **ike** command.

In the business partner scenario, the networks are not trusted, and the network administrator may want to restrict access to a smaller number of hosts behind the security gateway. In this case, the tunnel between the hosts will carry traffic protected by IP Sec for use between two particular hosts. This would set the protocol of the phase 2 tunnel to be AH or ESP.

In the remote access case, the tunnels would be setup on demand, and a high level of security would be applied. The IP addresses may not be meaningful, therefore, fully qualified domain names or *user at* fully qualified domain names would be preferred. Optionally, you can use KEYID to relate a key to a host ID.

Invoke WebSM to configure the tunnel. Use **wsm network** to fast path to the TCP /IP section. Under Virtual Private Networking, IP Security can be started by clicking on the "Selected" option at the top. This will load the IP Security kernel extensions and start the isakmpd and tmd daemons. A tunnel is created by defining the phase 1 and phase 2 endpoints and their associated security policies. Phase 1 is the authentication phase which is computational expensive and is needed less often. It sets up a secure channel between the negotiating parties needed before the final IP Security parameters and keys are computed. Phase 2 identities describes the type of traffic that will be using a particular tunnel, and can be configured for a single host or group of hosts (with the use of subnets or IP ranges) along with specified protocol and port numbers. The same phase 1 tunnel can be used to protect multiple phase 2 negotiations and key refreshes, as long as they take place between the same two endpoints, for example between two gateways.

For defining the phase 1 tunnel endpoints, click on "Internet Key Exchange (IKE) Tunnels". Here information is entered to describe the identities of the systems taking part in the negotiations. In most cases, IP addresses will be used, and a policy compatible with the remote side must be selected. IBM_low_prekey and IBM_med_prekey uses DES encryption and preshared keys in aggressive mode. IBM_high_prekey uses Triple DES Encryption in Main mode, also known as identity protect mode. Use matching policies on both sides, or contact the administrator on the remote end to define a matching policy. A policy containing several choices can be created to allow flexibility when proposing or matching on a policy.

The preshared key must also be entered under the key tab, and this value must match on both the remote and local machines. Precede the value with 0X for a hex value.

"Key Management (Phase 1) Protection Policies" is used to select the particular policy with the tunnel. Many predefined policies exist to from which to select, or a new policy can be created.

Similarly, the phase 2 tunnel endpoints and policies are created by defining phase 2 tunnels and phase 2 policies. From there the tunnel can be started either through the WebSM graphical user interface or on the command line with the **ike** command.

To start a tunnel negotiation or to allow the incoming system to act as a responder (depending on the role that is specified), the **ike** command can be used with a tunnel number, as shown in the following example:

```
ike cmd=activate numlist=1
```

The IP addresses can also be used. For example:

```
ike cmd=activate ipaddr=9.3.97.100,9.3.97.256
```

Since it may take several seconds for the commands to complete, the command will return once the negotiation is started. To ensure the command was successful, use the list option to display the tunnel status.

```
ike cmd=list
Phase 1 Tunnel ID [1]
Phase 2 Tunnel ID [1]
```

This will show the phase 1 and phase 2 tunnels that are currently active. To do a verbose listing of the tunnel,

```
ike cmd=list verbose
Phase 1 Tunnel ID
Local ID Type:
                         Fully_Qualified_Domain_Name
                         bee.austin.ibm.com
Local ID:
Remote ID Type: Fully_Qualified_Domain_Name ipsec.austin.ibm.com
Mode:
                          Aggressive
Moae:
Security Policy:
                         BOTH_AGGR_3DES_MD5
                          Initiator
Encryption Alg:
Auth Alg:
                          3DES-CBC
                        Preshared Key
Hash Alg:
                         MD5
Key Lifetime: 28800 Seconds
Key Lifesize: 0 Kbytes
Key Rem Lifetime: 28737 Seconds
Key Rem Lifesize: 0 Kbytes
Key Refresh Overlap: 5%
Tunnel Lifetime:
Tunnel Lifesize:
                         2592000 Seconds
                         0 Kbytes
Tun Rem Lifetime: 2591937 Seconds
Status:
                          Active
Phase 2 Tunnel ID 1
Local ID Type: IF
                          IPv4_Address
                          10.10.10.1
Local ID:
Local Subnet Mask:
                         N/A
Local Port:
                           any
```

Local Protocol:

Remote ID Type:

Remote ID:

Remote Subnet Mask:

Remote Port:

Remote Port:

Remote Portocol:

Mode:

Security Policy:

ESP_3DES_MD5_SHA_TUNNEL_NO_PFS
Role:

Initiator

Encryption Alg:

ESP_3DES

AH Transform:

N/A

Auth Alg:

HMAC-MD5

PFS:

No

SA Lifetime:

SA Lifesize:

SA Rem Lifetime:

SA Rem Lifetime:

SA Rem Lifetime:

SA Rem Lifetime:

Tunnel Lifetime:

Tunnel Lifetime:

Assoc P1 Tunnel:

ESP_tunnel

ESP_tunnel

ESP_tunnel

Status:

Active

Activating the IKE tunnel will cause filter rules for the new tunnel to be inserted into the dynamic filter table. These entries can be viewed using the Isfilt command with the –d option for dynamic filter rules:

```
1 permit 0.0.0.0 0.0.0.0 0.0.0.0 0.0.0.0 no udp eq 4001 eq 4001
both both no all packets 0 all
2 *** Dynamic filter placement rule *** no
0 permit 0.0.0.0 0.0.0.0 0.0.0.0 0.0.0.0 yes all any 0 any 0 both
both no all packets 0 all
*** Dynamic table ***
0 permit 0.0.0.0 0.0.0.0 0.0.0.0 0.0.0.0 no udp eq 500 eq 500
local both no all packets 0
0 permit 0.0.0.0 0.0.0.0 0.0.0.0 0.0.0.0 no ah any 0 any 0 both
inbound no all packets 0
0 permit 0.0.0.0 0.0.0.0 0.0.0.0 0.0.0.0 no esp any 0 any 0 both
inbound no all packets 0
1 permit 10.10.10.1 255.255.255.255 10.10.10.4 255.255.255.255 no
all any 0 any 0 both outbound yes all packets 1
1 permit 10.10.10.4 255.255.255.255 10.10.10.1 255.255.255.255 no
all any 0 any 0 both inbound yes all packets 1
```

This example shows a machine that has one IKE tunnel and no other tunnels. The dynamic filter rule (rule #2 in this example of the static table) can be moved by the user to control placement relative to all other user—defined rules. The rules in the dynamic table are constructed automatically as tunnels are negotated and corresponding rules are inserted into the filter table. These rules can be displayed, but not edited.

To turn on logging of the dynamic filter rules, the logging option for rule #2 is set to yes:

```
chfilt -v 4 -n 2 -l y
```

For more details on logging of IKE traffic, see the section on Logging Facilities, on page 4-23.

To deactivate the tunnel, use the remove option.

ike cmd=remove numlist=1

Predefined Filter Rules

There are several predefined filter rules that are autogenerated with certain events. When the <code>ipsec_v4</code> or <code>ipsec_v6</code> device is loaded, a predefined rule is inserted into the filter table and then activated. By default, this predefined rule is to permit all packets, but it is user configurable and you can set it to deny all packets.

Note: When configuring remotely, ensure that the deny rule is not enabled before the configuration is complete. This will keep your session from getting shut out of the machine. The situation can be avoided either by setting the default action to permit or by configuring a tunnel to the remote machine before activating ipsec.

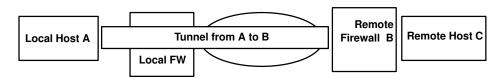
There is a predefined rule for both IPv4 and IPv6 filter tables. Either may be independently changed to deny all. This will keep traffic from passing unless that traffic is specifically defined by additional filter rules. The only other option to change on the predefined rules is **chfilt** with the **–I** option, which allows packets matching that rule to be logged.

To support IKE tunnels, a dynamic filter rule is placed in the IPv4 filter table. This is the position at which dynamic filter rules are inserted into the filter table. This position can be controlled by the user by moving its position up and down the filter table. Once the tunnel manager daemon and **isakmpd** daemon are initialized that will allow IKE tunnels to be negotiated, rules are automatically created in the dynamic filter table to handle IKE messages as well as AH and ESP packets.

Host-Firewall-Host

The host–firewall–host configuration option for tunnels (see figure) allows you to create a tunnel between your host and a firewall, then automatically generate the necessary filter rules for correct communication between your host and a host behind the firewall. The auto–generated filter rules permit all rules between the two non–firewall hosts over the tunnel specified. The default rules—for user datagram protocol (UDP), Authentication Headers (AH), and Encapsulating Security Payload (ESP) headers—should already handle the host to firewall communication. The firewall will have to be configured appropriately to complete the setup. You should use the export file from the tunnel you created to enter the SPI values and keys that the firewall needs.

Host - FW - Host



Logging Facilities

This section describes the configuration and format of system logs relating to IP Sec. As hosts communicate with each other, the transferred packets may be logged to the system log daemon, **syslogd**. Other important messages about IP Sec will appear as well. An administrator may choose to monitor this logging information for traffic analysis and debugging assistance. The following are the steps for setting up the logging facilities.

1. Edit the /etc/syslog.conf file to add the following entry:

local4.debug var/adm/ipsec.log

Use the local4 facility to record traffic and IP Sec events. Standard AIX priority levels apply. You should set the priority level of debug until traffic through IP Security tunnels and filters show stability and proper movement.

Note: The logging of filter events can create significant activity at the IP Sec host and can consume large amounts of storage.

- 2. Save /etc/syslog.conf.
- 3. Go to the directory you specified for the log file and create an empty file with the same name. In the case above, you would change to /var/adm directory and issue the command:

```
touch ipsec.log
```

4. Issue a **refresh** command to the **syslogd** subsystem:

```
refresh -s syslogd
```

- 5. While creating filter rules for your host, if you would like packets matching a specific rule to be logged, set the –I parameter for the rule to Y (yes) using the **genfilt** or the **chfilt** commands.
- 6. Finally, turn on packet logging and start the ipsec_logd daemon using the following command:

```
mkfilt -g start
```

You can stop packet logging by issuing the following command:

```
mkfilt -q stop
```

The sample log file below contains traffic entries and other IP Security log entries:

- 1. Aug 27 08:08:40 host1 : Filter logging daemon ipsec_logd (level 2.20) initialized at 08:08:40 on 08/27/97A
- 2. Aug 27 08:08:46 host1 : mkfilt: Status of packet logging set to Start at 08:08:46 on 08/27/97
- 3. Aug 27 08:08:47 host1 : mktun: IBM tunnel 1, 9.3.97.244, 9.3.97.130 activated.
- 4. Aug 27 08:08:47 host1 skeyd: Inserted new context for tunnel ID 1 local SPI: 1336 remote SPI: 1336 .
- 5. Aug 27 08:08:47 host1 : mkfilt: #:1 permit 0.0.0.0 0.0.0.0 0.0.0.0 dudp eq 4001 eq 4001 both both l=n f=y t=0 e= a=
- 6. Aug 27 08:08:47 host1 : mkfilt: #:2 permit 0.0.0.0 0.0.0.0 0.0.0.0 ah any 0 any 0 both both 1=n f=y t=0 e= a=
- 7. Aug 27 08:08:47 host1 : mkfilt: #:3 permit 0.0.0.0 0.0.0.0 0.0.0.0 esp any 0 any 0 both both l=n f=y t=0 e= a=
- 8. Aug 27 08:08:47 host1 : mkfilt: #:4 permit 10.0.0.1 255.255.255.255 10.0.0.2 255.255.255 icmp any 0 any 0 local outbound l=y f=y t=1 e= a=
- 9. Aug 27 08:08:47 host1 : mkfilt: #:4 permit 10.0.0.2 255.255.255.255 10.0.0.1 255.255.255.255 icmp any 0 any 0 local inbound l=y f=y t=1 e= a=
- 10. Aug 27 08:08:47 host1 : mkfilt: #:6 permit 0.0.0.0 0.0.0.0 0.0.0.0 all any 0 any 0 both both l=y f=y t=0 e= a=
- 11. Aug 27 08:08:47 host1 : mkfilt: Filter support (level 1.00) initialized at 08:08:47 on 08/27/97
- 12. Aug 27 08:08:48 host1 : #:6 R:p o:10.0.0.1 s:10.0.0.1 d:10.0.0.20 p:udp sp:3327 dp:53 r:l a:n f:n T:0 e:n l:67

- 13. Aug 27 08:08:48 host1 : #:6 R:p i:10.0.0.1 s:10.0.0.20 d:10.0.0.1 p:udp sp:53 dp:3327 r:l a:n f:n T:0 e:n 1:133
- 14. Aug 27 08:08:48 host1 : #:6 R:p i:10.0.0.1 s:10.0.0.15 d:10.0.0.1 p:tcp sp:4649 dp:23 r:l a:n f:n T:0 e:n 1:43
- 15. Aug 27 08:08:48 host1 : #:6 R:p o:10.0.0.1 s:10.0.0.1 d:10.0.0.15 p:tcp sp:23 dp:4649 r:l a:n f:n T:0 e:n l:41
- 16. Aug 27 08:08:48 host1 : #:6 R:p i:10.0.0.1 s:10.0.0.15 d:10.0.0.1 p:tcp sp:4649 dp:23 r:l a:n f:n T:0 e:n 1:40
- 17. Aug 27 08:08:51 host1 : #:4 R:p o:10.0.0.1 s:10.0.0.1 d:10.0.0.2 p:icmp t:8 c:0 r:l a:n f:n T:1 e:n l:84
- 18. Aug 27 08:08:51 host1 : #:5 R:p i:10.0.0.1 s:10.0.0.2 d:10.0.0.1 p:icmp t:0 c:0 r:l a:n f:n T:1 e:n l:84
- 19. Aug 27 08:08:52 host1 : #:4 R:p o:10.0.0.1 s:10.0.0.1 d:10.0.0.2 p:icmp t:8 c:0 r:l a:n f:n T:1 e:n l:84
- 20. Aug 27 08:08:52 host1 : #:5 R:p i:10.0.0.1 s:10.0.0.2 d:10.0.0.1 p:icmp t:0 c:0 r:l a:n f:n T:1 e:n l:84
- 21. Aug 27 08:32:27 host1 : Filter logging daemon terminating at 08:32:27 on 08/27/971

The following paragraphs explain the log entries.

- Filter logging daemon activated.
- 2. Filter packet logging set to on with **mkfilt** –g start.
- 3. IBM tunnel activation, showing tunnel ID, source address, destination address, and time stamp.
- 4. The **skeyd** daemon inserted the tunnel context, meaning that the IBM tunnel is ready for traffic.
- 5–10. Filters have been activated. Logging shows all loaded filter rules.
- Message showing activation of filters.
- 12–13. These entries show a DNS lookup for a host.
- 14–16. These entries show a partial Telnet connection (the others have been removed from this example for space reasons).
- 17–20. These entries show two pings.
- 21. Filter logging daemon shutting down.

The example below shows two hosts negotiating a phase 1 and phase 2 tunnel, then passing Telnet traffic through the phase 2 tunnel just created.

- 1. Aug 9 12:32:26 host 1 Tunnel Manager: 0: TM is processing a
 Connection_request_msg
- 2. Aug 9 12:32:26 host 1 Tunnel Manager: 1: Creating new P1 tunnel object (tid)
- 3. Aug 9 12:32:26 host 1 Tunnel Manager: 0: Built a P1 init request msq
- 4. Aug 9 12:32:41 host 1 Tunnel Manager: 1: TM is processing a P1_sa_created_msg (tid)

- 5. Aug 9 12:32:41 host 1 Tunnel Manager: 1: Received good P1 SA, updating P1 tunnel (tid)
- 6. Aug 9 12:32:41 host 1 Tunnel Manager: 0: Checking to see if any tunnels P2 tunnels need to start
- 7. Aug 9 12:32:56 host 1 Tunnel Manager: 0: TM is processing a Connection_request_msg
- 8. Aug 9 12:32:57 host 1 Tunnel Manager: 0: Connection object contains a P2 request
- 9. Aug 9 12:32:57 host 1 Tunnel Manager: 0: Received a connection object for an active P1 tunnel
- 10. Aug 9 12:32:57 host 1 Tunnel Manager: 1: Created blank P2 tunnel (tid)
- 11. Aug 9 12:32:57 host 1 Tunnel Manager: 0: Added reference of new P2 to the P1 list
- 12. Aug 9 12:32:57 host 1 Tunnel Manager: 0: Checking to see if any P2 tunnels need to start
- 13. Aug 9 12:32:57 host 1 Tunnel Manager: 1: Starting negotiations for P2 (P2 tid)
- 14. Aug 9 12:33:11 host 1 Tunnel Manager: 0: TM is processing a P2_sa_created_msg
- 15. Aug 9 12:33:11 host 1 Tunnel Manager: 1: received P2_sa_created for an existing tunnel as initiator (tid)
- 16. Aug 9 12:33:11 host 1 Tunnel Manager: 0: Writing filter rules
- 17. Aug 9 12:35:54 host 1 Tunnel Manager: 0: TM is processing a List_tunnels_msg
- 18. Aug 9 13:01:31 host 1 $\,$ mkfilt: Status of packet logging set to Start at 13:01:31 on 08/09/98
- 19. Aug 9 13:01:32 host 1 ipsec_logd: Filter logging daemon ipsec_logd (level 2.20) initialized at 13:01:32 on 08/09/98
- 20. Aug 9 13:01:32 host 1 ipsec_logd: TC_LOG6: Tunnel interface module for IPv6 was started at 12:31:39 on 08/09/98
- 21. Aug 9 13:01:32 host 1 ipsec_logd: TC_LOG4: Tunnel cache for IPv4 was cleared at 12:31:39 on 08/09/98
- 22. Aug 9 13:01:32 host 1 ipsec_logd: TC_LOG4: Tunnel 1 with ESP SPI 300 and AH SPI 0 for IPv4 was activated at 12:33:11 on 08/09/98
- 23. Aug 9 13:03:14 host 1 mkfilt: Filter rules updated at 13:03:14 on 08/09/98
- 25. Aug 9 13:03:14 host 1 mkfilt: #:2 permit 0.0.0.0.0.0.0.0.0.0.0.0.0.0.0.0.0 all any 0 any 0 both both l=y f=n t=0

- 26. Aug 9 13:03:14 host 1 $\,$ mkfilt: Filter rules updated at 13:03:14 on 08/09/98
- 27. Aug 9 13:03:14 host 1 $\,$ mkfilt: Filter support (level 1.00) initialized at 13:03:14 on 08/09/98
- 28. Aug 9 13:03:26 host 1 ipsec_logd: #:7001 R:p 0:10.10.10.1 S:10.10.10.1 D:10.10.10.4 P:tcp SP:32793 DP:23 R:r I:tr0 F:n T:1 L:41
- 29. Aug 9 13:03:26 host 1 ipsec_logd: #:7001 R:p I:10.10.10.1 S:10.10.10.4 D:10.10.10.1 P:tcp SP:23 DP:32793 R:r I:tr0 F:n T:1 L:41
- 30. Aug 9 13:03:26 host 1 ipsec_logd: #:7001 R:p 0:10.10.10.1 S:10.10.10.1 D:10.10.10.4 P:tcp SP:32793 DP:23 R:r I:tr0 F:n T:1 L:40
- 31. Aug 9 13:03:26 host 1 ipsec_logd: #:7001 R:p 0:10.10.10.1 S:10.10.10.1 D:10.10.4 P:tcp SP:32793 DP:23 R:r I:tr0 F:n T:1 T::41

The following paragraphs explain the log entries.

- 1–4 The **ike cmd=activate phase=1** command initiates a connection.
- The Tunnel Manager receives a valid phase 1 security association from the responder.
- The Tunnel Manager checks whether **ike cmd=activate** has a phase 2 value for more work. It does not.
- 7–13 The **ike cmd=activate phase=2** command initiates a phase 2 tunnel.
- 14–15 The Tunnel Manager receives a valid phase 2 security association from responder.
- The Tunnel Manager writes the dynamic filter rules.
- 17 The **ike cmd=list** command views the IKE tunnels.
- 18–19 Packet logging is turned on.
- 20–22 The phase 2 tunnel is inserted into the tunnel cache.
- 23–27 The filter rules are updated.
- 28–31 Packets from a Telnet session are sent over the phase 2 tunnel.

Labels in Field Entries

The fields in the log entries are abbreviated to reduce DASD space requirements:

- # The rule number that caused this packet to be logged.
- **R** Rule Type.
 - **p** Permit.
 - **d** Deny.
- i/o Direction the packet was traveling when it was intercepted by the filter support code. Identifies IP address of the adapter associated with the packet:
 - For inbound (i) packets, this is the adapter that the packet arrived on.
 - For outbound (o) packets, this is the adapter that the IP layer has determined should handle the transmission of the packet.

- **s** Specifies the IP address of the sender of the packet (extracted from the IP header).
- **d** Specifies the IP address of the intended recipient of the packet (extracted from the IP header).
- p Specifies the high-level protocol that was used to create the message in the data portion of the packet. May be a number or name, for example: udp, icmp, tcp, tcp/ack, ospf, pip, esp, ah, or all.
- sp/t Specifies the protocol port number associated with the sender of the packet (extracted from the TCP/UDP header). When the protocol is ICMP or OSPF, this field is replaced with t, which specifies the IP type.
- dp/c Specifies the protocol port number associated with the intended recipient of the packet (extracted from the TCP/UDP header). When the protocol is ICMP, this field is replaced with **c** which specifies the IP code.
- Specifies that no information is available
- **r** Indicates whether the packet had any local affiliation.
 - f Forwarded packets
 - I Local packets
 - Outgoing
 - **b** Both
- I Specifies the length of a particular packet in bytes.
- **f** Identifies if the packet is a fragment.
- T Indicates the tunnel ID.
- i Specifies what interface the packet came in on.

Coexistence of IP Security and IBM Secured Network Gateway 2.2/IBM Firewall 3.1/3.2

If you installed AIX 4.3 with the IP Sec function, the IP Security function is disabled for IP Version 6 if you install an IBM Firewall product afterwards. The IBM Firewall product overrides this IP Sec function with its own implementation.

For related information, see Interoperability Notes, on page 4-32.

IP Security Problem Determination

This section includes some hints and tips that may assist you when you encounter a problem. We recommend that you set up logging from the start. Logs are very useful in determining what is going on with the filters and tunnels. (See Logging Facilities, on page 4-23 for detailed log information.)

Error: Issuing **mktun** command results in the following error:

```
insert_tun_man4(): write failed : The requested
resource is busy.
```

Problem: The tunnel you requested to activate is already active or you have colliding SPI values.

To fix: Issue the **rmtun** command to deactivate, then issue the mktun command to activate. Check to see if the SPI values for the failing tunnel match any other active tunnel. Each tunnel should have its own unique SPI values.

Error: Issuing **mktun** command results in the following error:

```
Device ipsec_v4 is in Defined status.

Tunnel activation for IP Version 4 not performed.
```

Problem: You have not made the IP Security device available.

To fix: Issue the following command:

```
mkdev -l ipsec -t 4
```

You may have to change **–t** option to 6 if you are getting the same error for Version 6 tunnel activation. The devices must be in available state. To check the IP Security device state, issue the following command:

```
lsdev -Cc ipsec
```

Error: Issuing a **chfilt** command results in the following error:

```
Cannot modify the first rule.
or
Cannot modify a pre_defined filter rule.
```

Problem: You are not allowed to modify these filter rules. You may however change whether they log or not.

To fix: If you want these rules to log, just issue the command:

```
chfilt -v (4 or 6) -n (filter number) -l y
```

If you want to set up the default rules to pass Authentication Header (AH) or Encapsulating Security Payload (ESP) header packets to specific hosts only, then you may prevent the autogeneration of rules by using the **–g** parameter with the **gentun** command. Then you may add in the same rules for the AH and ESP packets with the specific host's IP address for source and the partner host's IP address for destination. Make sure these rules are placed before the actual tunnel traffic rules.

Error: Issuing a **gentun** command results in the following error:

```
Invalid Source IP address
```

Problem: You have not entered a valid IP address for the source address.

To fix: For IP Version 4 tunnels, please check to see that you have entered an available IP Version 4 address for the local machine. You cannot use host names for the source when generating tunnels, you may only use host names for the destination.

For IP Version 6 tunnels, please check to see that you entered an available IP Version 6 address. If you type netstat -in and no IP Version 6 addresses exist, run /usr/sbin/autoconf6 (interface) for a link local auto-generated address (using MAC address) or use ifconfig to manually assign an address.

Error: Issuing **mktun** command results in the following error:

```
insert_tun_man4(): write failed : A system call
received a parameter that is not valid.
```

Problem: Tunnel generation occurred with invalid ESP and AH combination or without the use of the new header format when necessary.

To fix: Check to see what authentication algorithms are in use by the particular tunnel in question. Remember that the HMAC_MD5 and HMAC_SHA algorithms require the new header format. The new header format can be changed using the SMIT fast path <code>ips4_basic</code> or the <code>-z</code> parameter with the <code>chtun</code> command. Also remember that <code>DES_CBC_4</code> cannot be used with the new header format.

Troubleshooting IKE Tunnel Errors

IKE Tunnel Process Flow

The IKE tunnels are set up by the communication of the **ike** command or the Web-based System Manager VPN panels with two daemons:

tmd The Tunnel Manager daemon

isakmpd The ISAKMP daemon

For IKE tunnels to be properly set up, both daemons must be running. If IP Security is set to start at reboot, these daemons start automatically. Otherwise, they must be started manually.

The Tunnel Manager gives requests to **isakmpd** to start a tunnel. If the tunnel already exists or is not valid (for instance, has an invalid remote address), it reports an error. If negotiation has started, it may take some time, depending on network latency, for the negotiation to complete. The **ike cmd=list** command can list the state of the tunnel to determine if the negotiation was successful. Also, the Tunnel Manager logs events to **syslog** to the levels of **debug**, **event**, and **information**, which can be used to monitor the progress of the negotiation.

The sequence is:

- 1. Use Web-based System Manager or the **ike** command to initiate a tunnel.
- 2. The **tmd** daemon gives the **isakmpd** daemon a connection request for key management (phase 1).
- 3. The isakmpd daemon responds with SA created or an error.
- 4. The **tmd** daemon gives the **isakmpd** daemon a connection request for a data management tunnel (phase 2).
- 5. The **isakmpd** daemon responds with SA created or an error.

- 6. Tunnel parameters are inserted into the kernel tunnel cache.
- 7. Filter rules are added to the kernel dynamic filter table.

When the machine is acting as a responder, the **isakmpd** daemon notifies the Tunnel Manager **tmd** daemon that a tunnel has been negotiated successfully and a new tunnel is inserted into the kernel. In such cases, the process starts with step 3 and continues until step 7, without the **tmd** daemon issuing connection requests.

isakmpd Logging

The **isakmpd** daemon logs to a separate log because of the number and size of the logging messages. Logging is enabled using the **ike cmd=log** command. The /etc/isakmpd.conf configuration file can be set up to specify the output files for each logging level. Levels may be set to **none**, **errors**, **events**, and **information**.

The **isakmpd** daemon code will either initiate or respond by sending or evaluating a proposal. If the proposal is accepted, a security association is created and the tunnel will be set up. If the proposal is not accepted or the connection times out before the negotation completes, the daemon indicates an error. The entries in **syslog** from the **tmd** indicate whether the negotiation succeeded. To find out the exact cause of a failed negotiation, the **isakmpd** log needs to be checked.

Tracing facilities

Tracing is a debug facility for tracing kernel events. Traces can be used to get more specific information about events or errors occurring in the kernel filter and tunnel code.

SMIT has an IP Security trace facility available through the Advanced IP Security Configuration menu. The information captured by this trace facility includes information on Error, Filter, Filter Information, Tunnel, Tunnel Information, Capsulation/Decapsulation, Capsulation Information, Crypto, and Crypto Information. By design, the error trace hook provides the most critical information. The info trace hook can generate a lot of information and may have an impact on system performance. This tracing will provide clues to you as to what a problem may be. Tracing information will also be required when speaking with an IBM IP Security Technician. To access the tracing facility, use the SMIT fast path **smit ips4_tracing** (for IP Version 4) or **smit ips6_tracing** (for IP Version 6).

ipsecstat

You can issue the **ipsecstat** command to generate the following sample report. This sample report shows that the IP Security devices are in the available state, that there are three authentication algorithms installed, three encryption algorithms installed, and that there is a current report of packet activity. This information may be useful to you in determining where a problem exists if you are troubleshooting your IP Security traffic.

```
IP Security Devices:
ipsec_v4 Available
ipsec_v6 Available
Authentication Algorithm:
HMAC_MD5 -- Hashed MAC MD5 Authentication Module
HMAC SHA -- Hashed MAC SHA Hash Authentication Module
KEYED MD5 -- Keyed MD5 Hash Authentication Module
Encryption Algorithm:
CDMF -- CDMF Encryption Module
DES_CBC_4 -- DES CBC 4 Encryption Module
DES_CBC_8 -- DES CBC 8 Encryption Module
3DES_CBC -- Triple DES CBC Encryption Module
IP Security Statistics -
Total incoming packets: 1106
Incoming AH packets:326
Incoming ESP packets: 326
Srcrte packets allowed: 0
Total outgoing packets:844
Outgoing AH packets:527
Outgoing ESP packets: 527
Total incoming packets dropped: 12
 Filter denies on input: 12
  AH did not compute: 0
  ESP did not compute:0
  AH replay violation:0
 ESP replay violation:
Total outgoing packets dropped:0
 Filter denies on input:0
Tunnel cache entries added: 7
Tunnel cache entries expired: 0
Tunnel cache entries deleted: 6
```

Interoperability Notes

The following sections describe interoperability solutions. For related information, see Coexistence of IP Security and IBM Secured Network Gateway 2.2/IBM Firewall 3.1 or 3.2, on page 4-28.

IBM Firewall 3.1/3.2, IBM Secured Network Gateway (SNG) 2.2

The IBM Firewall 3.1/3.2, and IBM SNG 2.2 products operate as a tunnel partner with the IP Security feature of AIX 4.3. The tunnel may be created on the firewall and exported, then imported into an AIX 4.3 host running IP Security by using the —n option with the **imptun** command. There is, however, a script called **ipsec_convert** that is shipped as a sample shell script that transforms an IP Security tunnel export file into the necessary files needed by the IBM Firewall 3.1/3.2, or IBM SNG 2.2 to import.

There are several items to note when exporting a tunnel that will have the IBM Firewall 3.1/3.2 or the IBM SNG 2.2 as a tunnel partner. They are as follows:

- The IBM Firewall 3.1/3.2, and IBM SNG 2.2 only use the KEYED_MD5 algorithm for AH, and will not be able to import a tunnel specifying HMAC_MD5 nor HMAC_SHA1 for AH.
- The IBM Firewall 3.1/3.2 and IBM SNG 2.2 do not support replay prevention.
- The IBM Firewall 3.1/3.2 and IBM SNG 2.2 do not accept a tunnel lifetime of zero (unlimited lifetime).
- Also, IP Security will create tunnels with its own ordered numbers for tunnel IDs; when importing into the firewall, make sure these numbers are not already in use.
- The IBM Firewall 3.1/3.2 and IBM SNG 2.2 do not support IP Version 6 tunnels.

Make sure that SNG 2.2 has been been updated with the correct service level.

FTP Software's IP Security

FTP Software's TCP/IP stack and IP Security function will operate as a tunnel partner with the IP Security feature of AIX 4.3. Follow the instructions from FTP Software to add IP Security. From the FTP Software's IP Security configuration table, you can choose to add an address for setting up secure communication. After that, a page comes up with the IP Security configuration entry fields. The source AH SPI and shared secret key (for AH) have been generated for you, but you may enter the destination AH SPI and shared secret key in the fields provided. The page also contains autogenerated source ESP SPI and source ESP key. When the box for encryption is selected, the source ESP SPI and source ESP key are shown.

For interoperability, follow these steps:

- On the AIX 4.3 host, add a tunnel using FTP Software's IP Security parameters for the destination AH SPI and key and the destination ESP SPI and key.
- Note that FTP Software's IP Security configuration page allows only hexadecimal numbers into their entry fields. You will have to convert the AH and ESP SPI values generated by FTP Software into decimal numbers before giving them to AIX 4.3 IP Security.
- When entering SPI values and keys into FTP's IP Security configuration page, leave off the 0x. FTP Software will also drop any leading zeros.
- Note that the policy can only be authentication after encryption (encr/auth), authentication only (auth), or encryption only (encr).
- Only DES CBC 4 and DES CBC 8 can be used for encryption.
- Only Keyed MD5 should be used for authentication.
- Be careful entering the key values, if they are not entered correctly, the tunnel will not function.
- You will have to reboot the Windows 95 box that has FTP Software's IP Security function when a new tunnel is added or when an existing tunnel is changed.

IP Security Reference

List of Commands

ike cmd=activate Starts an Internet Key Exchange (IKE) negotiation (AIX versions

4.3.2 and later)

ike cmd=remove Deactivates IKE tunnels (AIX versions 4.3.2 and later)

ike cmd=list Lists IKE tunnels (AIX versions 4.3.2 and later)

Creates a tunnel definition aentun mktun Activates tunnel definition(s) chtun Changes a tunnel definition rmtun Removes a tunnel definition Istun Lists tunnel definition(s) exptun Exports tunnel definition(s) imptun Imports tunnel definition(s) Creates a filter definition genfilt mkfilt Activates filter definition(s)

mvfilt Moves a filter rule

chfiltChanges a filter definitionrmfiltRemoves a filter definitionIsfiltLists filter definition(s)expfiltExports filter definition(s)impfiltImports filter definition(s)ipsec_convertLists status of IP securityipsecstatLists status of IP security

ipsectrcbuf Lists the contents of IP security tracing buffer

unloadipsec Unloads a crypto module

List of methods

defipsec Defines an instance of IP Security for IP Version 4 or IP Version 6

cfgipsec Configures and loads ipsec_v4 or ipsec_v6

ucfgipsec Unconfigures ipsec_v4 or ipsec_v6

Chapter 5. TTY Devices and Serial Communications

This chapter contains information on managing tty terminal devices. Topics discussed are:

- TTY Overview, on page 5-2
- Managing TTY Devices, on page 5-4
- Dynamic Screen Utility, on page 5-6
- Modems, on page 5-13
- ATE Overview, on page 5-38
- Setting Up ATE, on page 5-40
- TTY Troubleshooting, on page 5-41

TTY Overview

A tty terminal device is a character device that performs input and output on a character—by—character basis. The communication between terminal devices and the programs that read and write to them is controlled by the tty interface. Examples of tty devices are:

- Modems
- ASCII terminals
- System console (LFT)
- aixterm under AlXwindows

The tty devices can be added, deleted, listed, and changed like any other device on your system by using the Web-based System Manager Devices application, the SMIT tool, or device—specific commands.

TERM Values for Different Displays and Terminals

Information about terminal capabilities is stored in the **terminfo** database. The value of the **TERM** environment variable identifies the specific terminal description in the **terminfo** database. This provides all information that a program needs for communicating effectively with the current tty device.

TERM Values for Various Terminals		
Display/Terminal	Value	
3161 ASCII Terminal	ibm3161	
3163 ASCII Terminal	ibm3161	
DEC VT100 (terminal)	vt100	
DECVT220	vt220	
3151 ASCII Display Station with Cartridge or 3161 ASCII Display Station with Cartridge	ibm3161–C	
3162 ASCII Display Station	ibm3161	
3162 ASCII Display Station with Cartridge	ibm3162	
6091 Display	lft	
AlXwindows	aixterm	

For information about the entries in the **terminfo** database, see the **terminfo** file format. To convert **termcap** entries into **terminfo** entries, see the **captoinfo** command. (The **termcap** file contains the terminal descriptions for older Berkeley systems.)

Setting TTY Characteristics

The *line discipline* provides the hardware—independent user interface for communicating between the computer and an asynchronous device. For example, a user is able to erase a single line or to interrupt a currently running process by typing a particular sequence of characters. You can define the meaning of these character sequences as well as set other terminal characteristics, such as the communication speed, by using the Web-based System Manager Devices application, the **chdev** command, the System Management Interface Tool (SMIT), or the **stty** command.

Most applications (including shells and editors) are designed to interface with terminals using the POSIX line discipline (default). You can change to the Berkeley line discipline by using the **stty** command.

Setting Attributes on the Attached TTY Device

For proper communication between the host and an attached tty device, the following are required:

- A properly wired communications cable
- Matching communications values (line speed, word length, parity, stop bit, and interface) between the host and the attached tty device

Managing TTY Devices

To do any task in the following table except List, a tty device must be installed.

Managing TTY Devices Tasks				
Web-based System Manager: wsm devices fast path (Devices application) -OR-				
Task	SMIT Fast Path	Command or File		
List Defined TTY Devices	smit Isdtty	Isdev -C -c tty -H		
Add a TTY	smit mktty	mkdev –t tty ^{1,2}		
Move a TTY to Another Port ³	smit movtty	chdev –I Name –p ParentName –w ConnectionLocation ^{2,4}		
Change/Show Characteristics of a TTY	smit chtty	Isattr –I Name –E (to show); chdev –I Name (to change) ⁴		
Remove a TTY ³	smit rmtty	rmdev –I Name		
Configure a Defined TTY (Make Available for Use)	smit mktty	mkdev –I Name		

Note:

1. Other flags may be used to further specify the new tty device. For example, to define and configure an RS-232 tty device connected to port 0 on the 8-port asynchronous adapter sa3 with the speed attribute set to 19200 and other attributes set to values retrieved from the foo file:

```
mkdev -t tty -s rs232 -p sa3 -w 0 -a speed=19200 -f foo
```

- The **mkdev** and **chdev** commands support options that are not possible with Web-based System Manager or SMIT.
- 3. Disable the tty before doing this task. Refer to the **pdisable** command.
- 4. Use flags to change specific characteristics about a tty from the command line.

If adding or changing a tty from the command line, consult the following list to find out the *Attribute* name you should specify in the **–a** *Attribute*= *Value* flag for the characteristic you want to set. For example, specify –a speed= Value to set the baud rate of a tty device.

Characteristic	Attribute Name
Enable LOGIN	login
BAUD rate speed	speed
PARITY	parity
BITS per character	bpc
Number of STOP BITS	stops
TIME before advancing to next port setting	timeout
XON-XOFF handshaking	xon
TERMINAL type	term
FLOW CONTROL to be used	flow_disp
OPEN DISCIPLINE to be used	open_disp
STTY attributes for RUN time	runmodes

STTY attributes for LOGIN	logmodes
RUN shell activity manager	shell
LOGGER name	logger
STATUS of device at BOOT time	autoconfig
TRANSMIT buffer count	tbc
RECEIVE trigger level	rtrig
STREAMS modules to be pushed at open time	modules
INPUT map file	imap
OUTPUT map file	omap
CODESET map file	csmap
INTERRUPT character	intr
QUIT character	quit
ERASE character	erase
KILL character	kill
END OF FILE character	eof
END OF LINE character	eol
2nd END OF LINE character	eol2
DELAY SUSPEND PROCESS character	dsusp
SUSPEND PROCESS character	susp
LITERAL NEXT character	Inext
START character	start
STOP character	stop
WORD ERASE character	werase
REPRINT LINE character	reprint
DISCARD character	discard

Dynamic Screen Utility

The dynamic screen utility, or **dscreen** command, is a utility that allows a single physical terminal to be connected to several virtual terminal sessions (screens) at one time. It is mainly intended for use with terminals that have two or more pages of screen memory (for example, the IBM 3151 Models 310 or 410 display with the Cartridge for Expansion). With such terminals, switching between virtual screens also switches between physical terminal screen pages allowing each virtual screen's image to be saved and restored. On terminals without multiple pages of screen memory, **dscreen** can still be used to switch among virtual screen sessions although the appearance of the screen will not be maintained.

Note: For full support of **dscreen**, the terminal must be able to switch internal screen pages on command and must remember the cursor position for each page. While **dscreen** will work on both smart and dumb terminals, screen images are not saved during screen changes on dumb terminals.

dscreen Terminal Configuration Information File

The **dscreen** terminal configuration information file (or **dsinfo** file) is used to define a different set of keys to be used with **dscreen**. This might be done, for example, when the originally defined **dscreen** keys conflict with a software application in use on the system.

The **dsinfo** file terminal type assumes a single page of screen memory. Therefore, if a terminal supports additional pages of screen memory, the **dsinfo** file must be customized to use the appropriate sequence for page memory control. Consult the appropriate terminal reference guide for the specific control sequence.

The default **dsinfo** file is /**usr/lbin/tty/dsinfo**. Use the **–i** flag to specify a different **dsinfo** file. This remainder of this section will refer to the default file. However, the same information applies to any customized **dsinfo** file you create.

For more information concerning the dsinfo file, refer to "Dynamic Screen Assignment", on page 5-8.

Key Action Assignments

When **dscreen** is run, it starts a virtual screen. Some of the keys on the terminal keyboard will not be passed through to the virtual screen; instead, **dscreen** will intercept these keys and perform certain actions when they are pressed. The actions include:

Select Selects a specified screen

Block Blocks all input and output.

New Starts a new screen session.

End Ends dscreen.

Quit Quits dscreen.

Previous Switches to previous screen.

List Lists the **dscreen** assigned keys and their actions.

The function of each key is dependent on the terminal and the terminal description in the /usr/lbin/tty/dsinfo file.

Select Keys

When a new virtual screen is created, it is assigned a select key. Pressing the select key causes the following actions:

- A switch from the physical terminal to the video page associated with the particular virtual screen.
- Input and output is directed appropriately between the physical terminal and the virtual screen.

Once all of the select keys defined in the **dsinfo** file have virtual screens assigned to them, no more screens may be created. Individual screen sessions end when the original shell process exits. This frees the associated select key for use with another virtual screen. The **dscreen** utility is ended when there are no more active screens.

Block Keys

Block keys are used to stop output in a fashion similar to Ctrl–S key when using IXON flow control. The purpose of these keys is to allow for transparently setting up terminal sessions on two computers using a terminal that has two serial ports.

New Keys

Pressing a new screen key creates a new logical screen and assigns it to one of the select keys. Each new screen requires:

- A select key as defined in the dsinfo file
- A dscreen pseudo-terminal device
- Enough memory for the various structures used in screen tracking
- · A process to run the shell from

If any of these are not available, the new screen operation will fail with a message indicating the reason for the failure.

End and Quit Keys

Pressing an end key will cause the following to occur:

- Send a SIGHUP signal to all the screen sessions
- · Clean up
- Exit with a status of 0

Pressing a guit key will perform the same actions but will exit with status of 1.

Previous Kev

Pressing a previous key will switch the terminal to the screen that was last displayed.

Note:

- 1. Do not switch screens when the current screen is being written to; an escape sequence may be truncated and leave the terminal in an unknown state.
- Some terminal displays may save the cursor position for individual screens but may not save other states such as insert mode, inverse video, etc. If this is the case, users should avoid these modes while switching screens.

List Key

Pressing a list key will display a list of keys and their actions on the terminal display. Only those keys recognized by **dscreen** will be shown. When a new screen is created using **dscreen**, the message Press KEY for help, where KEY is the name of the list key displayed on the terminal. Note that the message is displayed *only* if there is a list key defined.

Dynamic Screen Assignment

The terminal description entry in the /usr/lbin/tty/dsinfo file will have the same number of screen selection keys as the terminal has physical screen pages. If more screen selection keys are defined than the number of physical screen pages, dscreen will dynamically assign physical screen pages to virtual screens.

When a virtual screen is selected that does not have an associated page of screen memory, **dscreen** assigns the least recently used physical screen to the virtual screen. Depending on the specifications maintained in the /usr/lbin/tty/dsinfo description file, an indication that the physical screen is connected to a different virtual screen may be noticeable; for example, the screen is cleared.

dsinfo File

The **dsinfo** file is a database of terminal descriptions used by the **dscreen** multiple screen utility. The file contains the following information:

- dscreen keys and the functions they perform
- · Number of screen memory pages for the terminal
- · Code sequences sent or received to use the above features

The terminal type entries in the default **dsinfo** file resemble the following 3151 ASCII terminal values:

```
# The Cartridge for Expansion (pn: 64F9314) needed for this entry
ibm3151|3151|IBM 3151,
dsks=\E!a^M|Shift-F1|,  # Selects first screen dsks=\E!b^M|Shift-F2|,  # Selects second screen dsks=\E!c^M|Shift-F3|,  # Selects third screen dsks=\E!d^M|Shift-F4|,  # Selects fourth screen dskc=\E!e^M|Shift-F5|,  # Creates a new screen
dske=\E:f^M|Shift-F6|\E pA\EH\EJ, # Go to screen 1 and end
dskq=\E!i^M|Shift-F9|\E pA\EH\EJ, # Go to screen 1 and quit
dsp=\E pA|\EH\EJ,
                                         # Terminal sequence for screen 1
dsp=\E pB|\EH\EJ,
                                        # Terminal sequence for screen 2
dsp=\E pC|\EH\EJ,
                                       # Terminal sequence for screen 3
                                        # Terminal sequence for screen 4
dsp=\E pD|\EH\EJ,
dst=10,
                                        # Allow 1 second timeout buffer
```

Entry Format for dsinfo

Entries in the **dsinfo** file consist of comma–separated fields. The first field is a list of alternate names for the terminal, each name is separated by a pipe (|) character. Any text preceded by a pound (#) character is regarded as a comment and ignored by **dscreen**. The remaining fields are strings describing the capabilities of the terminal to the **dscreen** utility. Within these strings, the following escape codes are recognized:

Escape Sequence	Description
\E,\e	escape character
\n,\I	newline (or linefeed) character
\r	carriage return
\t	tab character
\p	backspace character
\f	formfeed character
\s	space character
\nnn	character with octal value nnn
^ <i>X</i>	Ctrl-x for any appropriate x value

Any other character preceded by a backslash will yield the character itself. The strings are entered as *type=string*, where *type* is the type of string as listed below, and *string* is the string value.

It is important that the entry fields in the **dsinfo** file be separated by commas. If a comma is omitted or truncated from the end of a **dsinfo** file entry, the file will become unreadable by the **dscreen** utility and an error will be returned to the display.

String Types

The string types are as follows:

dskx

A string type that starts with dsk describes a key. The type must be four letters long, and the fourth letter *x* indicates what action is taken when the key is received. The key types are:

Туре	Action
dsks	Switch Screens
dskb	Block Input and Output
dske	End dscreen
dskq	Quit dscreen (exit status=1)
dskc	Create New Screen
dskp	Switch to Previous Screen
dskl	List Keys and Actions

Any other key type (that is, a string type dskx that does not end in s, b, e, q, p, or I) will cause no internal **dscreen** action, but will show up in the key listing and will be recognized and acted on. A type of dskn (n for No Operation) should be used when no internal **dscreen** action is desired.

The value string for each key has three substrings, which are separated by pipe (|) characters.

Note: Use \| to include the | character in one of the substrings.

The first substring is the sequence of characters that the terminal sends when the key is pressed. The second substring is a label for the key that is printed when a list of keys is displayed. The third substring is a sequence of characters that dscreen sends to the terminal when this key is pressed before performing the action this key requests.

A string type of dsp describes a physical screen in the terminal. One dsp string should be present for each physical screen in the terminal. The value string for each physical screen has two substrings, which are separated by a pipe (|) character.

The first substring is the sequence of characters to send to the terminal to display and output to the physical page on the terminal.

The second substring is sent to the terminal when the page is used for something new. This second substring is often set to the clear screen sequence. It is sent under the following two conditions:

- 1. When a new virtual terminal session is being created.
- 2. When there are more virtual terminals than there are physical screens. If a virtual terminal is selected that requires **dscreen** to reuse one of the physical screens, it will send this sequence to the screen to indicate that the screen contents do not match the output of the virtual terminal connected.

Note: Running with more virtual terminals than physical screens can be confusing and is not recommended; it can be avoided by defining no more screen selection keys (dsks=) than physical screens (dsp=) in the dsinfo entry.

String with a type of dst adjusts **dscreen**'s input timeout. The value of the string is a decimal number. The timeout value is in tenths of seconds and has a maximum value of 255 (default=1 [or 0.1 seconds]).

When **dscreen** recognizes a prefix of an input key sequence but does not have all the characters of the sequence, it will wait for more characters to be sent until it is recognizable. If the timeout occurs before more characters are received, the characters are sent on to the virtual screen and dscreen will not consider these characters as part of an input key sequence.

It may be necessary to raise this value if one or more of the keys dscreen is to trigger on is actually a number of keystrokes (that is assigning Ctrl-Z 1, Ctrl-Z 2, Ctrl-Z 3, etc., for screen selection and Ctrl-Z N for new screen and so on).

dsp

dst A

Example 1

The following example /usr/lbin/tty/dsinfo entry is for a Wyse-60 with three screen sessions:

```
wy60|wyse60|wyse model 60,
dsks=^A \^M|Shift-F1|,
dsks=^Aa^M|Shift-F2|,
dsks=^Ab^M|Shift-F3|,
dskc = \200|Ctrl-F1|,
dske=\201|Ctrl-F2|\Ew0\E+,
dskl=\202|Ctrl-F3|,
dsp=\Ew0|\E+,
dsp=\left| Ew1 \right| \left| E+, \right|
dsp=\langle Ew2 | \langle E+,
```

With this entry:

- Shift–F1 through Shift–F3 are used for selecting screens 1 through 3.
- Ctrl–F1 creates a new screen.
- Ctrl-F2 sends: Esc w 0 Esc + to the screen (switching to window 0 and clearing the screen) and then ends dscreen.
- Ctrl-F3 lists the keys and their functions.

Each time a physical screen is used for a new screen, the sequence Esc + will be sent to the terminal, which will clear the screen.

Example 2

This example is for a Wyse-60 with three screen sessions, but one of the screens is on a second computer communicating through the second serial port on the terminal:

```
wy60-1|wyse60-1|wyse model 60 - first serial port
dsks=^A \^M|Shift-F1|,
dsks=^Aa^M|Shift-F2|,
dsks=^Ab^M|Shift-F3|\Ed\#^Ab\r^T\Ee9,
dskc=\200|Ctrl-F1|,
dske=\201|Ctrl-F2|\Ed\#\201^T\Ew0\E+,
dskl=\202|Ctrl-F3|,
dsp=\Ew0|\E+, dsp=\Ew1|\E+,
wy60-2|wyse60-2|wyse model 60 - second serial port
dsks=^A ^M|Shift-F1|\Ed\#^A ^r^T\Ee8,
dsks=^Aa^M|Shift-F2|\Ed\#^Aa\r^T\Ee8,
dsks=^Ab^M|Shift-F3|,
dskc=\200|Ctrl-F1|,
dske=\201|Ctrl-F2|\Ed\#\201^T\Ew0\E+
dskl=\langle 202|Ctrl-F3|,
dsp=\langle Ew2 | \langle E+ ,
```

dscreen must be run on both computers, with terminal type wy60–1 on the first computer and terminal type wy60-2 on the second computer (using the -t option to dscreen). The wy60-1 entry will be examined first.

The first two key entries are unchanged from the original wy60 entry. The third key. however, has type dskb, which means block both input and output. When this key is pressed, the sequence:

```
Esc d # Ctrl-A b CR Ctrl-T Esc e 9
```

is sent to the terminal; after this output is blocked and dscreen continues scanning input for key sequences but discards all other input.

The sequence Esc d # puts the terminal in transparent print mode, which echoes all characters up to a Ctrl-T out through the other serial port.

The characters Ctrl-A b CR are sent out the other serial port, informing the dscreen process on the other computer that it should activate the window associated with the Shift-F3 key.

The Ctrl-T key sequence exits the transparent print mode. The Esc 9 key sequence causes the terminal to switch to the other AUX serial port for data communications.

At this point, the other computer takes over, sends an Esc w 2 to switch to the third physical screen, and then resumes normal communication.

The wy60–2 entry follows the same general pattern for keys Shift–F1 and Shift–F2:

- Switch to transparent print mode
- · Send function key string to other computer
- Switch transparent print off
- Switch to the other serial port

The end key, Ctrl-F2, works the same for both computers; it sends the end key sequence to the other computer through the transparent print mechanism, switches the terminal to window 0, clears the screen, then exits.

Modems

Modems provide serial communications across ordinary telephone lines. This section discusses modem standards, general modem setup, and specific configuration tips for popular modems.

Modem Overview

A modem is a device that allows you to connect one computer to another across ordinary telephone lines. The current telephone system is incapable of carrying the voltage changes required for a direct digital connection. A modem overcomes this limitation by modulating digital information into audio tones for transmission across the phone line, and by demodulating those tones back into digital information on reception. Modems are commonly used with Basic Network Utilities (BNU) or other implementations of the UNIX-to-UNIX Copy Program (UUCP). A high-speed (14,400 bps or greater) modem can be used with Serial Line Interface Protocol (SLIP) to provide Transmission Control Protocol/Internet Protocol (TCP/IP) connectivity as well.

Often, the term baud is used to refer to a modem's speed instead of bps. Baud is actually a measurement of the modulation rate. In older modems, only 1 bit was encoded in each signal change, so a modem's baud rate was equal to the modem's speed. Modems that operate at higher speeds, however, still generally operate at 2400 (or even 1200) baud, and encode two or more bits per signal change. A modem's bps rate is calculated by multiplying the number of data bits per signal with the baud (for example, 2400 baud x 6 bits per signal change = 14,400 bits per second). Most modern modems can communicate at a variety of speeds (for example, 14,400, 9600, 7800, 4800, and 2400 bps).

Telecommunications Standards

The older speeds of 300, 1200, and 2400 bps were well defined. However, as modem manufacturers began to devise methods for gaining higher speeds, each modem manufacturer started to use a proprietary method incompatible with modems from other manufacturers. Today, the ITU-TSS (formerly the United Nations Consultative Committee for International Telephony and Telegraphy, abbreviated CCITT) defines standards for most high-speed communications.

Even high-speed modems are much slower than other methods of computer communication. A high-speed modern may operate at 28,800 bps, but an Ethernet connection operates at 10,000,000 bps. In order to boost data throughput, high-speed modems typically offer one or more data compression algorithms. These algorithms can boost the throughput of a high-speed modem to speeds of 57,600 bps (if the data rate is 14,400 bps) or 115,200 bps (if the data rate is 28,800 bps). Note that these compression algorithms are sensitive to the data being transmitted. If the data has already been compressed (for example, with the compress command), the data compression methods of high-speed modems will offer little or no benefit, and might even reduce data throughput. When using a modem with data compression technology, the speed of the data terminal equipment/data circuit-terminating equipment (DTE/DCE) connection between the computer and the modem should be greater than the nominal data rate of the connection between modems, and equal . For example, with a V.32bis modem with V.42bis data compression, the data rate of the modem (the speed at which the modem communicates across telephone lines) is 14,400 bps. When the V.42bis compression is active, actual data throughput can reach 57,600 bps. To accommodate the greater throughput offered by data compression, the speed of the DTE/DCE between the computer and the modem should be set to 57,600 bps.

Attention: Some modems implementing data compression and modern modulation schemes may yield a higher data throughput than some systems and asynchronous adapters can accommodate.

Today, the ITU-TSS (formerly, the CCITT) defines standards for high-speed communications, including data compression algorithms. ITU-TSS standards are usually named V.nn, where nn is a number. Another, slightly less common standard is the Microcom Networking Protocol (MNP). Available in versions (called classes) 1-9, MNP is a high-performance, high-speed protocol that was available relatively early, and became something of a de facto standard before the advent of the CCITT standards.

ITU-TSS Communications Standards

Following is a list of some common communications standards defined by the ITU-TSS. Note that this only a partial list.

V.29 ITU-TSS standard for half-duplex 9600 bps

communications.

V.32 ITU-TSS standard for full-duplex 9600 bps

communications.

V.32bis ITU-TSS standard for 14,400 communications. V.32bis is a

revision to the V.32 standard.

V.34 ITU-TSS standard for 28,800 bps communications. Note

> that this standard achieves 28,800 bps data rates using multiple bit encoding, instead of the data compression scheme used by MNP Class 9. This standard was

previously referred to as V.fast.

V.42 ITU-TSS data compression standard.

V.42bis Revised ITU-TSS data compression standard.

MNP Communications Standards

MNP Class 1 An asynchronous, half-duplex, byte-oriented method of

transferring data realizing about 70% efficiency. Uncommon

in modern modems.

MNP Class 2 A full-duplex counterpart to MNP Class 1. Uncommon in

modern modems.

MNP Class 3 A synchronous, bit-oriented full-duplex method of

> transferring data realizing about 108% efficiency. (Efficiency greater than 100% is realized because the start/stop bits required for an asynchronous connection are eliminated. The DTE/DCE between the modem and the system are still

asynchronous).

MNP Class 4 An enhancement to MNP Class 3 including a mechanism

> for varying the packet size (adaptive packet assembly) and a means of eliminating redundant administrative overhead (data phase optimization). An MNP Class 4 modem offers

approximately 120% efficiency.

MNP Class 5 Class 5 includes data compression along with Class 4

features. An MNP Class 5 modem offers 200% efficiency.

MNP Class 6 MNP Class 6 allows incorporation of multiple, incompatible

> modulation techniques into one modem (universal link negotiation). This allows MNP Class 6 modems to begin communication at a slower speed and negotiate a transition to a higher speed. Class 6 also includes a statistical duplexing scheme that dynamically allocates utilization of half-duplex modulation to simulate full-duplex service. All

features of MNP Class 5 are supported.

MNP Class 7 Incorporates enhanced data compression. Combined with

Class 4, efficiencies of 300% can be realized.

MNP Class 8 N/A

MNP Class 9 Combines enhanced data compression with V.32

technology to allow data rates up to 28,800 bps.

Generic Modems

To set up a modem:

1. Attach the modem with the appropriate cables.

2. Add a tty for the modem.

3. Configure the modem.

Attaching the Modem with Appropriate Cables

The first step in setting up a modem is to attach the modem with the appropriate cables. Refer to the documentation of the system unit for information about the cables you should use.

Adding a TTY for the Modem

First, ensure that the system is turned on and that the modem is turned off. Use the Web-based System Manager fast path, wsm devices, or the SMIT fast path smit mktty. Example port settings for the tty are given below:

Parent Adapter SA₀ *Port Number {0}

STATE to be configured at {available}

boot time

Enable LOGIN disable **BAUD** rate {9600} **PARITY** {8} BITS per character {8} Number of STOP BITS {1} TIME before advancing to {0}

next port setting

XON-XOFF handshaking yes **TERMINAL** type {dumb}

Note that the port number, the baud rate, and the enabling of login depend on your setup:

- PORT number is dependent on the adapter/port to which the modem is attached.
- BAUD Rate is dependent on the capabilities of the modem.
- Enable LOGIN is dependent on port usage.

The Enable LOGIN parameter can have the following settings:

Setting Description

DISABLE The tty is set to off in the /etc/inittab file. No getty is started

on this port at system startup.

ENABLE The port starts a getty immediately. This setting is used

most often for dial-in-only modems, and terminal logins.

Description Setting

SHARE The port starts a getty when a carrier signal is asserted by a

> DCE. (A modem is a DCE.) A shared port can be used as both dial-out and dial-in without user intervention, but the DCE device must not assert carrier until true carrier is

established.

DELAY The port does not start a getty until a character is received

in the input buffer, that is, the DCE sends a character to the

DTE or passes a character through from the remote

DCE/DTE. (A delayed port can be used for both dial-in and dial-out but the user dialing into a delayed port will need to send a character or two before a login herald will be issued

to the port.)

Configuring the Modem

Two methods for configuring the modem are presented in this section. You only need to use one method.

Sending AT Commands with the cu Command

If you have the Basic Network Utilities (BNU) installed, you can use the cu command to configure a modem as follows:

1. Add the following line to your /usr/lib/uucp/Devices file. Do not add the line if it is already in the file. This should work with most Hayes-compatible modems. (Replace # with the number for your port.)

```
Direct tty# - Any direct
```

2. Enter the following commands (read the notes about each command before performing the command itself):

pdisable tty# cu -1 tty# Use cu -ml tty# if ix24051 is installed for **cu**. AT&F Restore the factory default configuration. In command state, echo characters from the keyboard to the screen. ATE1 (Make sure carrier is not ON on the port or modem.) **Note:** You need issue only one of the following commands. AT&D2 Monitor the data terminal ready (DTR) signal. When an on-to-off transition of the DTR signal occurs, the modem hangs up and enters the command state. Monitor the DTR signal. When an on-to-off transition of the DTR signal AT&D3 occurs, the modem hangs up and resets. AT&W Write the storable parameters of the current configuration to the modem's memory. Track the status of the carrier detect signal. (The modem may AT&C1 disconnect.) Note: Issue the following command only if the previous command used the modem to disconnect. cu -1 tty#

Set the carrier–detect response time. Default is 6. Possible values are 1

Put the modem in autoanswer mode.

to 255.

ATS0=1

ATS9=12

Write the storable parameters of the current configuration to the AT&W modem's memory.

Terminate the **cu** session.

3. Enter *one* of the following commands:

```
penable tty#
pshare tty#
pdelay tty#
pdisable tty#
```

The modem now has the basic configuration needed for most system communications. If you have problems, invoke cu with the cu -dl command to start a diagnostic trace on the connection.

Sending AT Commands Using a C Program

If the previous method failed, or if you do not have BNU installed, try running the following C program. Create a file called motalk.c containing the following code. Save the file. Compile and run it according to the instructions in the program comments.

```
/* MoTalk - A "C" program for modem setup.
                                                      * /
/*
          This program is meant as an aid only and is
                                                      * /
/*
                                                      * /
           not supported by IBM.
           compile: cc -o motalk motalk.c
/*
                                                      */
                 Usage: motalk /dev/tty? [speed]
                                                      * /
/*
#include <errno.h>
#include <stdio.h>
#include <signal.h>
#include <fcntl.h>
#include <termio.h>
FILE *fdr, *fdw;
int fd;
struct termio term save, stdin save;
void Exit(int sig)
{
  if (fdr) fclose(fdr);
  if (fdw) fclose(fdw);
  ioctl(fd, TCSETA, &term_save);
  close (fd);
  ioctl(fileno(stdin), TCSETA, &stdin_save);
  exit(sig);
}
main(int argc, char *argv[])
  char *b, buffer[80];
  int baud=0, num;
  struct termio term, tstdin;
  if (argc < 2 || !strcmp(argv[1], "-?"))
     fprintf(stderr, "Usage: motalk /dev/tty? [speed]\n");
     exit(1);
  if ((fd = open(argv[1], O_RDWR | O_NDELAY)) < 0)</pre>
  {
     perror(arqv[1]);
     exit (errno);
  if (argc > 2)
     switch(atoi(argv[2]))
        case 300: baud = B300;
```

```
break;
      case 1200: baud = B1200;
                  break;
      case 2400: baud = B2400;
                  break;
      case 4800: baud = B4800;
                  break;
      case 9600: baud = B9600;
                  break;
      case 19200: baud = B19200;
                  break;
      case 38400: baud = B38400;
                  break;
      default:
                  baud = 0;
                  fprintf(stderr, "%s: %s is an unsupported
                  baud\n'', argv[0], argv[2]);
                  exit(1);
      }
   }
/* Save stdin and tty state and trap some signals */
ioctl(fd, TCGETA, &term_save);
ioctl(fileno(stdin), TCGETA, &stdin_save);
signal(SIGHUP, Exit);
signal(SIGINT, Exit);
signal(SIGQUIT, Exit);
signal(SIGTERM, Exit);
/* Set stdin to raw mode, no echo */
ioctl(fileno(stdin), TCGETA, &tstdin);
tstdin.c_iflag = 0;
tstdin.c_lflag &= ~(ICANON | ECHO);
tstdin.c\_cc[VMIN] = 0;
tstdin.c_cc[VTIME] = 0;
ioctl(fileno(stdin), TCSETA, &tstdin);
/* Set tty state */
ioctl(fd, TCGETA, &term);
term.c_cflag |= CLOCAL|HUPCL;
if (baud > 0)
   term.c cflaq &= ~CBAUD
   term.c_cflag |= baud;
term.c_lflag &= ~(ICANON | ECHO); /* to force raw mode */
term.c_iflag &= ~ICRNL; /* to avoid non-needed blank lines */
term.c_cc[VMIN] = 0;
term.c_cc[VTIME] = 10;
ioctl(fd, TCSETA, &term);
fcntl(fd, F_SETFL, fcntl(fd, F_GETFL, 0) & ~O_NDELAY);
/* Open tty for read and write */
if ((fdr = fopen(argv[1], "r")) == NULL)
   perror(argv[1]);
  exit (errno);
if ((fdw = fopen(argv[1], "w")) == NULL)
   perror(argv[1]);
   exit (errno);
/* Talk to the modem */
puts ("Ready... ^C to exit");
while (1)
   if ((num = read(fileno(stdin), buffer, 80)) > 0)
      write(fileno(fdw), buffer, num);
```

```
if ((num = read(fileno(fdr), buffer, 80)) > 0)
         write(fileno(stdout), buffer, num);
      Exit (0);
   }
}
```

Hayes and Hayes-Compatible Modems

1. Change the tty settings using the Web-based System Manager fast path, wsm devices, or the SMIT fast path, smit chtty.

Example settings for the modem follow. (tty0 is the tty for this example and the baud rate may be changed depending on your modem. Generally, you want to keep the baud rate at the highest speed that the modem can allow.)

TTY tty0 TTY type tty **TTY** interface rs232

Description Asynchronous Terminal

Status Available Location 00-00-s1-00

Parent adapter sa0 **PORT** number s1 **Enable LOGIN** share **BAUD** rate 2400 **PARITY** none BITS per character 8 **Number of STOP BITS** 1 TIME before advancing to 0 next port setting

XON-XOFF handshaking no **TERMINAL** type dumb **INPUT** map file none **OUTPUT** map file none **CODESET** map file sbcs

STTY attributes for RUN TIME should be:

[hupcl,cread,brkint,icrnl,opost,tab3,onlcr,icanon,echo,echoe,echo k,echoctl, echoke,imaxbel,iexten]

(No ixon/ixoff needed)

STTY attributes for LOGIN should be:

[hupcl, cread, echoe, cs8]

(No **ixon/ixoff** needed)

2. Add the following line to /usr/lib/uucp/Systems file:

```
hayes Nvr HAYESPROG 2400
```

3. Add this to /usr/lib/uucp/Devices file:

```
# For programming the hayes modem only:
HAYESPROG tty0 - 2400 HayesProgrm2400
#regular ACU entry:
ACU tty0 - Any hayes
```

4. Add this to /usr/lib/uucp/Dialers file:

```
# This Entry is used to PROGRAM the modem ONLY:
# the next 3 lines should be made into one:
HayesProgrm2400 =,-, "" \d\dAT\r\c OK AT&F\r\c OK
ATM1\r\c OK
AT&D3\r\c OK AT&K3&C1\r\c OK ATL0E0Q2\r\c OK ATS0=1\r\c OK
AT&W\r\c
OK
                  "" \dAT\r\c OK ATDT\T\d\r\c CONNECT
hayes
```

5. To program the modem, enter the command cu -d hayes. This command uses cu to program the modem. Since no connection is made to another system, the command will fail. The modem will be programmed if sendthem AT&W and then OK got it appear in the output.

If you are not doing binary file transfers or using BNU, leave out the &K3 command, and set XON-XOFF handshaking to Yes in SMIT or checked in Web-based System Manager. However, it is more efficient to use hardware flow control (as opposed to XON-XOFF handshaking). To do that, use the settings and the Dialers entries from step 5.

After the modem is programmed, you must set request to send (RTS) line discipline to allow the system device driver to use hardware flow control. To do that, the modem must be connected to another system (to bring carrier high) and the following command must be run:

```
stty add rts < /dev/tty0
```

Be sure to substitute tty0 with your tty number. This works until the next reboot, when it must be done again. For the addition of RTS to be done automatically at system startup, type in and compile the addrts.c program that follows, and place a line in /etc/rc that calls the program on system startup. (This program does not require carrier to be high.) The command is:

```
addrts /dev/tty0
```

The full path to the command must be specified if the addrts program is not in one of the directories in the PATH. (Substitute your tty number for tty0 above.)

Adding RTS Discipline to TTY Ports

```
C Program to add RTS discipline to tty port(s).
   Tips:
         To make the program more permanent, insert the file name
         of the compiled version (complete with path) at the end of
         your "/etc/rc" file and the changes will take effect again
         at next reboot.
         NOTE: This program is supplied "as is" and is NOT supported
              by IBM. It is intended as an aid to administrators
              only.
      To create: vi addrts.c <enter>
      To compile: cc -o addrts addrts.c Usage is: addrts /dev/tty##
                            Program starts now
#include <stdio.h>
#include <fcntl.h>
#include <termios.h>
#include <sys/tty.h> main (argc,argv)
   int argc;
   char *argv[];
   int fd;
   if ( (fd = open(argv[1], O_NDELAY|O_RDWR)) <0 )</pre>
      printf("%s: could not open %s\n", argv[0], argv[1]);
      exit (22);
   ioctl(fd, TXADDCD, "rts");
   /* adds rts to the tty in the argument */
   close(fd);
```

Specific Modem Configuration Tips

Following are some specific configuration tips and setup information for a variety of popular modems.

Setup Information for the IBM 7855 Modem

Note: This guide uses tty0 as an example. Substitute the tty to which the modem is attached for tty0.

To use the IBM 7855 modem as a dial-in or dial-out modem on the system, the modem must be set to accept AT commands. To do this, press the right and left arrow keys on the front of the modem at the same time. The front panel should then read:

```
<Exit
          Enter>
```

Press the right arrow to enter, and the down arrow until you see First Setup; then the right arrow. Press the up arrow until you see Asynchronous AT Press the left arrow until you exit out of the setup menu.

- 1. Change the tty settings using the Web-based System Manager fast path, wsm devices, or the SMIT fast path, smit chtty.
- 2. Select the tty for the modem, and make the appropriate changes. Example settings follow:

Note: Options that are not needed are not included in the following list.

TTY tty0 TTY type tty **TTY** interface rs232

Description Asynchronous Terminal

Status Available 00-00-s1-00 Location

Parent adapter sa0 **PORT** number s1 **Enable LOGIN** share **BAUD** rate 19200 **PARITY** none BITS per character 8 **Number of STOP BITS** 1 TIME before advancing to 0

next port setting

XON-XOFF handshaking no **TERMINAL** type dumb **INPUT** map file none **OUTPUT** map file none **CODESET** map file sbcs

STTY attributes for RUN TIME should be:

[hupcl, cread, brkint, icrnl, opost, tab3, onlcr, icanon, echo, echoe, echo k, echoctl,echoke,imaxbel,iexten]

(No ixon/ixoff needed)

STTY attributes for LOGIN should be:

[hupcl, cread, echoe, cs8]

(No ixon/ixoff needed)

3. Add the following two lines to /usr/lib/uucp/Systems file (be sure to start in the leftmost column):

ibm7855 Nvr IBMPROG 19200 slip Nvr SLIPPROG 19200

4. Add the following lines to /usr/lib/uucp/Devices file (be sure to start in the leftmost column):

```
IBMPROG tty0 - Any IBMProgrm
SLIPPROG tty0 - Any SlipProgrm
# Usual ACU entry
ACU tty0 - Any ibm
# ACU entry to quietly dial the modem.
ACUQ tty0 - Any ibmq
# IBM7855 ACU entry to negotiate for error correction and compression
ACUECLC tty0 - Any ibmeclc
# IBM7855 ACU entry to negotiate for error correction and compression
# using large (256 byte) data block size
ACUECLB tty0 - Any ibmeclclb
# IBM 7855 ACU entry to use ECL without compression
ACUECL tty0 - Any ibmecl
```

5. Add the following lines to /usr/lib/uucp/Dialers file (be sure to start in the leftmost column):

```
# IBM 7855 program. Use ibmeclcbc for 256 byte block size. However,
# Noisey lines may obtain better throughput by setting the block
# size to be 64 bytes; compressed files will show better throughput # by
turning off compression.
# The following 3 lines should be joined into one line:
IBMProgrm =,-, "" \dATQ0\r\c OK AT&F\r\c OK ATM1\r\c OK
AT&D3\r\c OK AT&C1\R2\Q2\M14\r\c OK AT&B8N1L0E0\r\c OK ATS0=1\r\c
OK ATQ1&W0&Y0\r\c ""
# The following 3 lines should be joined into one line:
SlipProgrm =,-, "" \dATQ0\r\c OK AT&F\r\c OK ATM1\r\c OK
AT&D3\r\c OK AT&C0\R2\Q2\r\c OK AT&B8N1L0E0\r\c OK
ATS0=1\r\c OK ATQ1&W\r\c ""
      ibm
ibmq
# IBM 7855 with ECLC and Compression
\# (Note two backslashes are needed before the N; cu won't send \N
# without preceding it with a \)
ibmeclc = , -,
              "" \dAT\\N3%C1Q0\r\c OK ATDT\T\d\r\c CONNECT
# IBM 7855 ECLC Compression with 256 byte block size
ibmeclclb
         =,-, "" \dAT\\N3%C1Q0\A3\r\c OK ATDT\T\d\r\c CONNECT
# IBM 7855 ECL without compression
ibmecl = , -,
              "" \dAT\\N3%C0Q0\r\c OK ATDT\T\d\r\c CONNECT
```

Programming the IBM 7855

- 1. Disable the port (temporarily) with this command: pdisable tty0.
- 2. On the command line enter cu -d ibm7855.

The command fails because no one is connected, but if the following is in the output:

```
Sendthem (ATQ1&W0&Y0^MNO CR)
expect: ("")
got it
```

it is successfully programmed. The LED on the front should read ASYN8N A 9600 a. You can then pshare tty0.

After the modem is programmed, you can set RTS line discipline to allow hardware flow control. To do that, the modem must be connected to another system (to bring carrier high) and the following command must be run:

```
stty add rts < /dev/tty0
```

Be sure to substitute tty0 with your tty number. This will be good until the next system startup, then it will need to be done again.

Alternatively, you could type in and compile the addrts.c program, listed below, and place a line in /etc/rc that will call the program on system startup to add RTS to the line. This program does not require carrier to be high.

```
C Program to add RTS discipline to tty port(s).
     Tips:
          To make the program more permanent, insert the file
          name of the compiled version (complete with path) at
          the end of your "/etc/rc" file and the changes will
          take effect again at next reboot. Usage is addrts
          /dev/tty##.
          NOTE: This program is supplied "as is" and is NOT
                 supported by IBM. It is intended as an aid to
                 administrators only.
     To create: vi addrts.c <enter>
To compile: cc -o addrts addrts.c
                           Program starts now
#include <stdio.h>
#include <fcntl.h>
#include <termios.h>
#include <sys/tty.h>
main (argc, argv)
        int argc;
        char *argv[];
{
        int fd;
        if ( (fd = open(argv[1], O_NDELAY|O_RDWR)) <0 ) {
                 printf("%s: could not open %s\n", argv[0],
                 argv[1]);
                 exit (22);
        ioctl(fd, TXADDCD, "rts"); /* adds rts to the tty in the
        argument */
        close(fd);
}
```

Setup Information for the MultiTech MULTIMODEM II

1. Use the **L5**, **L6**, and **L7** AT commands to view the following modern parameters.

```
B1 E1 M1 Q0 R0 V1 X0 &E1 &E4 &E6 &E8 &E10 &E13 &E15
$MB9600 $SB9600 $BA0 &W0
ATL6
S0 S2 S3 S4 S5 S6 S7 S8 S9 S10 S11 S13 S24 S25 S30
001 043 013 010 008 002 045 002 025 007 070 037 020 001 000
OK
ATL7
$A0 &A0 #A0 &B0 &BS1 &C1 $D0 &D2 #DB0 $EB0 %E1 #F0 $F1 &G0 #L0 $MI0 &M0
&PO #P2 &QO $RO &R1 &RF1 &SO $SPO &SFO #T1 $VO $V6 $VDO &XO YO
$MB9600 $SB9600 $BA0 &W0
```

Note: In the example above, the s9 register has been set to 25. The default is 6. The normal setting should be between 15 and 20.

- 2. Modems that are set for baud rates at or above 9600 should have RTS/CTS (request to send/clear to send) hardware handshaking added to the tty port. There are two options for accomplishing this:
 - a. Use the following procedure:
 - Add clocal to the stty Attributes for Runtime for your tty using smit chtty.
 - On the command line, enter: pdisable tty N (N is your tty number).
 - Issue the command stty add rts < /dev/tty N
 - Issue the command stty get < /dev/tty N. RTS should appear on line.
 - b. Use the following C program to add RTS to the port.

```
/* C Program to add RTS discipline to tty port(s).
 NOTE: This program is supplied "as is" and is NOT supported by
       IBM. It is intended as an aid to administrators only.
                        Program starts now
#include <stdio.h>
#include <fcntl.h>
#include <termios.h>
#include <sys/tty.h>
main()
   int fd;
     fd = open("/dev/tty0", O_NDELAY|O_RDWR);
     ioctl(fd, TXADDCD, "rts"); /* adds rts to tty0
                                                               */
     close(fd);
                                                               * /
                            END Program
```

Setup Information for a Practical Peripherals Modem

- 1. Change the tty settings using the Web-based System Manager fast path, wsm devices, or the SMIT fast path, smit chtty.
- 2. Select the tty for the modem and make the appropriate changes. (tty0 is the tty for this example and the baud rate may be changed depending on your modem. Generally, you want to keep the baud rate at the highest speed that the modem can allow.) Example settings follow:

Note: Options that are not necessary are not included in the following list.

TTY ttv0 TTY type tty **TTY** interface rs232

Description Asynchronous Terminal

Status Available Location 00-00-s1-00

Parent adapter sa0 **PORT** number s1 **Enable LOGIN** share **BAUD** rate 2400 **PARITY** none **BITER** per character 8 **Number of STOP BITS** 1 TIME before advancing to 0

next port setting

XON-XOFF handshaking no **TERMINAL** type dumb **INPUT** map file none **OUTPUT** map file none **CODESET** map file sbcs

STTY attributes for RUN TIME should be:

[hupcl, cread, brkint, icrnl, opost, tab3, onlcr, icanon, echo, echoe, echok, echoctl, echoke, imaxbel, iexten]

(No ixon/ixoff needed)

STTY attributes for LOGIN should be:

[hupcl, cread, echoe, cs8]

(No ixon/ixoff needed)

3. Add the following line to /usr/lib/uucp/Systems file:

```
pracper Nvr PRACPERPROG 2400
```

4. Add this to /usr/lib/uucp/Devices file:

```
# For programming the Practical Peripherals modem only:
PRACPERPROG tty1 - 2400 PracPerProgram2400
```

```
# regular ACU entry:
ACU tty1 - Any hayes
```

5. Add this to /usr/lib/uucp/Dialers file:

```
# This Entry is used to PROGRAM the modem ONLY:
# the next 3 lines should be made into one:
                      =,-, "" \d\dAT\r\c OK AT&F\r\c OK
PracPerProgram2400
ATM1\r\c OK AT&D3\r\c OK AT&C1\r\c OK ATLOE0\r\c OK
ATS0=1S9=20\r\c OK AT&W\r\c
```

```
"" \dAT\r\c OK ATDT\T\d\r\c CONNECT
haves
```

6. To program the modem, enter the command cu -d pracper. **cu** is used to program the modem. Because there is no connection to another computer, the command fails.

The modem will be programmed if you see sendthem AT&W and then OK got it in the output.

After the modem is programmed, you can set RTS line discipline to allow hardware flow control. To do that, the modem must be connected to another system (to bring carrier high) and the command stty add rts < /dev/tty0 must be run. (Be sure to substitute tty0 with your tty number). This works until the next system restart, when it must be done again. Alternately, type in and compile the addrts.c program, and place a line in /etc/rc that calls the program on system startup. This program does not require carrier to be high.

Setup Information for the Telebit T1600 Modem

1. Add this to /usr/lib/uucp/Systems file:

```
telebit Nvr TELEPROG 19200
```

2. Add this to /usr/lib/uucp/Devicesca file:

```
#
 Devices entry to use to program the T1600
TELEPROG tty0 - 19200 TelebitProgram
 19200, 9600 baud (fast) Telebit Trailblazer modem
#
        tty0 - 19200 tbfast \D
ACU
ACUTB
       tty0 - 19200 tbfast \D
       tty0 - 9600 tbfast \D
ACU96
  2400 baud Telebit Trailblazer modem
ACU
       tty0 - 2400 tb2400 \D
       tty0 - 2400 tb2400 \D
ACU24
#
#
 1200 baud Telebit Trailblazer modem
#
         tty0 - 1200 tb1200 \D
ACU
         tty0 - 1200 tb1200 \D
ACU12
```

3. Add this to /usr/lib/uucp/Dialers file:

```
# Entry to program the modem ONLY
# The next 3 lines should be made into one long line:
TelebitProgram =,-, "" \dAT&F\r\c OK
ats0=1s2=255s7=60s11=50s41=2s45=255s51=252s63=1s58=2s64=1\r\c OK
ATs69=2s105=0s111=30s255=0M0&C1Q2&D3&Q0&R3&S1&T5\r\c OK
ATEOX12&W\r\c OK
# Telebit T1600 Dialers entries
#
tbfast =,-, "" \dATs50=255s7=60\r\c OK\r ATDT\T\r\c
CONNECT-\d\c-CONNECT
tb2400 =,-, "" \dATs50=3\r\c OK\r ATDT\T\r\c CONNECT
               "" \dATs50=2\r\c OK\r ATDT\T\r\c CONNECT
tb1200 = , -,
```

4. To program the modem, enter the command cu -d telebit. The cu command will think it is connecting to a system, so it will fail. Watch the command's output to make sure that an OK is received.

Configuration Information for the Telebit T2000 Modem

1. Add this to /usr/lib/uucp/Systems file:

```
telebit Nvr TELEPROG 19200
```

Add this to /usr/lib/uucp/Devices file:

```
Devices entry to use to program the T2000
TELEPROG tty0 - 19200 TelebitProgram
# 19200, 9600 baud (fast) Telebit Trailblazer modem
        tty0 - 19200 tbfast \D
ACU
ACUTB tty0 - 19200 tbfast \D
       tty0 - 9600 tbfast \D
# 2400 baud Telebit Trailblazer modem
         tty0 - 2400 tb2400 \D
ACU
ACU24
        tty0 - 2400 tb2400 \D
# 1200 baud Telebit Trailblazer modem
        tty0 - 1200 tb1200 \D
ACII
ACU12
        tty0 - 1200 tb1200 \D
# For sites with MNP modems
         tty0 - 1200 tb12mnp \D
ACUM12
ACUM24
         tty0 - 2400 tb24mnp \D
```

3. Add this to /usr/lib/uucp/Dialers file:

```
# Entry to program the modem ONLY
  # The next 4 lines will be made into one:
TelebitProgram =,-, "" \dAT&F\r\c OK
 \mathtt{ats2} = 255\$7 = 60\$11 = 50\$41 = 2\$45 = 255\$51 = 252\$52 = 2\$54 = 3\$58 = 2\$64 = 1 \\ \texttt{T} \\ \texttt{C} \\ \texttt{OK} \\ \texttt{C} \\
ATs69=1s92=1s96=0s105=0s110=1s111=30s130=3s131=1F1M0Q6TV1W0X3Y0&P0&T5\r\c
 OK ATEO&W\r\c OK 00
  # Telebit T2000 dialers Entries:
 tbfast =,-, "" \dATs50=255s7=60\r\c OK\r ATDT\T\r\c
 CONNECT-\d\c-CONNECT
tb2400 =,-, "" \dATs50=3\r\c OK\r ATDT\T\r\c CONNECT tb1200 =,-, "" \dATs50=2\r\c OK\r ATDT\T\r\c CONNECT
                                                                                                               "" \dAT\r\c OK ATS0=0S95=2S50=3S41=0\r\c OK ATDT\T\r\c
 tb24mnp =, -,
 CONNECT
tb12mnp =, -,
                                                                                                                  "" \dAT\r\c OK ATS0=0S95=2S50=2S41=0\r\c OK ATDT\T\r\c
 CONNECT
```

4. To program the modem, enter the command cu -d telebit. The cu command will think it is connecting to a system, so it will fail. Watch the command's output to make sure that an OK is expected it is received.

Setup Information for the Telebit T3000 Modem

1. Add this to /usr/lib/uucp/Systems file:

```
telebit Nvr TELEPROG 19200
```

2. Add this to /usr/lib/uucp/Devices file:

```
# Devices entry to use to program the T3000
TELEPROG tty1 - 19200 TelebitProgram
# 19200, 9600 baud (fast) Telebit Trailblazer modem
#
        ttv1 - 19200 tbfast \D
ACU
        tty1 - 19200 tbfast \D
ACUTB
ACU96
        tty1 - 9600 tbfast \D
# 2400 baud Telebit Trailblazer modem
         tty1 - 2400 tb2400 \D
ACU
ACU24
        tty1 - 2400 tb2400 \D
# 1200 baud Telebit Trailblazer modem
ACU
        tty1 - 1200 tb1200 \D
        tty1 - 1200 tb1200 \D
ACU12
# For sites with MNP modems
#
ACUM12
         tty1 - 1200 tb12mnp \D
ACUM24
         tty1 - 2400 tb24mnp \D
```

3. Add this to /usr/lib/uucp/Dialers file:

```
# Entry to program the modem ONLY
# The following 4 lines should be made into one long line:
TelebitProgram =,-, "" \dAT&F\r\c OK
\verb|ats0=1s2=255s7=60s11=50s41=2s45=255s51=252s63=1s58=2s64=1\r\c OK|
ATs69 = 2s105 = 0s111 = 30s255 = 0M0\&C1Q2\&D3\&Q0\&R3\&S1\&T5\r\c OK\ ATEOX12\&W\r\c
OK 00
# Telebit T2000 dialers Entries: The tbfast won't work for the T3000,
# but the rest should.
               "" \dATs50=255s7=60\r\c OK\r ATDT\T\r\c
tbfast =,-,
CONNECT-\d\c-CONNECT
tb2400 =,-, "" \dATs50=3\r\c OK\r ATDT\T\r\c CONNECT tb1200 =,-, "" \dATs50=2\r\c OK\r ATDT\T\r\c CONNECT
tb1200 =,-, "" \dATs50=2\r\c OK\r ATDT\T\r\c CONNECT
tb24mnp =,-, "" \dAT\r\c OK ATS0=0S95=2S50=3S41=0\r\c OK ATDT\T\r\c
tb12mnp = , -,
                 "" \dAT\r\c OK ATS0=0S95=2S50=2S41=0\r\c OK ATDT\T\r\c
CONNECT
```

4. To program the modem, enter the command cu -d telebit. The command will fail because no connection is established. Examine the debug output of the command to verify that ATEOX12&W was sent and that OK was received. This indicates the modem is programmed.

Setup Information for UDS Modems

- 1. Change the tty settings using the Web-based System Manager fast path, wsm devices, or the SMIT fast path, **smit chtty**.
- 2. Select the tty for the modem, and make the appropriate changes. Example changes follow (tty0 is the tty for this example):

Note: Options that are not necessary are not included.

TTY tty0 TTY type tty **TTY interface** rs232

Description Asynchronous Terminal

Status Available Location 00-00-s1-00

Parent adapter sa0 **PORT** number s1 **Enable LOGIN** share **BAUD** rate 19200 **PARITY** none BITS per character 8 **Number of STOP BITS** 1 TIME before advancing to 0 next port setting

XON-XOFF handshaking no **TERMINAL** type dumb **INPUT** map file none **OUTPUT** map file none CODESET map file sbcs

STTY attributes for RUN TIME should be:

[hupcl, cread, brkint, icrrnl, opost, tab3, onlcr, icanon, echo, echoe, echok, echoctl, echoke, imaxbel, iexten]

(No ixon/ixoff needed)

STTY attributes for LOGIN should be:

[hupcl, cread, echoe, cs8]

(No ixon/ixoff needed)

3. Add the following line to /usr/lib/uucp/Systems file:

uds Nvr UDSPROG Any

4. Add this to /usr/lib/uucp/Devices file:

```
-----begin text to insert-----
# For programming the UDS modem only:
UDSPROG ttv0 - 9600 udsmodemPROGRAM
# regular ACU entry:
ACU tty0 - Any uds
 -----end text to insert-----
** Add this to /usr/lib/uucp/Dialers file: **
 -----begin text to insert-----
# the next 2 lines should be made into one:
udsmodemPROGRAM =,-, "" \dAT&FQ2\r\c OK ATEOYO&C1&D2&S1%B5%E0*LC\r\c
OK AT&W\r\c OK
uds =,-, "" \dAT\r\c OK\r ATDT\T\d\r\c CONNECT
```

5. To program the modem, enter the command cu -d uds. The cu command is used to program the modem. Because there is no connection to another computer, the command fails. Watch the output of the command to see this.

Setup Information for the Telebit T3000 90202–01 WordBlazer Modem

When using the T3000 modem on a 128-port concentrator (operating at AIX level 3.2.3 Extended), change the tty settings using the Web-based System Manager fast path, wsm devices, or the SMIT fast path, smit chtty and change the following settings, both in the user interface and on the modem. Do not use the 38400 baud rate shown below unless your serial hardware can support this speed. Make changes to baud speeds accordingly.

Example settings follow (only entries that must be changed are shown):

TERMINAL type: dumb **Enable LOGIN** share

BAUD rate 38400 (may be 19200, or slower)

PARITY none 8 BITS per character

XON-XOFF handshaking no (at 9600 baud and over user RTS/CTS)

tty settings for modem on 128-port adapter:

Force Carrier: disable Perform Cooked disable

Processing in Adapter:

User Alternate RJ-45 disable

Pinouts:

Suggested settings for modem registers and commands:

Notes:

- 1. Use AT&V with this modem to view current values.
- 2. Change register 51 to reflect your hardware DTE speed.
- 3. Register 92 is very important if you will be connecting to modems operating at lower speeds. The suggested setting is 1.

AT&C1 When the local modem detects a carrier signal from the

remote modem, the CD signal goes ON after the

CONNECT result code is sent to the DTE.

AT&D3 Resets the modem and enters command mode when the

DTR signal goes from ON to OFF.

AT&S1 DSR is ON after the modem detects an answer tone, and it

stays ON throughout the connection.

ATS0=1 Sets the modem for Autoanswer mode. The modem will

answer the call after one ring.

ATS9=15 Sets the carrier redetect time to 15.

ATS50=0 Sets the modulation speed. A value of 0 indicates automatic

speed determination.

ATS51=253 Sets the DTE interface speed (the speed at which the

> computer talks to the modem, not the speed at which the modem talks to other modems). 253 is used for 38400 bps.

ATS59=15 Enable CONNECT suffixes. ATS92=1 Answering Sequence Selection.

ATQ2 Causes the modem to report result codes when originating

a call, but suppresses result codes when answering a call.

ATX12 Enables verbose result codes.

Troubleshooting Modem Problems

This section attempts to identify common problems when using a modem with your computer.

Symptom	Cause	Solution
The modem (or other device attached to the serial port) causes the system to gradually slow down and eventually hang. Turning off the device usually lets the system function normally again.	An intelligent modem has CD always ON. The system senses this and sends a login herald, which the modem tries to interpret as a command. The modem fails to recognize the login herald as a valid command, and echoes back to the tty port on the system. This cycle repeats continuously.	Set the tty port to delay on the system so no login herald will be sent. With this setting, only a valid carriage return character from the host logging in will cause a login herald to be sent. You can also change the modem's AT set profile to set CD to ON only when a valid carrier is sensed on the telephone line.

Software Services Modem Questionnaire

Before calling for assistance with modem problems, please collect the following information:

- · System Xtra?
- Level of the operating system. How long have you been at this level of the operating system?
- · Has the modem ever worked before?
- What type of modem are you using? What type of modem is on the other end of the telephone connection?
- To what adapter type (64–port. 128–port, S1,...) is the modem attached?
- To which port number is the modem attached?
- To which tty number is the modem attached?

- What type of cabling are you using?
- What is the login setting (share, delay, enable)?
- Can the modem connect to other modems?
- · Can other modems connect to your modem?
- What are the following values in Web-based System Manager, SMIT, modem, or port?
 - XON/XOFF?
 - RTS/CTS?
 - BPS rate?
- Include the following in your problem description:
 - Does the port lock intermittently?
 - Can you dial out? Can others dial in?
 - Any other specific and descriptive error conditions.
- Are there errors on the console? What are they?
- Are there errors in the error report? (errpt or errpt -a)
- What command are you using to dial out?
- What software is involved on the system?

AT Command Summary

The following is a summary of the Hayes Smartmodern command set. These commands comprise the AT command set used by many popular modems. This information comes from the Hayes Smartmodem 2400 Quick Reference Card, published by Hayes Microcomputer Products, Inc.

AT	Command prefix – precedes command line.
<cr></cr>	Carriage return (newline) character – terminated the command line.
A	Go off-hook, remain in command mode.
A /	Repeat previous command line. This command is not preceded with AT or followed by <cr></cr> /.
B0	Select CCITT V.22 standard for 1200 bps communications.
B1	Select Bell 212A standard for 1200 bps communications.
D	Enter originate mode, dial the number that follows, and attempt to go online. D is usually followed by T for tone, P for pulse may also be used.
DS = <i>n</i>	Dial the number stored in location n.
E0	Disable character echo in the command state.
E1	Enable character echo in the command state.
H0	Go on-hook (hang up the phone).
H1	Operate switch-hook and auxiliary relay.
10	Return product identification code.
I1	Perform checksum on firmware ROM; return checksum.
12	Perform checksum on firmware ROM; returns OK or ERROR as the result.
L0	Speaker off.

L1 Low speaker volume. L2 Medium speaker volume. L3 High speaker volume.

MO Speaker off.

М1 Speaker on until carrier detected.

M2 Speaker always on.

M3 Speaker on until carrier detected, except during dialing.

00 Enter online state.

01 Enter online state and initiate equalizer retrain.

Q0 Modem returns result codes.

Q1 Modem does not return result codes.

Sr Set pointer to register r. Sr=nSet register r to value n.

V0 Display result codes in numeric form.

V1 Display result codes in verbose form (as words). **X0** Enable features represented by result codes 0-4. **X1** Enable features represented by result codes 0-5, 10. **X2** Enable features represented by result codes 0-6, 10. **X3** Enable features represented by result codes 0-5, 7, 10. **X4** Enable features represented by result codes 0–7, 10.

Y0 Disable long space disconnect. **Y1** Enable long space disconnect.

Ζ Reset modem

&C0 Assume data carrier always present.

&C1 Track presence of data carrier.

&D0 Ignore DTR signal.

&D1 Assume command state when an on-to-off transition of

DTR occurs.

&D2 Hang up and assume command state when an on-to-off

transition of DTR occurs.

&D3 Reset when an on-to-off transition of DTR occurs. &F Recall the factory settings as the active configuration.

&G0 No guard tone. &G1 500 Hz guard tone. &G2 1800 Hz guard tone.

&J0 RJ-11/RJ41/RJ45S telco jack.

&J1 RJ-11/RJ-13 telco jack.

&P0 Pulse dial with make/break ratio 39/61. &P1 Pulse dial with make/break ratio 33/67.

&Q0 Operate in asynchronous mode. &Qn Operate in synchronous mode n. &R0 Track CTS according to RTS.

&R1 Ignore RTS; always assume presence of CTS.

&S0	Assume presence of DSR signal.
&S1	Track presence of DSR signal.
&T0	Terminate test in progress.
&T1	Initiate local analog loopback.
&T3	Initiate digital loopback.
&T4	Grant request from remote modem for remote data link (RDL).
&T5	Deny request from remote modem for RDL.
&T6	Initiate remote digital loopback.
&T7	Initiate remote digital loopback with self-test.
&T8	Initiate local analog loopback with self-test.
&V	View active configuration, user profiles, and stored numbers.
&W n	Save storable parameters of active configuration as user profile <i>n</i> .
&X0	Modem provides transmit clock signal.
&X1	Data terminal provides transmit clock signal.
&X2	Receive carrier provides transmit clock signal.
&Y n	Recall user profile <i>n</i> .
&Z n= <i>x</i>	Store phone number <i>x</i> in location <i>n</i> .

S–Register Summary

Register	Range	Description
S0	0–255	Select number of rings before answer.
S1	0–255	Ring count (incremented with each ring).
S2	0–127	Define escape sequence character (ASCII).
S3	0–127	Define carriage return character (ASCII).
S4	0–127	Define line feed character (ASCII).
S5	0–32, 127	Define backspace character (ASCII).
S6	2–255	Select wait-time in seconds before blind dialing.
S7	1–55	Select wait-time in seconds for carrier/dial tone.
S8	0–255	Select duration in seconds of comma.
S9	1–255	Carrier detect response time in .1 second increments (10 = 1 second).
S10	1–255	Delay between carrier loss and hangup in .1 second increments.

S11	50–255	Duration/spacing of tones in milliseconds.
S12	50–255	Escape sequence guard time in .02 second intervals.
S13	<u> </u>	Reserved.
S14	_	Reserved.
S15	_	Reserved.
S16	_	Reserved – functions for this register are controlled by the &T commands).
S17	_	Reserved.
S18	0–255	Test timer duration in seconds.
S19	_	Reserved.
S20	_	Reserved.
S21	_	Reserved.
S22	_	Reserved.
S23	_	Reserved.
S24	_	Reserved.
S25	0–255	Select DTR change detect time in .01 second intervals.
S26	0–255	RTS to CTS delay in .01 second intervals.
S27	_	Reserved.

Result Codes Summary

Number	Word	Description
0	OK	Command executed.
1	CONNECT	Connection established at 0–300 bps.
2	RING	Ring signal detected.
3	NO CARRIER	Carrier signal lost or not detected.
4	ERROR	Invalid command, checksum, error in command line, or command line too long.
5	CONNECT 1200	Connection established at 1200 bps.
6	NO DIALTONE	No dial tone detected.
7	BUSY	Busy signal detected.
8	NO ANSWER	No response when dialing a system.
9	CONNECT 2400	Connection established at 2400 bps.

Dial modifiers

The following lists and describes dial modifiers.

0–9 # * A–D	Digits and characters for dialing.	
Р	Pulse dial.	
Т	Tone dial.	
,	Delay processing of next character.	
!	Hookflash.	
@	Wait for silence.	
W	Wait for dial tone.	
;	Return to command state after dialing.	
R	Reverse mode.	
S = <i>n</i>	Dial number stored at location n.	

ATE Overview

The Asynchronous Terminal Emulation (ATE) program, an optional software product, enables a system to emulate a terminal on a remote system. Using ATE, you can log in to most systems that support asynchronous terminals, including any system that supports RS–232C or RS–422A connections. You can set ATE parameter values so the remote system recognizes your terminal as either an attached workstation terminal or a DEC VT100 terminal.

Setting Up ATE Overview

Before you run ATE, you must install the software and set up the ports and connections. ATE uses both direct (cabled) connections and modem connections. Local RS–232C connections allow a maximum distance of 15 meters (50 feet) between machines, and RS–422A connections allow up to 1200 meters (4000 feet) between machines.

Before you use ATE to call a remote system, be sure that the remote system's tty device is ready to accept a call. If another user on the remote system uses ATE to call your terminal, be sure your tty device is ready to accept the call.

See "Setting Up ATE", on page 5-40 for more information about installing and setting up ATE.

Note: You must be a member of a UNIX-to-UNIX Copy Program (UUCP) group to use ATE. A user with root authority can use Web-based System Manager or the System Management Interface Tool (SMIT) to set up a UUCP group.

Customizing ATE

The first time you run ATE, the program creates an **ate.def** default file in the current directory. The **ate.def** file contains parameters the ATE program uses for:

- Data transmission characteristics
- Local system features
- Dialing directory file
- · Control keys

To change the defaults, edit the ate.def file.

If the you need to run ATE with different settings, you can maintain **ate.def** files in different directories. You can then run ATE from the appropriate directory depending on the settings needed for specific sessions. However, running ATE from many directories requires multiple copies of the **ate.def** file, which uses system storage.

See "How to Edit the ATE Default File" in AIX 4.3 System User's Guide: Communications and Networks for details about editing the ate.def file.

You can temporarily change settings without modifying the default file. To do this, use the **alter** and **modify** subcommands. Settings you change with the **alter** or **modify** subcommand remain in effect until you exit the program with the **quit** subcommand. When you exit ATE, the settings return to the defaults set in the **ate.def** file.

When installed, ATE uses the /usr/lib/dir system—wide dialing directory file. You can temporarily change settings in the dialing directory file for a specific modem connection. Settings changed in this way revert to the default when the connection ends, rather than when you exit ATE. A user with root authority can modify the /usr/lib/dir file to include numbers for modems used by everyone on the system. Individual users can also create their own dialing directory files and modify their copies of the ate.def file to make ATE use those directories.

"How to Set up an ATE Dialing Directory" in AIX 4.3 System User's Guide: Communications and Networks explains how to set up ATE to use a customized dialing directory.

You can edit the dialing directory file to include frequently used phone numbers Additionally, you can change the baud rate, data character length, stop bits, parity, echoing, and line-feeds for a phone number if these characteristics differ from the defaults. If a number is not in the directory file, you can complete the connection by using the connect subcommand.

Note: A dialing directory file can contain up to 20 lines (one entry per line). ATE ignores subsequent lines.

Changing ATE Characteristics

The following table identifies the ATE characteristics that the user can change and the appropriate methods for changing each characteristic.

Note: All ATE characteristics can be changed in the ate.def file.

Changing ATE Characteristics		
Characteristic Change with		
Control keys	ate.def file	
Data character length	alter subcommand or dialing directory entry	
Dialing directory file name	directory subcommand	
Echoing (on or off)	modify subcommand or dialing directory entry	
File name for incoming data (capture file)	modify subcommand	
Final dial suffix for the modem	alter subcommand	
Initial dial prefix for the modem	alter subcommand	
Line feeds	modify subcommand or dialing directory entry	
Number of redialing attempts	alter subcommand	
Number of stop bits	alter subcommand or dialing directory entry	
Parity (even or odd)	alter subcommand or dialing directory entry	
Port name (device)	alter subcommand	
Rate (bits per second)	alter subcommand or dialing directory entry	
Telephone number	dialing directory entry	
Transfer protocol (pacing or xmodem)	alter subcommand	
Type of pacing (character or interval)	alter subcommand	
VT100 emulation (on or off)	modify subcommand	
Wait time between redialing attempts	alter subcommand	
Write (capture) incoming data to a file	modify subcommand	
Xon/Xoff protocol (on or off)	modify subcommand	

Setting Up ATE

This article provides information on setting up Asynchronous Terminal Emulation (ATE).

Prerequisites

- The ATE program must be installed on the system. ATE is an optional program product.
- The user must have root user authority to set up the port for the communications device.

Procedure

To prepare ATE to run on the system:

- 1. Install an asynchronous adapter card in an appropriate slot in the system unit, unless the system has a built-in serial port.
- 2. Plug the RS-232C or RS-422A cable into the adapter card or the built-in serial port.
- 3. Add a tty device for the communications port. To do this, use the Web-based System Manager fast path, wsm devices, or enter:

```
smit mktty
```

- Select Add a TTY.
- Select the tty type.
- 6. Select a parent adapter
- 7. Select a port.
- 8. Set the Enable LOGIN field to disable.
- Set Terminal Type to HFT or dumb.
- 10. Make the necessary adjustments for the environment. The most common changes are line speed, parity settings, number of bits per character, and whether the line is to be driven as a remote or local line. Use BPC 8 and no parity if National Language Support (NLS) is required.
- 11. Set up the port for the device.
 - To set up a port to call out with ATE, use the pdisable command. For example, to set up port tty1, enter:

```
pdisable tty1
```

 To set up a port so that others can call in, use the penable command. For example, to let other systems call in to the tty2 port, enter:

```
penable tty2
```

12. Ensure the device has previously been defined to the remote system. Once the device is defined, the ATE program must be customized to reflect the device settings on the remote system. Customize the default settings with the alter and modify subcommands or by editing the ate.def default file. To change the default settings for a telephone connection, use a dialing directory file entry.

TTY Troubleshooting

This section discusses troubleshooting the tty subsystem:

- Respawning Too Rapidly Errors, on page 5-41
- Error Log Information and tty Log Identifiers, on page 5-42

Respawning Too Rapidly Errors

The system records the number of getty processes spawned for a particular tty in a short time period. If the number of getty processes spawned in this time frame exceeds five, then the Respawning Too Rapidly error is displayed on the console and the port is disabled by the system.

The tty stays disabled for about 19 minutes or until the system administrator enables the port again. At the end of the 19 minutes, the system will automatically enable the port, resulting in the spawning of a new getty process.

Possible Causes

- Incorrect modem configuration
- A port is defined and enabled but no cable or device is attached to it
- Bad cabling or loose connection
- Noise on communication line
- Corruption of, or tampering with, /etc/environment or /etc/inittab files
- tty configuration is corrupted
- · Hardware is defective

Procedures for Recovery

Correct modem configuration:

Ensure that the modem's carrier detect is *not* forced high.

Note: The following applies to Hayes–compatible modems.

- a. Connect to the modem and examine the active profile.
- b. The modem's carrier detect should be set to &C1 rather than &C0 (forced high). Use the following AT modem commands to set and change the carrier attribute:

```
AT&C1
AT&W
```

Notes:

- a. Refer to "Sending AT Commands with the cu Command", on page 5-16
- b. Refer to your modem documentation for further information.
- Disable the tty, remove the tty definition, or attach a device to the port:
 - To disable the tty definition use the chdev command as follows:

```
chdev -l ttyName -a Login=disable
```

After running this command, the tty will not become enabled after a system restart.

- To remove the tty definition:
 - a. Disable the tty port, use the **pdisable** command, enter:

```
pdisable ttyName
```

- b. Remove the tty definition from the system. Refer to "Managing TTY Devices", on page 5-4 for further information.
- Check for bad cables or loose connections:
 - a. Check cabling. Tighten loose connections and replace damaged or inappropriate connectors.
 - b. Verify that the suspected cabling or that the cable meets the same standard. Replace damaged or inappropriate cables.
- Eliminate noise on communication line:
 - a. Verify that cabling is correct length and impedance.
 - b. Ensure that toroid rings are in place where needed on longer cables.
 - c. Check routing of cables; they should not be close to fluorescent lights or motors.
- Check for corruption of, or tampering with, the /etc/environment or the /etc/inittab files:
 - a. If possible, compare these files against known good copies.
 - b. Copy the files as a backup and make changes as needed.
 - c. In the /etc/environment file, remove any lines that are not:
 - blank lines
 - comment lines
 - variable=value
 - d. In the /etc/inittab file, examine the tty devices lines.If the tty is set to off, it is likely that the tty port is not being used. If it is not being used, remove the tty definition or attach a device to the port.
- Remove corrupted tty configuration:
 - a. Remove the tty definition. Use the Web-based System Manager Devices application or see "Managing TTY Devices", on page 5-4 for further information.
 - b. If you want a hard copy record of the tty definition before removing it, press the Image key (F8 or Esc+8). This will capture the current screen image and copy it to the smit.log file in your \$HOME directory.
 - c. Read the tty definition. See the instructions for Adding a TTY under "Managing TTY Devices", on page 5-4.
- Locate defective hardware:
 - a. Run diagnostics using the diag command.
 - b. If any hardware problems are detected, follow local problem solving procedures.

Error Log Information and TTY Log Identifiers

The following sections discuss important error logging files and commands and common error report messages relating to ttys.

Important Error Logging Files and Commands

Command: errclear

This command deletes entries from the error log. The entire log can be erased with errclear 0 or entries with specified error ID numbers, classes, or types can be removed.

Command: errpt

This command generates an error report from entries in the system error log. The most used format for this command is errpt -a | pg, which generates a detailed report starting with the most current errors.

File: /var/adm/ras/errlog

This file stores instances of errors and failures encountered by system. The errlog file tends to become quite lengthy. If not cleared on a regular basis, it can occupy quite a bit of space on your hard disk. Use the errclear command mentioned above to clean out this file.

File: /usr/include/sys/errids.h

The **errids.h** header file correlates error IDs with error labels.

Common Error Report Messages

Message	Description	Comments
Core Dump	Software program abnormally terminated	This error is logged when a software program abnormally terminates and causes a core dump. Users may not be exiting applications correctly, the system may have been shut down while users were working in application, or the user's terminal may have locked up and the application stopped.
Errlog On	Errdaemon turned on	This error is logged by the error daemon when the error logging is started. The system automatically turns off error logging during shutdown.
Lion Box Died	Lost communication with 64–port concentrator	This error is logged by the 64–port concentrator driver if communications with the concentrator are lost. If you receive this error, check the date and time stamp to see if user may have caused this message to occur. A series of these errors can indicate a problem with the 64–port adapter or its associated hardware.
Lion Buffero	Buffer overrun: 64-port concentrator	This error occurs when the hardware buffer in a 64–port concentrator is overrun. If device and cabling allow, try adding request to send (RTS) handshaking to the port and device. Also try lowering the baud rate.

Lion Chunknumc	Bad chunk count: 64-port controller	This error occurs when the value for the number of characters in a chunk does not match the actual values in the buffer. This error may indicate a problem with the hardware; try running diagnostics on devices.
Lion Hrdwre	Cannot access memory on 64–port controller	This error is logged by the 64–port concentrator driver if it is unable to access memory on the 64–port controller.
Lion Mem ADAP	Cannot allocate memory: ADAP structure	This error is logged by the 64–port concentrator driver if the malloc routine for the adap structure fails.
Lion Mem List	Cannot allocate memory: TTYP_T List	This error is logged by the 64–port concentrator driver if the malloc routine for the <i>ttyp_t</i> list structure fails
Lion Pin ADAP	Cannot pin memory: ADAP structure	This error is logged by the 64–port concentrator driver if the pin routine for the adap structure fails.
SRC	Software program error	This error is logged by the System Resource Controller (SRC) daemon in the event of some abnormal condition. Abnormal conditions are divided in three areas: failing subsystems, communication failures, and other failures.
Lion Unkchunk	Unknown error code from the 64-port concentrator	Error Code: Number of characters in the chunk received.
TTY Badinput	Bad ttyinput return	Error Code: System error code (see sys/errno.h). This is logged by the tty driver if ttyinput routine returns an error.
TTY Overrun	Receiver overrun on input	The sending device is ignoring flow control and overloading the hardware buffer on the adapter. This occurs prior to the driver accessing the hardware FIFO. Try adding rts handshaking to the port and to device (if possible).

TTY TTYHOG	TTYHOG overrun	The sending device is ignoring flow control. This error occurs after the hardware first in/first out (FIFO) has been accessed and has written to the software buffer. In other words, the tty hog error is logged because the buffer in which incoming characters are placed is overflowing. When the buffer is about three—fourths full, the device driver notifies the hardware (in this case, tty) to send the XOFF char to the sender to stop sending data. If the sender continues to send data and the buffer is not emptied, the device driver flushes the buffer and logs the error message in the error log.
TTY Parerr	Parity/Framing error on input	This error indicates parity errors on incoming data to asynchronous ports on a character—by—character basis.
TTY Prog PTR	Software error: T_HPTR field invalid	This error is logged by the tty driver if <i>t_hptr</i> pointer is null.

Chapter 6. Micro Channel, ISA, and PCI Adapters

This chapter provides installation and configuration information for Micro Channel, ISA, and PCI adapters. Topics discussed are support and configuration for the following:

- Micro Channel adapters (Multiport/2, on page 6-2 and Portmaster, on page 6-2)
- ISA adapters (Multiport Model 2, on page 6-4)
- ISA/PCI Wide Area Network (WAN) Adapters (Multiport Model 2 (ISA), on page 6-4, 2—Port Multiprotocol (PCI), on page 6-8, and ARTIC960HX(PCI)).

Note: Not all adapters are described in this chapter. Please refer also to the documentation delivered with the system unit or with the adapters ordered separately.

Micro Channel Wide Area Network (WAN) Adapters

This section explains the configuration requirements for the following Micro Channel adapters:

- Multiport/2, on page 6-2
- Portmaster, on page 6-2
- Configuring Multiport/2 and Portmaster Adapters, on page 6-2

This section also includes information on the following:

- Supported Multiport/2 Adapters, on page 6-2
- Supported Portmaster Adapters, on page 6-2
- Device Driver Support, on page 6-2

Supported Multiport/2 Adapters

Realtime Interface Co-Processor Multiport/2 adapters:

- 4-Port EIA 232D Multiport/2 adapter
- 4-Port EIA 232D/4-Port EIA 422A Multiport/2 adapter
- 8-Port EIA 232D Multiport/2 adapter
- 8-Port EIA 422A Multiport/2 adapter
- 6-Port Synchronous EIA 232D Multiport/2 adapter

Supported Portmaster Adapters

In Web-based System Manager, SMIT, and in the following information, these adapters are referred to generically as *Portmaster Adapter/A*:

- 4-Port Multiprotocol communications controller
- Realtime Interface Co

 Processor Portmaster adapters:
 - 8-Port EIA 232D Portmaster Adapter/A
 - 8-Port EIA 422A Portmaster Adapter/A
 - 6-Port V.35 Portmaster Adapter/A
 - 6-Port X.21 Portmaster Adapter/A

Device Driver Support

The 4–Port Multiprotocol communications controller device driver is supplied with the operating system.

To use a Realtime Interface Co–Processor Multiport/2 adapter or a Realtime Interface Co–Processor Portmaster adapter, you will need additional software to provide device driver configuration support. One of the following is required:

- Realtime Interface Co–Processor licensed program OR
- A device driver from a business partner or a customer–written device driver

Configuring Multiport/2 and Portmaster Adapters

The following procedures explain how to configure device drivers and ports for a Multiport/2 or Portmaster communications controller (adapter).

Prerequisites

Install the communications controller (adapter).

Configuring Multiport/2 and Portmaster Adapters Tasks	
Web-based System Manager: wsm network fast path (wsm network application)	
-OR-	
Task	SMIT Fast Path
Add a Device Driver	smit commodev ^{1,2}
Add Ports on the Adapter	smit commodev, select adapter, then Manage Ports, then Add a Multiprotocol Port
Reconfigure Ports on the Adapter	smit commodev, select adapter, then Manage Ports, then Change/Show Characteristics of a Port
Remove a Port on the Adapter	smit commodev, select adapter, then Manage Ports, then Remove a Port
Make a Defined Port Available	smit commodev, select adapter, then Manage Ports, then Configure a Defined Port
Transfer Adapter's Device Driver from Defined to Available	smit commodev, select adapter, then Manage Ports, then Configure a Defined Device Driver
Remove Adapter's Device Driver	smit commodev, select adapter, then Manage Ports, then Remove a Device Driver

Note:

- 1. This menu varies depending on the software you have installed. The device drivers for the 4–Port Multiprotocol communications controller are included with the basic operating system. To use another type of Multiport/2 adapter or Portmaster Adapter/2, you must have installed the Realtime Interface Co–Processor licensed program, or a device driver from a business partner or a customer–written device driver.
- 2. If you are using a device driver other than the 4–Port Multiprotocol communications controller device driver, consult the documentation for that device driver.

ISA/PCI Wide Area Network (WAN) Adapters

This section explains the installation and configuration requirements for the Multiport Model 2 adapter (ISA), the 2–Port Multiprotocol adapter (PCI), and the ARTIC960HX PCI adapter.

Multiport Model 2 Overview

The Multiport Model 2 Adapter (MM2) device driver is a component of the communication I/O subsystem. This device driver provides support for the Multiport Model 2 at a maximum speed of 64K bps. The modems used must provide the clocking, since only external clocking is supported.

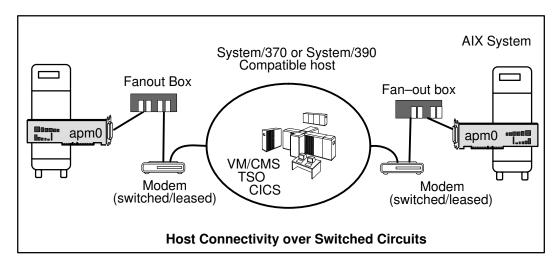
The following options provide access to the Multiport Model 2 device driver:

- Systems Network Architecture (SNA)
- The generic data link control (GDLC) Programming Interface
- User-written applications compatible with the MPQP-API (Multiprotocol Quad Port-Application Programming Interface).

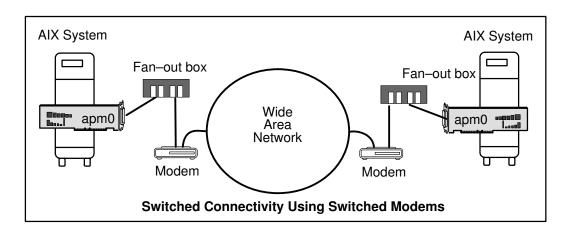
These options require use of the **mpq***i* special file, which allows access to the Multiport Model 2 adapter through the Multiport Model 2 device driver. The **mpq***i* special file resides in the /**dev** directory.

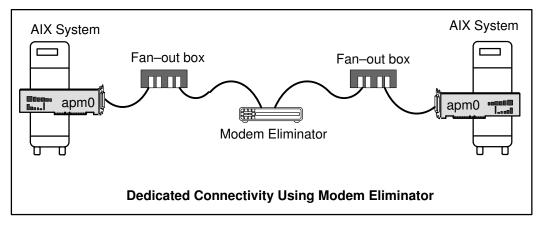
Note: The *i* in **mpq***i* specifies the instance of the device driver; for example, mpq0.

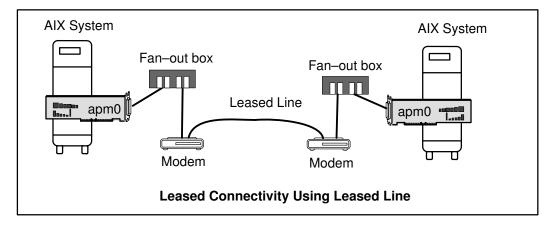
The Multiport Model 2 device driver allows connectivity to remote host systems using the Multiport Model 2 adapter, either directly over a leased line or over switched circuits, as shown in the **Host Connectivity** figure. The device driver can provide a gateway between work group environments and remote data processing facilities.



Direct connectivity to an AIX system using the Multiport Model 2 adapter is also possible over a wide area network using switched modems, through a modem eliminator, or over leased lines (see the figures).







Configuring the Multiport Model 2 Adapter

The following procedures explain how to configure a Multiport Model 2 adapter.

Configuring the Multiport Model 2 Adapter Tasks		
Web-based System Manager: wsm network fast path (wsm network application)		
-OR-		
Task	SMIT Fast Path	Command or File
Add a Multiport Model 2 Adapter ¹	smit makmm2	mkdev ²
List All Multiport Model 2 Adapters	smit Ismm2	Isdev ²
Change/Show a Multiport Model 2 Adapter	smit chgmm2	chdev ²
Configure a Defined Multiport Model 2 Adapter	smit cfgmm2	mkdev ²
Manage Device Drivers for Multiport Model 2 Adapters	smit mm2isa_dd	
Generate an Error Report	smit errpt	
Trace a Multiport Model 2 Adapter	smit mm2trace	
Remove a Multiport Model 2 Adapter	smit rmvmm2	rmdev ²
Manage DLC Services	smit mpaserv	
Add User-Written Applications	smit mm2apps	

Note:

- 1. For information on adding ports, see "Configuring Multiport/2 and Portmaster Adapters", on page 6-2.
- 2. Details about command line options are provided in the command descriptions for **mkdev**, **Isdev**, **chdev**, or **rmdev** in *AIX Commands Reference*.

Multiport Model 2 Adapter Object Information and Attributes

The following sections provide information on Multiport Model 2 adapter object classes and attributes.

Predefined Device Object Information

The **Predefined Devices (PdDv)** object class contains entries for all known device types supported by the system. This information is modified only when a new device is added. The following tables contain information for the Multiport Model 2 adapter.

Device type:	portmaster
Device class:	adapter
Device subclass:	isa
Unique Type:	/adapter/isa/portmaster
Prefix name:	арт

Flag	Setting
Base Device	no
Vital Product Data (VPD)	no
Detectable	yes
Change Status	set to NEW when defining device
Bus Extender	no
Field Replaceable Unit (FRU)	yes

Predefined Connection Object Information

The **Predefined Connection (PdCn)** object class contains connection and dependency information for devices. The following table contains predefined connection information for the Multiport Model 2 adapter.

Connection Location:	0
Unique Type:	/adapter/isa/portmaster
Connection Key:	portmaster

Predefined Attribute Object Information

The **Predefined Attribute (PdAt)** object class contains an entry for each existing attribute for a device, beyond what is defined in the **Predefined Device** object class. The **PdAt** information for the Multiport Model 2 adapter includes the following definitions:

Attribute	Description	Possible Values
bus_intr_lvl	ISA Interrupt Level	3,4,7,9,10,11,12
bus_io_addr	Bus IO Address	0x2A0-0x3EA0, skip by 0x400
bus_mem_addr	Bus Memory Address	0xC0000-0xDE0000, skip by 0x200
window_size	Window Size	0x2000 (nonchangeable)
intr_priority	Interrupt Priority	2 (nonchangeable)

Subtypes:

- 4-Port Selectable Multiport Model 2 Adapter
- Multiport Model 2 Adapter
- 6-Port Multiport Model 2 Adapter (V.35)
- 6-Port Multiport Model 2 Adapter (X.21)
- 8-Port Multiport Model 2 Adapter (RS-232)
- 8-Port Multiport Model 2 Adapter (RS-422)

Note: Currently, only the 4-Port Selectable Multiport Model 2 Adapter is supported.

Multiport Model 2 Power Management

The device driver for the Multiport Model 2 adapter provides for Power Management (PM) support. PM is a technique that enables hardware and software to minimize system power consumption. PM is generally only important to low–end models, such as Notebooks and systems operating on battery.

When Power Management is enabled, the system enters a power–saving mode under one of the following conditions:

- Expiration of the idle timer
- A direct command from the user

- A low battery
- Closing of the Notebook lid

PM state transitions include the following:

- enable
- standby
- suspend
- hibernation
- shutdown

Each transition implies a further decrease in the power supply to the various system components.

Refer to "Using Power Management" in AIX 4.3 System Management Guide: Operating System and Devices for detailed information on Power Management and its configuration and function.

Impact to External Connection – Network Provider

Bringing a system to the **suspend**, **hibernation**, or **shutdown** states results in loss of power to the adapters. All network connections are lost. Externally, this can be viewed as pulling the physical connection (cable), since it results in loss of signal at the physical layer. A data set ready (DSR) error is generated. Any application that has the device driver opened should be stopped before suspending, hibernating, or shutting down the system.

Once power is restored, all previously configured ports are restored to an **available** state. At that time, any applications previously using these connections can be restarted.

2-Port Multiprotocol HDLC Network Device Driver Overview

The 2–Port Multiprotocol Adapter high–level data link control (HDLC) device driver is a component of the communication I/O subsystem. This device driver provides support for the HDLC operation over the 2–Port Multiprotocol adapter at speeds up to 64K bps.

The following options provide access to the 2–Port Multiprotocol HDLC Network device driver:

- Systems Network Architecture (SNA)
- The synchronous data link control (SDLC) version of the GDLC Programming Interface
- User-written applications compatible with the SDLC MPQP-API (Multiprotocol Quad Port-Application Programming Interface)

Note: The above options require use of the **mpq**n special file, which allows access to the 2–Port Multiprotocol adapter's HDLC device driver through the SDLC COMIO device driver emulation subsystem. This subsystem must be installed and configured for each HDLC network device.

 User-written applications compatible with the HDLC Common Data Link Interface (CDLI) API

Configuring the 2-Port Multiprotocol Adapter

The following table explains how to configure a 2-Port Multiprotocol adapter.

Configuring the 2-Port Multiprotocol Adapter Tasks	
Web-based System Manager: wsm network fast path (wsm network application)	
-OR-	
Task	SMIT Fast Path
Add a Device Driver to the Adapter	smit mkhdlcdpmpdd
Reconfigure the Device Driver on the Adapter	smit chhdlcdpmpdd
Remove a Device Driver on the Adapter	smit rmhdlcdpmpdd
Make a Defined Device Driver Available	smit cfghdlcdpmpdd
Add an SDLC COMIO Emulator on the Adapter	smit mksdlcsciedd
Reconfigure the SDLC COMIO Emulator on the Adapter	smit chsdlcsciedd
Remove an SDLC COMIO Emulator on the Adapter	smit rmsdlcsciedd
Make a Defined SDLC COMIO Emulator Available	smit cfgsdlcsciedd

ARTIC960HX PCI Adapter Overview

The ARTIC960HX PCI Adapter MPQP COMIO device driver emulator is a component of the communication I/O subsystem. This device driver provides support for the ARTIC960HX PCI adapter at a maximum speed of 2M bps. The modems used must provide the clocking, since only external clocking is supported.

The following options provide access to the Multiport Model 2 device driver:

- Systems Network Architecture (SNA)
- The generic data link control (GDLC) Programming Interface
- User—written applications compatible with the MPQP—API (Multiprotocol Quad Port—Application Programming Interface), such as bisync applications.

These options require use of the **mpq**x special file, which allows access to the ARTIC960HX PCI adapter through the MPQP COMIO emulation device driver. This device driver must be installed and configured for each port on the ARTIC960HX PCI adapter. The **mpq**x special file resides in the /**dev** directory.

Note: The x in **mpq**x specifies the instance of the device driver; for example, mpq0.

The MPQP COMIO emulation device driver allows connectivity to remote host systems using the ARTIC960HX PCI adapter, either directly over a leased line. The device driver can provide a gateway between work group environments and remote data processing facilities.

Configuring the MPQP COMIO Emulation Driver over the ARTIC960HX PCI Adapter

The following table explains how to configure the MPQP COMIO Emulation Driver over the ARTIC960HX PCI Adapter.

Configuring the MPQP COMIO Emulation Driver Tasks	
Web-based System Manager: wsm network fast path (wsm network application)	
-OR-	
Task	SMIT Fast Path
Add a Device Driver	smit mktsdd
Reconfigure the MPQP COMIO Emulation Driver	smit chtsdd
Remove a Device Driver	smit rmtsdd
Configure a Defined Device Driver	smit cfgtsdd
Add a Port	smit mktsdports
Reconfigure an MPQP COMIO Emulation Port	smit chtsdports
Remove a Port	smit rmtsdports
Configure a Defined Port	smit cfgtsdports
Trace the MPQP COMIO Emulation Driver	smit trace_link

Chapter 7. Data Link Control

Generic data link control (GDLC) is a generic interface definition that allows application and kernel users a common set of commands to control data link control (DLC) device managers within the operating system. Topics discussed in this chapter are:

- Generic Data Link Control Environment Overview, on page 7-2
- Implementing the GDLC Interface, on page 7-5
- Installing GDLC Data Link Controls, on page 7-6
- GDLC Interface ioctl Entry Point Operations, on page 7-7
- GDLC Special Kernel Services, on page 7-10
- GDLC Problem Determination, on page 7-11
- Managing DLC Device Drivers, on page 7-14

Generic Data Link Control Environment Overview

Generic data link control (GDLC) is a generic interface definition that provides application and kernel users a common set of commands to control DLC device managers within the operating system.

For more information about the GDLC environment, see:

- Implementing GDLC Interface, on page 7-5
- Installing GDLC Data Link Controls, on page 7-6
- GDLC Interface ioctl Entry Point Operations, on page 7-7
- GDLC Special Kernel Services, on page 7-10
- GDLC Problem Determination, on page 7-11

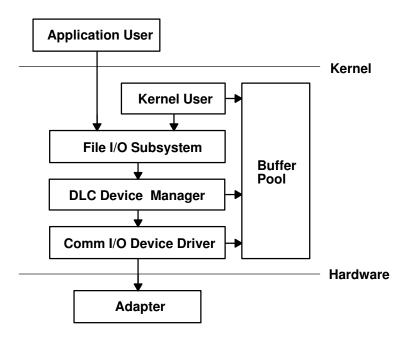
The GDLC interface specifies requirements for entry point definitions, functions provided, and data structures for all DLC device managers. DLCs that conform to the GDLC interface include:

- 8023 (IEEE 802.3 for Ethernet)
- ETHER (Standard Ethernet)
- SDLC (Synchronous Data Link Control)
- TOKEN (Token-Ring)
- FDDI (Fiber Distributed Data Interface)

DLC device managers perform higher—layer protocols and functions beyond the scope of a kernel device driver. However, the managers reside within the kernel for maximum performance and use a kernel device driver for their I/O requests to the adapter. A DLC user is located above or within the kernel.

Synchronous data link control (SDLC) and IEEE 802.2 Data Link Control are examples of DLC device managers. Each DLC device manager operates with a specific device driver or set of device drivers. SDLC, for example, operates with the multiprotocol device driver for the system's product and its associated adapter.

The basic structure of a DLC environment is shown in the DLC Device Manager Environment figure. Users within the kernel have access to the communications memory buffers (mbufs) and call the add entry points through the **fp** kernel services. Users above the kernel access the standard interface—to—kernel device drivers, and the file system calls the **dd** entry points. Data transfers require a move of data between user and kernel space.



DLC Device Manager Environment

The components of the DLC device manager environment are:

ve the kernel as an application or access

method.

Kernel User Resides within the kernel as a kernel process or device

manager.

File I/O Subsystem Converts the file-descriptor and file-pointer subroutines to

file pointer accesses of the switch table.

Buffer Pool Provides data—buffer services for the communications

subsystem.

Comm I/O Device Driver Controls hardware adapter I/O and direct memory access

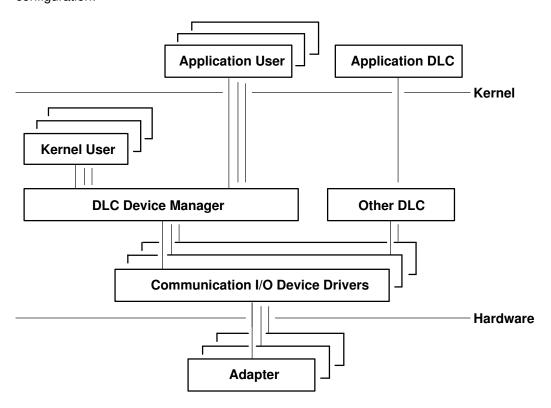
(DMA) registers, and routes receive packets to multiple

DLCs.

Adapter Attaches to the communications media.

A device manager written in accordance with GDLC specifications runs on all the operating system hardware configurations containing a communications device driver and its target adapter. Each device manager supports multiple users above as well as multiple device drivers and adapters below. In general, users operate concurrently over a single adapter, or each user operates over multiple adapters. DLC device managers vary based on their protocol constraints.

The Multiple User and Multiple Adapter Configuration figure illustrates a multiple user configuration:



Multiple User and Multiple Adapter Configuration

Meeting the GDLC Criteria

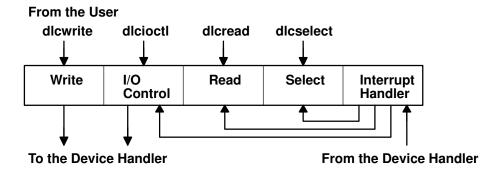
A GDLC interface must meet the following criteria:

- Be flexible and accessible to both application and kernel users.
- Have multiple user and multiple adapter capability, allowing protocols to take advantage of multiple sessions and ports.
- Support connection—oriented and connectionless DLC device managers.
- For special requirements beyond the scope of the DLC device manager in use, must allow transparent data transfer.

Implementing the GDLC Interface

Each DLC device manager is a standard /dev entry operating in the kernel as a multiplexed device manager for a specified protocol. For an adapter not in use by DLC, each open subroutine to a DLC device manager creates a kernel process. An open subroutine is also issued to the target adapter's device handler. If needed, issue additional open subroutines for multiple DLC adapter ports of the same protocol. Any open subroutine targeting the same port does not create additional kernel processes, but links the open subroutine with the existing process. There is always one kernel process for each port in use.

The internal structure of a DLC device manager has the same basic structure as a kernel device handler, except that a kernel process replaces the interrupt handler in asynchronous events. The read, write, I/O control, and select blocks function as shown in the Standard Kernel Device Manager figure.



Standard Kernel Device Manager

Installing GDLC Data Link Controls

You can install DLCs separately or in a group. A DLC device manager is automatically added into the kernel and set to the "Available" state for each type of DLC installed. Installation can be verified by issuing the **Islpp** command as follows:

lslpp -h dlctype

where dlctype is one of the following:

bos.dlc.8023 IEEE Ethernet (802.3) Data Link Control **bos.dlc.ether** Standard Ethernet Data Link Control

bos.dlc.fddiFDDI Data Link Controlbos.dlc.sdlcSDLC Data Link Control

bos.dlc.token Token-Ring Data Link Control

Information about an installed DLC can be displayed through Web-based System Manager, the System Management Interface Tool (SMIT), or the command line. On heavily used systems or communications ports, it may be necessary to change the DLC attributes to fine—tune the DLC's performance. If receive performance is slow, and the system error log indicates that ring queue overflow is occurring between the DLC and its device handler, increase the DLC's queue depth for incoming data. Finally, it is advisable to remove an installed DLC from the kernel when it is not needed for a lengthy period of time. This removal does not deinstall the DLC from the system, but allows kernel resources to be freed for other tasks until the DLC is needed again. Instructions for all of these tasks are in "Managing DLC Device Drivers", on page 7-14.

GDLC Interface ioctl Entry Point Operations

The generic data link control (GDLC) interface supports the following **ioctl** subroutine operations:

DLC_ENABLE_SAP Enables a service access point (SAP).

DLC_DISABLE_SAP Disables a SAP.

DLC START LS Starts a link station on a particular SAP as a caller or

listener.

DLC HALT LS Halts a link station.

DLC_TRACETraces a link station's activity for short or long activities.DLC_CONTACTContacts a remote station for a particular local link station.DLC_TESTTests the link to a remote for a particular local link station.

DLC ALTER Alters a link station's configuration parameters.

DLC_QUERY_SAP Queries statistics of a particular SAP.

DLC QUERY LSQueries statistics of a particular link station.

DLC_EXIT_LBUSY
Enters local—busy mode on a particular link station.

DLC_EXIT_LBUSY
Exits local—busy mode on a particular link station.

DLC_EXIT_SHOLD
Exits short—hold mode on a particular link station.

DLC_EXIT_SHOLD
Exits short—hold mode on a particular link station.

DLC_GET_EXCEP
Returns asynchronous exception notifications to the

application user.

Note: This **ioctl** subroutine operation is not used by the kernel user since all exception conditions are passed to the

kernel user through their exception handler.

DLC_ADD_GRP Adds a group or multicast receive address to a port.

DLC_ADD_FUNC_ADDR Adds a group or multicast receive functional address to a

port.

DLC_DEL_FUNC_ADDR Removes a group or multicast receive functional address

from a port.

IOCINFO Returns a structure that describes the GDLC device

manager. See the /usr/include/sys/devinfo.h file format

for more information.

Service Access Point

A SAP identifies a particular user service that sends and receives a specific class of data. This allows different classes of data to be routed separately to their corresponding service handlers. Those DLCs that support multiple concurrent SAPs have addresses known as Destination SAP and Source SAP imbedded in their packet headers. DLCs that can only support a single SAP do not need or use SAP addressing, but still have the concept of enabling the one SAP. In general, there is a SAP enabled for each DLC user on each port.

Most SAP address values are defined by IEEE standardized network-management entities or user-defined values as specified in the Token-Ring Network Architecture Reference. Some of the common SAP addresses are:

Null SAP (0x00) Provides some ability to respond to remote nodes even when no

> SAP has been enabled. This SAP supports only connectionless service and responds only to exchange identification (XID) and

TEST Link Protocol Data Units (LPDUs).

(0x04)

SNA Path Control Denotes the default individual SAP address used by Systems

Network Architecture (SNA) nodes.

PC Network NETBIOS (0xF0) Used for all DLC communication that is driven by Network Basic

Input/Output System (NetBIOS) emulation.

Discovery SAP

(0xFC)

Used by the local area network (LAN) name-discovery services.

Global SAP (0xFF) Identifies all active SAPs.

Link Station

A link station (LS) identifies an attachment between two nodes for a particular SAP pair. This attachment can operate as a connectionless service (datagram) or connection-oriented service (fully sequenced data transfer with error recovery). In general, there is one LS started for each remote attachment.

Local-Busy Mode

When an LS is operating in a connection-oriented mode, it needs to stop the remote station's sending of information packets for reasons such as resource outage. Notification can then be sent to the remote station to cause the local station to enter local-busy mode. Once resources are available, the local station notifies the remote that it is no longer busy and that information packets can flow again. Only sequenced information packets are halted with local-busy mode. All other types of data are unaffected.

Short-Hold Mode

Use the short-hold mode of operation when operating over data networks with the following characteristics:

- · Short call-setup time
- Tariff structure that specifies a relatively small fee for the call setup compared to the charge for connect time.

During short-hold mode, an attachment between two stations is maintained only while there is data available for transfer between the two stations. When there is no data to send, the attachment is cleared after a specified time-out period and only reestablished when there is new data to transfer.

Testing and Tracing a Link

To test an attachment between two stations, instruct an LS to send a test packet from the local station. This packet is echoed back from the remote station if the attachment is operating correctly.

Some data links are limited in their support of this function due to protocol constraints. SDLC, for example, only generates the test packet from the host or primary station. Most other protocols, however, allow test packets to be initiated from either station.

To trace a link, line data, and special events (such as station activation, termination, and time outs) can be logged in the generic trace facility for each LS. This function helps determine the cause of certain communications attachment problems. The GDLC user can select either short or long entries to be traced.

Short entries consist of up to 80 bytes of line data, while long entries allow full packets of data to be traced.

Tracing can be activated when an LS is started, or it can be dynamically activated or terminated at any time afterward.

Statistics

Both SAP and LS statistics can be queried by a GDLC user. The statistics for a SAP consist of the current SAP state and information about the device handler. LS statistics consist of the current station states and various reliability, availability, and serviceability counters that monitor the activity of the station from the time it is started.

GDLC Special Kernel Services

Generic data link control (GDLC) provides special services for a kernel user. However, a trusted environment must exist within the kernel. Instead of the DLC device manager copying asynchronous event data into user space, the kernel user must specify function pointers to special routines called function handlers. Function handlers are called by the DLC at the time of execution. This allows maximum performance between the kernel user and the DLC layers. Each kernel user is required to restrict the number of function handlers to a minimum path length and use the communications memory buffer (mbuf) scheme.

A function handler must never call another DLC entry directly. This is because direct calls are made under lock, causing a fatal sleep. The only exception to this rule is that a kernel user may call the **dlcwritex** entry point during its service of any of the four receive data functions. Calling the **dlcwritex** entry point allows immediate responses to be generated without an intermediate task switch. Special logic is required within the DLC device manager to check the process identification of the user calling a write operation. If it is a DLC process and the internal queuing capability of the DLC has been exceeded, the write is sent back with a bad return code (**EAGAIN** return value) instead of putting the calling process (DLC) to sleep. It is then up to the calling user subroutine to return a special notification to the DLC from its receive data function to ensure a retry of the receive buffer at a later time.

The user-provided function handlers are:

Datagram Data Received Routine	Called any time a datagram packet is received for the kernel user.
Exception Condition Routine	Called any time an asynchronous event occurs that must notify the kernel user, such as SAP Closed or Station Contacted.
I–Frame Data Received Routine	Called each time a normal sequenced data packet is received for the kernel user.
Network Data Received Routine	Called any time network-specific data is received for the kernel user.
XID Data Received Routine	Called any time an exchange identification (XID) packet is received for the kernel user.

The dicread and dicselect entry points for DLC are not called by the kernel user because the asynchronous functional entries are called directly by the DLC device manager. Generally, any queuing of these events must occur in the user's function handler. If, however, the kernel user cannot handle a particular receive packet, the DLC device manager may hold the last receive buffer and enter one of two special user-busy modes:

User-Terminated Busy Mode (I-frame only)

If the kernel user cannot handle a received I-frame (due to problems such as queue blockage), a DLC_FUNC_BUSY return code is given back, and DLC holds the buffer pointer and enters local-busy mode to stop the remote station's I-frame transmissions. The kernel user must call the Exit Local Busy function to reset local-busy mode and start the reception of I-frames again. Only normal sequenced I-frames can be stopped. XID, datagram, and network data are not affected by local-busy mode.

Timer-Terminated Busy Mode (all frame types)

If the kernel user cannot handle a particular receive packet and wants DLC to hold the receive buffer for a short period and then re-call the user's receive function, a DLC FUNC RETRY return code is given back to DLC. If the receive packet is a sequenced I-frame, the station enters local-busy mode for that period. In all cases, a timer is started; once the timer expires, the receive data functional entry is called again.

GDLC Problem Determination

Each generic data link control (GDLC) provides problem determination data that can be used to isolate network problems. Four types of diagnostic information are provided:

- DLC Status Information, on page 7-11
- DLC Error Log, on page 7-12
- GDLC Link Station Trace Facility, on page 7-12
- LAN Monitor Trace, on page 7-13

DLC Status Information

Status information can be obtained for a service access point (SAP) or a link station (LS) using the DLC QUERY SAP and DLC QUERY LS ioctl subroutines to call the specific DLC kernel device manager in use.

The DLC QUERY SAP ioctl subroutine obtains individual device driver statistics from various devices:

- Token–Ring
- Ethernet
- Multiprotocol

The DLC QUERY LS ioctl subroutine obtains LS statistics from various DLCs. These statistics include data link protocol counters. Each counter is reset by the DLC during the **DLC START LS** ioctl subroutine, and generally runs continuously until the LS is terminated and its storage is freed. If a counter reaches the maximum count, the count is frozen and no wraparound occurs.

The suggested counters provided by a DLC device manager are listed as follows. Some DLCs can modify this set of counters based on the specific protocols supported. For example, the number of rejects or receive-not-ready packets received might be meaningful.

Test Commands Sent	Contains a binary count of the test commands sent to the remote station by GDLC, in response to test commands issued by the user.
Test Command Failures	Contains a binary count of the test commands that did not complete properly due to problems such as the following:
	Invalid response
	Bad data compare
	Inactivity
Test Commands Received	Contains a binary count of valid test commands received, regardless of whether the response is completed properly.
Sequenced Data Packets Transmitted	Contains a binary count of the total number of normal sequenced data packets transmitted to the remote LS.
Sequenced Data Packets Transmitted	Contains a binary count of the total number of normal sequenced data packets retransmitted to the remote LS.
Maximum Contiguous Retransmissions	Contains a binary count of the maximum number of times a single data packet has been retransmitted to the remote LS before acknowledgment. This counter is reset each time a valid acknowledgment is received.
Sequenced Data Packets Received	Contains a binary count of the total number of normal sequenced data packets correctly received.

Invalid Packets Received Contains a binary count of the number of invalid commands

or responses received, including invalid control bytes.

invalid I-fields, and overflowed I-fields.

Adapter Detected Receive

Errors

Contains a binary count of the number of receive errors

reported back from the device driver.

Errors

Adapter Detected Transmit Contains a binary count the number of transmit errors

reported back from the device driver.

Receive Inactivity

Time-outs

Contains a binary count of the number of receive time-outs

that have occurred.

Command Polls Sent Contains a binary count of the number of command packets

sent that requested a response from the remote LS.

Command Repolls Sent Contains a binary count of the total number of command

packets retransmitted to the remote LS due to a lack of

response.

Command Contiguous

Repolls

Contains a binary count of the number of times a single command packet was retransmitted to the remote LS due to a lack of response. This counter is reset each time a valid

response is received.

DLC Error Log

Each DLC provides entries to the system error log whenever errors are encountered.

The user can obtain formatted errorlog data by issuing the errpt command. When used with the -N DLCName flag, the errpt command produces a summary report of all the error log entries for the resource name indicated by the DLCName parameter. Valid values for the DLCName parameter include:

SYSXDLCE Standard Ethernet datalink **SYSXDLCI** IEEE 802.3 Ethernet datalink

SYSXDLCT Token-ring datalink **SYSXDLCS** SDLC datalink

The format of each required alert vector can be found in SNA Format and Protocol Reference Manual: Management Services.

For more information on the error log facility, refer to "Error Logging Overview" AIX Version 4.3 Problem Solving Guide and Reference.

GDLC Link Station Trace Facility

GDLC provides optional entries to a generic system trace channel as required by the system product Reliability/Availability/Serviceability. The default is trace disabled. This provides maximum performance and reduces the number of system resources used. For information on additional ways to use trace facilities, see "Managing DLC Device Drivers", on page 7-14.

Trace Channels

The operating system supports up to seven generic trace channels in operation at the same time. A user must allocate a channel before starting an LS trace, with the **DLC START LS** ioctl operation or the DLC TRACE ioctl operation. This is accomplished with the trcstart and trcon subroutines.

Trace activity in the LS must be stopped either by halting the LS or by issuing an ioctl (DLC_TRACE, flags=0) operation to that station. Once the LS has stopped tracing, the channel is disabled with the trcoff subroutine and returned to the system with the trcstop subroutine.

Trace Reports

The user can obtain formatted tracelog data by issuing the trcrpt command with the appropriate file name, such as:

```
trcrpt /tmp/link1.log
```

This example produces a detailed report of all link trace entries in the /tmp/link1.log file, if a prior trcstart subroutine specified the /tmp/link1.log file as the (-o) name for the trace log.

Trace Entries

For each trace entry, GDLC generates the trcgenkt kernel service to the kernel generic trace.

LAN Monitor Trace

Each of the local area network data link controls (DLCETHER, DLC8023, DLCFDDI, and DLCTOKEN) provides an internal monitor trace capability that can be used to identify the execution sequence of pertinent entry points within the code. This is useful if the network is having problems that indicate the data link is not operating properly, and the sequence of events may indicate the cause of the problems. This trace is shared among the LAN data link controls, and inactive is the default.

The LAN monitor trace can be enabled by issuing the following command:

```
trace -j 246
```

where 246 is the hook ID to be traced.

Tracing can be stopped with the **trcstop** command and a report can be obtained with the following command:

```
trcrpt -d 246
```

where 246 is the hook ID of the trace for which you want a report.

Note: Exercise caution when enabling the monitor trace, since it directly affects the performance of the DLCs and their associates.

For information on additional ways to use trace facilities, see "Managing DLC Device Drivers", on page 7-14.

Managing DLC Device Drivers

A DLC must be added to the system prior to use. Each installed DLC is automatically added after installation and at each system reboot. If a DLC has been removed without a subsequent reboot, it may be re-added.

Managing DLC Device Drivers Tasks		
Web-based System Manager: wsm devices fast path (Devices application)		
-OR-		
Task	SMIT Fast Path	Command or File
Add an Installed DLC	Choose one (by device driver name): smit cmddlc_sdlc smit cmddlc_token smit cmddlc_qllc smit cmddlc_ether ¹ smit cmddlc_fddi then select Add	mkdev ²
Change DLC Attributes ^{3,4}	Choose one (by device driver name): smit cmddlc_sdlc_ls smit cmddlc_token_ls smit cmddlc_qllc_ls smit cmddlc_ether_ls smit cmddlc_fddi_ls	chdev ²
Start DLC Local Area Network Monitor Trace ⁵	smit trace	trace –j <i>nnn</i> where the value <i>nnn</i> is the hook ID to be traced
Stop DLC Local Area Network Monitor Trace	smit trestop	trcstop ²
Generate DLC Local Area Network Monitor Trace Report	smit trcrpt	trcrpt –d <i>nnn</i> where the value <i>nnn</i> is the hook ID to be reported
List Current DLC Information ³	Choose one (by device driver name): smit cmddlc_sdlc_ls smit cmddlc_token_ls smit cmddlc_qllc_ls smit cmddlc_ether_ls¹ smit cmddlc_fddi_ls	Isdev ² or Isattr ²
Remove a DLC ^{3,6}	Choose one (by device driver name): smit cmddlc_sdlc_rm smit cmddlc_token_rm smit cmddlc_qllc_rm smit cmddlc_ether_rm ¹ smit cmddlc_fddi_rm	rmdev ²

Notes:

- 1. The SMIT fast path for an Ethernet device manager includes both Standard Ethernet and IEEE 802.3 Ethernet device managers.
- 2. Details about command line options are provided in the command descriptions for mkdev, chdev, trace, trcstop, trcrpt, Isdev, Isattr, or rmdev in AIX Commands Reference.
- 3. A DLC must be installed and added before you can list, show, change, or remove its attributes. An attribute change is only successful if there are no active opens against the target DLC. Before issuing the change action, the user may have to stop services such as SNA, OSI, or NetBIOS from using the DLC.
- 4. Changing the receive-queue size directly affects system resources. Make this change only if the DLC is having receive-queue problems, such as sluggish performance or overflows between the DLC and its device handler.
- 5. Exercise caution when enabling the monitor trace since it directly affects the performance of the DLCs and their associates.
- 6. Removing a DLC is only successful if there are no active opens against the target DLC. Before issuing the remove action, the user may have to stop services such as SNA, OSI, or NetBIOS from using the DLC.

Chapter 8. Basic Networking Utilities

This chapter presents information on installing, configuring, and maintaining Basic Network Utilities (BNU). The following topics are discussed:

- BNU Overview, on page 8-2
- Configuring BNU, on page 8-10
- Maintaining BNU, on page 8-18
- BNU Configuration Files, on page 8-29
- BNU Files, Commands, and Directories Reference, on page 8-36

BNU Overview

The Basic Networking Utilities (BNU) are a group of programs, directories, and files that can be used to communicate with any UNIX system on which a version of the UNIX-to-UNIX Copy Program (UUCP) is running. BNU is one of the Extended Services programs that can be installed with the base operating system.

A group of commands related to UUCP, a UNIX-to-UNIX communication program developed by AT&T and modified as part of the Berkeley Software Distribution (BSD) are contained in BNU. BNU provides commands, processes, and a supporting database for connections to local and remote systems. Communication networks such as Token-Ring and Ethernet are used to connect systems on local networks. A local network can be connected to a remote system by hardwire or telephone modem. Commands and files can then be exchanged between the local network and the remote system.

Topics discussed in this section are:

- How BNU Works
- BNU File and Directory Structure
- BNU Security
- BNU Daemons

Before users on your system can run BNU programs, BNU must be installed and configured.

BNU is controlled by a set of configuration files that determine whether remote systems can log in to the local system and what they can do after they log in. These configuration files must be set up according to the requirements and resources of your system.

To maintain BNU, you must read and remove log files periodically and check the BNU queues to ensure jobs are transferring to remote systems properly. You must also periodically update the configuration files to reflect changes in your system or remote systems.

For more information, see:

- Configuring BNU
 - Information to Collect Before Configuring BNU
- Maintaining BNU
 - Working with BNU Log Files
 - Using BNU Maintenance Commands

How BNU Works

BNU uses a set of hardware connections and software programs to communicate between systems. A structure of directories and files tracks BNU activities. This structure includes a set of public directories, a group of administrative directories and files, configuration files, and lock files. Most of the directories for BNU are created during the installation process. Some of the administrative directories and files are created by various BNU programs.

With the exception of the remote login commands, BNU works as a batch system. When a user requests a job sent to a remote system, BNU stores the information needed to complete the job. This is known as *queuing* the job. At scheduled times, or when a user instructs it to do so, BNU contacts various remote systems, transfers queued work, and accepts jobs. These transfers are controlled by the configuration files on your system and those of the remote system.

National Language Support for BNU Commands

All BNU commands, except **uucpadm**, are available for National Language Support. User names need not be in ASCII characters. However, all system names must be in ASCII characters. If a user attempts to schedule a transfer or a remote command execution involving non–ASCII system names, BNU returns an error message.

BNU File and Directory Structure

BNU uses a structure of directories and files to keep track of their activities. This structure includes:

- · Public directories
- · Configuration files
- Administrative directories and files
- Lock files

Most of the directories for BNU are created during the installation process. Some of the administrative directories and files are created by various BNU programs as they run.

BNU Public Directories

The BNU public directory, /var/spool/uucppublic, stores files that have been transferred to the local system from other systems. The files wait in the public directory until users claim them with the uupick command. The public directory is created when BNU is installed. Within the public directory, BNU creates a subdirectory for each remote system that sends files to the local system.

BNU Configuration Files

The BNU configuration files, also known as the BNU supporting database, reside in the /etc/uucp directory. The files must be configured specifically for your system. They are owned by the uucp login ID and can be edited only with root authority. The configuration files contain information about:

- · Accessible remote systems
- Devices for contacting the remote systems
- Times to contact the remote systems
- What the remote systems are allowed to do on your system

Some configuration files also specify limits on BNU activities that prevent your system from becoming overloaded.

The BNU configuration files include:

Devices Contains information about available devices, including both

modems and direct connections.

DialcodesContains dialing code abbreviations, which allow you to shorten

phone numbers in the Systems file.

Dialers Specifies calling command syntax for a specific modem type

("dialer").

Maxuuscheds Limits simultaneous scheduled jobs.

Maxuuxqts Limits simultaneous remote command executions.

Permissions Contains access permission codes. This file is the primary file for

determining the security for BNU.

Poll Specifies when the BNU program should poll remote systems to

initiate tasks.

Sysfiles Lists files that serve as the Systems, Devices, and Dialers files

for the BNU configuration. If this file is not used, the default files

are /etc/uucp/Systems, /etc/uucp/Devices, and

/etc/uucp/Dialers.

Systems Lists accessible remote systems and information needed to

contact them, including the device to use and the user name and password combinations you need to log in. Also specifies the

times when the systems can be contacted.

The configuration files cross—reference each other when BNU is in use. For example:

The Devices file contains a Token field that refers to entries in the Dialers file.

- The **Systems** file contains an entry for a *Class* of device. A device of each *Class* referred to in the **Systems** file must be defined in the **Devices** file.
- The **Poll** file contains entries for systems your system calls. Each of these systems must be defined in the **Systems** file.

Entries in the BNU configuration files depend on the types of connections between your system and each remote system. For example, special entries must be made if Transmission Control Protocol/Internet Protocol (TCP/IP) or direct connections are used to contact other systems. If modems are used to contact other systems, the modems must be defined in the **Dialers** file.

The **Systems**, **Devices**, and **Permissions** files must be configured on your system before you can contact remote systems using BNU. Other configuration files enable you to use BNU capabilities, such as automatic polling. Many of the configuration files must be modified periodically to reflect changes to your system or the systems you contact. The **Sysfiles** file can be used to specify files other than the default **Systems**, **Devices**, and **Dialers** files to fulfill the same role.

BNU Administrative Directories and Files

The BNU administrative directories and files are in subdirectories of the /var/spool/uucp directory. These directories and files contain two types of information:

- · Data waiting to be transferred to other systems
- Log and error information about BNU activities

Under the /var/spool/uucp directory, BNU creates the following directories:

.Admin Contains four administrative files:

audit

Foreign

errors

xferstats

These files contain error and log information about BNU activities.

.Corrupt Contains copies of files that cannot be processed by the BNU

program.

.Log and **.Old** Contain log files from BNU transactions.

.Status Stores the last time the **uucico** daemon tried to contact remote

systems.

.Workspace Holds temporary files that the file transport programs use

internally.

.Xqtdir Contains execute files with lists of commands that remote

systems can run.

SystemName Contains files used by file transport programs. These files are:

• Command (C.*)

• Data (D.*)

• Execute (X.*)

Temporary (TM.*)

BNU creates a *SystemName* directory for each remote system it contacts.

The directories whose names begin with a dot are *hidden*. They cannot be found with an **Is** or **Ii** command unless the **–a** flag is used. When the **uucico** daemon is started, it searches the /**var**/spool/uucp directory for work files and transfers the files from any directory that is not hidden. The **uucico** daemon sees only the *SystemName* directories, not the other administrative directories.

The files in the hidden directories are owned by the uucp login ID. These files can be accessed only with root authority or with a login ID with a UID of 5.

For further information about maintaining the BNU administrative directories, see "Maintaining BNU", on page 8-18.

BNU Lock Files

The BNU lock files are stored in the /etc/locks directory. When BNU uses a device to connect to a remote computer, it places a lock file for that device in the /etc/locks directory. When another BNU program or any other program needs the device, that program checks the /etc/locks directory for a lock file. If a lock file exists, the program waits until the device is available or uses another device for the communication.

In addition, the **uucico** daemon places lock files for remote systems in the /**etc/locks** directory. Before contacting a remote system, the **uucico** daemon checks the /**etc/locks** directory for a lock file for that system. These files prevent other instances of the **uucico** daemon from establishing duplicate connections to the same remote system.

Note: Other software besides BNU, such as Asynchronous Terminal Emulation (ATE) and TCP/IP, uses the /etc/locks directory.

BNU Security

Because other systems contact your system to log in, transfer files, and enter commands, BNU provides a means to establish security. BNU security enables you to restrict what users of remote systems can do on the local system (users of remote systems can also restrict what you can do on their systems). BNU runs several daemons to complete its activities and uses administrative directories to store the files it needs. BNU also keeps a log of its own activities.

BNU security works on several levels. When you configure BNU, you can determine:

- Who on your system has access to BNU files.
- Which remote systems your system can contact.
- How users on remote systems log in to your system.
- What users on remote systems can do on your system once they log in.

uucp Login ID

When BNU is installed, all of the configuration files, daemons, and many of the commands and shell procedures are owned by the uucp login ID. The uucp login ID has a user ID (UID) of 5 and a group ID (GID) of 5. The **cron** daemon reads the /var/spool/cron/crontabs/uucp file to schedule automatic jobs for BNU.

Usually, logging in as user uucp is not allowed. To change files that are owned by the uucp login ID, log in with root authority.

Attention: Allowing remote systems to log in to the local system with the uucp login ID seriously jeopardizes the security of the local system. Remote systems logged in with the uucp ID can display and possibly modify the local **Systems** and **Permissions** files depending on the other permissions specified in the LOGNAME entry. It is strongly recommended that you create other BNU login IDs for remote systems and reserve the uucp login ID for the person responsible for administering BNU on the local system. For the best security, each remote system that contacts the local system should have a unique login ID with a unique UID number.

BNU Login IDs

The startup shell for BNU login IDs is the **uucico** daemon (/usr/sbin/uucp/uucico). When remote systems call your system, they automatically start the **uucico** daemon on your system. Login IDs for BNU have a uucp group ID of 5.

Login IDs used by remote systems need passwords. In order to prevent security from prompting a new BNU login ID for a new password when the remote system logs in, you must set the password as soon as you create the account. To do this, use the **passwd** command followed by the **pwdadm** command. For example, to set a password for the login ID nuucp, log in as the root user and enter the following commands:

```
passwd nuucp
pwadm -f NOCHECK
nuucp
```

The system prompts you for a password for the nuucp login ID. Completing these steps allows the remote system to log in without being immediately prompted for a new password (which the batch–oriented nuucp login ID can not provide).

After creating the login ID for a remote system, notify that system's BNU administrator of the login ID and password to access your system.

Creating a BNU Administrative Login ID

A user with root authority can set up a BNU administrative login ID. This is useful if you want to delegate BNU administration duties to a user without root authority. The BNU administrative login ID should have password security, a UID of 5, and be in a uucp group ID 5. The login shell for the administrative login should be the /usr/bin/sh program (instead of the uucico daemon). Giving the BNU administrative login a UID of 5 causes it to have the same privileges as the uucp login ID. Thus, for security, remote systems should not be allowed to log in as the BNU administrator.

Adding BNU Login Shells to the login.cfg File

User configuration stanzas in the **login.cfg** and user files provide configuration information for programs that change user attributes or add new users. The stanza in the **login.cfg** file is labeled usw. The stanzas in the user files are labeled with the individual user names.

The shells attribute in the usw stanza defines the valid shells on the system. The value is a list of comma–separated full path names. The default is:

/usr/bin/sh,/usr/bin/bsh,/usr/bin/csh,/usr/bin/ksh

Before using Web-based System Manager or the System Management Interface Tool (SMIT) to add a new BNU user, add the program name /usr/sbin/uucp/uucico to the usw shells stanza. The new program name should be separated from the last entry by a comma and no blanks; for example:

/usr/bin/sh,/usr/bin/bsh,/usr/bin/csh,/usr/bin/ksh,/usr/sbin/uucp/uucico

Attention: Web-based System Manager or SMIT will fail when specifying /usr/sbin/uucp/uucico as a user's login shell if the program name is not added to the login.cfg file.

Security and the Systems and remote.unknown Files

On most BNU systems, only remote systems listed in the /etc/uucp/Systems file or one of its substitutes (specified in the Sysfiles file) can log in to the local system. The /usr/sbin/uucp/remote.unknown script is executed whenever an unknown system attempts to call the local system. This script refuses to let the unknown system log in and makes an entry in the /var/spool/uucp/.Admin/Foreign file recording the time of the login attempt.

With root authority, or as a BNU administrator, you can modify the **remote.unknown** shell procedure to log more information about the remote system or to store the information in a different file. For example, you can modify the shell procedure to send mail to the BNU administrator whenever an unknown system tries to log in.

By taking away execute permissions on the **remote.unknown** shell procedure, you enable unknown machines to log in. In this case, you should add a MACHINE=OTHER entry to the /etc/uucp/Permissions file to establish permissions for the unknown machines.

Your system can contact only remote systems listed in the **Systems** file. This prevents users on your system from contacting unknown systems.

Security and the Permissions File

The /etc/uucp/Permissions file determines:

- Remote login user names for logging in to the local system.
- Approved commands and privileges for remote systems logging in to the local system.

The /etc/uucp/Permissions file contains two types of entries:

LOGNAME Defines login names and the privileges associated with them.

LOGNAME entries take effect when a remote system calls the

local system and attempts to log in.

MACHINE Defines machine names and the privileges associated with them.

MACHINE entries take effect when the remote system attempts to

carry out commands on the local system.

Options in the **Permissions** file enable you to establish various levels of security for each remote system. For example, if many remote systems share one login ID on the local system, use the VALIDATE option to require each remote system to use a unique login ID. The SENDFILES, REQUEST, and CALLBACK options specify which system has control, keeping the local system in control of transactions if necessary.

The READ, WRITE, NOREAD, and NOWRITE options define access to specific directories on the local system. These options also control where on your system remote users can place data. The COMMANDS option limits the number of commands users on remote systems can execute on the local system. The COMMANDS=ALL option allows total privileges to systems closely associated with your system.

Attention: The COMMANDS=ALL option can seriously jeopardize the security of your system.

BNU Daemons

The BNU software includes four daemons stored in the /usr/sbin/uucp directory:

uucico Facilitates file transfers.

uusched Facilitates work request scheduling of files queued in the local

spooling directory.

uuxqt Facilitates remote command executions.uucpd Facilitates communications using TCP/IP.

The **uucico**, **uusched**, and **uuxqt** daemons are started by the **cron** daemon according to a schedule set by the BNU administrator. With root authority, you can also start these daemons manually. The **uucpd** daemon should be started by the TCP/IP **inetd** daemon.

Using the uucico Daemon

The **uucico** daemon transports the files required to send data from one system to another. The **uucp** and **uux** commands start the **uucico** daemon to transfer command, data, and execute files to the designated system. The **uucico** daemon is also started periodically by the BNU scheduler, the **uusched** daemon. When started by the **uusched** daemon, the **uucico** daemon attempts to contact other systems and execute the instructions in the command files.

How the Daemon Process Begins

To run the instructions in the command files, the **uucico** daemon first checks the /etc/uucp/Systems file (or one or more other files specified by /etc/uucp/Sysfiles) for the system to be called. The daemon then checks the **Systems** file entry for a valid time to call. If the time is valid, the **uucico** daemon checks the *Type* and *Class* fields and accesses the /etc/uucp/Devices file (or one or more other files specified by /etc/uucp/Sysfiles) for a device that matches.

After finding a device, the **uucico** daemon checks the /etc/locks directory for a lock file for the device. If one exists, the daemon checks for another device of the requested type and speed.

When no device is available, the daemon returns to the **Systems** files for another entry for the remote system. If one exists, the daemon repeats the process of searching for a device. If another entry is not found, the daemon makes an entry in the

/var/spool/uucp/.Status/SystemName file for that remote system and goes on to the next request. The command file remains in the queue. The uucico daemon attempts the transfer again at a later time. The later attempt is called a *retry*.

When the Daemon Reaches the Remote System

When the **uucico** daemon reaches the remote system, it uses the instructions in the **Systems** files to log in. This causes an instance of the **uucico** daemon to be invoked on the remote system as well.

The two **uucico** daemons, one on each system, work together to make the transfer. The **uucico** daemon on the calling system controls the link, specifying the requests to be performed. The **uucico** daemon on the remote system checks the local permissions for whether they allow the request to be performed. If so, the file transfer starts.

After the **uucico** daemon on the calling system has finished transferring all requests it has for the remote system, it sends a hangup request. When the remote **uucico** daemon has transactions to send to the calling system, it denies the hangup request, and the two daemons reverse roles.

Note: Either the /etc/uucp/Permissions file on the local system or the /etc/uucp/Permissions file on the remote system can forbid the daemons to reverse roles. In this case, the remote system must wait to transfer files until it calls the local system.

When nothing is left to be transferred in either direction, the two **uucico** daemons hang up. At this point, the **uuxqt** daemon is called to execute remote command requests.

Throughout the transfer process, the **uucico** daemons on both systems log messages in the BNU log and error files.

Using the uusched Daemon

The **uusched** daemon schedules the transfer of files that are queued in the spooling directory on the local system. The spooling directory is /var/spool/uucppublic. When the **uusched** daemon is invoked, it scans the spooling directory for command files, then randomizes the files and starts the **uucico** daemon. The **uucico** daemon transfers the files.

Using the uuxqt Daemon

When a user issues the **uux** command to run a specified command on a designated system, the **uuxqt** daemon runs the command. After creating the necessary files, the **uux** command starts the **uucico** daemon, which transfers those files to the public spooling directory on the specified system.

The **uuxqt** daemon periodically searches the spooling directory for command–execution requests on every connected system. When it locates such a request, the **uuxqt** daemon checks for necessary files and permissions. Then, if permitted, the daemon runs the specified command.

Using the uucpd Daemon

The **uucpd** daemon must be able to run on the remote system before BNU can establish communications with a remote computer with Transmission Control Protocol/Internet Protocol (TCP/IP). The **uucpd** daemon is a subserver of the TCP/IP **inetd** daemon and is started by the **inetd** daemon.

By default, the **inetd** daemon is configured to start the **uucpd** daemon. However, if this has been changed on your system, you may need to reconfigure the **inetd** daemon to start the **uucpd** daemon.

Configuring BNU

The following procedures detail how to configure Basic Network Utilities (BNU) for various types of connections, including hardwired, modem, and Transmission Control Protocol/Internet Protocol (TCP/IP) connections.

Prerequisites

- BNU must be installed on your system.
- You must have root user authority to edit the BNU configuration files.
- If you are using direct connections for BNU communications, the appropriate hardwired connections between your system and the remote systems must be set up.
- If you are using modems for BNU communications, you must have installed and configured each modem.
- If one or more of your connections uses TCP/IP, then TCP/IP must be running between your system and the apropriate remote systems.
- Collect the information you need to configure BNU. This information should include a list
 of remote systems and lists of devices and modems to use to connect to the systems.

Information to Collect before Configuring BNU

Before configuring BNU, gather the information listed:

- For each remote system your system will call, collect the following information:
 - System name
 - Login name your system should use on the remote system
 - Password for the login name
 - Login and password prompts on the remote system
 - Type of connection you will use to reach the remote system (TCP/IP, direct, or telephone)
- · If the connection is direct, collect:
 - The bit rate of the connection
 - The port on the local system to which the connection is attached
- If the connection is a telephone connection, collect:
 - The telephone number of the remote system
 - The speed of your modem that is compatible with that of the remote system

Note: If any of the remote systems will call your system, ensure the BNU administrator on each of the remote systems has all the preceding information about your system.

- For each local modem that you will use for BNU connections, collect the following information:
 - The chat script for the modem (consult the modem's documentation)

Note: For some modems, the chat script is already in the /etc/uucp/Dialers file.

- The local port for the modem

Using the information you collect, make a list of each device you need to connect to a remote system. Following is a sample list for local system morgan:

```
direct:
hera 9600 tty5
zeus& 2400 tty2
ariadne 2400 tty1
hayes modem (tty3): apollo, athena
TCP/IP: merlin, arthur, percy
```

To connect to system hera, a direct connection at a speed of 9600 from port tty5 is used. To connect to system apollo, the hayes modem connected to port tty3 is used. TCP/IP is used to connect to systems merlin, arthur, and percy.

Procedure

For BNU to function correctly at your site, you must configure the remote communications facilities to:

- List the devices used to establish a hardwired, telephone, or modem communications link.
- List the modems used to contact remote systems over the telephone network.
- List the accessible remote systems.
- List the alphabetic abbreviations representing the prefixes of telephone numbers used to contact the specified remote systems (optional).
- Set access permissions specifying the ways in which local and remote systems may communicate.
- Schedule monitoring for the networked remote systems (optional).

To create these lists, permissions, schedules, and procedures:

- · Modify the BNU configuration files
- Edit the /var/spool/cron/crontabs/uucp file to remove the comment characters (#) from the beginnings of the lines that schedule the automatic maintenance routines.

You must configure the **Systems**, **Devices**, and **Permissions** files before BNU will run correctly at your site. However, it is not necessary to modify the BNU configuration files in any particular order.

To configure BNU on your system:

1. Make sure that BNU is installed on your system by running the command:

```
lslpp -h bos.net.uucp
```

If BNU is installed, you will see bos.net.uucp in the output. If you do not see it, install bosext1 from the install tape.

 Set up appropriate login IDs and passwords for remote systems that will call your system, and tell the person responsible for administering BNU or UNIX-to-UNIX Copy Program (UUCP) on each remote system the login and password you have provided. This is done by editing the /etc/passwd, /etc/group, /etc/security/login.cfg, and /etc/security/passwd files.

Attention: Allowing remote systems to log into the local system with the uucp login ID seriously jeopardizes the security of your system. Remote systems logged in with the uucp ID can display and possibly modify (depending on the permissions specified in the LOGNAME entry of the **Permissions** file) the local **Systems** and **Permissions** files. It is strongly recommended that you create other BNU login IDs for remote systems and reserve the uucp login ID for the person administering BNU on the local system. For the best security, each remote system that contacts the local system should a have unique login ID with a unique UID number. These login IDs should have GIDs of 5.

Note: User configuration stanzas in the **login.cfg** and user files provide configuration information for Web-based System Manager and the System Management Interface Tool (SMIT). The stanza in the **login.cfg** file is labeled usw. The stanzas in the user file are labeled with the individual user names.

Before using Web-based System Manager or SMIT to add a new BNU user, add the program name /usr/sbin/uucp/uucico to the usw shells stanza. The new program name should be separated from the last entry by a comma and no blanks; for example:

/usr/bin/sh,/usr/bin/bsh,/usr/bin/csh,/usr/bin/ksh,/usr/sbin/uucp/uucico

Attention: Web-based System Manager or SMIT will fail when specifying /usr/sbin/uucp/uucico as a user's login shell if the program name is not added to the login.cfg file.

a. In the /etc/passwd file, add a user uucpadm to the system with the same group and user ID as uucp. This user ID will be helpful for several administration tasks as well as system debugging. Make sure that the uucpadm's shell is not /usr/sbin/uucp/uucico, and make the home directory /usr/sbin/uucp. The suggested entry in /etc/passwd is:

```
uucpadm:!:5:5::/usr/sbin/uucp:/bin/ksh
```

b. You have the option of maintaining separate logins or having one login for all BNU connections. If you need to maintain complete control over access by each individual machine, you must create separate login IDs, as well as combine the MACHINE and LOGNAME entries in the **Permissions** file. A few example /etc/passwd entries are shown here:

```
Umicrtk:!:105:5:micrtk
uucp:/usr/spool/uucppublic:/usr/sbin/uucp/uucico
Ufloydl:!:106:5:floydl
uucp:/usr/spool/uucppublic:/usr/sbin/uucp/uucico
Uicus:!:107:5:icus
uucp:/usr/spool/uucppublic:/usr/sbin/uucp/uucico
Urisctkr:!:108:5::/usr/spool/uucppublic:/usr/sbin/uucp/uucico
```

c. If you want to have one set of permissions and do not want to maintain separate control for any of your UUCP connections, you can have a single login for all machines such as the following:

```
nuucp:!:6:5::/usr/spool/uucppublic:/usr/sbin/uucp/uucico
```

- d. The user ID (the third colon separated field) must be unique to avoid a security risk. The group ID (the fourth separated field) must be 5, the same group as uucp. You can change the home directory (the sixth field) to any valid directory, but the login shell (the seventh field) must be /usr/sbin/uucp/uucico.
- e. Make sure that the /etc/group file contains the new users. An example of such an entry is:

```
uucp:!:5:uucp, uucpadm, nuucp, Umicrtk, Uicus, Urisctakr
```

- f. You may want to add any users to group uucp who will be using modems to dial out with programs other than the **cu** command.
- g. After editing these files as root, set up a password for the new users with the command passwd UserName.
- h. The login shell of your BNU logins (/usr/sbin/uucp/uucico) must be added to the list of shells in login.cfg. Look for the shells = stanza and append, /usr/sbin/uucp/uucico to the end of the list.

i. Sometimes, the default herald with all of its Ctrl—J's, will cause a uucico process that is trying to login to give up. (You may see the message Enough already.) You can avoid that by commenting out (with asterisks) the default stanza, and defining a stanza for your tty something like this:

```
/dev/tty0:
    herald = "\nrisc001 login:"
```

j. If you change a password from the root login, the flags entry in the stanza for the user in /etc/security/passwd will contain the following:

```
flags = ADMCHG
```

Change it to:

```
flags =
```

Otherwise, when the remote uucico logs in, it will be prompted to enter a new password, which it cannot do. Hence the login will fail.

k. After logging in as uucpadm, run the following command to read the current crontab for BNU into a temporary file:

```
crontab -l > /tmp/cron.uucp
```

I. Next, edit /tmp/cron.uucp to uncomment entries. They should look like:

```
20,50 * * * * /bin/bsh -c "/usr/sbin/uucp/uudemon.poll > /dev/null"
25,55 * * * * /bin/bsh -c "/usr/sbin/uucp/uudemon.hour > /dev/null"
45 23 * * /bin/bsh -c "/usr/sbin/uucp/uudemon.cleanu > /dev/null"
48 8,12,16 * * * /bin/bsh -c "/usr/sbin/uucp/uudemon.admin >/dev/null"
```

Entries can be changed to suit your needs. Read the edited version into the BNUs crontab with this command:

```
crontab /tmp/cron.uucp
```

m. Check to make sure your changes took effect by running this command:

```
crontab -1
```

- n. Set up BNU's data files: Systems, Permissions, Devices, Dialers, and Sysfiles. You could use the /usr/sbin/uucp/uucpadm command to initially set up the files and then edit them to suit your exact needs. Note that the Sysfiles file allows you to specify files other than /etc/uucp/Systems, /etc/uucp/Devices, and /etc/uucp/Dialers for BNU configuration. See Sysfiles for more information.
- 3. Decide whether to use dial—code abbreviations for telephone numbers (see the **Dialcodes** file format). If you decide to use dial—code abbreviations in the **Systems** files, set up the **Dialcodes** entry for each abbreviation. Refer to Dialcodes File Format for BNU in *AIX Files Reference* for details.

If you are using TCP/IP for your BNU connections, use the **netstat** command to see whether the **uucpd** daemon is runnable, by entering:

```
netstat -a
```

The **uucpd** daemon is started by the **inetd** daemon. If the **uucpd** daemon is not able to run, reconfigure the **inetd** daemon to start the **uucpd** daemon.

4. Using your list of devices, modify the **Devices** file on your system. Make an entry for each modem and each direct connection. If you are using TCP/IP, make sure you uncomment the TCP/IP entry in the **Devices** file. You can configure the /etc/uucp/Sysfiles file to specify other files to use for Devices configuration. Refer to the Devices File Format for BNU in AIX Files Reference for details on the Devices file. Refer to Sysfiles File Format for BNU for details on the Sysfiles file in AIX Files Reference.

Also, if you are using TCP/IP, check to see whether the /etc/services file includes:

```
uucp 540/tcp uucpd
```

If not, add the line.

- 5. Using your information about each remote system, modify the **Systems** file on your system. Use the commented examples in the **Systems** file as a guide when specifying your configuration. See the "BNU Systems File Format" in *AIX Files Reference* for details. If you are using TCP/IP, ensure the host–name table in the /etc/hosts file includes the name of the remote computer with which you want to connect. You can configure the /etc/uucp/Sysfiles file to specify other files to use for Systems configuration. Refer to Sysfiles File Format for BNU in *AIX Files Reference* for more information.
- 6. Using your information about devices and modems, make sure the **Dialers** file on your system contains an entry for each modem. If you are using TCP/IP and direct connections, make sure the TCP/IP entry and direct entries are present in the file. Refer to Dialers File Format for BNU in AIX Files Reference for details. You can configure the /etc/uucp/Sysfiles file to specify other files to use for Dialers configuration. Refer to Sysfiles File Format for BNU in AIX Files Reference for more information.
- 7. Decide how much access to your system you want to provide to each remote system you call and to each remote system that calls you. Set up appropriate entries for each system and each login name in the **Permissions** file. Refer to Permissions File Format for BNU in *AIX Files Reference* for details.
- 8. Issue the **uucheck** command to verify that everything is in place:

/usr/sbin/uucp/uucheck -v

The **uucheck** command verifies that the directories, programs, and support files are set up properly and that the **Permissions** file entries are consistent. If the **uucheck** command reports any errors, fix the errors.

9. If you wish, set up automatic monitoring of BNU operations and automatic polling of remote systems.

Setting Up Automatic Monitoring of BNU

Prerequisites

- Complete the steps in "Configuring BNU", on page 8-10.
- You must have root user authority to edit the /var/spool/cron/crontabs/uucp file.

Procedure

BNU uses the **cron** daemon to start BNU daemons and to monitor BNU activity. The **cron** daemon reads the /**var/spool/cron/crontabs/uucp** file for instructions about when to start BNU procedures.

- 1. Log in as a user with root user authority.
- 2. Using an ASCII text editor, edit the /var/spool/cron/crontabs/uucp file.
- 3. Uncomment the lines for the BNU maintenance procedures, **uudemon.admin** and **uudemon.cleanup**. You can change the times these procedures are run if your system needs maintenance at more or less frequent intervals. It is best, however, to run the **uudemon.admin** command at least once a day and the **uudemon.cleanup** command at least once a week.
- 4. You can use the crontabs/uucp file to schedule other BNU maintenance commands, such as the uulog, uuclean, or uucleanup commands. In addition, you can use the crontabs/uucp file to instruct the cron daemon to start the uucico, uuxqt, or uusched daemons at specific times.

Setting Up BNU Polling of Remote Systems

Prerequisites

- 1. Complete the steps in "Configuring BNU", on page 8-10.
- 2. You must have root authority to edit the /var/spool/cron/crontabs/uucp file and the /etc/uucp/Poll file.

Procedure

To enable BNU to poll remote systems for jobs, list the systems in the /etc/uucp/Poll file. In addition, run the uudemon.hour and uudemon.poll commands periodically.

- 1. Decide which remote systems to automatically poll. Decide how often you want to poll each one. Specify times for each system with the **Poll** file as seldom as once a day or as often as you wish.
- 2. Log in as a user with root authority.
- 3. Using an ASCII text editor or the **uucpadm** command, edit the **PoII** file. Add an entry for each system your system will poll.

Note: The systems listed in the **PolI** file must also be listed in the /etc/uucp/Systems file.

4. Using an ASCII text editor, edit the /var/spool/cron/crontabs/uucp file. Remove the comment characters (#) from the lines that run the uudemon.hour and uudemon.poll commands. You can change the times these commands are run. However, be sure to schedule the uudemon.poll command approximately five minutes before you schedule the uudemon.hour command.

BNU will now automatically poll the systems listed in the **Poll** file at the times you have specified.

Using the /etc/uucp/Systems File

The remote systems accessible with BNU commands are identified when the BNU program is installed. They are listed in the /etc/uucp/Systems files. The /etc/uucp/Systems file is the default Systems file. The system administrator can specify additional files in the /etc/uucp/Sysfiles file.

Each entry in a **Systems** file contains:

- Name of the remote system
- Times when users can connect to the remote system
- Type of link (direct line or modem)
- Speed of transmission over the link
- Information needed to log in to the remote system

Each entry in a **Systems** file represents one remote system. To establish communications, the remote system must be listed in the local **Systems** file. A **Systems** file must be present on every system that uses the BNU facility. Normally, only the root user can read the **Systems** files. Any user, however, can list the names of remote BNU systems using the **uuname** command.

Editing Devices Files for Hardwired Connections

Prerequisites

You must have root authority to edit the /etc/uucp/Devices file or another file specified in /etc/uucp/Sysfiles as a Devices file.

Procedure to Set Up a System Name Entry

To set up a hardwired connection specifying a port and a remote system, make an entry as follows:

- 1. Enter the name of the remote system to which you want to connect the local computer over the hardwired line in the *Type* field in the second line of the entry.
- 2. Enter the device name appropriate for the hardwired connection used at your site in the *Line* field in both lines of the entry.
- 3. Enter a (hyphen) for a placeholder in the *Line2* field in both lines of the entry.
- 4. Enter the transmission rate appropriate for the hardwired connection used at your site in the *Class* field in both lines of the entry.
- 5. Enter direct (all lowercase) in the Dialer-Token Pairs field in both lines of the entry.

Continue adding entries to the **Devices** file until you have listed each hardwired device connecting the local system to a remote system.

Procedure to Set Up a Direct Entry

To set up a hardwired connection between two systems that use a permanent asynchronous serial connection, make a one–line entry as follows:

- 1. Enter the name of the remote system in the first (*Type*) field.
- 2. Enter the name of the tty device in the second (Line) field.
- 3. Enter a (hyphen) for a placeholder in the third (Line2) field.
- 4. Enter the transmission rate appropriate for the hardwired connection used at your site in the fourth (*Class*) field.
- 5. Enter direct (all lowercase) in the fifth (Dialer-Token Pairs) field.

Continue adding entries to the **Devices** file until you have listed each hardwired device connecting the local system to a remote system.

Editing Devices File for Autodialer Connection

Prerequisites

You must have root authority to edit the /etc/uucp/Devices file or another file specified in /etc/uucp/Sysfiles as a Devices file.

Procedure

In telephone–connection entries, the *Type* field is specified as an automatic calling unit (ACU). Type ACU as the *Type* field entry in all remote connections established over a phone line. To set up **Device** file entries for autodialer connections, make a one–line entry for each modem:

- 1. Enter ACU in the first (*Type*) field.
- 2. The second (*Line*) field contains the name of the device that is attached to the modem. Enter the device name appropriate for your site.
- 3. Enter a (hyphen) as a placeholder in the third (*Line2*) field, unless the autodialer is a standard 801 dialer. If the autodialer is a standard 801 dialer, enter 801.

4. In the fourth (*Class*) field, enter the baud rate appropriate for your modem and line (this can be 300, 1200, 2400, or higher, depending on the modem) or the class of your modem (for example, D2400).

Note: If the modem can be used at more than one specific rate, make a separate entry in the **Devices** file for each rate. If the modem can be used at any rate, enter the word Any in the *Class* field.

5. Enter the name of the modem as the *Dialer* field entry in the fifth (*Dialer–Token Pair*) field. If you are planning to include complete phone numbers in the /etc/uucp/Systems file or another Systems file specifies in /etc/uucpSysfiles, leave the *Token* field blank. (A blank instructs the BNU program to use the default \D token.) If you are planning to use dialing–code abbreviations specified in the /etc/uucp/Dialcodes file, enter the token \T.

Continue adding entries to the **Devices** file until you have listed each connection between the local system and a remote system that uses a telephone line and a modem.

Editing Devices File for TCP/IP

Prerequisites

You must have root authority to edit the /etc/uucp/Devices file or another file specified in /etc/uucp/Sysfiles as a Devices file.

Procedure

If your site is using the TCP/IP system, include the relevant TCP/IP entry in the **Devices** file. To set up the file for use with the TCP/IP system, enter the following line in the **Devices** file:

TCP - - TCP

Maintaining BNU

BNU must be maintained to work properly on your system. To maintain BNU:

- Read and remove log files periodically.
- Use the uuq and uustat commands to check the BNU queues to ensure jobs are transferring to remote systems properly.
- Schedule automatic commands that poll remote systems for jobs, return unsent files to users, and send you periodic messages about BNU status.
- Periodically update the configuration files to reflect changes in your system.

In addition, occasionally check with administrators of remote systems to keep up with changes on their systems that may affect your configuration. For example, if the supervisor of system venus changes your system's password, you will need to put the new password in the /etc/uucp/Systems file (or the appropriate Systems file specified by /etc/uucp/Sysfiles) before your system can log in to system venus.

See BNU Files, Commands, and Directories Reference, on page 8-36 for a list of commands used to maintain BNU.

Working with BNU Log Files

BNU creates log files and error files to track its own activities. These files must be checked and removed periodically to keep them from filling the storage space on your system. BNU provides several commands for use in cleaning log files:

- uulog
- uuclean
- uucleanup
- uudemon.cleanu

Run these commands manually or use entries in the /var/spool/cron/crontabs/uucp file to run the commands by the cron daemon.

Log Files in the .Log and .Old Directories

BNU creates individual log files in the /var/spool/uucp/.Log directory. BNU creates these log files for each accessible remote system, using the uucp, uuto, or uux command. BNU places status information about each transaction in the appropriate log file each time someone on the system uses BNU. When more than one BNU process is running the system cannot access the log file. Instead, it places the status information in a separate file with a .LOG prefix.

The **uulog** command displays a summary of **uucp** or **uux** requests, by user or by system. The **uulog** command displays the files. However, you can also have BNU automatically combine the log files into a primary log file. This is called *compacting* the log files and can be done with the **uudemon.cleanu** command, usually run by the **cron** daemon.

The **cron** daemon runs the **uudemon.cleanu** command. The **uudemon.cleanu** command combines the **uucico** and **uuxqt** log files on the local system and stores them in the /var/spool/uucp/.Old directory. At the same time, the command removes old log files previously stored in the .Old directory. By default, the **uudemon.cleanu** command saves log files that are two days old.

If storage space is a problem, consider reducing the number of days that files are kept. To track BNU transactions over a longer period of time, consider increasing the number of days that files are kept. To change the default time for saving log files, modify the shell procedure for the **uudemon.cleanu** command. This script is stored in the /**usr/sbin/uucp** directory and can be modified with root authority.

Other BNU Log Files

BNU also collects information and stores it in the /var/spool/uucp/.Admin directory. This directory contains the errors, xferstats, Foreign, and audit files. These files must be checked and removed occasionally to save storage space. BNU creates each file when it is needed.

When another system contacts your system with the **uucico** daemon's debugging mode on, it invokes the **uucico** daemon on your system with debugging turned on. The debugging messages generated by the daemon on the local system are stored in the **audit** file. This file can get quite large. Check and remove the **audit** file often.

The **errors** file records errors encountered by the **uucico** daemon. Checking this file can help you correct problems such as incorrect permissions on BNU work files.

The **xferstats** file contains information about the status of every file transfer. Check and remove this file occasionally.

The **Foreign** file is important to the security of your system. Whenever an unknown system attempts to log in to the local system, BNU calls the **remote.unknown** shell procedure. This shell procedure logs the attempt in the **Foreign** file. The **Foreign** file contains the names of the systems that have attempted to call the local system and been refused. If a system has been attempting frequent calls, use this information when considering whether to allow that system access.

Systemwide Log Files used by BNU

Because many BNU processes need root authority to complete their tasks, BNU creates frequent entries in the /var/spool/sulog log file. Similarly, using the cron daemon to schedule BNU tasks creates multiple entries in the /var/spool/cron/log file. When using BNU, check and clean these files.

BNU Maintenance Commands

The Basic Networking Utilities contain several commands for monitoring BNU activities and cleaning BNU directories and files.

Cleanup Commands

BNU contains three commands that clean directories and remove files that have not been sent:

uuclean	Deletes all files older than a specified number of hours, from	n tha
uuciean	- Deletes all lifes older man a specified number of nours. Ifor	ппе

BNU administrative directories. Use the **uuclean** command to specify a directory to be cleaned or a type of file to be deleted. You can also instruct the command to notify the owners of the deleted files. The **uuclean** command is the Berkeley equivalent of

the uucleanup command.

uucleanup Performs functions similar to the **uuclean** command. However,

the **uucleanup** command checks the age of files based on *days* rather than hours. Use the **uucleanup** command to send a warning message to users whose files have not been transferred, notifying them that the files are still in the queue. The **uucleanup** command also removes files relating to a specified remote

system.

uudemon.cleanu A shell procedure that issues the **uulog** and **uucleanup**

commands to compress the BNU log files and remove log and work files over three days old. The **uudemon.cleanu** command is

run by the **cron** daemon.

Status-checking Commands

BNU also provides commands for checking the status of transfers and log files:

uuq Displays jobs currently in the BNU job queue. Use the **uuq**

command to display the status of a specified job or of all jobs. With root authority, you can use the **uuq** command to delete a job

from the queue.

uustat Provides information similar to that provided by the **uuq**

command, in a different format. Use the **uustat** to check the status of jobs and delete jobs you own. With root authority, you

can also delete jobs belonging to other users.

uulog Displays a summary of **uucp** or **uux** requests, by user or by

system. The uulog command displays the file names. See

"Working with BNU Log Files", on page 8-18.

uupoll Forces a poll of a remote system. This is helpful when work for

that system is waiting in the queue and needs to be transferred,

before the system is scheduled to be called automatically.

uusnap Displays a very brief summary of BNU status. For each remote

system, this command shows the number of files awaiting transfer. However, it does not show how long they have been waiting. The **uusnap** command is the Berkeley equivalent of the

uustat command.

Shell Procedures

BNU is delivered with two shell procedures used for maintenance:

uudemon.cleanu Discussed under "Cleanup Commands", on page 8-19

uudemon.admin Issues the **uustat** command The **uustat** command reports the

status of BNU jobs. It sends the results to the uucp login ID as mail. You can modify the **uudemon.admin** shell procedure to send the mail elsewhere, or use a mail program to reroute all mail

for the uucp login ID to the user responsible for BNU

administration.

These shell procedures are stored in the /usr/sbin/uucp directory. Copy the procedures and modify the copy, if you want to change what they do. Run the procedures from the command line or schedule them to be run by the **cron** daemon.

To automatically run the **uudemon.cleanu** and **uudemon.admin** commands, remove the comment characters (#) from the beginning of the relevant lines in the /var/spool/cron/crontabs/uucp file.

Monitoring a BNU Remote Connection

Prerequisites

- The BNU program must be installed on your system.
- A link (hardwired, modem, or TCP/IP) must be set up between your system and the remote system.
- The BNU configuration files, including the **Systems** file, **Permissions** file, **Devices** file, and **Dialers** file (and **Sysfiles** file, if applicable), must be set up for communications between your system and the remote system.

Note: You must have root user authority to modify the BNU configuration files.

Procedure

The **Uutry** command can help you monitor the **uucico** daemon process if users at your site report file—transfer problems.

1. Issue the **uustat** command to determine the status of all the transfer jobs in the current queue as follows:

```
uustat -q
```

The system displays a status report like the following:

```
venus 3C (2) 05/09-11:02 CAN'T ACCESS DEVICE hera 1C 05/09-11:12 SUCCESSFUL merlin 2C 5/09-10:54 NO DEVICES AVAILABLE
```

This report indicates that three command (**C**.*) files intended for remote system venus have been in the queue for two days. There could be several reasons for this delay. For example, perhaps system venus has been shut down for maintenance or the modem has been turned off.

2. Before you begin more extensive troubleshooting activities, issue the **Uutry** command as follows to determine whether your local system can contact system venus now:

```
/usr/sbin/uucp/Uutry -r venus
```

This command starts the **uucico** daemon with a moderate amount of debugging and the instruction to override the default retry time. The **Uutry** command directs the debugging output to a temporary file, /tmp/venus.

3. If your local system succeeds in establishing a connection to system venus, the debugging output contains a good deal of information. However, the final line in this script, which follows, is the most important:

```
Conversation Complete: Status SUCCEEDED
```

If the connection is successful, assume that the temporary file—transfer problems are now resolved. Issue the **uustat** command again to make certain that the files in the spooling directory have been transferred successfully to the remote system. If they have not, use the steps in "Monitoring a BNU File Transfer", on page 8-22 to check for file—transfer problems between your system and the remote system.

4. If your local system cannot contact the remote system, the debugging output generated by the **Uutry** command contains the following type of information (the exact form of the output may vary):

```
mchFind called (venus)
conn (venus)
getto ret -1
Call Failed: CAN'T ACCESS DEVICE
exit code 101
Conversation Complete: Status FAILED
```

First, check the physical connections between the local and remote systems. Make sure that the remote computer is turned on and all cables are properly connected, that the ports are enabled or disabled (as appropriate) on both systems, and that the modems (if applicable) are working.

If the physical connections are correct and secure, then verify all the relevant configuration files on both the local and remote systems, including the following:

- Make certain that the entries in the Devices, Systems, and Permissions files (and Sysfiles file, if applicable) in the /etc/uucp directory are correct on both systems.
- If you are using a modem, make sure that the /etc/uucp/Dialers file (or an alternate file specified in /etc/uucp/Sysfiles) contains the proper entry. If you are using dial—code abbreviations, be sure the abbreviations are defined in the /etc/uucp/Dialcodes file.

- If you are using a TCP/IP connection, make sure that the uucpd daemon can be run
 on the remote system and that the configuration files contain the correct TCP entries.
- 5. Once you have checked the physical connections and configuration files, issue the **Uutry** command again. If the debugging output still reports that the connection failed, you may need to confer with a member of your systems support team. Save the debugging output produced by the **Uutry** command. This may prove helpful in diagnosing the problem.

Monitoring a BNU File Transfer

Prerequisites

- 1. The BNU program must be installed on and configured for your system.
- 2. Establish a connection to a remote system using the steps given in "Monitoring a BNU Remote Connection", on page 8-20.

Monitoring a File Transfer

Use this procedure to monitor a file transfer to a remote system. Monitoring a file transfer is useful when file transfers to the remote system in question are failing for unknown reasons. The debugging information produced by the **uucico** daemon (called by the **Uutry** command) can help you find out what is working incorrectly.

The **Uutry** command enables you to monitor file transfers, as follows:

1. Prepare a file for transfer using the **uucp** command with the **-r** flag by entering:

```
uucp -r test1 venus!~/test2
```

The **-r** flag instructs the BNU program to place the test1 file in the queue but *not* to start the **uucico** daemon.

2. Issue the **Uutry** command with the **-r** flag to start the **uucico** daemon with debugging turned on by entering:

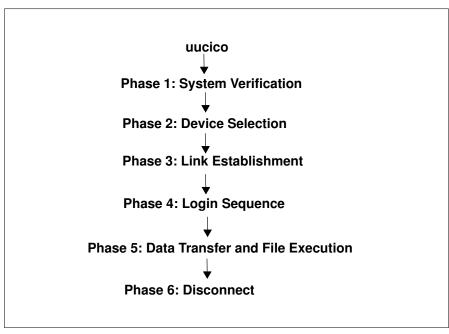
```
/usr/sbin/uucp/Uutry -r venus
```

This instructs the **uucico** daemon to contact remote system <code>venus</code> overriding the default retry time. The daemon contacts system <code>venus</code>, logs in, and transfers the file, while the **Uutry** command produces debugging output that enables you to monitor the **uucico** process. Press the Interrupt key sequence to stop the debugging output and return to the command prompt.

The **Uutry** command also stores the debugging output in the /**tmp**/SystemName file. If you break out of the debugging output before the connection is complete, you can page through the output file to see the outcome of the connection.

Debugging BNU Problems

BNU error messages can be linked to a specific phase in the conversation flow. Use the BNU Conversation Flow Diagram and the following error descriptions to help diagnose your BNU problems. Some of the following messages may not be sent from BNU, but are included in case another uucp version is in use.



BNU Conversation Flow Diagram

PHASE 1 Status Messages

Assert Error	The local system unit is having problems. Check the error report for possible causes by issuing the command errpt – a pg .
System not in Systems	If you supply a remote system name that is not found in the Systems files, this status message is created. BNU will terminate. Use the uuname command to check the system name again.
Wrong time to call	The Systems file has restrictions on times to allow outgoing calls. BNU will keep trying until the time is right. Check the Systems file.
Callback required	The network has restricted usage either for security or economic reasons, and access is denied at this time.
Cannot call No Call	These errors mean BNU recently tried to call the remote system and failed. It will not immediately try again. They can also be caused by an old system status file being retained thus keeping uucico from trying again.

PHASE 2 Status Messages

Dialer Script Failed No Device Available Can't Access Device Your **Dialers** file script did not complete successfully.

The modem or the outgoing phone line from your system is busy. Check for an error in the device entry of the **Systems** file. Also, check the **Devices** and **Dialers** files to be sure logical devices have physical devices associated with them. The file /etc/uucp/Sysfiles may be specifying an alternate **Systems**, **Devices**, or **Dialers** file that is not correctly configured. Is the device in use by some other program? Check the /etc/locks directory for lock on port. If a lock file exists (for example, **LCK..TTY0**), check to see if the process identified by the number in the lock file is still active. If not, you can remove it (for example, rm /etc/locks/LCK..TTY0). Also check the permissions on the port.

The

Dial Failed Failed (call to system) These errors appear when your system dials another successfully but the other system does not answer. It may also indicate a problem in the **Devices** files. Enter the command $\mathtt{uucico-r1-x6-s}$ SystemName. It could be that BNU is expecting some string that it is not receiving. Make the connection by hand to find out what needs to be incorporated into the **Systems** files entry to satisfy the request. Please keep "timing" in mind; perhaps some delays in the modem's dial string are needed. This could also mean that the port is busy, you dialed an incorrect number, or BNU lost ownership of the port.

OK Auto Dial These are informative messages only and do not indicate an error.

PHASE 3 Status Messages

Handshake Failed
(LCK)

The device is being used by someone else; the process could not create the **LCK** file. Sometimes **LCK** files must be manually removed by the administrator. After a number of retries, see your system administrator. See if another process has control of the port (for example, another instance of **uucico**).

Login Failed

The login failed due to a bad connection or possibly a slow machine.

Timeout

The remote system did not respond within a set period of time. This could also indicate a problem with the chat script.

Succeeded (Call to

The call was completed.

System)

BNU (continued)

These are informative messages only and do not indicate an error.

PHASE 4 Status Messages

Startup Failed Remote reject after login After login, **uucico** is started on the remote system. If there is a problem initiating a conversation between the two systems, these messages are created. You may have also logged into the incorrect BNU account or the initial handshake failed.

Wrong machine name

A machine was called incorrectly or the machine's name was changed.

Bad login/machine combination

Remote has a LCK file for me

OK Talking LOGIN:

PASSWORD:

The login to the remote system failed. The problem could be an incorrect phone number, an incorrect login or password, or an error in the chat script.

Both systems were simultaneously trying to call each other. The local request will fail temporarily.

These are informative messages only and do not indicate an error.

If the login or password prompt is in all capital letters, the modem may be in echo mode (E1 on Hayes compatibles). This causes the modem to echo back, or send, a RING to your system when an incoming call is received. The **getty** command receives the string and accordingly changes the login: or password: into all caps. Change the echo mode on the modem to **off** (use ATE0 for Hayes compatibles).

Note: Keep in mind that once this change is made, you should use $\mathtt{ATE1}$ in the chat script of your **Dialers** files, or you will not get the expected \mathtt{OK} back from the modem.

If the remote port is set for <code>delay</code> or <code>getty -r</code> and the chat script expects key input, then the ports set for <code>delay</code> are expecting one or more carriage returns before proceeding with the login. Try beginning the chat script on the dialing system with the following:

Interpreted, the above reads as follows: expect nothing, send return, delay, return, delay, return, delay, return.

PHASE 5 Status Messages

Alarm

Remote access to path/file denied copy (failed)

Bad read

Conversation failed

Requested Copy (succeeded)

uucico is having trouble with the connection. Either the connection is bad or "xon/xoff" is set to yes on the modem.

These messages indicate a permission problem; check file and path permissions.

The remote system ran out of space, most likely in the spool area, or **uucico** could not read or write to device.

The modem's carrier detect was lost. Possibly the modem was turned off, the cable is loose or disconnected, or the remote system crashed or is shut down. Telephone disconnection can also cause this error.

These are informative messages only and do not indicate an error.

PHASE 6 Status Messages

OK (Conversation Complete)

Conversation succeeded

The remote system can deny the hangup request and reverse the roles (meaning the remote system has work for the local system to do). Once the two **uucico**s agree that no more work exists, they hang up.

This is an informative message only and does not indicate an error.

Debugging BNU Login Failures Using the uucico Daemon

Prerequisites

- BNU must be installed on your system.
- A link (hardwired, modem, or TCP/IP) must be set up between your system and the remote system.
- The BNU configuration files, including the Sysfiles file (if applicable), the Systems file, Permissions file, Devices file, and Dialers file, must be set up for communications between your system and the remote system.

Note: You must have root user authority to modify the BNU configuration files.

• You must have root user authority to invoke the **uucico** daemon in debugging mode.

Procedure

1. To produce debugging information about a local—to—remote system connection that is not working, start the **uucico** daemon with the **–x** flag as follows:

```
/usr/sbin/uucp/uucico -r 1 -s venus -x 9
```

where -r 1 specifies the master, or caller mode; -s venus, the name of the remote system to which you are trying to connect; and -x 9, the debug level that produces the most detailed debugging information.

If the expect–send sequence entry in a Systems file in the format of /etc/uucp/Systems is:

```
venus Any venus 1200 - "" \n in:--in: uucp1 word:
mirror
```

the uucico daemon connects the local system to the remote system venus. The debugging output is similar to:

```
expect: ""
got it
sendthem (^J^M)
expect (in:)^
M^Jlogin:got it
sendthem (uucp1^M)
expect (word:)^
M^JPassword:got it
sendthem (mirror^M)
imsg >^M^JPShere^@Login Successful: System=venus
```

where:

expect: ""	Specifies that the local system will not wait for any information from the remote system.		
got it	Acknowledges that the message has been received.		
sendthem (^J^M)	Specifies that the local system will send the remote system a carriage return and a new line.		
expect (in:)	Specifies that the local system will expect to receive the remote system login prompt, which ends in the $in:$ character string.		
^M^Jlogin:got it	Confirms that the local system will receive the remote login prompt.		
sendthem (uucp1^M)	Specifies that the local system will send the ${\tt uucp1}$ login ID to the remote system.		

expect (word:) Specifies that the local system will expect to receive the

remote system password prompt, which ends in the

word: character string.

^M^JPassword:got it Confirms the local system will receive the remote

password prompt.

sendthem (mirror^M) Specifies that the local system will send the password for

the uucp1 login ID to the remote system.

imsg >^M^J^PShere^@Login Successful: System=venus

Confirms the local system is successfully logged in to remote system venus.

Note:

1. The expect—send debugging output produced by the uucico command can come either from information in the /etc/uucp/Dialers file or from information in the /etc/uucp/Systems file. Information about communication with the modem comes from the Dialers file, while information about communication with the remote system comes from the Systems file. (Note that /etc/uucp/Systems and /etc/uucp/Dialers are default BNU configuration files. Other files can be specified in /etc/uucp/Sysfiles to serve the same role.)

2. To set up a connection with a remote system, you must be familiar with the login sequence of that system.

Contacting Connected UNIX Systems Using the tip Command

Use the **tip** command to contact any connected system running the UNIX operating system. The **tip** command is installed with the Basic Networking Utilities (BNU) and can use the same asynchronous connections used by BNU.

The **tip** command uses variables and escape signals, as well as flags, to control its operations. The flags can be entered at the command line. The escape signals can be used over a connection with a remote system to start and stop file transfers, change the direction of a file transfer, and exit to a subshell.

tip Command Variables

The **tip** command variables define settings such as the end–of–line character, the break signal, and the mode of file transfers. Variable settings can be initialized at run time using a **.tiprc** file. Variable settings can also be changed during execution using the ~**s** escape signal. Some variables, such as the end–of–line character, can be set for an individual system in that system's entry in the **remote** file.

The **tip** command reads three files, the **phones** file, **remote** file, and **.tiprc** file, to determine initial settings for its variables. The **.tiprc** file must always be in the user's home directory. The names and locations of the **remote** and **phones** files can vary. The names of the **remote** file and the **phones** file can be determined by environment variables:

PHONES Specifies the name of the user's phone file. The file can have any

valid file name and must be set up in the format of the file /usr/lib/phones-file. The default file is etc/phones. If a file is specified with the PHONES variable, it is used in place of (not in

addition to) the /etc/phones file.

REMOTE Specifies the name of the user's remote system definition file. The

file can have any valid file name and must be set up in the format of the /usr/lib/remote-file file. The default file is /etc/remote. If a file is specified with the **REMOTE** variable, it is used in place of

(not in addition to) the /etc/remote file.

To use an environment variable, set it before starting the **tip** command. As an alternative, the names of the **phones** and **remote** files can be determined using the **tip** command **phones** variable and **remote** variable, respectively, in the **.tiprc** file.

Note: The **tip** command reads only the *last* **remote** or **phones** file specified. Thus, if you specify a **remote** or **phones** file with a variable, the new file is used in place of (not in addition to) any previous files you specified.

The tip command uses variable settings in the following order:

- The command checks the settings of the PHONES and REMOTE environment variables for the files to use as the phones and remote files.
- The command reads the .tiprc file and sets all variables accordingly. If the phones or remote variable is set in the .tiprc file, this setting overrides the environment variable setting.
- When a connection to a remote system is initiated, the command reads the remote file entry for that system. The settings in the remote file entry override settings made in the .tiprc file.
- 4. If the *-BaudRate* flag is used with the **tip** command, the specified rate overrides all previous baud rate settings.
- 5. A setting made with the ~s escape signal overrides all previous settings of a variable.

Note: Any **tip** user can create a **.tiprc** file and use this file to specify initial settings for **tip** variables. The **.tiprc** file must be placed in the user's **\$HOME** directory.

tip Command Configuration Files

Before the **tip** command can connect to a remote system, the /**etc/remote** and /**etc/phones** files must be established.

/etc/remote Defines attributes of remote systems such as the port and type of

device to use to reach the system, as well as the signals to use to

indicate the beginnings and endings of transmissions.

/etc/phones Lists telephone numbers used to contact remote systems over a

modem line.

To establish one of these files, copy a sample file to the correct name and modify it to suit the needs of your site. Sample **remote** and **phones** files are delivered with the bos.net.uucp package. The sample **remote** file is named /usr/lib/remote-file. The sample **phones** file is named /usr/lib/phones-file.

Note: You must have root authority to create files in the /usr/lib directory.

A tip user can also create customized remote and phones files. An individual remote file must be in the format of the /usr/lib/remote-file file and specified with the remote variable or the REMOTE environment variable. An individual phones file must be in the format of the /usr/lib/phones-file file and specified with the phones variable or the PHONES environment variable. If an individual phones or remote file is specified with one of the variables, that file is read *in place of* (not in addition to) the /etc/phones or /etc/remote file.

Users of **tip** can use combinations of individual **phones** and **remote** files. For example, a user could read the default **remote** file, /**usr/lib/remote-file**, but use an individual **phones** file named with the **phones** variable.

BNU Configuration Files

Basic Network Utilities (BNU) uses the following configuration files:

/etc/uucp Contains all the configuration files for BNU.

/var/spool/uucppublic Contains files that have been transferred.

/etc/uucp/Systems Contains a list of systems to which the uucico program can

connect.

/etc/uucp/Devices Defines the device type, location, speed, and other basic

communication parameters for many system dial—out programs. Only dial—out connections use this file.

/etc/uucp/Permissions Creates security control, with limitations, over machines

attempting to communication with your machine.

/etc/uucp/Dialers Specifies the dialer types. Each dialer uses a specific

command set when attempting to dial the modem. The most common dialer types are hayes, direct, and TCP

(Transmission Control Protocol).

/etc/uucp/Dialcodes Makes standardized names for certain parts of a phone

number. For example, if you frequently make calls to a certain area code in San Francisco, you could create the

following entry: SF09, 1415.

/etc/uucp/Sysfiles Enables a BNU administrator to specify files to fill the role of

BNU configuration files other than /etc/uucp/Systems, /etc/uucp/Devices, and /etc/uucp/Dialers. Distinctions can be made between what files are used for uucico traffic

versus cu-related (cu, ct, slattach) activity.

/usr/sbin/uucp/remote.unk

nown

Defines a shell script. It is run by the BNU program when a remote computer that is not listed in the local permissions

file attempts to communicate with that local system.

/etc/uucp/Poll Schedules polling of passing systems. Its format is similar

to the crontab file. Poll format is SiteName, a tab, and the

hours to poll (0-23), separated by spaces.

Correlation of Files

Systems file: SystemName Any v32ibm 9600 555-1111

Devices file: v32ibm tty0 - Any ibm \D

Dialers file: $ibm = -, \#'' \setminus d$

ATSFI\r\c#OK#AFE1SD3L2MICOSCI\r\c#OK...

BNU Configuration for a TCP/IP Connection Example

The following files are set up for a Transmission Control Protocol/Internet Protocol (TCP/IP) connection between systems <code>zeus</code> and <code>hera</code>, where <code>zeus</code> is considered the local system and <code>hera</code> the remote system.

Entries in the Local System's Files

Files entries on local system venus include the following:

Systems File

The **Systems** file on system zeus should contain the following entry to allow zeus to contact system hera:

hera Any TCP, t - - in:--in: uzeus word: birthday

This specifies that system zeus can call system hera at any time, using the t protocol for communications with system hera. System zeus logs in to system hera as uzeus with the password birthday.

Note: The **t** protocol supports the **tcp** protocol. Therefore, always use the **t** protocol for BNU communications over TCP/IP connections. However, the **t** protocol cannot be used when the *Type* field is <code>ACU</code> (automatic calling unit) or when a modem connection is being used.

BNU uses the *Type* and *Class* fields in the **Systems** file to find the appropriate device for the connection. Accordingly, it checks the **Devices** file for an entry of type TCP.

Devices File

A **Devices** file used by **uucico** on system zeus should contain the following entry for TCP/IP connections:

```
TCP - - TCP
```

Because the device type is TCP, there are no *Class*, *Line*, or *Line2* entries. The *Dialer* is also specified as TCP. Accordingly, BNU looks in the **Dialers** files for a TCP entry.

Dialers File

The **Dialers** file used by **uucico** on system zeus should contain a TCP/IP entry as follows:

```
TCF
```

This entry specifies that no dialer configuration is required.

Note: Dialer configuration is never required over a TCP/IP connection.

Permissions File

The **Permissions** file on system zeus contains the following entry specifying system hera's access to system zeus:

```
LOGNAME=uhera SENDFILES=yes REQUEST=yes \
MACHINE=zeus:hera VALIDATE=uhera /
READ=/var/spool/uucppublic:/home/hera \
WRITE=/var/spool/uucppublic:/home/hera COMMANDS=ALL
```

This combined LOGNAME and MACHINE entry provides the following permissions to system hera on system zeus:

- System hera can request and send files regardless of who initiated the call.
- System hera can read and write to the public directory and the /home/hera directory on system zeus.
- System hera can execute all commands on system zeus.
- System hera must log in to system zeus as user uhera and cannot use any other login ID for BNU transactions.

Note: Because the permissions are the same regardless of which system initiates the call, the preceding LOGNAME and MACHINE entries are combined. Separately, they are:

```
LOGNAME=uhera VALIDATE=hera SENDFILES=yes REQUEST=yes& \
READ=/var/spool/uucppublic:/home/hera \
WRITE=/var/spool/uucppublic:/home/hera

MACHINE=zeus:hera REQUEST=yes COMMANDS=ALL\
READ=/var/spool/uucppublic:/home/hera \
WRITE=/var/spool/uucppublic:/home/hera
```

Entries in the Remote System's Files

Files containing telephone connection entries on remote system merlin include the following:

Systems File

A **Systems** file on system hera should contain the following entry to allow hera to contact system zeus:

```
zeus Any TCP, t - - ogin:--ogin: uhera ord: lightning
```

This specifies that system hera can call system zeus at any time, using the t protocol for communications with system zeus. System hera logs in to system zeus as user uhera with the password lightning. Again, BNU next checks the **Devices** files for an entry of type TCP.

Note: The **t** protocol supports the **tcp** protocol. Therefore, always use the **t** protocol for BNU communications over TCP/IP connections. However, the **t** protocol cannot be used when the *Type* field is ACU or when a modem connection is being used.

Devices File

The **Devices** file used by **uucico** on system hera should contain the following entry for TCP/IP connections:

```
TCP - - TCP
```

Because the device type is TCP, there are no *Class*, *Line*, or *Line2* entries. The *Dialer* is also specified as TCP. Accordingly, BNU looks in the **Dialers** files for a TCP entry.

Dialers File

The **Dialers** file used by **uucico** on system hera should contain a TCP/IP entry as follows:

TCP

This entry specifies that no dialer configuration is required.

Note: Dialer configuration is never required over a TCP/IP connection.

Permissions File

The **Permissions** file on system hera contains the following entry specifying system zeus's access to system hera:

```
LOGNAME=uzeus SENDFILES=yes REQUEST=yes \
MACHINE=hera:zeus VALIDATE=zeus COMMANDS=rmail:who:uucp
```

This combined LOGNAME and MACHINE entry provides the following permissions to system zeus on system hera:

- System zeus can request and send files regardless of who initiated the call.
- System zeus can read and write only to the public directory (the default).
- System zeus can run only the rmail, who, and uucp commands.
- System zeus must log in to system hera as user uzeus and cannot use any other login ID for BNU transactions.

Note: Separately, the LOGNAME and MACHINE entries are:

```
LOGNAME=uzeus VALIDATE=zeus SENDFILES=yes REQUEST=yes MACHINE=hera:zeus COMMANDS=rmail:who:uucp REQUEST=yes
```

BNU Configuration for a Telephone Connection Example

The following sample files are set up to connect systems venus and merlin over a telephone line using modems. System venus is considered the local system and system venus in the remote system.

On both systems, the device tty1 is hooked to a Hayes modem at 1200 baud. The login ID used for system venus to log into system merlin is uvenus, and the associated password is mirror. The login ID for system merlin to log into system venus is umerlin, and the associated password is oaktree. The phone number for the modem attached to venus is 9=3251436; the number of the merlin modem is 9=4458784. Both computers include partial phone numbers in their **Systems** files and dial—codes in their **Dialcodes** files.

Entries on the Local System

Files containing telephone connection entries on local system venus include the following:

Systems File

The **Systems** file on venus should contain the following entry for merlin, including a phone number and a dialing prefix:

```
merlin Any ACU 1200 local8784 "" in:--in: uvenus word: mirror
```

System venus can call system merlin at any time, using an ACU device at 1200 baud and logging in as uvenus with the password mirror. The telephone number is expanded based on the code local in the **Dialcodes** file, and the device to be used is determined based on the *Type* and *Class* entries. Accordingly, BNU checks the **Devices** files for a device of type ACU and class 1200.

Dialcodes File

The **Dialcodes** file on system venus contains the following dial—code prefix for use with the number in the **Systems** file:

```
local 9=445
```

Given this code, the telephone number for system merlin in the **Systems** file is expanded to 9=4458784.

Devices File

The **Devices** file on system <code>venus</code> should contain the following entry for the connection to system <code>merlin</code>:

```
ACU tty1 - 1200 hayes \T
```

The port to be used is tty1, and the *Dialer* entry in the *Dialer–Token Pairs* field is hayes. The *Token* entry, T , indicates that the telephone number is to be expanded using a code from the **Dialcodes** file. BNU checks the **Dialers** files for a hayes dialer type.

Dialers File

A **Dialers** file used by **uucico** on system venus should contain the following entry for the hayes modem:

```
hayes =,-, "" \dAT\r\c OK \pATDT\T\r\c CONNECT
```

Note: The expect-send characters are defined in the **Dialers** file format.

Permissions File

The **Permissions** file on system <code>venus</code> contains the following entries specifying the ways in which system <code>merlin</code> can conduct **uucico** and **uuxqt** transactions with system <code>venus</code>:

```
LOGNAME=umerlin REQUEST=yes SENDFILES=yes \
READ=/var/spool/uucppublic:/home/merlin \
WRITE=/var/spool/uucppublic:/home/merlin
MACHINE=venus:merlin VALIDATE=umerlin REQUEST=yes SENDFILES=yes \
COMMANDS=ALL \
READ=/var/spool/uucppublic:/home/merlin \
WRITE=/var/spool/uucppublic:/home/merlin
```

System merlin logs in to system venus as umerlin, which is a unique login for system merlin. It can request and send files regardless of who initiated the call. Also, system merlin can read and write to the /var/spool/uucppublic directory and the /home/merlin directory on system venus. It can issue all commands in the default command set on system venus.

Entries on the Remote System

Files containing telephone connection entries on remote system merlin include the following:

Systems File

A **Systems** file on merlin should contain the following entry for venus, including a phone number and a dialing prefix:

```
venus Any ACU 1200 intown4362 "" in:--in: umerlin word: oaktree
```

System merlin can call system venus at any time, using an ACU device at 1200 baud and logging in as user umerlin with the password oaktree. The telephone number is expanded based on the code intown in the **Dialcodes** file, and the device to be used is determined based on the *Type* and *Class* entries. Accordingly, BNU checks the **Devices** file(s) for a device of type ACU and class 1200.

Dialcodes File

The **Dialcodes** file on system merlin contains the following dial—code prefix for use with the number in the **Systems** file:

```
intown 9=325
```

Therefore, the expanded telephone number to reach system venus is 9=3254362.

Devices File

A **Devices** file on system merlin should contain the following entry for the connection to venus:

```
ACU tty1 - 1200 hayes \T
```

The ACU is attached to port tty1, and the dialer is hayes. The telephone number is expanded with information from the **Dialcodes** file. BNU checks the **Dialers** files for an entry for a hayes modem.

Dialers File

A **Dialers** file used by **uucico** on system merlin should contain the following entry for its modem:

```
hayes =,-, "" \dAT\r\c OK \pATDT\T\r\c CONNECT
```

Permissions File

The **Permissions** file on system merlin contains the following entries specifying system venus's access to merlin:

```
LOGNAME=uvenus SENDFILES=call REQUEST=no \
WRITE=/var/spool/uucppublic:/home/venus \
READ=/var/spool/uucppublic:/home/venus
MACHINE=merlin:venus VALIDATE=uvenus \
READ=/ WRITE=/ COMMANDS=ALL REQUEST=yes \
NOREAD=/etc/uucp:/usr/etc/secure \
NOWRITE=/etc/uucp:/usr/etc/secure
```

BNU Configuration for a Direct Connection Example

The following files are set up for a hardwired connection between systems <code>zeus</code> and <code>hera</code>, where <code>zeus</code> is considered the local system and <code>hera</code> the remote system. The hardwired device on system <code>zeus</code> is <code>tty5</code>; on system <code>hera</code> it is <code>tty1</code>. The speed of the connection is <code>1200</code> bps. The login ID for system <code>zeus</code> on system <code>hera</code> is <code>uzeus</code>, and the associated password is <code>thunder</code>. The login ID for system <code>hera</code> on system <code>zeus</code> is <code>uhera</code>, and the associated password is <code>portent</code>.

Entries in the Local System's Files

Files containing telephone connection entries on local system venus include the following:

Systems File

A Systems file on zeus should contain the following entry for the remote system hera:

```
hera Any hera 1200 - "" \r\d\r\d\r in:--in: uzeus word: thunder
```

This entry specifies that system hera can log in to system zeus at any time, using a direct connection specified in the **Devices** files. To find the entry in the **Devices** files, BNU uses the third and fourth fields of the **Systems** entry. Thus, BNU looks for an entry in the **Devices** files with a *Type* of hera and a *Class* of 1200. System zeus logs in to system hera as user uzeus with the password thunder.

Devices File

A **Devices** file on zeus should contain the following entry in order to connect to the remote system hera:

```
hera tty5 - 1200 direct
```

This entry specifies that system <code>zeus</code> uses the device <code>tty5</code> at 1200 bps to communicate with system <code>hera</code>. Note that the <code>Dialer</code> in both <code>Dialer-Token Pairs</code> fields is <code>direct</code>. When connecting to system <code>hera</code>, BNU checks the <code>Dialers</code> file for a <code>direct</code> entry.

Dialers File

A Dialers file on system zeus must contain the following entry for direct connections:

```
direct
```

This specifies that no handshaking is required on the direct connection.

Permissions File

The **Permissions** file on the local system zeus contains the following entry specifying the ways in which the remote system hera can conduct **uucico** and **uuxqt** transactions with zeus:

```
LOGNAME=uhera MACHINE=hera VALIDATE=uhera REQUEST=yes \
SENDFILES=yes MACHINE=hera READ=/ WRITE=/ COMMANDS=ALL
```

This entry specifies that system hera logs in as uhera. Since the VALIDATE=uhera option is included, system hera cannot log in to system zeus with any other login ID, nor can any other remote system use the uhera ID. System hera can read and write to any directory on system zeus, and can send and request files regardless of who initiated the call. System hera can also initiate any commands on system zeus.

Note: Since the permissions that are granted are the same regardless of which system initiated the connection, the LOGNAME and MACHINE entries have been combined. Separately, they are:

```
LOGNAME=uhera REQUEST=yes SENDFILES=yes READ=/ WRITE=/ MACHINE=zeus:hera VALIDATE=uhera READ=/ WRITE=/ REQUEST=yes \ COMMANDS=ALL
```

Attention: Providing the permissions in the preceding example is equivalent to giving any user on the remote system a login ID on the local system. Such liberal permissions can jeopardize your security and should usually be given only to well–trusted remote systems at the same site.

Entries in the Remote System's Files

Files containing telephone connection entries on remote system merlin include the following:

Systems File

A Systems file on system hera must contain the following entry for zeus:

```
zeus Any zeus 1200 - "" \r\d\r\d\r in:--in: uhera word: portent
```

This entry specifies that system hera can log in to system zeus at any time, using a direct connection specified in the **Devices** files. To find the entry in the **Devices** files, BNU uses the third and fourth fields of the **Systems** entry. Thus BNU looks for an entry in the **Devices** files with a *Type* of zeus and a *Class* of 1200. System hera logs in to system zeus as user uhera with the password portent.

Devices File

A **Devices** file on system hera must contain the following entry for communications with zeus:

```
zeus ttyl - 1200 direct
```

This entry specifies that system hera uses the device tty1 at 1200 bps to communicate with system zeus. Since the *Dialer* is specified as direct, BNU checks the **Dialers** files for a direct entry.

Dialers File

A **Dialers** file on system hera must contain the following entry for direct connections:

```
direct
```

This specifies that no dialer configuration is required on the direct connection.

Permissions File

The **Permissions** file on system hera contains the following entries specifying the ways zeus can conduct **uucico** and **uuxqt** transactions with hera:

```
LOGNAME=uzeus REQUEST=yes SENDFILES=yes READ=/ WRITE=/ MACHINE=hera:zeus VALIDATE=uzeus REQUEST=yes COMMANDS=ALL READ=/\WRITE=/
```

These entries specify that system <code>zeus</code> logs in to system <code>hera</code> as <code>uzeus</code>. Because the <code>VALIDATE=uzeus</code> option is included, system <code>zeus</code> cannot log in to system <code>hera</code> with any other login ID, nor can any other remote system use the <code>uzeus</code> ID. System <code>zeus</code> can read and write to any directory on system <code>hera</code>, and can send and request files regardless of who initiated the call. System <code>zeus</code> can also initiate any commands on system <code>hera</code>.

Attention: Providing the permissions in the preceding example is equivalent to giving any user on the remote system a login ID on the local system. Such liberal permissions can jeopardize your security and normally should be given only to remote systems at the same site.

BNU Files, Commands, and Directories Reference

BNU Directories

/etc/uucp Contains all Basic Network Utilities (BNU) configuration

files.

/etc/locks Contains lock files for system devices. Used by other

subsystems in addition to BNU.

/var/spool/uucppublic Contains files that have been transferred by BNU.

/var/spool/uucp Contains BNU administrative files.

/var/spool/uucp/.Workspace Holds temporary files that the file transport programs

use internally.

/var/spool/uucp/.Xqtdir Contains execute files with lists of commands that

remote systems can run.

/var/spool/uucp/SystemName Contains files used by file transport programs.

BNU Files

/etc/uucp/Systems A list of systems to which uucico can connect.

/etc/uucp/Devices Defines basic communication parameters for

dial-out connections.

/etc/uucp/Permissions Defines permissions for remote machines

contacting the local machine through BNU.

Maxuuscheds Limits simultaneous scheduled jobs.

Maxuuxqts Limits simultaneous remote command executions.

/etc/uucp/Dialers Specifies dialer and modem type.

/etc/uucp/Dialcodes Contains the initial digits of telephone numbers

used to establish remote connections over a phone

line.

/usr/sbin/uucp/remote.unknown A shell script executed when an unknown remote

computer attempts to communicate.

/usr/sbin/uucp/Sysfiles Assigns alternate or additional Systems, Devices,

and Dialers files.

/etc/uucp/Poll Determines when a remote system is called.

uudemon.adminSends a BNU status report to a specified login ID.uudemon.cleanuCleans BNU spooling directories at prescheduled

times.

uudemon.hour Initiates file transport calls to remote systems.

uudemon.poll Polls remote systems listed in the /etc/uucp/Poll

ile.

/var/spool/uucp/auditContains audit information from BNU activities./var/spool/uucp/ForeignContains error information about BNU activities./var/spool/uucp/errorsContains error information about BNU activities.

/var/spool/uucp/xferstats Contains statistical information about BNU

activities.

/var/spool/uucp/Corrupt Contains copies of files that cannot be processed

by the BNU program.

/var/spool/uucp/.Log Contains log files from current BNU transactions.

/var/spool/uucp/.Old Contains log files from old BNU transactions.

/var/spool/uucp/.Status Stores the last time the uucico daemon tried to

contact remote systems.

/var/spool/uucp/SystemName/C.* These files are the commands allowed when

connected to SystemName.

/var/spool/uucp/SystemName/D.* These files are data files associated with

SystemName.

/var/spool/uucp/SystemName/X.* Executable files on SystemName.

/var/spool/uucp/SystemName/TM.* Temporary files used when connected to

SystemName.

BNU Commands

ct Connects to another system over a telephone line.

cu Connects to another system.

tip A variation of **cu** that requires special configuration.

uucp Copies files from one system to another running BNU or a version of the

UNIX-to-UNIX Copy Program (UUCP).

uudecode Reconstructs a binary file encoded with **uuencode**.

uuencode Encodes a binary file into ASCII form for transmission using BNU.

uuname Provides information about accessible systems.

uupoll Forces a call to a remote system.uuq Displays the BNU job queue.

uusend Sends a file to a remote host running BNU or UUCP.

uusnap Displays a brief summary of BNU's status.uustat Reports the status of BNU operations.

uuto Copies files to another system using BNU or UUCP.

uux Runs a command on a remote system.

uucheck Checks the /etc/uucp/Permissions file for correct configuration.uuname Shows the names of all systems that can be reached through BNU.

uucleanuucleanupCleans BNU spooling directories.

uukick Contacts a remote system with debugging enabled.

uulog Displays BNU log files.

uutry Contacts a remote system with debugging enabled; allows override of retry

time.

uucpadm Administers the BNU system.

uupick
 uucp
 Login ID with full administrative authority over the BNU subsystem.
 Uutry
 Contacts a remote system with debugging turned on, and saves the

debugging output in a file.

BNU Daemons

uucico Contacts remote systems and transfers files.

uucpd Allows BNU to run on top of Transmission Control Protocol/Internet

Protocol (TCP/IP).

uusched Schedules BNU jobs.

uuxqt Runs command requests from remote systems.

Chapter 9. Network Management

The Network Management facility provides comprehensive management of system networks through the use of Simple Network Management Protocol (SNMP) enabling network hosts to exchange management information. SNMP is an internetworking protocol designed for use with TCP/IP—based internets. The following sections provide information to assist the network manager in understanding and working with SNMP:

- SNMP for Network Management, on page 9-2
- SNMP Access Polices, on page 9-3
- SNMP Daemon, on page 9-4
- Configuring the SNMP Daemon, on page 9-5
- SNMP Daemon Processing, on page 9-6
- SNMP Daemon Support of EGP Family of MIB Variables, on page 9-11
- SNMP Daemon RFC Conformance, on page 9-26
- SNMP Daemon Implementation Restrictions, on page 9-27
- SNMP Daemon Logging Facility, on page 9-28
- Problem Determination for the SNMP Daemon, on page 9-31

You may also want to consult the information in "SNMP Overview for Programmers" in *AIX Communications Programming Concepts*.

SNMP for Network Management

SNMP network management is based on the familiar client/server model that is widely used in TCP/IP—based network applications. Each host that is to be managed runs a process called an *agent*. The agent is a server process that maintains the Management Information Base (MIB) database for the host. Hosts that are involved in network management decision—making may run a process called a manager. A *manager* is a client application that generates requests for MIB information and processes responses. In addition, a manager may send requests to agent servers to modify MIB information.

For more information on SNMP, you may want to read "Simple Network Management Protocol (SNMP)" in *AIX Communications Programming Concepts*. The following RFCs may also be of use:

RFC 1155	Structure and Identification of Management Information (SMI) for TCP/IP-based internets.
RFC 1157	Simple Network Management Protocol (SNMP).
RFC 1213	Management Information Base (MIB) for Network Management of TCP/IP-based internets.
RFC 1227	Simple Network Management Protocol (SNMP) single multiplexer (SMUX) protocol and Management Information Base (MIB).
RFC 1228	Simple Network Management Protocol–Distributed Program Interface (SNMP–DPI).
RFC 1229	Extensions to the generic interface Management Information Base (MIB).
RFC 1231	IEEE 802.5 token-ring Management Information Base (MIB).
RFC 1398	Definitions of Managed objects for the Ethernet. For example, interface type.
RFC 1512	Fiber Distributed Data Interface (FDDI) Management Information Base (MIB).

SNMP Access Policies

The **snmpd** agent uses a simple authentication scheme to determine which Simple Network Management Protocol (SNMP) manager stations can access its Management Information Base (MIB) variables. This authentication scheme involves the specification of SNMP access policies. An SNMP access policy is an administrative relationship involving an association among an SNMP community, an access mode, and an MIB view.

An *SNMP community* is a group of one or more hosts and a community name. A community name is a string of octets that an SNMP manager must embed in an SNMP request packet for authentication purposes.

The access mode specifies the access the hosts in the community are allowed with respect to retrieving and modifying the MIB variables from a specific SNMP agent. The access mode must be one of: none, read—only, read—write, or write—only.

A *MIB view* defines one or more MIB subtrees that a specific SNMP community can access. The MIB view can be the entire MIB tree or a limited subset of the entire MIB tree.

When the SNMP agent receives a request, the agent verifies the community name with the requesting host's Internet Protocol (IP) address to determine if the requesting host is a member of the SNMP community identified by the community name. If the requesting host is indeed a member of the SNMP community, the SNMP agent then determines if the requesting host is allowed the specified access for the specified MIB variables as defined in the access policy associated with that community. If all criteria are met, the SNMP agent will attempt to honor the request. Otherwise, the SNMP agent will generate an *authenticationFailure* trap or return the appropriate error message to the requesting host.

The SNMP access policies for the **snmpd** agent are user–configurable and are specified in the /**etc/snmpd.conf** file. To configure the SNMP access policies for the **snmpd** agent, see the /**etc/snmpd.conf** file.

SNMP Daemon

The Simple Network Management Protocol (SNMP) daemon is a background server process that can be run on any Transmission Control Protocol/Internet Protocol (TCP/IP) workstation host. The daemon, acting as SNMP agent, receives, authenticates, and processes SNMP requests from manager applications. Read "Simple Network Management Protocol," "How a Manager Functions," and "How an Agent Functions" in *AIX Communications Programming Concepts* for more detailed information on agent and manager functions.

Note: The terms SNMP daemon, SNMP agent, and agent are used interchangeably.

The **snmpd** daemon requires the loopback TCP/IP interface to be active for minimal configuration. Enter the following command before starting TCP/IP:

ifconfig lo0 loopback up

Configuring the SNMP Daemon

The Simple Network Management Protocol (SNMP) daemon will attempt to bind sockets to certain well–known User Datagram Protocol (UDP) and Transmission Control Protocol (TCP) ports, which must be defined in the /etc/services file as follows:

```
snmp 161/udp
snmp-trap 162/udp
smux 199/tcp
```

The snmp service must be assigned port 161, as required by RFC 1157. The /etc/services file assigns ports 161, 162, and 199 to these services. If the /etc/services file is being serviced off another machine, these assigned ports must be made available in the served /etc/services file on the server before the SNMP daemon can run.

The SNMP daemon reads a configuration file, /etc/snmpd.conf, on startup and when a refresh command (if the snmpd daemon is invoked under System Resource Controller control) or kill -1 signal is issued. This configuration file specifies the community names and associated access privileges and views, hosts for trap notification, logging attributes, snmpd—specific parameter configurations, and single multiplexer (SMUX) configurations for the SNMP daemon. See the /etc/snmpd.conf file for more information.

SNMP Daemon Processing

The Simple Network Management Protocol (SNMP) daemon processes SNMP requests from manager applications. Read "Simple Network Management Protocol (SNMP)," "How a Manager Functions," and "How an Agent Functions" in *AIX Communications Programming Concepts* for more detailed information on agent and manager functions.

Message Processing and Authentication

All requests, traps, and responses are transmitted in the form of ASN.1–encoded messages. A message, as defined by RFC 1157, has the following structure:

Version Community PDU

where *Version* is the SNMP version (currently version 1), *Community* is the community name, and *PDU* is the protocol data unit that contains the SNMP request, response, or trap data. A PDU is also encoded according to ASN.1 rules.

The SNMP daemon receives and transmits all SNMP protocol messages through the Transmission Control Protocol/Internet Protocol (TCP/IP) User Datagram Protocol (UDP). Requests are accepted on well–known port 161. Traps are transmitted to the hosts listed in the trap entries in the /etc/snmpd.conf file that are listening on well–known port 162.

When a request is received, the source IP address and the community name are checked against a list containing the IP addresses, community names, permissions, and views as specified in the community and view entries in the /etc/snmpd.conf file. The snmpd agent reads this file at startup and on a refresh command or a kill -1 signal. If no matching entry is found, the request is ignored. If a matching entry is found, access is allowed according to the permissions specified in the community and view entries for that IP address, community, and view name association in the /etc/snmpd.conf file. Both the message and the PDU must be encoded according to the ASN.1 rules.

This authentication scheme is not intended to provide full security. If the SNMP daemon is used only for get and get—next requests, security may not be a problem. If set requests are allowed, the set privilege may be restricted.

See the /etc/snmpd.conf file for further information. See "Management Information Base (MIB)" in AIX Communications Programming Concepts for further information.

Request Processing

There are three types of request PDUs that may be received by the SNMP daemon. The request types are defined in RFC 1157, and the PDUs all have the following format:

Request PDU Format					
request-ID	error-status	error-index	variable-bindings		
GET	0	0	VarBindList		
GET-NEXT	0	0	VarBindList		
SET	0	0	VarBindList		

The request–ID field identifies the nature of the request; the error–status field and error–index field are unused and must be set to 0 (zero); and the variable–bindings field contains a variable–length list of numeric–format instance IDs whose values are being requested. If the value of the request–ID field is SET, the variable–bindings field is a list of pairs of instance IDs and values.

Read "Using the Management Information Base (MIB) Database" in *AIX Communications Programming Concepts* for a discussion of the three request types.

Response Processing

Response PDUs have nearly the same format as request PDUs:

Response PDU Format			
request-ID	error-status	error-index	variable-bindings
GET-RESPONSE	ErrorStatus	ErrorIndex	VarBindList

If the request was successfully processed, the value for both the error–status and error–index field is 0 (zero), and the variable–bindings field contains a complete list of pairs of instance IDs and values.

If any instance ID in the variable—bindings field of the request PDU was not successfully processed, the SNMP agent stops processing, writes the index of the failing instance ID into the error—index field, records an error code in the error—status field, and copies the partially completed result list into the variable—bindings field.

RFC 1157 defines the following values for the error–status field:

Values for the Error–Status Field			
Value	Value	Explanation	
noError	0	Processing successfully completed (error–index is 0).	
tooBig	1	The size of the response PDU would exceed an implementation—defined limit (error—index is 0).	
noSuchName	2	An instance ID does not exist in the relevant MIB view for GET and SET request types or has no successor in the MIB tree in the relevant MIB view for GET–NEXT requests (nonzero error–index).	
badValue	3	For SET requests only, a specified value is syntactically incompatible with the type attribute of the corresponding instance ID (nonzero error-index).	
readOnly	4	Not defined.	
genErr	5	An implementation—defined error occurred (nonzero error—index); for example, an attempt to assign a value that exceeds implementation limits.	

Trap Processing

Trap PDUs are defined by RFC 1157 to have the following format:

Trap PDU Format					
enterprise	agent-	generic-	specific-	time-stamp	variable-
	address	trap	trap		bindings
Object ID	Integer	Integer	Integer	TimeTicks	VarBindList

The fields are used as follows:

enterprise	The object identifier assigned to the vendor implementing the agent. This is the value of the sysObjectID variable, and it is unique for each implementer of an SNMP agent. The value assigned to this implementation of the agent is 1.3.6.1.4.1.2.3.1.2.1.1.3 , or risc6000snmpd.3 .	
agent-address	IP address of the	ne object generating the trap.
generic-trap	Integer, as follo	ws:
	0	coldStart
	1	warmStart
	2	linkDown
	3	linkUp
	4	authenticationFailure
	5	egpNeighborLoss
	6	enterpriseSpecific
specific-trap	Unused, reserved for future development.	
time-stamp	Elapsed time, in hundredths of a second, from the last reinitialization of the agent to the event generating the trap	
variable-bindings	Extra information, dependent on generic-trap type.	

The following generic-trap values indicate that certain system events have been detected:

coldStart	The agent is reinitializing. Configuration data or MIB variable values, or both, may have changed. Measurement epochs should be restarted.
warmStart	The agent is reinitializing but configuration data or MIB variable values have not changed. In this implementation of the SNMP agent, a <i>warmStart</i> trap is generated when the /etc/snmpd.conf file is reread. The configuration information in the /etc/snmpd.conf file is for agent configuration that has no side effects on SNMP manager databases. Measurement epochs should not be restarted.
linkDown	The agent has detected that a known communications interface has been disabled.
linkUp	The agent has detected that a known communications interface has been enabled.
authenticationFailure	A message was received that could not be authenticated.
egpNeighborLoss	An Exterior Gateway Protocol (EGP) neighbor was lost. This value is only generated when the agent is running on a host that runs the gated daemon using EGP.
enterpriseSpecific	Not implemented; reserved for future use.

The *linkDown* and *linkUp* traps contain a single instance ID/value pair in the variable—bindings list. The instance ID identifies the **ifIndex** of the adapter that was disabled

or enabled, and the value is the **ifIndex** value. The trap for *egpNeighborLoss* also contains a binding consisting of the instance ID and value of *egpNeighAddr* for the lost neighbor.

Generation of linkUp and linkDown Traps

Note: In the following sections, interfaces applies to both a TCP/IP interface and a Common Data Link Interface (CDLI) token—ring, Ethernet, or Fiber Distributed Data Interface (FDDI) device. CDLI allows **snmpd** to monitor the Ethernet, token—ring, and FDDI devices even if they are not running TCP/IP. The SNMP daemon stills needs TCP/IP loopback to run, but it is no longer a requirement on interfaces.

The *linkUp* and *linkDown* traps are generated when the **snmpd** agent determines there is a change in the state for a known interface. If the current state of a known interface is down and the interface state changes to up, a *linkUp* trap will be generated. Likewise, if the current state of a known interface is up and the interface changes to down, a *linkDown* trap is generated.

There is no concept of up or down to a CDLI device if there is not a TCP/IP interface running on that device. CDLI devices with a TCP/IP interface layer attached is always considered to be up unless the device is removed from the system. The other values of the interfaces table are also effected by whether or not a CDLI device has a TCP/IP interface or not. If a CDLI device has a TCP/IP interface, then all the statistics in the interfaces table pertain to the TCP/IP interface and the device specific MIB should be used to retrieve statistics about the device. If there is not a TCP/IP interface layer running on the CDLI device, the iftable entries for the device are retrieved from the CDLI device itself.

If an interface is known to the **snmpd** agent, there is an entry for that interface in the **snmpd** interfaces table. When an interface is attached or detached by the **ifconfig** command, the entries in the **snmpd** agent's interfaces table change. A *coldStart* trap is generated to indicate the configuration change. The purpose of generating a *coldStart* trap is to guarantee the receiving host understands that crucial MIB variable values may have changed. In particular, measurement epochs should be restarted. Even though links may be enabled or disabled, no *linkUp* or *linkDown* traps are generated.

For instance, a **netstat –in** command may provide the following information for a host's network configuration:

Name	Mtu	Network	Address	Ipkts	Ierrs	Opkts	0errs	Coll	
100	1536	127	127.0.0.1	6228	0	6228	0	0	
en0	1500	192.100.154	192.100.154.7	585287	0	666636	0	0	
tr0	1500	129.35.32	129.35.42.141	397632	:3 0	2414030	0 (0	

In this example, network configuration, the **snmpd** agent has three entries in its interface table from TCP/IP interfaces: one for *lo0*, one for *en0*, and one for *tr0*. In the interface index table, ifIndex.1 might refer to lo0, ifIndex.2 might refer to *en0*, and ifIndex.3 might refer to *tr0*. The term "might refer to" implies that since interfaces are dynamic, the actual entry number may vary. In the example, there is an additional token–ring adapter in the workstation that the TCP/IP kernel does not know about; the interface for this adapter is named *tr1*. Since, the token–ring adapter is a CDLI device, there is an entry in the interfaces table, for example, ifIndex.4. The workstation has a serial optical device. The serial optical device has not been configured for TCP/IP and it is not a CDLI device. So, neither TCP/IP nor **snmpd** recognize this device. Its interface name is so0.

If you issue the **ifconfig tr1** command, TCP/IP attaches *tr1*, but does not mark the interface as up. The **snmpd** agent changes the method for reporting interface table statistics to the TCP/IP interface layer from the CDLI device level statistics. The **snmpd** agent then generates a *coldStart* trap. This action has not added a new entry, because an entry for *tr1* already existed.

If you issue the **ifconfig so0** command, TCP/IP attaches *so1*, but does not mark the interface as up. The **snmpd** agent then adds a fifth entry for this new interface to its interfaces table and sets ifIndex.5 to refer to *so0*. The SNMP agent then generates a *coldStart* trap. Since, *so0* is not a CDLI device, it was not in the interface table as a result of device configuration, so it had to be added when the TCP/IP interface layer was configured.

The crucial interface index (ifIndex) values that an SNMP manager may have stored for the four original entries will not change. But the *coldStart* trap is a signal to SNMP managers that their MIB databases need to be updated. Upon refreshing its database, an SNMP manager learns about this new entry in the **snmpd** agent's interface table.

If the system administrator issued an **ifconfig tr1 up** command, the new interface is marked up and the method of obtaining statistics is changed. The **snmpd** agent would send a *coldStart* trap; a *linkUp* trap would not be sent because the state of the device did not change. A CDLI device is always considered up, until it gets a TCP/IP interface.

The *coldStart* trap indicates to the SNMP managers that the **snmpd** agent configuration had changed, so the managers should update their MIB databases. A *linkUp* trap generated after the *coldStart* trap is meaningless to the SNMP manager because the manager's database already has the information from its database refresh.

The system administrator elected to detach the *en0* interface from the above network configuration. Once detached, the **snmpd** agent updates its statistics gathering method to use the CDLI device statistics (Ethernet is a CDLI device). The agent generates a *coldStart* trap to inform the SNMP managers that something in the interfaces is different. In this case, the crucial ifIndex values have not changed, just the method of statistic retrieval.

The system administrator has elected to detach the *so0* interface from the above network configuration. Upon the detachment of an interface, the **snmpd** agent updates its interfaces table. In this example, all the indices remain the same except that the fifth entry has been removed. If the third entry had been removed, the fourth and fifth entry would have been renumbered to be the third and fourth entries. In either case, the **snmpd** agent generates a *coldStart* trap.

In this case, the crucial ifIndex values that an SNMP manager may have stored for the original entries will change. The SNMP manager needs to refresh its MIB database to reflect the changes in the **snmpd** agent's interfaces table. In this situation, a *linkDown* trap is not generated. An SNMP manager cannot take immediate action upon receipt of a *linkDown* trap because the ifIndex values in its database are no longer valid.

In order for the **snmpd** agent to catch all interface status changes, the **snmpd** agent periodically checks the TCP/IP kernel and CDLI device list to determine the status of the interfaces. The time interval at which these checks are run is user—configurable.

In the event that the **snmpd** agent has received a request for a MIB variable in the interfaces table and the **snmpd** agent determines that an interface has changed state to the degree that a *coldStart* should be generated, the **snmpd** agent returns a *genErr* and issues the *coldStart* trap.

Refer to "User Datagram Protocol (UDP)", on page 3-28, "Exterior Gateway Protocol (EGP)", on page 3-32, and "TCP/IP Addressing", on page 3-52 for more information on protocols and Internet addresses.

SNMP Daemon Support for the EGP Family of MIB Variables

If the agent host is running the **gated** daemon with the Exterior Gateway Protocol (EGP) enabled, there are several Management Information Base (MIB) variables in the EGP group supported by the **gated** daemon which the **snmpd** agent can access.

The following EGP MIB variables have a single, unique instance:

egpInMsgs
Number of EGP messages received without error.

Number of EGP messages received in error.

Total number of EGP messages transmitted by the gated daemon running on the agent's host.

Pumber of EGP messages that could not be sent by the agent host's gated daemon due to resource limitations.

Autonomous system number of the agent host's gated

daemon.

The following EGP MIB variables have an instance for each EGP peer or neighbor acquired by the agent host's **gated** daemon:

egpNeighState	The state of th	is EGP peer:
	1	idle
	2	acquisition
	3	down
	4	up
	5	cease
egpNeighAddr	IP address of t	his EGP peer.
egpNeighAs		ystem number of this EGP peer. Zero (0) utonomous system number of this peer is not
egpInNeighMsgs	Number of EG EGP peer.	P messages received without error from this
egpNeighInErrs	Number of EGP messages received in error from this EGP peer.	
egpNeighOutMsgs	Number of loca peer.	ally generated EGP messages to this EGP
egpNeighOutErrs		ally generated EGP messages not sent to this to resource limitations.
egpNeighInErrMsgs	Number of EGP-defined error messages received from this EGP peer	
egpNeighOutErrMsgs	Number of EGP-defined error messages sent to this EGP peer.	
egpNeighStateUp	Number of EGP state transitions to the UP state with this EGP peer.	
egpNeighStateDowns	Number of EGP state transitions from the UP state to any other state with this EGP peer.	
egpNeighIntervalHello	Interval between hundredths of	en EGP Hello command retransmissions in a second.
egpNeighIntervalPoll	Interval between hundredths of	en EGP poll command retransmissions in a second.

egpNeighMode Polling mode of this EGP peer. The mode can be either

active (1) or passive (2).

egpNeighEventTrigger Control variable triggers operator—initiated start and stop

events to this EGP peer. This MIB variable can be set to

start (1) or stop (2).

If the **gated** daemon is not running, or if the **gated** daemon is running but is not configured to communicate with the **snmpd** agent, or if the **gated** daemon is not configured for EGP, get and set requests for the values of these variables will return the *noSuchName* error response code.

The **gated** daemon configuration file, /etc/gated.conf, should contain the following statement:

snmp yes;

The **gated** daemon is internally configured to be an Simple Network Management Protocol (SNMP) single multiplexer (SMUX) protocol peer, or proxy agent, of the **snmpd** daemon. When the **gated** daemon starts up, it registers the *ipRouteTable* MIB variable tree with the **snmpd** agent. If the **gated** daemon is configured for EGP, then the **gated** daemon also registers the EGP MIB variable tree. After this registration is complete, an SNMP manager can successfully make requests to the **snmpd** agent for the *ipRouteTable* an EGP MIB variables supported by this agent host's **gated** daemon. Thus, when the **gated** daemon is running, all MIB routing information is obtained using the **gated** daemon. In this case, set requests to the *ipRouteTable* are not allowed.

The SMUX communication between the **gated** daemon and the **snmpd** daemon takes place over the well–known Transmission Control Protocol (TCP) port 199. If the **gated** daemon should terminate, **snmpd** immediately unregisters the trees the **gated** daemon previously registered. If the **gated** daemon is started before the **snmpd** daemon, the **gated** daemon periodically checks for the **snmpd** daemon until the SMUX association can be established.

To configure the **snmpd** agent to recognize and allow the SMUX association with the **gated** daemon client, the user must add a SMUX entry to the /**etc/snmpd.conf** file. The client object identifier and password specified in this SMUX entry for the **gated** daemon must match those specified in the /**etc/snmpd.peers** file.

The **snmpd** agent supports set requests for the following read–write MIB I and MIB II variables:

sysContact

The textual identification of the contact person for this agent's host. This information includes the name of this person and how to contact this person: for example, "Bob Smith, 555–5555, ext 5." The value is limited to 256 characters. If, for a set request, the string for this MIB variable is greater than 256 characters, the **snmpd** agent will return the error <code>badValue</code>, and the set operation will not be performed. The initial value of <code>sysContact</code> is defined in <code>/etc.snmp.conf</code>. If nothing is defined, the value is null string.

Instance	Value	Action
0	"string"	The MIB variable is set to "string".

sysName

The host name for this agent's host. Typically this is the node's fully qualified domain name. The value is limited to 256 characters. If, for a set request, the string for this MIB variable is greater than 256 characters, the **snmpd** agent will return the error *badValue*, and the set operation will not be performed.

Instance	Value	Action
0	"string"	The MIB variable is set to
		"string".

sysLocation

A textual string stating the physical location of the machine on which this snmpd agent resides: for example, "Austin site, building 802, lab 3C-23." The value is limited to 256 characters. If, for a set request, the string for this MIB variable is greater than 256 characters, the **snmpd** agent will return the error badValue, and the set operation will not be performed. The initial value of sysLocation is defined in /etc/snmp.conf. If nothing is defined, the value is null string.

Instance	Value	Action
0	•	The MIB variable is set to "string".

ifAdminStatus The desired state of an interface adapter on the agent's host. Supported states are up and down. The state can be set to testing but such an action has no effect on the operational state of the interface.

Instance	Value	Action
f		The interface adapter with ifIndex f is enabled.

Note: It is possible that the *ifAdminStatus* value can be set to up or down, yet the actual operational change of the interface failed. In such a case, a get request of the ifAdminStatus may reflect up while an ifOperStatus for that interface may reflect down. If such a situation should occur, the network administrator would have to issue another set request to set ifAdminStatus to up to attempt the operational change again.

atPhysAddress The hardware address portion of an address table binding on the agent's host (an entry in the Address Resolution Protocol, or ARP, table). This is the same MIB variable as ipNetToMediaPhysAddress.

Instance	Value	Action
f.1.n.n.n.n	hh:hh:hh:hh:hh	For the interface with iflndex f, any existing ARP table binding for IP address n.n.n.n is replaced with the binding (n.n.n.n, hh:hh:hh:hh:hh). If a binding did not exist, the new binding is added. hh:hh:hh:hh:hh:hh is a twelve—hexadecimal—digit hardware address.

atNetAddress The IP address corresponding to the hardware or physical address specified in atPhysAddress. This is the same MIB variable as ipNetToMediaNetAddress.

Instance	Value	Action
f.1.n.n.n.n		For the interface with ifIndex f, an existing ARP table entry for IP address n.n.n.n is replaced with IP address m.m.m.m.

ipForwarding Indicates whether this agent's host is forwarding datagrams. See "SNMP Daemon Implementation Restrictions", on page 9-27 for more information on this MIB variable.

Instance	Value	Action
0	1	If the agent's host has more than one active interface, then the TCP/IP kernel is configured to forward packets. If the agent's host has only one active interface, the set request fails.
	2	The TCP/IP kernel on the agent's host is configured to not forward packets.

ipDefaultTTL The default time—to—live (TTL) value inserted into IP headers of datagrams originated by the agent's host.

Instance	Value	Action
0		The default time—to—live value used by IP protocol support is set to the integer n.

ipRouteDest The destination IP address of a route in the route table.

Instance	Value	Action
n.n.n.n		The destination route for route n.n.n.n is set to the IP address m.m.m.m.

ipRouteNextHop

The gateway by which a destination IP address can be reached from the agent's host (an entry in the route table).

Instance	Value	Action
n.n.n.n	m.m.m.m	A route table entry to reach network n.n.n.n using gateway m.m.m.m is added to the route table. The host portion of the IP address n.n.n.n must be 0 to indicate a network address.

ipRouteType The state of a route table entry on the agent's host (used to delete entries).

Instance	Value	Action
h.h.h.h	1	Any route to host IP address h.h.h.h is deleted.
n.n.n.n	2	Any route to host IP address n.n.n.n is deleted.

ip Net To Media Phys Address

The hardware address portion of an address table binding on the agent's host (an entry in the ARP table). This is the same MIB variable as atPhysAddress.

Instance	Value	Action
f.1.n.n.n.n	hh:hh:hh:hh:hh	For the interface with ifIndex f, any existing ARP table binding for IP address n.n.n.n is replaced with the binding (n.n.n.n, hh:hh:hh:hh:hh). If a binding did not exist, the new binding is added. hh:hh:hh:hh:hh:hh is a 12–hexadecimal–digit hardware address.

ipNetToMediaNetAddress

The IP address corresponding to the hardware or physical address specified in *ipNetToMediaPhysAddress*. This is the same MIB variable as *atNetAddress*.

Instance	Value	Action
f.1.n.n.n.n		For the interface with ifIndex f, an existing ARP table entry for IP address n.n.n.n is replaced with IP address m.m.m.m.

ipNetToMediaType

The type of mapping from the IP address to the physical address.

Instance	Value	Action
f.1.n.n.n.n	1	For the interface with ifIndex f, for an existing ARP binding from IP address to physical address, the mapping type is set to 1, or other.
	2	For the interface with iflndex f, for an existing ARP binding from IP address to physical address, the mapping type is set to 2, or invalid. As a side effect, the corresponding entry in the ipNetMediaTable is invalidated; that is, the interface is disassociated from this ipNetToMediaTable entry.
	3	For the interface with iflndex f, for an existing ARP binding from IP address to physical address, the mapping type is set to 3, or dynamic.
	4	For the interface with ifIndex f, for an existing ARP binding from IP address to physical address, the mapping type is set to 4, or static.

snmp Enable Authen Traps

Indicates whether the **snmpd** agent is configured to generate *authenticationFailure* traps.

Instance	Value	Action
0	1	The snmpd agent will not generate authentication failure traps.
	2	The snmpd agent will generate authentication failure traps.

smuxPstatus The status of an SMUX protocol peer (used to delete SMUX peers).

Instance	Value	Action
n	1	snmpd agent does nothing.
	2	snmpd agent will stop communicating with SMUX peer n.

smuxTstatus The status of a SMUX MIB tree (used to delete MIB tree mounts).

Instance	Value	Action
I.m.m.mp	1	snmpd agent does nothing.
	2	Unmounts SMUX mounting of MIB tree m.m.m where <i>I</i> is the length of MIB tree instance and <i>p</i> is the smuxTpriority.

The variables listed below are the settable variables as defined in RFC 1229. The **snmpd** daemon allows the user to set these variables. The underlying device might not allow the setting of such variables. You should check with each device to see what is and is not supported.

ifExtnsPromiscuous

The status of the promiscuous mode on a given device. This is used to enable and disable promiscuous mode on a given device. The **snmpd** action is final and complete. When **snmpd** is told to turn off, promiscuous mode is turned completely off regardless of the other applications on the machine.

Instance	Value	Action
n	1	Turns on the promiscuous mode for device n.
	2	Turns off the promiscuous mode for device n.

ifExtnsTestType

The test initiation variable. When this variable is set the appropriate test is run for that device. An Object Identifier is the value of the variable. The specific value is dependent on the device type and the test wished to be run. Currently, the only define test that **snmpd** knows to run is testFullDuplexLoopBack test.

Instance	Value	Action
n		Start the test specified by oid.

ifExtnsRcvAddrStatus

The address status variable. When this variable is set, the specified address comes into existence with the appropriate level of duration. **snmpd**

only allows the setting of temporary addresses because it is not able to set device Object Data Manager (ODM) records and it is only allowed to set multicast/broadcast addresses.

Instance	Value	Action
n.m.m.m.m.m	1	Add the address as something other than a temporary or permanent address.
	2	Remove the address from usage.
	3	Add the address as a temporary address.
	4	Add the address as a permanent address.

The variables listed below are the settable variables as defined in RFC 1231. The **snmpd** daemon allows the user to set these variables. The underlying device might not allow the setting of such variables. You should check with each device to see what is and is not supported.

dot5Commands

The command the token-ring device should perform.

Instance	Value	Action
n	1	Does nothing. Returned.
	2	Tells the token–ring device to open.
	3	Tells the token-ring to reset.
	4	Tells the token–ring device to close.

dot5RindSpeed

The current ring speed or bandwidth.

Instance	Value	Action
n	1	An unknown speed.
	2	1 megabit ring speed.
	3	4 megabit ring speed.
	4	16 megabit ring speed.

dot5ActMonParticipate

The object specifies whether the device should participate in the active monitor selection process.

Instance	Value	Action
n	1	Should participate.
	2	Should not participate.

dot5Functional

The functional mask that allows the token–ring device to specify from what addresses it will receive frames.

Instance	Value	Action
n	m.m.m.m.m	Functional mask to be set.

The following variables are defined in the RFC as read—only but you are encouraged to make them read—write. They are complex timer manipulations. You should look them up in the RFC to gain a full understanding of their interactions. **snmpd** allows the requestor to set them, but the device may not. Check the device driver documentation for more information. The variables are:

- dot5TimerReturnRepeat
- dot5TimerHolding
- dot5TimerQueuePDU
- dot5TimerValidTransmit
- dot5TimerNoToken
- dot5TimerActiveMon
- dot5TimerStandbyMon
- dot5TimerErrorReport
- dot5TimerBeaconTransmit
- dot5TimerBeaconReceive

The variables listed below are the settable variables as defined in RFC 1512. The SNMP daemon allows the user to set these variables. It uses the FDDI Station Management (SMT) 7.2 protocol standard to get the information. This is determined at the microcode level. Check the microcode on the FDDI documentation to ensure that the SMT 7.2 microcode is being used.

fddimibSMTUserData

A variable holding 32 bytes of user information.

Instance	Value	Action
n	_	Stores 32 bytes of user information.

fddimibSMTConfigPolicy

The status of the configuration policies, specifically the hold policy usage.

Instance	Value	Action
n	0	Do not use the hold policy.
	1	Use the hold policy.

fddimibSMTConnectionPolicy

The status of the connection policies in the FDDI node. See RFC 1512 for more information about the specific settable values.

Instance	Value	Action
n		Defines the connection policies.

fddimibSMTTNotify

The timer, expressed in seconds, used in the Neighbor Notification protocol. It has a range of 2 seconds to 30 seconds, and its default value is 30 seconds.

Instance	Value	Action
n	k	Defines the timer value.

fddimibSMTStatRptPolicy

The status of the status reporting frame generation.

Instance	Value	Action
n	1	Indicates that the node will generate status reporting frames for implemented events.
	2	Indicates that the node will not create status reporting frames.

fddimibSMTTraceMaxExpiration

This variable defines the maximum timer expiration value for trace.

Instance	Value	Action
n		Defines the maximum timer expiration in milliseconds.

fddimibSMTStationAction

This variable causes the SMT entity to take a specific action. Refer to the RFC to get specific information about this variable.

Instance	Value	Action
n		Defines an action on the SMT entity. Values range from 1 to 8.

fddimibMACRequestedPaths

Defines the paths the medium access control (MAC) should be inserted.

Instance	Value	Action
n.n	k	Defines the requested path
		for the MAC.

fddimibMACFrameErrorThreshold

Threshold for when a MAC status report should be generated. Defines the number of error that must occur before a report is generated.

Instance	Value	Action
n.n		Defines the number of errors that must be observed before a MAC status report is generated.

fddimibMACMAUnitdataEnable

This variable determines the value of the **MA_UNITDATA_Enable** flag in RMT. The default and initial value of this flag is true (1).

Instance	Value	Action
n.n	1	Marks the MA_UNITDATA_Enable flag true.
	2	Marks the MA_UNITDATA_Enable flag false.

fddimibMACNotCopiedThreshold

A threshold for determining when a MAC condition report will be generated.

Instance	Value	Action
n.n	k	Defines the number of errors that must be observed before a MAC condition report is generated.

The following three variables are timer variables that are interactive among themselves. Before changing any of these variables, you should have a good understanding of their meaning as defined in **RFC 1512**.

- fddimibPATHTVXLowerBound
- fddimibPATHTMaxLowerBound
- fddimibPATHMaxTReq

fddimibPORTConnectionPolicies

Specifies the connection policies for the specified port.

Instance	Value	Action
n.n		Defines the connection policies for the specified port.

fddimibPORTRequestedPaths

This variable is a list of permitted paths where each list element defines the port's permitted paths. The first octet corresponds to 'none', the second octet to 'tree', and the third octet to 'peer'.

Instance	Value	Action
n.n	CCC	Defines the port's paths.

fddimibPORTLerCutoff

The link error rate estimate at which a link connection will be broken. It ranges from 10**–4 to 10**–15 and is reported as the absolute value of the base 10 logarithm (default of 7).

Instance	Value	Action
n.n	k	Defines the port's LerCutoff.

fddimibPORTLerAlarm

The link error rate estimate at which a link connection will generate an alarm. It ranges from 10**–4 to 10**–15 and is reported as the absolute value of the base 10 logarithm of the estimate (default is 8).

Instance	Value	Action
n.n	k	Defines the port's LerAlarm.

fddimibPORTAction

This variable causes the port to take a specific action. Refer to the RFC to get specific information about this variable.

Instance	Value	Action
n		Defines an action on the defined port. The values range from 1 to 6.

Note: RFC 1213 describes all variables in the atEntry and ipNetToMediaEntry tables as read—write. Set support is implemented only for the atEntry variables atPhysAddress and atNetAddress, and the ipNetToMediaEntry variables ipNetToMediaPhysAddress, ipNetToMediaNetAddress, and ipNetToMediaType. To accept set requests that may specify the remaining unsupported attributes in these two tables, set requests for the remaining variables are accepted: atlfIndex and ipNetToMediaIfIndex. No error response is returned to the set request originator, but a subsequent get request will show that the original values are retained.

RFC 1213 describes all variables in the *ipRouteEntry* table as read—write except *ipRouteProto*. As mentioned above, set support is implemented only for the variables *ipRouteDest*, *ipRouteNextHop*, and *ipRouteType*. To accept set requests that may specify several unsupported route attributes, set requests for the remaining variables in the *ipRouteEntry* table are accepted: *ipRouteIfIndex*, *ipRouteMetric1*, *ipRouteMetric2*, *ipRouteMetric3*, *ipRouteMetric4*, *ipRouteMetric5*, *ipRouteAge*, and *ipRouteMask*. No error response is returned to the set request originator, but a subsequent get request will show that the original values are retained. The **snmpd** daemon does not coordinate routing with the **routed** daemon. If the **gated** daemon is running and has registered the *ipRouteTable* with the **snmpd** daemon, set requests to the *ipRouteTable* are not allowed.

RFC 1229 describes settable variables; **snmpd** allows for these to be set. Refer to the previous entries for actual deviations.

Examples

The following examples use the **snmpinfo** command. It is assumed that the **snmpinfo** default community name, public, has read—write access for the respective MIB subtree.

```
snmpinfo -m set sysContact.0="Primary contact: Bob Smith, office phone: 555-5555, beeper: 9-123-4567. Secondary contact: John Harris, phone: 555-1234."
```

This command sets the value of sysContact.0 to the specified string. If an entry for sysContact.0 already exists, it is replaced.

```
snmpinfo -m set sysName.0="bears.austin.ibm.com"
```

This command sets the value of sysName.0 to the specified string. If an entry for sysName.0 already exists, it is replaced.

```
snmpinfo -m set sysLocation.0="Austin site, building 802, lab 3C-23, southeast corner of the room."
```

This command sets the value of sysLocation.0 to the specified string. If an entry for sysLocation.0 already exists, it is replaced.

```
snmpinfo -m set ifAdminStatus.2=2
```

This command disables the network interface adapter which has the ifIndex of 2. If the assigned value is 1, the interface adapter is enabled.

```
snmpinfo -m set atPhysAddress.2.1.192.100.154.2=02:60:8c:2e:c2:00
snmpinfo -m set
ipNetToMediaPhysAddress.2.1.192.100.154.2=02:60:8c:2e:c2:00
```

These two commands change the hardware address in the ARP table entry for 192.100.154.2 to 02:60:8c:2e:c2:00. These two commands affect the same ARP table entry. The MIB variable *atPhysAddress* is a deprecated variable and is being replaced with the MIB variable *ipNetToMediaPhysAddress*. Thus, *atPhysAddress* and *ipNetToMediaPhysAddress* access the same structure in the TCP/IP kernel ARP table.

```
\label{lem:snmpinfo} \mbox{ snmpinfo -m set atNetAddress.2.1.192.100.154.2=192.100.154.3 snmpinfo -m set ipNetToMediaNetAddress.2.1.192.100.154.2=192.100.154.3 }
```

These commands change the IP address in the ARP table entry for 192.100.154.2 to 192.100.154.3. These two commands affect the same ARP table entry. The MIB variable atNetAddress is a deprecated variable and is being replaced with the MIB variable ipNetToMediaNetAddress. Thus, atNetAddress and ipNetToMediaNetAddress access the same structure in the TCP/IP kernel ARP table.

```
snmpinfo -m set ipForwarding.0=1
```

This command sets the TCP/IP kernel so that it can forward packets if the agent's host has more than one interface that is up. If the host has only one active interface, then the set request fails and the **snmpd** agent returns the error, *badValue*.

```
snmpinfo -m set ipDefaultTTL=50
```

This command allows an IP datagram using default time—to—live (TTL) to pass through up to 50 gateways before being discarded. When each gateway processes a datagram, the gateway subtracts 1 from the time—to—live field. In addition, each gateway decrements the time—to—live field by the number of seconds the datagram waited for service at that gateway before passing the datagram on to the next destination.

```
snmpinfo -m set ipRouteDest.192.100.154.0=192.100.154.5
```

This command sets the destination IP address of the route associated with 192.100.154.0 to the IP address 192.100.154.5, assuming route 192.100.154 already existed.

```
snmpinfo -m set ipRouteNextHop.192.100.154.1=129.35.38.47
```

This command sets a route to host 192.100.154.1 using the gateway host 129.35.38.47, assuming route 192.100.154.1 already existed.

```
snmpinfo -m set ipRouteNextHop.192.100.154.0=192.100.154.7
```

This command sets a route to the class C network 192.100.154 using the gateway host 192.100.154.7, assuming route 192.100.154.0 already existed. Note that the host part of the address must be 0 to indicate a network address.

```
snmpinfo -m set ipRouteType.192.100.154.5=2
```

This command deletes any route to host 192.100.154.5.

This command creates a new route from host 129.35.128.90 to 129.35.128.1 as a gateway.

```
snmpinfo -m set ipNetToMediaType.2.1.192.100.154.11=4
```

This command sets the ARP table entry for 192.100.154.11 to static.

```
snmpinfo -m set snmpEnableAuthenTraps=2
```

This command causes the **snmpd** agent on the specified host to not generate *authenticationFailure* traps.

```
snmpinfo -m set smuxPstatus.1=2
```

This command invalidates the SMUX peer 1. The side effect is that the connection between the **snmpd** agent and this SMUX peer is terminated.

```
snmpinfo -m set smuxTstatus.8.1.3.6.1.2.1.4.21.0=2
```

This command invalidates or removes the mounting of the SMUX tree 1.3.6.1.2.1.4.21, the *ipRoute* Table. The first number in the instance indicates the number of levels in the SMUX tree identifier. The final number in the instance indicates the smuxTpriority. In this example, there are 8 levels in the SMUX tree identifier:

```
1.3.6.1.2.1.4.21. The priority, 0, is the highest priority.
```

```
\verb|snmpinfo-m| set if Extns Promiscuous.1=1 if Extns Promiscuous.2=2|
```

This command turns on promiscuous mode for the first device in the interfaces table and turns off promiscuous mode for the second device in the interfaces table.

```
snmpinfo -m set ifExtnsTestType.1=testFullDuplexLoopBack
```

This command starts the testFullDuplexLoopBack test on interface 1.

```
snmpinfo -m set ifExtnsRcvAddrStatus.1.129.35.128.1.3.2=2
```

This command tells interface 1 to remove physical address 129.35.128.1.3.2 from its list of acceptable addresses.

```
snmpinfo -m set dot5Commands.1=2
```

This command tells the first interface to do an open.

```
snmpinfo -m set dot5RingSpeed.1=2
```

This command tells the first interface to set it ring speed to 1 megabit.

```
snmpinfo -m set dot5ActMonParticipate.1=1
```

This command tells the first interface to participate in the active monitor selection process.

```
snmpinfo -m set dot5Functional.1=255.255.255.255.255.255
```

This command sets the functional address mask to allow everything.

```
snmpinfo -m set fddimibSMTUserData.1="Greg's Data"
```

This command set the user data on the first SMT entity to "Greg's Data".

```
snmpinfo -m set fddimibMACFrameErrorThreshold.1.1=345
```

This command set the threshold for frame errors to 345 on the first SMT entity's first MAC.

Note: All of the variables described earlier fall into one of the methods used above to set the variable.

Refer to *Address Resolution Protocol*, on page 3-23 and *Internet Addresses*, on page 3-52 for more information on protocols and Internet addresses.

SNMP Daemon RFC Conformance

RFC 1157 requires that each set request variable assignment "should be effected as if simultaneously set with respect to all other assignments specified in the same message" (on pages 25 and 26 of the RFC). This means that a set request with multiple instance ID/value pairs should be processed in an *all-or-none* fashion. That is, either all the new values are assigned without error, or else none of the variables in the request have modified values. This requirement is also known as "atomic commit with rollback."

Note: RFC 1157 does not address the problems of order dependency or consistency. Ordering dependencies may exist, such as in:

```
snmpinfo -m set -h host1
ipNetToMediaPhysAddress.f.1.n.n.n=hh:hh:hh:hh:hh:hh:hh
ifAdminStatus.f=1
snmpinfo -m set -h host1 iproutenexthop.n.n.n.n=m.m.m.m ifAdminStatus.f=1
```

The adapter with **ifIndex f** must be enabled before an address table entry can be associated with it or before a route can be established that reaches a gateway by using that adapter. In these cases, the order of the variables is significant and must be the reverse of that shown above. If the RFC 1157 atomic commit policy were followed exactly, ordered set requests would have undefined effects.

For the Management Information Base (MIB) variables supported, the **snmpd** agent prechecks the specified values for the MIB variables in the set request. If a specified value does not meet the precheck requirements, the set request is rejected. The actual implementation for the set request is technically a best effort, not a true atomic commit and rollback.

During the set process when the **snmpd** agent is actually modifying MIB variable values, if a failure occurs, the original values of the MIB variables already set are not restored.

RFC 1213 describes all variables in the ipRouteEntry table as read—write. As described above, set support is only implemented for **ipRouteDest**, **ipRouteType**, and **ipRouteNextHop** (see note). To accept set requests that may specify several unsupported route attributes (such as the **ipRouteMetric1** or the **ipRouteProto** attributes), set requests for these unsupported variables are accepted. No error response is returned to the request originator, but a subsequent get request will show that the original values have been retained.

SNMP Daemon Implementation Restrictions

The current implementation of the Simple Network Management Protocol (SNMP) agent does *not* have:

- User interface for proxy agent support
- Nontrivial authentication support

Not all values defined in RFC 1213 are supported, as illustrated in the following note:

Note: A set request for **ipRouteType 2** (invalid) causes deletion of a route. A get request will never return this value for an existing route.

The following implementation restrictions should be observed:

- An interface must be enabled by a locally issued ifconfig command before it can be
 enabled or disabled by the ifAdminStatus variable. There is no way to associate an
 Internet Protocol (IP) address with an interface except with a locally issued ifconfig
 command.
- The **snmpd** agent does not provide a way to alter the IP address associated with an interface adapter. This must be done with a locally issued **ifconfig** command.
- The maximum message size that can be sent is 9KB (9216 bytes). The maximum message that can be received is 40KB (40960 bytes).

If you set the value of the MIB variable **ipForwarding** to 1, indicating that the TCP/IP kernel should forward packets, the **snmpd** agent performs some interface checks before it will allow the set operation to succeed. If there is more than one active interface, then the MIB variable **ipForwarding** is set to 1 and the Transmission Control Protocol/Internet Protocol (TCP/IP) kernel configurable option **ipforwarding** is set to 1. Otherwise, the **snmpd** agent does not perform the set operation and an error *badValue* is returned.

SNMP Daemon Logging Facility

Logging activities for the **snmpd** daemon can be handled in three ways:

- From the snmpd command line
- From the **snmpd** configuration file
- By the syslogd daemon

The debug level of the **snmpd** logging messages is user–configurable. The debug level can be one of the following:

- 0 All NOTICES, EXCEPTIONS, and FATAL messages
- 1 Level 0 plus DEBUG messages
- 2 Level 1 plus a hexadecimal dump of the request and response packets
- 3 Level 2 plus an English version of the request and response packets.

The default debug level is 0. Level 3 messages are not written by the **syslogd** daemon.

If the **snmpd** daemon is configured for logging, logging can be toggled on and off by issuing a **SIGUSR1** signal, or a **kill –30** to the **snmpd** daemon. If the **snmpd** daemon is invoked under System Resource Controller (SRC) control, the SRC **traceson** and **tracesoff** commands can also be used to start and stop the logging process. If logging is activated, the **tracesoff** command will stop logging. Likewise, if logging is temporarily disabled, the **traceson** command can be used to restart the logging. If the **snmpd** daemon is not configured for logging, these commands have no effect. These commands have no effect on logging by the **syslogd** daemon.

Should the log file reach the file size limit, the **snmpd** daemon will rotate the log file. The **snmpd** daemon will keep up to four levels of rotated log files. For example, if you name the log file LogFile, the files are rotated as follows:

- LogFile.3 is deleted.
- LogFile.2 is renamed LogFile.3.
- LogFile.1 is renamed LogFile.2.
- LogFile.0 is renamed LogFile.1.
- LogFile is renamed LogFile.0.
- Logging continues in LogFile.

If you are logging with **syslogd** at the same time, a log message will be logged by the **syslogd** daemon stating that the log files are rotating. The **syslogd** log file will not be rotated by the **snmpd** daemon.

If the **snmpd** daemon is invoked under SRC control, the **Issrc** command with the long status option will list the **snmpd** logging parameters. The **Issrc** command will not display information about the **syslogd** logging activities.

Logging Directed from the snmpd Command Line

To direct logging from the **snmpd** command line, the **-f** flag must be specified at the **snmpd** invocation. If the **-f** flag is not specified, logging cannot be directed from the command line. If the **-f** flag is specified, the full path name and file name of the logging file must be specified as the **snmpd** daemon forks and changes to the root directory at startup. If the **snmpd** daemon cannot open the file, the **snmpd** daemon will direct logging from the configuration file. If **syslogd** is also handling **snmpd** log messages, an EXCEPTIONS message will be logged to the **syslogd** log file stating the reason why the **snmpd** log file could not be opened. If the log file specified with the **-f** flag is successfully opened, the log file cannot be changed during execution of the **snmpd** daemon.

The debug level is specified on the **snmpd** command line with the **-d** flag. If the **-d** flag is not specified, the debug level will default to 0. The debug level cannot change during execution of the **snmpd** daemon.

The size of the log file will be **unlimited**; that is, it will be the system maximum file size.

If logging is directed from the command line, logging entries in the configuration file will be ignored both at startup and during any refresh of the **snmpd** daemon.

See the **snmpd** command in in AIX Commands Reference.

Logging Directed from the Configuration File

To direct logging from the configuration file, the **–f** flag must not be specified on the **snmpd** command line. If the **–d** flag is specified on the **snmpd** command line, the debug level specified with the **–d** flag becomes the default debug level. If the **–d** flag is not specified on the **snmpd** command line, the default debug level is 0.

The logging parameters are specified in logging entries in the **snmpd** configuration file. The configurable logging parameters include: log file name, maximum log file size, debug level, and enablement. If the log file name is not specified, logging is not enabled. Because the **snmpd** daemon forks and changes to the root directory at startup, the full path name of the log file must be specified. The default enablement is disabled. Thus, if the enablement parameter is not specified as *enabled*, logging to the log file will not take place.

The default file size is **unlimited**; that is, it defaults to the system maximum file size.

The **snmpd** configuration file is read at **snmpd** startup and upon a refresh. Thus, logging parameters do not need to be specified before the **snmpd** daemon is invoked. Because the file is reread when the **snmpd** daemon receives a refresh (if **snmpd** is started under SRC control) or a SIGHUP signal (**kill** –1), the logging parameters can be specified at any time. In addition, the logging parameters can be changed at any time during the running of the **snmpd** daemon.

See the **snmpd.conf** file for more information on how to configure the **snmpd** daemon for logging directed from the configuration file.

Logging by the syslogd Daemon

Logging by the **syslogd** daemon can take place alone or in conjunction with logging directed from either the **snmpd** command line or configuration file.

To configure the **syslogd** daemon to log messages for the **snmpd** daemon, you must be the root user. Edit the /**etc/syslog.conf** file and add an entry such as the following:

```
daemon.debug /var/tmp/snmpd.syslog
```

The /var/tmp/snmpd.syslog file must exist before the syslogd daemon rereads the /etc/syslog.conf configuration file in order for the syslogd daemon to log the snmpd daemon log messages to this file. To create this file, issue the following command:

```
touch /var/tmp/snmpd.syslog
```

Then issue the following command to force the **syslogd** daemon to reread its configuration file:

```
refresh -s syslogd
```

Note that the **syslogd** daemon will log *all* daemon messages to this log file, not just the **snmpd** log messages.

If the **syslogd** daemon is configured to log messages from the daemon facility at the **syslogd** LOG_DEBUG severity level and higher, all messages at **snmpd** debug level 2 or lower from the **snmpd** daemon can be logged into a **syslogd** configured file. If level 3 is

specified as the **snmpd** debug level, the **syslogd** daemon logs only **snmpd** level 2 messages.

If logging is to be handled solely by the **syslogd** daemon, first the **syslogd** daemon must be configured for logging as described above. The **snmpd** daemon debug level must be specified on the command line with the **-d** option, or else in the configuration file in a logging entry in the *level=Value* field. If no debug level is specified, the default level of 0 will be used. If logging is directed from the **snmpd** configuration file, the debug level can change during execution of the **snmpd** daemon.

The **SIGUSR1** signal (**kill –30**) and the SRC **traceson** and **tracesoff** commands have no effect on logging by the **syslogd** daemon.

No logging will take place if the **snmpd** daemon or the **syslogd** daemon are not configured to log messages as directed by the **snmpd** command line or configuration file.

Problem Determination for the SNMP Daemon

If the **snmpd** agent is not behaving as you expected, following are some hints to help you determine and correct the problem. If you are having problems, it is strongly recommended that you start up the **snmpd** agent with some type of logging. If you are having trouble with the invocation of the **snmpd** daemon, it is strongly recommended that you set up the **syslogd** daemon for logging at the daemon facility and DEBUG severity level. See "SNMP Daemon Logging Facility", on page 9-28, the **snmpd** command, and the **snmpd.conf** file for more information on **snmpd** logging.

Daemon Termination Problem

If the **snmpd** daemon terminates as soon as it is invoked, the following are possible reasons for failure and probable solutions:

 The reason the snmpd daemon terminated should be logged in the snmpd log file or the configured syslogd log file. Check the log file to see the FATAL error message.

Solution: Correct the problem and restart the **snmpd** daemon.

The snmpd command line usage was incorrect. If you invoked the snmpd command
without the System Resource Controller (SRC), the required usage statement should be
echoed to the screen. If you invoked the snmpd daemon under SRC control, the usage
message will not be echoed to the screen. You will need to check the log file to see the
usage message.

Solution: Invoke the **snmpd** command with the correct usage statement.

 The **snmpd** daemon must be invoked by the root user. The **snmpd** agent will not run if it is not invoked by the root user.

Solution: Switch to the root user and restart the **snmpd** daemon.

The snmpd.conf file must be owned by the root user. The snmpd agent checks the
ownership of the configuration file. If the file is not owned by the root user, the snmpd
agent terminates as this is considered a fatal error.

Solution: Make sure you are the root user, change the ownership of the configuration file to the root user, and restart the **snmpd** daemon.

• The snmpd.conf file must exist. If you did not specify the configuration file on the snmpd command line with the -c flag, the /etc/snmpd.conf file must exist. If you have accidentally removed the /etc/snmpd.conf file, reinstall the bos.net.tcp.client image or else reconstruct the file with the appropriate configuration entries as defined in the snmpd.conf file man page. If you did indeed specify the configuration file on the snmpd command line with the -c flag, make sure that the file exists and that the file is owned by the root user. You must specify the full path and file name of the configuration file if you do not wish to use the default /etc/snmpd.conf file.

Solution: Make sure the specified configuration file exists and that this file is owned by the root user. Restart the **snmpd** daemon.

 The udp port 161 is already bound. Make sure that the snmpd daemon is not already running. Issue the ps -eaf | grep snmpd command to determine if an snmpd daemon process is already executing. Only one snmpd agent can bind to udp port 161.

Solution: Either kill the existing **snmpd** agent or do not try to start up another **snmpd** daemon process.

Daemon Failure Problem

If the **snmpd** daemon fails when you issue a **refresh** or a **kill –1** signal, the following are possible reasons for failure and probable solutions:

- The reason the snmpd daemon terminated should be logged in the snmpd log file or the configured syslogd log file. Check the log file to see the FATAL error message.
 - Solution: Correct the problem and restart the **snmpd** daemon.
- Make sure that you specified the complete path and file name of the configuration file
 when you invoked the **snmpd** daemon. The **snmpd** daemon forks and changes to the
 root directory at invocation. If you did not specify the complete path name of the
 configuration file, the **snmpd** agent cannot find the file on a refresh. This is a fatal error
 and will cause the **snmpd** agent to terminate.
 - *Solution*: Specify the complete path and file name of the **snmpd** configuration file. Make sure the configuration file is owned by the root user. Restart the **snmpd** daemon.
- Make sure that the snmpd configuration file still exists. The file may have been
 accidentally removed after the snmpd agent was invoked. If the snmpd agent cannot
 open the configuration file, the snmpd agent terminates.
 - *Solution*: Recreate the **snmpd** configuration file, make sure the configuration file is owned by the root user, and restart the **snmpd** daemon.

MIB Variable Access Problem

If Management Information Base (MIB) variables cannot be accessed from the **snmpd** agent; if the **snmpd** agent is running, but the Simple Network Management Protocol (SNMP) manager application times out waiting for a response from the **snmpd** agent, try the following:

Check the network configuration of the host on which the snmpd agent is running using
the netstat -in command. Make sure the lo0, loopback, device is up. If the device is
down, there will be an * (asterisk) to the left of the lo0. The lo0 must be up for the snmpd
agent to service requests.

Solution: Issue the following command to start up the loopback interface:

```
ifconfig lo0 inet up
```

 Make sure that the snmpd daemon has a route to the host where you issued the requests.

Solution: On the host where the **snmpd** daemon is running, add a route to the host where the SNMP request issued the **route add** command. See the **route** command for more information.

- Check to see that the hostname and the host IP address are the same value.
 - Solution: Reset the hostname to correspond to the host IP address.
- Check to see that localhost is defined to be the lo0 IP address.

Solution: Define *localhost* to be the same address used by the lo0 IP address (usually 127.0.0.1).

MIB Variable Access in Community Entry Problem

If a community entry is specified in the configuration file with a MIB view name, but MIB variables cannot be accessed, check the following:

Make sure that you have correctly specified the community entry. If you have specified a
view name in the community entry, all fields in the community are absolutely required.

Solution: Specify all fields in the community entry in the configuration file. Refresh the **snmpd** agent and try your request again.

- Make sure the access mode in the community entry corresponds with your request type.
 If you are issuing a get or get—next request, make sure that the community has
 read—only or read—write permission. If you are issuing a set request, make sure that the
 community has read—write permission.
 - *Solution*: Specify the correct access mode in the community entry. Refresh the **snmpd** agent and try your request again.
- Make sure that you have specified a view entry for the specified view name in the
 community entry in the configuration file. If you have indeed specified a view name in the
 community entry, but you have no corresponding view entry, the **snmpd** agent will not
 allow access for that community. A view entry is absolutely required for a view name
 specified in a community entry in the configuration file.
 - *Solution*: Specify a view entry for the view name specified in the community entry. Refresh the **snmpd** agent and try your request again.
- If you have specified iso as the MIB subtree for your view entry, make sure that you
 have specified iso.3. The instance of 3 is required for the snmpd agent to access the
 org portion of the iso tree.
 - Solution: Specify the MIB subtree as iso.3 in the view entry. Refresh the **snmpd** agent and try your request again.
- Check the IP address and network mask in the community entry. Make sure that the host from which you are issuing the SNMP request is included in the community you are specifying with the community name.
 - Solution: Change the *IP address* and *network mask* fields in the community entry in the configuration file to include the host from which you are issuing the SNMP request.

No Response from Agent Problem

If the *IP address* in the community is specified as 0.0.0.0, but there is no response from the **snmpd** agent, try the following:

- Check the network mask field in the community entry. If you want general access for this
 community name, the network mask must be 0.0.0.0. If you have specified the network
 mask to be 255.255.255.255, you have configured the snmpd agent to not allow any
 requests with the specified community name.
 - *Solution*: Specify the *network mask* in the community entry to 0.0.0.0. Refresh the **snmpd** agent and try your request again.
- Make sure the access mode in the community entry corresponds with your request type.
 If you are issuing a get or get—next request, make sure that the community has
 read—only or read—write permission. If you are issuing a set request, make sure that the
 community has read—write permission.
 - *Solution*: specify the correct access mode in the community entry. Refresh the **snmpd** agent and try your request again.

noSuchName Problem

If in attempting to set an MIB variable that the **snmpd** agent is supposed to support, **noSuchName** error message is returned, the following may be the reason:

The set request you issued did not include a community name for a valid community with write access. The SNMP protocol dictates that a set request with a community with inappropriate access privileges be answered with the **noSuchName** error message.

Solution: Issue the **set** request with a community name for a community that has write privileges and includes the host from which the **set** request is issued.

Chapter 10. Network File System

This chapter provides information on the Network File System (NFS), a mechanism for storing files on a network. The following topics are discussed:

- Network File System Overview, on page 10-2
- NFS Installation and Configuration, on page 10-10
- PC-NFS, on page 10-19
- WebNFS, on page 10-21
- Network Lock Manager, on page 10-22
- Secure NFS, on page 10-25
- NFS Problem Determination, on page 10-34
- NFS Reference, on page 10-42

Network File System Overview

The Network File System (NFS) is a distributed file system that allows users to access files and directories located on remote computers and treat those files and directories as if they were local. For example, users can use operating system commands to create, remove, read, write, and set file attributes for remote files and directories.

The NFS software package includes commands and daemons for NFS, Network Information Service (NIS), and other services. Although NFS and NIS are installed together as one package, each is independent and each is configured and administered individually. See AIX 4.3 NIS/NIS+ Guide for details on NIS and NIS+.

AIX supports the latest NFS protocol update, NFS Version 3. AIX also provides an NFS Version 2 client and server and is, therefore, backward compatible with an existing install base of NFS clients and servers.

The following topics are discussed in this section:

- NFS Services, on page 10-2
- NFS Access Control Lists (ACL) Support, on page 10-3
- Cache File System Support, on page 10-3
- NFS Mapped File Support, on page 10-4
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NFS Services

NFS provides its services through a client–server relationship. The computers that make their *file systems*, or *directories*, and other resources available for remote access are called *servers*. The act of making file systems available is called *exporting*. The computers, or the processes they run, that use a server's resources are considered *clients*. Once a client *mounts* a file system that a server exports, the client can access the individual server files (access to exported directories can be restricted to specific clients).

The major services provided by NFS are:

Mount service	From the /usr/sbin/rpc.mountd daemon on the server and the /usr/sbin/mount command on the client.
Remote File access	From the /usr/sbin/nfsd daemon on the server and the /usr/sbin/biod daemon on the client.
Remote execution service	From the /usr/sbin/rpc.rexd daemon on the server and the /usr/bin/on command on the client.
Remote System Statistics service	From the /usr/sbin/rpc.rstatd daemon on the server and the /usr/bin/rup command on the client.
Remote User Listing service	From the /usr/lib/netsvc/rusers/rpc.rusersd daemon on the server and the /usr/bin/rusers command on the client.
Boot Parameters service	Provides boot parameters to SunOS diskless clients from the /usr/sbin/rpc.bootparamd daemon on the server.
Remote Wall service	From the /usr/lib/netsvc/rwall/rpc.rwalld daemon on the server and the /usr/sbin/rwall command on the client.

Spray service Sends a one—way stream of Remote Procedure Call (RPC)

packets from the /usr/lib/netsvc/spray/rpc.sprayd daemon on the server and the /usr/sbin/spray command

on the client.

PC authentication service Provides a user authentication service for PC–NFS from the

/usr/sbin/rpc.pcnfsd daemon on the server.

Note: A computer can be both an NFS server and an NFS client simultaneously.

An NFS server is *stateless*. That is, an NFS server does not have to remember any transaction information about its clients. In other words, NFS transactions are atomic: a single NFS transaction corresponds to a single, complete file operation. NFS requires the client to remember any information needed for later NFS use.

NFS Access Control Lists (ACL) Support

Although NFS supports access control lists (ACLs), they are no longer used as the default. If you want to use access control lists with NFS, use the **acl** option with the NFS **–o** flag, as shown in the following example:

```
mount -o acl
```

This support is handled by an RPC program that exchanges information about ACLs between clients and servers. The ACL support does not change the NFS protocol specification; it is a separate function.

The operating system adds ACLs to the regular file system. Since the normal NFS protocol does not support them, ACLs cannot be seen by normal NFS clients. Unexpected behavior results. A user on an NFS client might presume access to a file after looking at the permission bits, but the permissions could have been altered by the ACL associated with the file. Permissions on a server are enforced at the server according to the ACL on the server, so a user on the client machine could receive a permissions error.

When a client first attempts to access a remote mounted file system, it attempts to contact the ACL RPC program on the server.

If the server is a Version 3.2 server, the client consults the ACL associated with a file before granting access to the program on the client. This provides the expected behavior on the client when the request is sent over to the server. In addition, the **aclget**, **aclput**, and **alcedit** commands can be used on the client to manipulate ACLs.

Note: NFS no longer uses access control lists as the default.

Cache File System (CacheFS) Support

The Cache File System (CacheFS) is a general—purpose file system caching mechanism that improves NFS server performance and scalability by reducing server and network load. Designed as a layered file system, CacheFS provides the ability to cache one file system on another. In an NFS environment, CacheFS increases the client—per—server ratio, reduces server and network loads and improves performance for clients on slow links, such as Point—to—Point Protocol (PPP).

You create a cache on the client so file systems that you specify to be mounted in the cache can be accessed by the user locally instead of across the network. When a user first requests access to these files, they are placed in the cache. The cache does not get filled until the user requests access to a file or files. Initial file requests may seem slow, but subsequent uses of the same file(s) are faster.

Note:

- 1. You cannot cache the / (root) or /usr file systems.
- 2. You can mount only file systems that are shared. (See the **exportfs** command.)
- 3. There is no performance gain in caching a local Journaled File System (JFS) disk file system.

4. You must have root or system authority to do the tasks in the following table.

CacheFS Tasks			
Web-based System Manager: wsm network fast path (wsm network application)			
-OR-			
Task	SMIT Fast Path	Command or File	
Set up a cache	cachefs_admin_create	cfsadmin –c MountDirectoryName ¹	
Specifying Files for Mounting	cachefs_mount	mount –F cachefs –o backfstype=FileSysType, cachedir=CacheDirectory[,o ptions] BackFileSystem MountDirectoryName ² or edit /etc/filesystems	
Modify the Cache	cachefs_admin_change	remove the cache, then recreate it using appropriate mount command options	
Display Cache Information	cachefs_admin_change	cfsadmin –I MountDirectoryName	
Remove a Cache	cachefs_admin_remove	Unmount the file system: umount MountDirectoryName	
		2. Determine the cache ID: cfsadmin –I MountDirectoryName	
		3. Delete the file system: cfsadmin –d CachelD CacheDirectory	
Check File System Integrity	cachefs_admin_check	fsck_cachefsCacheDirector y ³	

Note:

- 1. After you have created the cache, do not perform any operations within the cache directory (**cachedir**) itself. This causes conflicts within the CacheFS software.
- 2. If you use the **mount** command option to specify files for mounting, the command must be reissued each time the system is rebooted.
- 3. Use the **-m** or **-o** options of the **fsck_cachefs** command to check the file systems without making any repairs.

NFS Mapped File Support

NFS mapped file support allows programs on a client to access a file as though it were memory. Using the **shmat** subroutine, users can map areas of a file into their address space. As a program reads and writes into this region of memory, the file is read into memory from the server or updated as needed on the server.

Mapping files over NFS is limited in three ways:

- Files do not share information well between clients.
- Changes to a file on one client using a mapped file are not seen on another client.

 Locking and unlocking regions of a file is not an effective way to coordinate data between clients.

If an NFS file is to be used for data sharing between programs on different clients, record locking and the regular **read** and **write** subroutines should be used.

Multiple programs on the same client can share data effectively using a mapped file. Advisory record locking can coordinate updates to the file on the client, provided that the entire file is locked. Multiple clients can share data—using mapped files only if the data never changes, as in a static database.

Three Types of Mounts

There are three types of NFS mounts: predefined, explicit, and automatic.

Predefined mounts are specified in the /etc/filesystems file. Each stanza (or entry) in this file defines the characteristics of a mount. Data such as the host name, remote path, local path, and any mount options are listed in this stanza. Predefined mounts should be used when certain mounts are always required for proper operation of a client.

Explicit mounts serve the needs of the root user. Explicit mounts are usually done for short periods of time when there is a requirement for occasional unplanned mounts. Explicit mounts can also be used if a mount is required for special tasks and that mount should not be generally available on the NFS client. These mounts are usually fully qualified on the command line by using the **mount** command with all needed information. Explicit mounts do not require updating the /etc/filesystems file. File systems mounted explicitly remain mounted unless explicitly unmounted with the **umount** command or until the system is restarted.

Automatic mounts are controlled by the **automount** command, which causes the **AutoFS** kernel extension to monitor specified directories for activity. If a program or user attempts to access a directory that is not currently mounted, then **AutoFS** intercepts the request, arranges for the mount of the file system, then services the request.

NFS Mounting Process

Clients access files on the server by first mounting a server's exported directories. When a client mounts a directory, it does not make a copy of that directory. Rather, the mounting process uses a series of remote procedure calls to enable a client to access the directories on the server transparently. The following describes the mounting process:

- When the server starts, the /etc/rc.nfs script runs the exportfs command, which reads
 the server /etc/exports file, and then tells the kernel which directories are to be exported
 and which access restrictions they require.
- 2. The **rpc.mountd** daemon and several **nfsd** daemons (8, by default) are then started by the **/etc/rc.nfs** script.
- When the client starts, the /etc/rc.nfs script starts several biod daemons (8, by default), which forward client mount requests to the appropriate server.
- 4. Then the /etc/rc.nfs script executes the mount command, which reads the file systems listed in the /etc/filesystems file.
- The mount command locates one or more servers that export the information the client wants and sets up communication between itself and that server. This process is called binding.
- 6. The **mount** command then requests that one or more servers allow the client to access the directories in the client /etc/filesystems file.
- 7. The server rpc.mountd daemon receives the client mount requests and either grants or denies them. If the requested directory is available to that client, the rpc.mountd daemon sends the client's kernel an identifier called a *file handle*.

8. The client kernel then ties the file handle to the mount point (a directory) by recording certain information in a *mount record*.

Once the file system is mounted, the client can perform file operations. When the client does a file operation, the **biod** daemon sends the file handle to the server, where the file is read by one of the **nfsd** daemons to process the file request. Assuming the client has access to perform the requested file operation, the **nfsd** daemon returns the necessary information to the client's **biod** daemon.

/etc/exports File

The /etc/exports file indicates all directories that a server exports to its clients. Each line in the file specifies a single directory. The server automatically exports the listed directories each time the NFS server is started. These exported directories can then be mounted by clients. The syntax of a line in the /etc/exports file is:

```
directory -options[,option]
```

The directory is the full path name of the directory. Options may designate a simple flag such as ro or a list of host names. See the specific documentation of the /etc/exports file and the exports command for a complete list of options and their descriptions. The /etc/rc.nfs script does not start the nfsd daemons or the rpc.mountd daemon if the /etc/exports file does not exist.

The following example illustrates entries from an /etc/exports file:

```
/usr/games -ro,access=ballet:jazz:tap
/home -root=ballet,access=ballet
/var/tmp
/usr/lib -access=clients
```

The first entry in this example specifies that the /usr/games directory can be mounted by the systems named ballet, jazz, and tap. These systems can read data and run programs from the directory, but they cannot write in the directory.

The second entry in this example specifies that the / home directory can be mounted by the system ballet and that root access is allowed for the directory.

The third entry in this example specifies that any client can mount the /var/tmp directory. (Notice the absence of an access list.)

The fourth entry in this example specifies an access list designated by the netgroup clients. In other words, these machines designated as belonging to the netgroup clients can mount the /usr/lib directory from this server. (A *netgroup* is a network—wide group allowed access to certain network resources for security or organizational purposes. Netgroups are controlled by using NIS or NIS+. For more information, see *AIX 4.3 NIS/NIS+ Guide*.)

/etc/xtab File

The /etc/xtab file has a format identical to the /etc/exports file and lists the currently exported directories. Whenever the exportfs command is executed, the /etc/xtab file changes. This allows you to export a directory temporarily without having to change the /etc/exports file. If the temporarily exported directory is unexported, the directory is removed from the /etc/xtab file.

Note: The /etc/xtab file is updated automatically, and should not be edited.

Implementation of NFS

NFS can be, and is, implemented on a wide variety of machine types, operating systems, and network architectures. NFS achieves this independence using the *Remote Procedure Call* (RPC) protocol.

Remote Procedure Call (RPC) Protocol

RPC is a library of procedures. The procedures allow one process (the client process) to direct another process (the server process) to execute procedure calls as if the client process had executed the calls in its own address space. Because the client and the server are two separate processes, they need not exist on the same physical system (although they can).

NFS implemented as a set of RPC calls in which the server services certain types of calls made by the client. The client makes such calls based on the file system operations that are done by the client process. NFS, in this sense, is an RPC application.

Because the server and client processes can reside on the two different physical systems which may have completely different architectures, RPC must address the possibility that the two systems may not represent data in the same manner. So, RPC uses data types defined by the eXternal Data Representation (XDR) protocol.

eXternal Data Representation (XDR) Protocol

XDR is the specification for a standard representation of various data types. By using a standard data type representation, a program can be confident that it is interpreting data correctly, even if the source of the data is a machine with a completely different architecture.

In practice, most programs do not use XDR internally. Rather, they use the data type representation specific to the architecture of the computer on which the program is running. When the program needs to communicate with another program, it converts its data into XDR format before sending the data. Conversely, when it receives data, it converts the data from XDR format into its own specific data type representation.

The portmap Daemon

Each RPC application has associated with it a program number and a version number. These numbers are used to communicate with a server application on a system. The client, when making a request from a server, needs to know what port number that server is accepting requests on. This port number is associated with the User Datagram Protocol (UDP) or Transmission Control Protocol (TCP) that is being used by the service. The client knows the program number, the version number, and the system name or host name where the service resides. The client needs a way to map the program number and version number pair to the port number of the server application. This is done with the help of the **portmap** daemon.

The **portmap** daemon runs on the same system as the NFS application. When the server starts running it, registers with the **portmap** daemon. As a function of this registration, the server supplies its program number, version number, and UDP or TCP port number. The **portmap** daemon keeps a table of server applications. When the client tries to make a request of the server, it first contacts the **portmap** daemon (on a well–known port) to find out which port the server is using. The **portmap** daemon responds to the client with the port of the server that the client is requesting. The client, upon receipt of the port number, is able to make all of its future requests directly to the server application.

Controlling NFS

The NFS, NIS, and NIS+ daemons are controlled by the System Resource Controller (SRC). This means you must use SRC commands such as **startsrc**, **stopsrc**, and **Issrc** to start, stop, and check the status of the NFS, NIS, and NIS+ daemons.

Some NFS daemons are not controlled by the SRC: specifically, **rpc.rexd**, **rpc.rusersd**, **rpc.rwalld**, and **rpc.rsprayd**. These daemons are started and stopped by the **inetd** daemon.

The following table lists the SRC-controlled daemons and their subsystem names.

Daemons and Their Subsystems		
File Path	Subsystem Name	Group Name
/usr/sbin/nfsd	nfsd	nfs
/usr/sbin/biod	biod	nfs
/usr/sbin/rpc.lockd	rpc.lockd	nfs
/usr/sbin/rpc.statd	rpc.statd	nfs
/usr/sbin/rpc.mountd	rpc.mountd	nfs
/usr/lib/netsvc/yp/ypserv	ypserv	ур
/usr/lib/netsvc/yp/ypbind	ypbind	ур
/usr/lib/netsvc/rpc.yppasswdd	yppasswdd	ур
/usr/lib/netsvc/rpc.ypupdated	ypupdated	ур
/usr/sbin/keyserv	keyserv	keyserv
/usr/sbin/portmap	portmap	portmap

NIS+ daemons are described in *AIX 4.3 NIS/NIS+ Guide*. Each of these daemons can be specified to the SRC commands by using either their subsystem name or the appropriate group name. These daemons support neither the long–listing facility of SRC nor the SRC trace commands.

For more information on using the SRC, see "System Resource Controller Overview" in AIX 4.3 System Management Concepts: Operating System and Devices.

Change the Number of biod and nfsd Daemons

To change the number of **biod** or **nfsd** daemons running on the system, use the **chnfs** command. For example, to set the number of **nfsd** daemons to 10 and the number **biod** daemons to 4, run the command:

This command temporarily stops the daemons currently running on the system, modifies the SRC database code to reflect the new number, and restarts the daemons.

Note: In the AIX NFS implementation, the number of **biod** daemons are controllable only per mount point using the **biod** –**o** option. Specification using **chnfs** is retained for compatibility purposes only and has no real effect on the number of threads performing I/O.

Change Command Line Arguments for Daemons Controlled by SRC

Many NFS, NIS, and NIS+ daemons have command–line arguments that can be specified when the daemon is started. Since these daemons are not started directly from the command line, you must update the SRC database so that the daemons can be started correctly. To do this, use the **chssys** command. The **chssys** command has the format:

```
chssys -s Daemon -a 'NewParameter'
For example:
```

chssys -s nfsd -a '10'

changes the **nfsd** subsystem so that when the daemon is started, the command line looks like nfsd 10. The changes made by the **chssys** command do not take effect until the subsystem has been stopped and restarted.

Start the NFS Daemons at System Startup

The NFS daemons, by default, are not started during installation. When installed, all of the files are placed on the system, but the steps to activate NFS are not taken. You can start the NFS daemons at system startup through:

- The Web-based System Manager fast path, wsm network
- The SMIT fast path, smit mknfs
- The mknfs command

All of these methods place an entry in the **inittab** file so that the /etc/rc.nfs script is run each time the system restarts. This script, in turn, starts all NFS daemons required for a particular system.

Start the NFS Daemons

The file size limit for files located on an NFS server is taken from the process environment when **nfsd** is started. To use a specific value, edit the file /**etc/rc.nfs** and add a ulimit, using the **ulimit** command, with the desired limit before the **startsrc** command for **nfsd**.

The NFS daemons can be started individually or all at once. To start NFS daemons individually:

```
startsrc -s Daemon
```

where <code>Daemon</code> is anyone of the SRC–controlled daemons. For example, to start the **nfsd** daemons:

```
startsrc -s nfsd
```

To start all of the NFS daemons:

```
startsrc -q nfs
```

Note: If the /etc/exports file does not exist, the nfsd and the rpc.mountd daemons will not be started. You can create an empty /etc/exports file by running the command touch /etc/exports. This will allow the nfsd and the rpc.mountd daemons to start, although no file systems will be exported.

Stop the NFS Daemons

The NFS daemons can be stopped individually or all at once. To stop NFS daemons individually:

```
stopsrc -s Daemon
```

where <code>Daemon</code> is anyone of the SRC–controlled daemons. For example, to stop the **rpc.lockd** daemon:

```
stopsrc -s rpc.lockd
```

To stop all NFS daemons at once:

```
stopsrc -g nfs
```

Get the Current Status of the NFS Daemons

You can get the current status of the NFS daemons individually or all at once. To get the current status of the NFS daemons individually:

```
lssrc -s Daemon
```

where <code>Daemon</code> is anyone of the SRC–controlled daemons. For example, to get the current status of the <code>rpc.lockd</code> daemon:

```
lssrc -s rpc.lockd
```

To get the current status of all NFS daemons at once:

```
lssrc -a
```

NFS Installation and Configuration

For information on installing the Network File System (NFS), see the AIX Installation Guide.

Checklist for Configuring NFS

Once the NFS software is installed on your systems, you are ready to configure NFS.

- 1. Determine which systems in the network will be servers and which will be clients (a system can be configured as both a server and a client).
- 2. For each system (whether client or server), follow the instructions in "Start the NFS Daemons at System Startup", on page 10-9.
- 3. For each NFS server, follow the instructions in "Configuring an NFS Server", on page 10-10.
- 4. For each NFS client, follow the instructions in "Configuring an NFS Client", on page 10-10.
- If you want personal computers on your network to have access to your NFS servers (beyond being able to mount file systems), configure PC–NFS by following the instructions in "PC–NFS", on page 10-19.

Configuring an NFS Server

To configure an NFS server:

- 1. Start NFS using the instructions in "Start the NFS Daemons Using SRC", on page 10-9.
- 2. Create the /etc/exports file.

Configuring an NFS Client

 Verify that NFS is the default remote file system. (If this is not done, you will need to specify the -v nfs flag when using the mount command.) Using a text editor, open the /etc/vfs file and search for the following entry:

```
#%defaultvfs jfs nfs
#nfs 2 /sbin/helpers/nfsmnthelp none remote
```

If pound signs (#) appear at the beginning of the lines as shown, delete them.

- 2. Start NFS using the instructions in "Start the NFS Daemons", on page 10-9.
- 3. Establish the local mount point using the **mkdir** command. For NFS to complete a mount successfully, a directory that acts as the mount point (or place holder) of an NFS mount must be present. This directory should be empty. This mount point can be created like any other directory, and no special attributes are needed for this directory.

Note: The mount points for all NFS mounts must exist on your system before you can mount a file system, with one exception. If the **automount** daemon is used it may not be necessary to create mount points. See the **automount** documentation for details.

4. Establish and mount the predefined mounts by following the instructions in "Establishing Predefined NFS Mounts", on page 10-15.

Exporting an NFS File System

You can export an NFS file system using the Web-based System Manager wsm network application, or you can use one of the following procedures.

- To export an NFS file system using the System Management Interface Tool (SMIT):
 - a. Verify that NFS is already running by entering the command lssrc -g nfs. The output should indicate that the **nfsd** and the **rpc.mountd** daemons are active. If they are not, start NFS using the instructions in "Start the NFS Daemons", on page 10-9.
 - b. Use

```
smit mknfsexp
```

- c. Specify appropriate values in the PATHNAME of directory to export, MODE to export directory, and EXPORT directory now, system restart or both fields.
- d. Specify any other optional characteristics you want, or accept the default values by leaving the remaining fields as they are.
- e. When you have finished making your changes, SMIT updates the /etc/exports file. If the /etc/exports file does not exist, then it will be created.
- f. Repeat steps 3 through 5 for each directory you want to export.
- To export an NFS file system using a text editor:
 - a. Open the /etc/exports file with your favorite text editor.
 - b. Create an entry for each directory to be exported, using the full path name of the directory. List each directory to be exported starting in the left margin. No directory should include any other directory that is already exported. See the /etc/exports file documentation for a description of the full syntax for entries in the /etc/exports file.
 - c. Save and close the /etc/exports file.
 - d. If NFS is currently running, enter:

```
/usr/sbin/exportfs -a
```

The **–a** option tells the **exportfs** command to send all information in the **/etc/exports** file to the kernel. If NFS is not running, start NFS using the instructions in "Start the NFS Daemons", on page 10-9.

• To export an NFS file system temporarily (in other words, without changing the /etc/exports file), enter:

```
exportfs -i /dirname
```

where <code>dirname</code> is the name of the file system you want to export. The **exports –i** command specifies that the <code>/etc/exports</code> file is not to be checked for the specified directory, and all options are taken directly from the command line.

Unexporting an NFS File System

You can unexport an NFS directory using the Web-based System Manager wsm network application, or you can use one of the following procedures.

- To unexport an NFS directory using SMIT:
 - a. Enter:

```
smit rmnfsexp
```

b. Enter the appropriate path name in the PATHNAME of exported directory to be removed field.

The directory is now removed from the /etc/exports file and is unexported.

- To unexport an NFS directory by using a text editor:
 - a. Open the /etc/exports file with your favorite text editor.
 - b. Find the entry for the directory you wish to unexport, and the delete that line.
 - c. Save and close the /etc/exports file.
 - d. If NFS is currently running, enter:

```
exportfs -u dirname
```

where *dirname* is the full path name of the directory you just deleted from the /etc/exports file.

Changing an Exported File System

You can change an exported NFS file system using the Web-based System Manager wsm network application, or you can use one of the following procedures.

- To change an exported NFS file system using SMIT:
 - a. Unexport the file system by entering:

```
exportfs -u /dirname
```

where dirname is the name of the file system you want to change.

b. Enter:

```
smit chnfsexp
```

- c. Enter the appropriate path name in the PATHNAME of exported directory field.
- d. Make whatever changes you want.
- e. Exit SMIT.
- f. Re-export the file system by entering:

```
exportfs /dirname
```

where dirname is the name of the file system you just changed.

- To change an exported NFS file system by using a text editor:
 - a. Unexport the file system by entering:

```
exportfs -u /dirname
```

where dirname is the name of the file system you want to change.

- b. Open the /etc/exports file with your favorite text editor.
- c. Make whatever changes you want.
- d. Save and close the /etc/exports file.
- e. Re-export the file system by entering:

```
exportfs /dirname
```

where dirname is the name of the file system you just changed.

Enabling Root User Access to an Exported File System

When you export a file system, you can enable the root user on a particular machine to have root access to that file system. By default, the root user is not granted root access to exported file systems. When a person logged in as the root user on one host requests access to a particular file from NFS, the user ID of the requester is mapped by NFS to the user ID of user nobody (nobody is one of the user names placed in the /etc/password file by default). The access rights of user nobody are the same as those given to the public (others) for a particular file. For example, if others only execute permission for a file, then user nobody can only execute the file.

To enable root user access to an exported file system, follow the instructions in "Changing an Exported File System", on page 10-12. If you use the Web-based System Manager or SMIT method, specify in the HOSTS allowed root access field the name of the host to which you want to grant root access. If you edit the file with a text editor, add the qualifier -root=hostname to the file system entry. For example,

```
/usr/tps -root=hermes
```

specifies that the root user on host hermes may access the /usr/tps directory with root privileges.

Mounting an NFS File System Explicitly

To mount an NFS directory explicitly, use the Web-based System Manager fast path, **wsm network**, or use the following procedure:

1. Verify that the NFS server has exported the directory:

```
showmount -e ServerName
```

where ServerName is the name of the NFS server. This command displays the names of the directories currently exported from the NFS server. If the directory you want to mount is not listed, export the directory from the server.

- Establish the local mount point using the **mkdir** command. For NFS to complete a mount successfully, a directory that acts as the mount point (or place holder) of an NFS mount must be present. This directory should be empty. This mount point can be created like any other directory, and no special attributes are needed for this directory.
- 3. Enter:

```
mount ServerName:/remote/directory /local/directory
```

where <code>ServerName</code> is the name of the NFS server, <code>/remote/directory</code> is the directory on the NFS server you want to mount, and <code>/local/directory</code> is the mount point on the NFS client.

4. On the client machine, enter the following SMIT fast path:

```
smit mknfsmnt
```

5. Make changes to the following fields that are appropriate for your network configuration. Your configuration may not require completing all of the entries on this screen.

Note: If you are using the ASCII SMIT interface, press the Tab key to change to the correct value for each field, but do *not* press Enter until you get to step 7.

- PATHNAME of mount point.
- PATHNAME of remote directory.
- HOST where remote directory resides.
- MOUNT now, add entry to /etc/filesystems or both?
- /etc/filesystems entry will mount the directory on system RESTART.
- MODE for this NFS file system.
- 6. Use the default values for the remaining entries or change them, depending on your NFS configuration.
- 7. When you finish making all the changes on this screen, SMIT mounts the NFS file system.
- 8. When the Command: field shows the OK status, exit SMIT.

The NFS file system is now ready to use.

Automatic Mounting of a File System Using AutoFS

AutoFS relies on the use of the **automount** command to propagate the automatic mount configuration information to the AutoFS kernel extension and start the **automountd** daemon. Through this configuration propagation, the extension automatically and transparently mounts file systems whenever a file or a directory within that file system is opened. The extension informs the **autmountd** daemon of mount and unmount requests, and the **autmountd** daemon actually performs the requested service.

Because the name—to—location binding is dynamic within the **autmountd** daemon, updates to a Network Information Service (NIS) map used by the **autmountd** daemon are transparent to the user. Also, there is no need to premount shared file systems for applications that have hard—coded references to files and directories, nor is there a need to maintain records of which hosts must be mounted for particular applications.

AutoFS allows file systems to be mounted as needed. With this method of mounting directories, all file systems do not need to be mounted all of the time; only those being used are mounted.

For example, to mount an NFS directory automatically:

1. Verify that the NFS server has exported the directory by entering:

```
showmount -e ServerName
```

where ServerName is the name of the NFS server. This command displays the names of the directories currently exported from the NFS server.

2. Create an AutoFS map file. AutoFS will mount and unmount the directories specified in this map file. For example, suppose you want to use AutoFS to mount the /usr/local/dir1 and /usr/local/dir2 directories as needed from the serve1 server onto the /usr/remote/dir1 and /usr/remote/dir2 directories, respectively. In this example, the map file name is /tmp/mount.map.

```
dir1 -rw serve1:/usr/local/dir1 dir2 -rw serve1:/usr/local/dir2
```

- Ensure that the AutoFS kernel entension is loaded and the automountd daemon is running. This can be accomplished in two ways:
 - a. Using **SRC**: Issue **Issrc –s automountd**. If the **automountd** subsystem is not running, issue **startsrc –s automountd**.
 - b. Using the automount command: Issue /usr/bin/automount -v.

Define the map file using the command line interface by entering:

```
/usr/sbin.automount /usr/remote /tmp/mount.map
```

where /usr/remote is the **AutoFS** mount point on the client. Now, if a user runs the **cd** /usr/remote/dir1 command, the **AutoFS** kernel extension will intercept access to the directory and will issue a remote procedure call to the **automountd** daemon, which will mount the /usr/remote/dir1 directory and then allow the **cd** command to complete.

```
/usr/sbin/automount /usr/remote /tmp/mount.map
```

where /usr/remote is the mount point on the NFS client. Now, if a user runs the cd /usr/remote/dir1 command, the automount daemon will mount the /usr/remote/dir1 directory and then allow the cd command to complete.

To stop the automount daemon, issue the stopsrc –s automountd command.

If, for some reason, the automountd daemon was started without the use of SRC, issue:

```
kill automountd_PID
```

where *automountd_PID* is the process ID of the **automountd** daemon. (Running the **ps**—**e** command will display the process ID of the **automountd** daemon.) The **kill** command sends a SIGTERM signal to the **automountd** daemon.

Establishing Predefined NFS Mounts

You can establish predefined NFS mounts using the Web-based System Manager wsm network application, or you can use one of the following procedures.

Attention: Define the **bg** (background) and **intr** (interruptible) options in the /etc/filesystems file when establishing a predefined mount that will be mounted during system startup. Mounts that are noninterruptible and running in the foreground can hang the client if the network or server is down when the client system starts up. If a client cannot access the network or server, the user must start the machine again in maintenance mode and edit the appropriate mount requests.

- To establish predefined mounts through SMIT:
 - a. Enter:

```
smit mknfsmnt
```

- b. Specify values in this screen for each mount you want to predefine. You must specify a value for each required field (those marked with an asterisk (*) in the left margin). You may specify values for the other fields or accept their default values. This method creates an entry in the /etc/filesystems file for the desired mount and attempts the mount.
- To establish the NFS default mounts by editing the /etc/filesystems file:
 - a. Open the /etc/filesystems file with your favorite text editor.
 - b. Add entries for each of the remote file systems that you want mounted when the system is started. For example:

```
/home/jdoe:
dev = /home/jdoe
mount = false
vfs = nfs
nodename = mach2
options = ro, soft
type = nfs_mount
```

This stanza directs the system to mount the <code>/home/jdoe</code> remote directory over the local mount point of the same name. The file system is mounted as read—only (ro). Because it is also mounted as <code>soft</code>, an error is returned in the event the server does not respond. By specifying the <code>type</code> parameter as <code>nfs_mount</code>, the system attempts to mount the <code>/home/jdoe</code> file (along with any other file systems that are specified in the <code>type = nfs_mount</code> group) when the <code>mount -t</code> nfs <code>mount</code> command is issued.

The example stanza below directs the system to mount the /usr/games file system at system startup time. If the mount fails, the system continues to attempt to mount in the background.

```
/usr/games:
dev = /usr/games
mount = true
vfs = nfs
nodename = gameserver
options = ro,soft,bg
type = nfs_mount
```

The following parameters are required for stanzas pertaining to NFS mounts:

dev=filesystem_name	Specifies the path name of the remote file system being mounted.
mount=[true false]	If $true$, specifies that the NFS file system will be mounted when the system boots. If $false$, the NFS file system will not be mounted when the system boots.

nodename=hostname Specifies the host machine on which the remote file system

resides.

vfs=nfs Specifies that the virtual file system being mounted is an

NFS file system.

The following parameters are optional for stanzas pertaining to NFS mounts:

type=type_name

Defines the file system being mounted as part of the <code>type_name</code> mount group. This parameter is used with the **mount –t** command, which mounts groups of specified file systems at the same time.

options=options

Specifies one or more of the following options parameters:

biods=*N* Specifies the number of biod daemons to start.

The default is 6. N is an integer.

bg Specifies to try the mount again in the background

if the first mount attempt fails.

fg Specifies to try the mount again in the foreground if

the first mount attempt fails.

noacl Disables, for this mount only, the Access Control

List (ACL) support provided by the NFS journaled

file system.

When used between two systems, NFS supports access control lists. If the <code>noacl</code> option is used when mounting a file system, NFS does not use ACLs. The effect of the <code>noacl</code> option equals what happens when an NFS client on a system mounts from an NFS server that does not support ACLs.

For more information about ACLs, refer to "NFS Access Control List (ACL) Support", on page 10-3.

retry=*n* Sets the number of times to try the mount.

rsize=n Sets the read buffer size to the number of bytes

specified by n.

wsize=n Sets the write buffer size to the number of bytes

specified by n.

timeo=n Sets the NFS time out to the tenths of a second

specified by n. Use this variable to avoid situations that can occur in networks where the server load

can cause inadequate response time.

retrans=n Sets the number of NFS retransmissions to the

number specified by n.

port=n Sets the server port to the number specified by n.

soft Returns an error if the server does not respond.

hard Continues to try the request until the server

responds.

Note: When you specify a hard mount, it is possible that the process can hang while waiting for a response. To be able to interrupt the process and end it from the keyboard, use

the intr variable in the mount variables.

intr Allows keyboard interrupts on hard mounts.

ro Sets the read-only variable.

variable along with this variable to avoid error conditions that can conflict with applications if a soft mount is attempted as read–write. See "NFS Problem Determination", on page 10-34 for

Problem Determination", on page 10-34 for information on hard— and soft—mounted problems.

secure Specifies to use a more secure protocol for NFS

transactions.

actimeo=n Extends flush time by n seconds for both regular

files and directories.

Note: The attribute cache retains file attributes on the client. Attributes for a file are assigned a time to be erased. If the file is modified before the flush time, then the flush time is extended by the time since the previous modification (under the assumption that recently changed files are likely to change again soon). There are minimum and maximum flush time extensions for regular files and for directories.

acregmin=n Holds cached attributes for at least n seconds after

file modification.

acregmax=n Holds cached attributes for no more than n

seconds after file modification.

acdirmin=n Holds cached attributes for at least n seconds after

directory update.

acdirmax=n Holds cached attributes for no more than n

seconds after directory update.

Note: If you do not set the following options, the kernel automatically sets them to these default values:

biods=6
fg
retry=10000
rsize=8192
wsize=8192
timeo=7
retrans=3
port=NFS_PORT
hard
secure=off
acregmin=3
acregmax=60
acdirmin=30
acdirmax=60

- c. Remove any directory entries that you do not want to mount automatically at system startup.
- d. Save and close the file.
- e. Run the **mount –a** command to mount all the directories specified in the /etc/filesystems file.

Unmounting an Explicitly or Automatically Mounted File System

To unmount an explicitly or automatically mounted NFS directory, enter:

umount /directory/to/unmount

Removing Predefined NFS Mounts

You can remove a predefined NFS mount using the Web-based System Manager wsm network application, or you can use one of the following procedures.

- To remove a predefined NFS mount through SMIT:
 - a. enter:

smit rmnfsmnt

- To remove a predefined NFS mount by editing the /etc/filesystems file:
 - a. Enter the command: umount /directory/to/unmount.
 - b. Open the /etc/filesystems file with your favorite editor.
 - c. Find the entry for the directory you just unmounted, and then delete it.
 - d. Save and close the file.

PC-NFS

PC-NFS is a program for personal computers that enables the personal computer to mount file systems exported by a Network File System (NFS) server. The personal computer can also request network addresses and host names from the NFS server. Additionally, if the NFS server is running the **rpc.pcnfsd** daemon, the personal computer can access authentication and print-spooling services.

You may want to configure the **rpc.pcnfsd** daemon on the following:

- Systems that perform user authentication services
- · Systems that offer print-spooling
- All Network Information Service (NIS) master and slave servers.

Note: Because NIS networks are typically configured so that PC–NFS can pick any NIS server as the default server, it is important that all servers have the **rpc.pcnfsd** program running. If running this program on all NIS servers is not practical, or if you want to limit requests to a specific server, add a **net pcnfsd** command to the **autoexec.bat** file on each personal computer to force it to use a specific NIS server. For more information, see *AIX 4.3 NIS/NIS+ Guide*.

PC-NFS Authentication Service

By default, PC–NFS presents itself to NFS servers as the <code>nobody</code> user. With <code>nobody</code> privileges, all personal computer user files appear as owned by <code>nobody</code>, and consequently you cannot distinguish between different personal computer users. The authentication capability of the <code>rpc.pcnfsd</code> daemon allow you to monitor system resources and security by recognizing individual users and assigning them different privileges.

With the **rpc.pcnfsd** daemon running, a PC–NFS user can issue the **net name** command from a personal computer to log in to PC–NFS in the same manner as a user can log in to AIX. The user's name and password are verified by the **rpc.pcnfsd** daemon. This authentication procedure does not make a server more secure, but it does provide more control over access to files that are available through NFS.

PC-NFS Print-Spooling Service

The print–spooling service of the **rpc.pcnfsd** daemon enables personal computers running PC–NFS to print to printers not directly attached to the personal computer. Specifically, PC–NFS redirects files intended for personal computer printers to a file on an NFS server. This file is placed in a spool directory on the NFS server. The **rpc.pcnfsd** daemon then invokes the server's printing facility. (The spooling directory must be in an exported file system so that PC–NFS clients can mount it.) When PC–NFS requests that the **rpc.pcnfsd** daemon print the file, it provides the following information:

- · Name of the file to be printed
- Login ID of the user on the client
- Name of the printer to be used.

Configuring the rpc.pcnfsd Daemon

To configure the **rpc.pcnfsd** daemon:

- 1. Install PC-NFS program on your personal computer.
- 2. Select a location for the spool directory on the NFS server. The default spool directory is /var/tmp. The spool directory must have at least 100K bytes of free space.
- 3. Export the spool directory. Do not put access restrictions on the exported directory that could cause access problems in your network. For details of this procedure, see "Exporting an NFS File System", on page 10-11.

- 4. Start the **rpc.pcnfsd** daemon by following the instructions in "Starting the rpc.pcnfsd Daemon", on page 10-20.
- 5. Verify that the **rpc.pcnfsd** daemon is accessible by following the instructions in "Verifying the rpc.pcnfsd Daemon Is Accessible", on page 10-20.

Note: Because printer–redirection requests sometimes cause file listings of zero length to be left in the PC–NFS spool directories, you should periodically clear spooling directories of these entries.

Starting the rpc.pcnfsd Daemon

To start the **rpc.pcnfsd** daemon using the default spooling directory:

1. Using your favorite text editor, uncomment the following entry in the /etc/inetd.conf file:

```
pcnfsd sunrpc_udp udp wait root /usr/sbin/rpc.pcnfsd pcnfsd 150001 1
```

2. Save the file and exit the text editor.

To start the **rpc.pcnfsd** daemon using a directory different from the default:

1. Using your favorite text editor, add the following entry to the /etc/rc.nfs file:

```
if [ -f /usr/sbin/rpc.pcnfsd ] ; then
/usr/sbin/rpc.pcnfsd -s spooldir ; echo ' rpc.pcnfsd\c'
fi
```

where spooldir specifies the full path name of the spool directory.

- 2. Save the file and exit the text editor.
- 3. Using your favorite text editor, comment the following entry in the /etc/inetd.conf file:

```
pcnfsd sunrpc_udp udp wait root /usr/sbin/rpc.pcnfsd pcnfsd 150001 1
```

by placing a pound sign (#) at the beginning of the line. This prevents the **inetd** daemon from starting the **rpc.pcnfsd** daemon using the default spool directory.

4. Start the **rpc.pcnfsd** daemon print spooler by entering the following at the command line:

```
/usr/sbin/rpc.pcnfsd -s spooldir
```

where *spooldir* specifies the full path name of the spool directory.

For more information on updating the **inetd** configuration database, see "Configuring the inetd Daemon", on page 3-94.

Note: The default directory that the **rpc.pcnfsd** daemon uses cannot be changed from the **inetd.conf** file.

Verifying the rpc.pcnfsd Daemon Is Accessible

To verify that the **rpc.pcnfsd** daemon is accessible, enter:

```
rpcinfo -u host 150001
```

where *host* specifies the host name of the system on which you are configuring **rpc.pcnfsd**, and 15001 is the RPC program number of the **rpc.pcnfsd** daemon. After you enter the command, you should receive the message that the program is ready and waiting.

WebNFS

AIX provides NFS server capability for WebNFS. Defined by Sun Microsystems, WebNFS is a simple extension of the NFS protocol that allows easier access to servers and clients through Internet firewalls.

A WebNFS-enhanced web browser can use an NFS universal resource locator (URL) to access data directly from the server. An example NFS URL is:

```
nfs://www.YourCompany.com/
```

WebNFS works in tandem with existing web-based protocols to provide data to clients.

WebNFS also takes advantage of the scalability of NFS servers.

Network Lock Manager

The network lock manager is a facility that works in cooperation with the Network File System (NFS) to provide a System V style of advisory file and record locking over the network. The network lock manager (**rpc.lockd**) and the network status monitor (**rpc.statd**) are network—service daemons. The **rpc.statd** daemon is a user level process while the **rpc.lockd** daemon is implemented as a set of kernel threads (similar to the NFS server). Both daemons are essential to the kernel's ability to provide fundamental network services.

Note: Mandatory or enforced locks are not supported over NFS.

Network Lock Manager Architecture

The network lock manager contains both server and client functions. The client functions are responsible for processing requests from the applications and sending requests to the network lock manager at the server. The server functions are responsible for accepting lock requests from clients and generating the appropriate locking calls at the server. The server will then respond to the client's locking request.

In contrast to NFS, which is stateless, the network lock manager has an implicit state. In other words, the network lock manager must remember certain information about a client, that is, whether the client currently has a lock. The network status monitor, **rpc.statd**, implements a simple protocol that allows the network lock manager to monitor the status of other machines on the network. By having accurate status information, the network lock manager can maintain a consistent state within the stateless NFS environment.

Network File Locking Process

When an application wants to obtain a lock on a local file, it sends its request to the kernel using the **lockf**, **fcntl**, or **flock** subroutine. The kernel then processes the lock request. However, if an application on an NFS client makes a lock request for a remote file, the Network Lock Manager client will generate a Remote Procedure Call (RPC) to the server to handle the request.

When the client receives an initial remote lock request, it registers interest in the server with the client's **rpc.statd** daemon. The same is true for the network lock manager at the server. On the initial request from a client, it will register interest in the client with the local network status monitor.

Crash Recovery Process

Each machine's **rpc.statd** daemon notifies every other machine's **rpc.statd** daemon of its activities. When a machine's **rpc.statd** daemon receives notice that another machine crashed or recovered, it notifies its **rpc.lockd** daemon.

If a server crashes, clients with locked files must be able to recover their locks. If a client crashes, its servers must hold the client locks while it recovers. Additionally, to preserve the overall transparency of NFS, the crash recovery must occur without requiring the intervention of the applications themselves.

The crash recovery procedure is simple. If the failure of a client is detected, the server releases the failed client locks, on the assumption that the client application will request locks again as needed. If the crash and recovery of a server is detected, the client lock manager retransmits all lock requests previously granted by the server. This retransmitted information is used by the server to reconstruct its locking state during a grace period. (The grace period, 45 seconds by default, is a time period within which a server allows clients to reclaim their locks.)

The **rpc.statd** daemon uses the host names kept in /**etc/sm** and /**etc/sm.bak** to keep track of which hosts must be informed when the machine needs to recover operations.

Starting the Network Lock Manager

By default, the /etc/rc.nfs script starts the rpc.lockd and rpc.statd daemons along with the other NFS daemons. If NFS is already running, you can verify that the rpc.lockd and rpc.statd daemons are running by following the instructions in "Get the Current Status of the NFS Daemons", on page 10-9. The status of these two daemons should be active. If the rpc.lockd and rpc.statd daemons are not active, and therefore not running, do the following:

- 1. Using your favorite text editor, open the /etc/rc.nfs file.
- 2. Search for the following lines:

3. If there is a pound sign (#) at the beginning of any of these lines, delete it, then save and exit the file. Then start the **rpc.statd** and **rpc.lockd** daemons by following the instructions in "Start the NFS Daemons", on page 10-9.

Note: Sequence is important. Always start the statd daemon first.

4. If NFS is running and the entries in the /etc/rc.nfs file are correct, stop and restart the rpc.statd and rpc.lockd daemons by following the instructions in "Stop the NFS Daemons", on page 10-9 and "Start the NFS Daemons", on page 10-9.

Note: Sequence is important. Always start the statd daemon first.

If the **rpc.statd** and **rpc.lockd** daemons are still not running, see "Troubleshooting the Network Lock Manager."

Troubleshooting the Network Lock Manager

If you receive a message on a client similar to:

```
clnttcp_create: RPC: Remote System error - Connection refused
rpc.statd:cannot talk to statd at {server}
```

then the machine thinks there is another machine which needs to be informed that it may have to take recovery measures. When a machine reboots, or when **rpc.lockd** and **rpc.statd** are stopped and restarted, machine names are moved from /**etc/sm** to /**etc/sm.bak** and the **rpc.statd** tries to inform each machine corresponding to each entry in /**etc/sm.bak** that recovery procedures are needed.

If the **rpc.statd** can reach the machine, then its entry in /**etc/sm.bak** is removed. If **rpc.statd** cannot reach the machine, then it will keep trying at regular intervals. Each time the machine fails to respond, the timeout generates the above message. In the interest of locking integrity, the daemon will continue to try; however, this can have an adverse effect on locking performance. The handling is different, depending on whether the target machine is just unresponsive or semi–permanently taken out of production. To eliminate the message:

- 1. Verify that the **statd** and **lockd** daemons on the server are running by following the instructions in "Get the Current Status of the NFS Daemons", on page 10-9. (The status of these two daemons should be *active*.)
- 2. If these daemons are not running, start the **rpc.statd** and **rpc.lockd** daemons on the server by following the instructions in "Start the NFS Daemons", on page 10-9.

Note: Sequence is important. Always start the statd daemon first.

After you have restarted the daemons, remember that there is a grace period. During this time, the **lockd** daemons allow reclaim requests to come from other clients that

previously held locks with the server, so you will not get a new lock immediately after starting the daemons.

Alternatively, you can eliminate the message by:

- Stop the rpc.statd and rpc.lockd daemons on the client by following the instructions in "Stop the NFS Daemons", on page 10-9.
- 2. On the client, remove the target machine entry from /etc/sm.bak by entering:

```
rm /etc/sm.bak/TargetMachineName
```

This action will keep the target machine from being aware that it may need to participate in locking recovery, so it should only be used when it can be determined that the machine does not have any applications running that are participating in network locking with the affected machine.

3. Start the **rpc.statd** and **rpc.lockd** daemons on the client by following the instructions in "Start the NFS Daemons", on page 10-9.

If you are unable to obtain a lock from a client, do the following:

- 1. Use the **ping** command to verify that the client and server can reach and recognize each other. If the machines are both running and the network is intact, check the host names listed in the /etc/hosts file for each machine. Host names must exactly match between server and client for machine recognition. If a name server is being used for host name resolution, make sure the host information is exactly the same as that in the /etc/hosts file.
- Verify that the rpc.lockd and rpc.statd daemons are running on both the client and the server by following the instructions in "Get the Current Status of the NFS Daemons", on page 10-9. The status of these two daemons should be active.
- 3. If they are not active, start the **rpc.statd** and **rpc.lockd** daemons by following the instructions in "Start the NFS Daemons", on page 10-9.
- 4. If they are active, you may need to reset them on both clients and servers. To do this, stop all the applications that are requesting locks.
- 5. Next, stop the **rpc.statd** and **rpc.lockd** daemons on both the client and the server by following the instructions in "Stop the NFS Daemons", on page 10-9.
- 6. Now, restart the **rpc.statd** and **rpc.lockd** daemons, first on the server and then on the client, by following the instructions in "Start the NFS Daemons", on page 10-9.

Note: Sequence is important. Always start the **statd** daemon first.

If the procedure does not alleviate the locking problem, run the **lockd** daemon in debug mode, by doing the following:

- 1. Stop the **rpc.statd** and **rpc.lockd** daemons on both the client and the server by following the instructions in "Stop the NFS Daemons", on page 10-9.
- 2. Start the **rpc.statd** daemon on the client and server by following the instructions in "Start the NFS Daemons", on page 10-9.
- 3. Start the **rpc.lockd** daemon on the client and server be entering:

```
/usr/sbin/rpc.lockd -d1
```

When invoked with the **-d1** flag, the **lockd** daemon provides diagnostic messages to standard output. At first, there will be a number of messages dealing with the grace period; wait for them to time out. After the grace period has timed out on both the server and any clients, run the application that is having lock problems and verify that a lock request is transmitted from client to server and server to client.

Secure NFS

In addition to the standard UNIX authentication system, the Network Information Services NIS and NIS+ and the Network File System (NFS) provide a means to authenticate users and machines in networks on a message—by—message basis. This additional authentication system uses Data Encryption Standard (DES) encryption and public key cryptography.

This section discusses the following topics:

- Secrecy, on page 10-25
- Secrecy in NFS, on page 10-27
- Naming Network Entities for DES Authentication, on page 10-29
- The /etc/publickey File, on page 10-29
- Booting Considerations of Public Key Systems, on page 10-30
- Performance Considerations, on page 10-30
- Checklist for Administering Secure NFS, on page 10-30
- How to Configure Secure NFS, on page 10-31
- How to Export a File System Using Secure NFS, on page 10-31
- How to Mount a File System Using Secure NFS, on page 10-32

Secrecy

Throughout history, various groups of people have sought to communicate in such a way that only the sender and receiver know the contents of a given message. To achieve this secrecy, senders and receivers use a *cipher*, a scheme for converting a *plaintext* message into *ciphertext* and back again. *Encryption* is the process of converting plain text into cipher text, and *decryption* is the process of converting cipher text into plain text.

One of the earliest ciphers, the *Caesar cipher*, is attributed to Julius Caesar. In this cipher, one letter is substituted for another. For example, 'A' becomes 'C', 'B' becomes 'D', ..., 'Y' becomes 'A', and 'Z' becomes 'B'. So, the Caesar cipher encrypts the phrase **ATTACK AT DAWN** as **CVVCEM CV FCYP**.

If the Carthaginians could *cryptanalyze* the Caesar cipher and break it, the Roman *cryptographers* would have to invent an entirely new cipher. Since cipher development is an expensive process, the Romans might use a cipher *key* in order to get more use out of their cipher. For example, instead of specifying a letter–for–letter substitute, the Romans might specify a key K, where K indicates the number of positions to shift a letter. That is, if K = 2, then 'A' becomes 'C'. If K = 4, then 'A' becomes 'E', and so on. With this scheme, if the Carthaginians break the cipher, all the Romans must do is change the key. Of course, the Carthaginians might realize what algorithm the Italians were using, and exhaustively try every value of K from 1 to 26. If the Carthaginians had a computer, their task would be a trivial programming exercise.

Data Encryption Standard

Modern ciphers are designed to address the fact that computers can be powerful tools for an intruder attempting to break a cipher. In 1977, the U.S. government adopted a cipher as its Data Encryption Standard. This cipher is widely used in industry. DES is a highly complex algorithm which converts 64—bit blocks of plain text into 64—bit blocks of cipher text using a 56—bit key. Because of the complexity of the algorithm and the size of the cipher key, DES is essentially unbreakable. For example, if an intruder had a computer which could compute the DES algorithm at a rate of one key per microsecond, the computer would need over two thousand years to try every possible key.

Public Key Cryptography

A significant weakness in any encryption algorithm is the key used. If the sender and receiver are to communicate securely using a cipher, both the sender and receiver must know the key. They must agree upon a key either using a separate communications link, which itself must be secure, or in person.

To address this problem, two researchers (Diffie and Hellman) developed a technique by which the sender and receiver can exchange keys publicly without compromising the security of their communication. Their technique has three requirements:

Decipher(Encipher(plaintext, E), D) = plaintext

where E is the encryption key (known to the public), and D is the decryption key (known only by the receiver).

That is, Encipher and Decipher functions are inverses of each other. Therefore, if you take the encrypted text string returned by Encipher(plaintext, E), and use it along with the key D for the Decipher function, Decipher will return the original plain text.

- An intruder cannot deduce Decipher() from Encipher().
- Encipher() is unbreakable.

The following outline describes how a sender sends a secret message to a receiver.

- 1. The sender obtains the receiver's public encryption key.
- 2. The sender converts the plain text message into cipher text by computing the result:

```
ciphertext = Encipher( plaintext, E )
```

- 3. The sender sends the cipher text message to the receiver.
- 4. The receiver receives the cipher text message and converts it into plaintext by computing the result:

```
plaintext = Decipher( ciphertext, D )
```

Even if an intruder intercepts the message, the intruder will not be able to decipher it because the intruder does not have the receiver's decipher key. (For that matter, the sender cannot decipher the cipher text message either.)

Authentication

A primary application of secrecy is *authentication*. One common method of authentication (which is the standard UNIX method of authentication) uses a password. In other words, when a user wants to login, the operating system requires the user to provide a password known only by the operating system and the user. If the user provides the correct password, the operating system concludes that the user is who the user claims to be. Note that this method requires the operating system to store the user's password in a file on the system, although in encrypted form. This means that two different entities know a single password.

Public key cryptography provides an alternative to password authentication. Suppose a sender wants to send a message, and the receiver wants to be certain that the message is from the sender and not from an intruder pretending to be the sender. The authentication process will occur in the following manner:

- 1. First the sender enciphers a "request to send" message using the receiver's public key, and then sends the request.
- 2. The receiver receives the "request to send" message and deciphers it using the receiver's private key.
- 3. The receiver enciphers a "token" message using the sender's public key, and then sends the token.
- 4. The sender receives the token and deciphers it with the sender's private key. Now, when the sender sends messages to the receiver, the sender will begin each message with the token, signifying that the sender is, in fact, the sender. If an intruder attempts to send

messages in the name of the sender, the receiver will reject them because the intruder does not know what the token is.

Note that, unlike password authentication, the receiver is able to authenticate the sender without knowing the sender's private key. For more information on authentication systems, see "Understanding RPC Authentication" in *AIX Communications Programming Concepts*.

Secrecy in NFS

NFS, NIS, and NIS+ use the DES algorithm for different purposes. NFS uses DES to encrypt a time stamp in the Remote Procedure Call (RPC) messages sent between NFS servers and clients. This encrypted time stamp authenticates machines just as the "token" authenticates the sender, as described in "Authentication", on page 10-26. (For information about NIS and NIS+ see *AIX 4.3 NIS/NIS+ Guide*.)

Because NFS can authenticate every single RPC message exchanged between NFS clients and servers, this provides an additional, optional, level of security for each file system. By default, file systems are exported with the standard UNIX authentication. To take advantage of this additional level of security, you can specify the **secure** option when you export a file system.

Public Key Cryptography for Secure NFS

Both the user's public and the secret keys are stored and indexed by net name in the **publickey.byname** map. The secret key is DES—encrypted with the user's login password. The **keylogin** command uses the encrypted secret key, decrypts it with the login password, then gives it to a secure local key server to save for use in future RPC transactions. Users are not aware of their public and secret keys because the **yppasswd** command, in addition to changing the login password, generates the public and secret keys automatically.

The **keyserv** daemon is an RPC service that runs on each NIS and NIS+ machine. For information on how NIS+ uses **keyserv**, see *AIX 4.3 NIS/NIS+ Guide*. Within NIS, **keyserv** executes the following three public key subroutines:

- key setsecret subroutine
- key_encryptsession subroutine
- key decryptsession subroutine

The **key_setsecret** subroutine tells the key server to store the user's secret key (SK_A) for future use; it is normally called by the **keylogin** command. The client program calls the **key_encryptsession** subroutine to generate the encrypted conversation key, which is passed in the first RPC transaction to a server. The key server looks up the server public key and combines it with the client secret key (set up by a previous **key_setsecret** subroutine) to generate the common key. The server asks the key server to decrypt the conversation key by calling the **key_decryptsession** subroutine.

Implicit in these subroutine calls is the name of the caller, which must be authenticated in some manner. The key server cannot use DES authentication to do this, since it would create a deadlock. The key server solves this problem by storing the secret keys by the user ID (UID) and only granting requests to local root processes. The client process then executes a root user owned **setuid** subroutine which makes the request on the part of the client, telling the key server the real UID of the client.

Authentication Requirements

Secure NFS authentication is based on the ability of a sender to encrypt the current time, which the receiver can then decrypt and check against its own clock. This process has two requirements:

- The two agents must agree on the current time.
- The sender and receiver must be using the same DES encryption key.

Agreeing on the Current Time

If the network uses time synchronization, then the **timed** daemon will keep the client and server clocks in sync. If not, the client will compute the proper time stamps based on the server's clock. To do this, the client determines the server time before starting the RPC session, and then computes the time difference between its own clock and that of the server's. The client then adjusts its time stamp accordingly. If, during the course of an RPC session, the client and server clocks get out of sync to the point where the server begins rejecting the client's requests, the client will redetermine the server time.

Using the Same DES Key

The client and server compute the same DES encryption key by using public key cryptography. For any client A and server B, there is a key that only A and B can deduce. This key is called the *common key*. The client derives the common key by computing the following formula:

$$K_{AB} = PK_{B}^{SK_{A}}$$

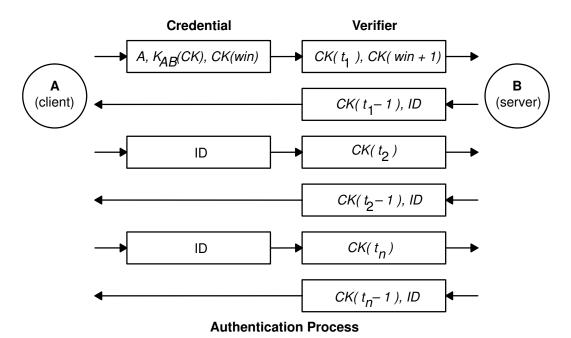
where *K* stands for the *common Key*, *PK* stands for *Public Key*, and *SK* stands for *Secret Key*, and each of these keys is a 128–bit number. The server derives the same common key by computing the following formula:

$$K_{AB} = PK_{A}^{SK_{B}}$$

Only the server and client can calculate this common key since doing so requires knowing one secret key or the other. Because the common key has 128 bits, and DES uses a 56-bit key, the client and server extract 56 bits from the common key to form the DES key.

Authentication Process

When a client wants to talk to a server, it randomly generates a key used for encrypting the time stamps. This key is known as the *conversation key* (*CK*). The client encrypts the conversation key using the DES common key (described in "Authentication Requirements", on page 10-27) and sends it to the server in the first RPC transaction. This process is shown in the Authentication Process illustration.



This figure shows client A connecting to server B. The term K(CK) means CK is encrypted with the DES common key K. In its first request, the client's RPC credential contains the client name (A), the conversation key (CK), and the variable called win (window) encrypted with CK. (The default window size is 30 minutes.) The client's verifier in the first request contains the encrypted time stamp and an encrypted verifier of the specified window, win + 1. The window verifier makes guessing the right credential much more difficult, and increases security.

After authenticating the client, the server stores the following items in a credential table:

- · Client name, A
- Conversation key, CK
- Window
- Time stamp

The server only accepts time stamps that are chronologically greater than the last one seen, so any replayed transactions are guaranteed to be rejected. The server returns to the client in the verifier an index ID into the credential table, plus the client's time stamp minus one, encrypted by *CK*. The client knows that only the server could have sent such a verifier, since only the server knows what time stamp the client sent. The reason for subtracting one from the time stamp is to ensure that it is not valid and cannot be reused as a client verifier. After the first RPC transaction, the client sends just its ID and an encrypted time stamp to the server, and the server sends back the client's time stamp minus one, encrypted by *CK*.

Naming Network Entities for DES Authentication

DES authentication does its naming by using net names. The following paragraphs describe how NIS handles DES authentication. For information on how NIS+ handles DES authentication, see *AIX 4.3 NIS/NIS+ Guide*.

A *net name* is a string of printable characters to authenticate. The public and secret keys are stored on a per–net–name rather than a per–user–name basis. The **netid.byname** NIS map maps the net name into a local UID and group–access list.

User names are unique within each domain. Net names are assigned by concatenating the operating system and user ID with the NIS and Internet domain names. A good convention for naming domains is to append the Internet domain name (com, edu, gov, mil) to the local domain name.

Network names are assigned to machines as well as to users. A net name of a machine is formed much like that of a user. For example, a machine named hal in the eng.ibm.com domain has the net name unix.hal@eng.ibm.com. Proper authentication of machines is important for diskless machines that need full access to their home directories over the network.

To authenticate users from any remote domain, you should make entries for them in two NIS databases. One is an entry for their public and secret keys; the other is for their local UID and group—access list mapping. Users in the remote domain can then access all of the local network services, such as the NFS and remote logins.

/etc/publickey File

The /etc/publickey file contains names and public keys, which NIS and NIS+ use to create the publickey map. The publickey map is used for secure networking. Each entry in the file consists of a network user name (which may refer to either a user or a host name), followed by the user's public key (in hexadecimal notation), a colon, and the user's encrypted secret key (also in hexadecimal notation). By default, the only user in the /etc/publickey file is the user nobody.

Do not use a text editor to alter the /etc/publickey file because the file contains encryption keys. To alter the /etc/publickey file, use either the chkey or newkey commands.

Booting Considerations of Public Key Systems

When restarting a machine after a power failure, all of the stored secret keys are lost, and no process can access secure network services, such as mounting an NFS. Root processes could continue if there were someone to enter the password that decrypts the root user's secret key. The solution is to store the root user's decrypted secret key in a file that the key server can read.

Not all **setuid** subroutine calls operate as they should. For example, if a **setuid** subroutine is called by owner dave, who has not logged into the machine since it booted, the subroutine cannot access any secure network services as dave. However, most **setuid** subroutine calls are owned by the root user, and the root user's secret key is always stored at startup time.

Performance Considerations

Secure NFS affects system performance in two ways. First, both the client and server must compute the common key. The time it takes to compute the common key is about 1 second. As a result, it takes about 2 seconds to establish the initial RPC connection, since both client and server have to perform this operation. After the initial RPC connection, the key server caches the results of previous computations, and so it does not have to recompute the common key every time.

Second, each RPC transaction requires four DES encryption operations:

- 1. The client encrypts the request time stamp,
- 2. The server decrypts it,
- 3. The server encrypts the reply time stamp,
- 4. And finally, the client decrypts it.

Since system performance is reduced by secure NFS, you will have to weigh the benefits of increased security and your system's performance requirements.

Administering Secure NFS Checklist

The following is a checklist of items to help ensure that secure NFS operates properly:

- When mounting a file system with the -secure option on a client, the server name must match the server host name in the /etc/hosts file. If a name server is being used for host name resolution, make sure the host information returned by the name server matches the entry in the /etc/hosts file. Authentication errors will result if these data do not match, because the net names for machines are based on the primary entries in the /etc/hosts file and keys in the publickey map are accessed by net name.
- Do not mix secure and nonsecure exports and mounts. Otherwise, file access may be determined incorrectly. For example, if a client machine mounts a secure file system without the **secure** option or mounts an nonsecure system with the **secure** option, users have access as nobody, rather than as themselves. This condition also occurs if a user unknown to NIS or NIS+ attempts to create or modify files on a secure file system.
- Because NIS must propagate a new map after each use of the chkey and newkey commands, only use these commands when the network is lightly loaded.
- Do not delete the /etc/keystore file or the /etc/.rootkey file. If you reinstall, move, or upgrade a machine, save the /etc/keystore and /etc/.rootkey files.
- Instruct your users to use the **yppasswd** command rather than the **passwd** command to change passwords. Doing so will keep passwords and private keys synchronized.
- Because the login command does not retrieve keys out of the publickey map for the
 keyserv daemon, the user must execute the keylogin command. You may want to place
 the keylogin command in each user's profile file to execute the command automatically
 during login. Note that the keylogin command requires the user to enter his or her
 password again.

- When you generate keys for the root user at each host with either the newkey –h or chkey command, you must execute the keylogin command to pass the new keys to the keyserv daemon. The keys are stored in the /etc/.rootkey file, which is read by the keyserv daemon each time the daemon is started.
- Periodically verify that the **yppasswdd** and **ypupdated** daemons are running on the NIS master server. These daemons are necessary for maintaining the **publickey** map.
- Periodically verify that that the keyserv daemon is running on all machines using secure NFS.

Configuring Secure NFS

To configure secure NFS on NIS master and slave servers, use the Web-based System Manager wsm network application or use the following procedure. For information on using NFS with NIS+ see AIX 4.3 NIS/NIS+ Guide.

1. On the NIS master server, create an entry for each user in the NIS /etc/publickey file using the newkey command. This command has two options. For a regular user, enter:

```
smit newkey
OR
newkey -u username
For a root user on a host machine, enter:
newkey -h hostname
```

Alternatively, users can establish their own public keys using the **chkey** or **newkey** commands.

- 2. Create the NIS **publickey** map by following the instructions in *AIX 4.3 NIS/NIS+ Guide*. The corresponding NIS **publickey.byname** map should reside only on the NIS servers.
- 3. Uncomment the following stanzas in the /etc/rc.nfs file:

```
#if [ -x /usr/sbin/keyserv ]; then
# startsrc -s keyserv
#fi
#if [ -x /usr/lib/netsvc/yp/rpc.ypupdated -a -d
/etc/yp/'domainname' ]; then
# startsrc -s ypupdated
#fi
#DIR=/etc/passwd
#if [ -x /usr/lib/netsvc/yp/rpc.yppasswdd -a -f $DIR/passwd ];
then
# startsrc -s yppasswdd
#fi
```

4. Start the keyserv, ypupdated, and yppasswdd daemons using the startsrc command.

To configure secure NFS on NIS clients, start the **keyserv** daemon using the startsrc command.

Exporting a File System Using Secure NFS

You can export a secure NFS using the Web-based System Manager wsm network application or by using one of the following procedures.

- To export a secure NFS file system using SMIT:
 - a. Verify that NFS is already be running by executing the command Issrc –g nfs. The output should indicate that the nfsd and the rpc.mountd daemons are active. If they are not, start NFS using the instructions in "Starting the NFS Daemons", on page 10-9.

- b. Verify that the **publickey** map exists and that the **keyserv** daemon is running. See "Configuring Secure NFS", on page 10-31 for more information.
- c. Execute the **smit mknfsexp** fast path.
- d. Specify the appropriate values for the PATHNAME of directory to export, MODE to export directory, and EXPORT directory now, system restart or both fields. Specify ves for the Use SECURE option field.
- e. Specify any other optional characteristics you wish, or simply accept the default values by leaving the remaining fields as they are.
- f. Exit SMIT. If the /etc/exports file does not exist, then it will be created.
- g. Repeat steps 3 through 6 for each directory you wish to export.
- To export a secure NFS file system using a text editor:
 - a. Open the /etc/exports file with your favorite text editor.
 - b. Create an entry for each directory to be exported, using the full path name of the directory. List each directory to be exported starting in the left margin. No directory should include any other directory that is already exported. See the /etc/exports file documentation for a description of the full syntax for entries in the /etc/exports file, including how to specify the secure option.
 - c. Save and close the /etc/exports file.
 - d. If NFS is currently running, enter:

```
/usr/sbin/exportfs -a
```

The **–a** option tells the **exportfs** command to send all information in the **/etc/exports** file to the kernel. If NFS is not running, start NFS using the instructions in "Start the NFS Daemons", on page 10-9.

To export an NFS file system temporarily (that is, without changing the /etc/exports file):

Enter:

```
exportfs -i -o secure /dirname
```

where <code>dirname</code> is the name of the file system you want to export. The **exports –i** command specifies that the <code>/etc/exports</code> file is not to be checked for the specified directory, and all options are taken directly from the command line.

Mounting a File System Using Secure NFS

To mount a secure NFS directory explicitly:

 Verify that the NFS server has exported the directory. Do this by executing the command:

```
showmount -e ServerName
```

where <code>ServerName</code> is the name of the NFS server. This command displays the names of the directories currently exported from the NFS server. If the directory you want to mount is not listed, export the directory from the server.

- 2. Establish the local mount point using the **mkdir** command. For NFS to complete a mount successfully, a directory that acts as the mount point (or place holder) of an NFS mount must be present. This directory should be empty. This mount point can be created like any other directory, and no special attributes are needed for this directory.
- 3. Verify that the **publickey** map exists and that the **keyserv** daemon is running. See "Configuring Secure NFS", on page 10-31 for more information.

4. Enter:

mount -o secure ServerName:/remote/directory /local/directory

where ServerName is the name of the NFS server, /remote/directory is the directory on the NFS server you want to mount, and /local/directory is the mount point on the NFS client.

Note: Only the root user can mount a secure NFS.

Refer to "Establishing Predefined NFS Mounts", on page 10-15 for more details on mounting with NFS.

NFS Problem Determination

As with other network services, problems can occur on machines that use the Network File System (NFS). Troubleshooting for these problems involves understanding the strategies for tracking NFS problems, recognizing NFS-related error messages, and selecting the appropriate solutions. When tracking down an NFS problem, isolate each of the three main points of failure to determine which is not working: the server, the client, or the network itself.

Note: See "Troubleshooting the Network Lock Manager", on page 10-23 for file lock problems.

Identifying Hard-Mounted and Soft-Mounted File Problems

When the network or server has problems, programs that access hard–mounted remote files fail differently from those that access soft–mounted remote files.

If a server fails to respond to a hard-mount request, NFS prints the message:

```
NFS server hostname not responding, still trying
```

Hard—mounted remote file systems cause programs to hang until the server responds because the client retries the mount request until it succeeds. You should use the **–bg** flag with the **mount** command when performing a hard mount so that if the server does not respond, the client will retry the mount in the background.

If a server fails to respond to a soft-mount request, NFS prints the message:

```
Connection timed out
```

Soft—mounted remote file systems return an error after trying unsuccessfully for a while. Unfortunately, many programs do not check return conditions on file system operations, so you do not see this error message when accessing soft—mounted files. However, this NFS error message will print on the console.

Identifying NFS Problems Checklist

If a client is having NFS trouble, do the following:

- 1. Verify that the network connections are good.
- 2. Verify that the **inetd**, **portmap**, and **biod** daemons are running on the client, by following the instructions in "Getting the Current Status of the NFS Daemons", on page 10-9.
- 3. Verify that a valid mount point exists for the file system being mounted. For more information, see "Configuring an NFS Client", on page 10-10.
- 4. Verify that the server is up and running by running the following command at the shell prompt of the client:

```
/usr/bin/rpcinfo -p server_name
```

If the server is up, a list of programs, versions, protocols, and port numbers is printed, similar to the following:

program	vers	proto	port	
100000	2	tcp	111	portmapper
100000	2	udp	111	portmapper
100005	1	udp	1025	mountd
100001	1	udp	1030	rstatd
100001	2	udp	1030	rstatd
100001	3	udp	1030	rstatd
100002	1	udp	1036	rusersd
100002	2	udp	1036	rusersd
100008	1	udp	1040	walld
100012	1	udp	1043	sprayd
100005	1	tcp	694	mountd
100003	2	udp	2049	nfs
100024	1	udp	713	status
100024	1	tcp	715	status
100021	1	tcp	716	nlockmgr
100021	1	udp	718	nlockmgr
100021	3	tcp	721	nlockmgr
100021	3	udp	723	nlockmgr
100020	1	udp	726	llockmgr
100020	1	tcp	728	llockmgr
100021	2	tcp	731	nlockmgr

If a similar response is not returned, log in to the server at the server console and check the status of the **inetd** daemon by following the instructions in "Get the Current Status of the NFS Daemons", on page 10-9.

5. Verify that the **mountd**, **portmap** and **nfsd** daemons are running on the NFS server by entering the following commands at the client shell prompt:

```
/usr/bin/rpcinfo -u server_name mount
/usr/bin/rpcinfo -u server_name portmap
/usr/bin/rpcinfo -u server_name nfs
```

If the daemons are running at the server, the following responses are returned:

```
program 100005 version 1 ready and waiting program 100000 version 2 ready and waiting program 100003 version 2 ready and waiting
```

The program numbers correspond to the commands, respectively, as shown in the example output above. If a similar response is not returned, log in to the server at the server console and check the status of the daemons by following the instructions in "Get the Current Status of the NFS Daemons", on page 10-9.

6. Verify that the /etc/exports file on the server lists the name of the file system that the client wants to mount and that the file system is exported. Do this by entering the command:

```
showmount -e server_name
```

This command will list all the file systems currently exported by the <code>server_name</code>.

Asynchronous Write Errors

When an application program writes data to a file in an NFS-mounted file system, the write operation is scheduled for asynchronous processing by the **biod** daemon. If an error occurs at the NFS server at the same time that the data is actually written to disk, the error is returned to the NFS client and the **biod** daemon saves the error internally in NFS data structures. The stored error is subsequently returned to the application program the next time it calls either the **fsync** or **close** functions. As a consequence of such errors, the application is not notified of the write error until the program closes the file. A typical example of this event is when a file system on the server is full, causing writes attempted by a client to fail.

NFS Error Messages

The following sections explain error codes that can be generated while using NFS.

nfs_server Error Message

Insufficient transmit buffers on your network can cause the following error message:

```
nfs_server: bad sendreply
```

To increase transmit buffers, use the Web-based System Manager fast path, **wsm devices**, or the System Management Interface Tool (SMIT) fast path, **smit commodev**. Then select your adapter type, and increase the number of transmit buffers.

mount Error Messages

A remote mounting process can fail in several ways. The error messages associated with mounting failures are as follows:

```
mount: ... already mounted

The file system that you are trying to mount is already mounted.

mount: ... not found in /etc/filesystems

The specified file system or directory name cannot be matched.
```

If you issue the **mount** command with either a directory or file system name but not both, the command looks in the /etc/filesystems file for an entry whose file system or directory field matches the argument. If the **mount** command finds an entry such as the following:

```
/dancer.src:
    dev=/usr/src
    nodename = d61server
    type = nfs
    mount = false
```

then it performs the mount as if you had entered the following at the command line:

```
/usr/sbin/mount -n dancer -o rw, hard /usr/src /dancer.src ... not in hosts database
```

On a network without Network Information Service, this message indicates that the host specified to the **mount** command is not in the /etc/hosts file. On a network running NIS, the message indicates that NIS could not find the host name in the /etc/hosts database or that the NIS ypbind daemon on your machine has died. If the /etc/resolv.conf file exists so that the name server is being used for host name resolution, there can be a problem in the **named** database. See "Name Resolution on an NFS Server", on page 10-40.

Check the spelling and the syntax in your **mount** command. If the command is correct, your network does not run NIS, and you only get this message for this host name, check the entry in the /etc/hosts file.

If your network is running NIS, make sure that the **ypbind** daemon is running by entering the following at the command line:

```
ps -ef
```

You should see the **ypbind** daemon in the list. Try using the **rlogin** command to log in remotely to another machine, or use the **rcp** command to remote—copy something to another machine. If this also fails, your **ypbind** daemon is probably stopped or hung.

If you only get this message for this host name, you should check the /etc/hosts entry on the NIS server.

```
mount: ... server not responding: port mapper failure - RPC timed out

Either the server you are trying to mount from is down or its port mapper is

stopped or hung. Try rebooting the server to restart the inetd, portmap,

and ypbind daemons.
```

If you cannot log in to the server remotely with the **rlogin** command but the server is up, you should check the network connection by trying to log in remotely to some other machine. You should also check the server's network connection.

```
mount: ... server not responding: program not registered

This means that the mount command got through to the port mapper, but the rpc.mountd NFS mount daemon was not registered.
```

```
mount: access denied ...
```

Your machine name is not in the export list for the file system you are trying to mount from the server.

You can get a list of the server's exported file systems by running the following command at the command line:

```
showmount -e hostname
```

If the file system you want is not in the list, or your machine name or netgroup name is not in the user list for the file system, log in to the server and check the /etc/exports file for the correct file system entry. A file system name that appears in the /etc/exports file, but not in the output from the showmount command, indicates a failure in the mountd daemon. Either the daemon could not parse that line in the file, it could not find the directory, or the directory name was not a locally mounted directory. If the /etc/exports file looks correct and your network runs NIS, check the server's ypbind daemon. It may be stopped or hung. For more information, see AIX 4.3 NIS/NIS+ Guide.

```
mount: ...: Permission denied
```

This message is a generic indication that some part of authentication failed on the server. It may be that in the previous example, you are not in the export list, the server could not recognize your machine's **ypbind** daemon, or that the server does not accept the identity you provided.

Check the server's /etc/exports file, and, if applicable, the ypbind daemon. In this case you can just change your host name with the hostname command and retry the mount command.

```
mount: ...: Not a directory

Either the remote path or the local path is not a directory. Check the spelling in your command and try to run on both directories.
```

```
mount: ...: You are not allowed
```

You must have root authority or be a member of the system group to run the **mount** command on your machine because it affects the file system for all users on that machine. NFS mounts and unmounts are only allowed for root users and members of the system group.

Identifying the Cause of Slow Access Times for NFS

If access to remote files seems unusually slow, ensure that access time is not being inhibited by a runaway daemon, a bad **tty** line, or a similar problem.

Checking Processes

At the server, enter the following at the command line:

If the server seems fine and other users are getting timely responses, make sure your **biod** daemons are running. Try the following steps:

- Run the ps –ef command and look for the biod daemons in the display.
 If they are not running or are hung, continue with steps 2 and 3.
- 2. Stop the biod daemons that are in use by issuing the following command:

```
stopsrc -x biod -c
```

3. Start the **biod** daemons by issuing the following command:

```
startsrc -s biod
```

To determine if the **biod** daemons are hung, run the **ps** command as above, copy a large file from a remote system, and then run the **ps** command again. If the **biod** daemons do not accumulate CPU time, they are probably hung.

Checking Network Connections

If the **biod** daemons are working, check the network connections. The **nfsstat** command determines whether you are dropping packets. Use the **nfsstat** –c and **nfsstat** –s commands to determine if the client or server is retransmitting large blocks. Retransmissions are always a possibility due to lost packets or busy servers. A retransmission rate of 5% is considered high.

The probability of retransmissions can be reduced by changing communication adapter transmit queue parameters. The System Management Interface Tool (SMIT) can be used to change these parameters.

The following values are recommended for NFS servers.

Communication Adapter Maximum Transmission Unit (MTU) and Transmit Queue Sizes			
Adapter	MTU	Transmit Queue	
Token Ring	1500	50	
4Mb	3900	40 (Increase if the nfsstat command times out.)	
16Mb	1500	40 (Increase if the nfsstat command times out.)	
	8500	40 (Increase if the nfsstat command times out.)	
Ethernet	1500	40 (Increase if the nfsstat command times out.)	

The larger MTU sizes for each token–ring speed reduce processor use and significantly improve read/write operations.

Note:

- 1. Apply these values to NFS clients if retransmissions persist.
- 2. All nodes on a network must use the same MTU size.

Setting MTU Sizes

To set MTU size, use the Web-based System Manager fast path, **wsm network**, or the SMIT fast path, **smit chif**. Select the appropriate adapter and enter an MTU value in the Maximum IP Packet Size field.

The **ifconfig** command can be used to set MTU size (and *must* be used to set MTU size at 8500). The format for the **ifconfig** command is:

```
ifconfig trn NodeName up mtu MTUSize
```

where trn is your adapter name, for example, tr0.

Another method of setting MTU sizes combines the ifconfig command with SMIT.

- 1. Add the **ifconfig** command for token rings, as illustrated in the previous example, to the /etc/rc.bsdnet file.
- 2. Enter the smit setbootup option fast path. Toggle the Use BSD Style field to yes.

Setting Transmit Queue Sizes

Communication adapter transmit queue sizes are set with SMIT. Enter the **smit chgtok** fast path, select the appropriate adapter, and enter a queue size in the Transmit field.

Fixing Hung Programs

If programs hang during file-related work, the NFS server could have stopped. In this case, the following error message may be displayed:

```
NFS server hostname not responding, still trying
```

The NFS server (hostname) is down. This indicates a problem with the NFS server, the network connection, or the NIS server.

Check the servers from which you have mounted file systems if your machine hangs completely. If one or more of them is down, do not be concerned. When the server comes back up, your programs continue automatically. No files are destroyed.

If a soft-mounted server dies, other work is not affected. Programs that time out trying to access soft-mounted remote files fail with the errno message, but you will still be able to access your other file systems.

If all servers are running, determine whether others who are using the same servers are having trouble. More than one machine having service problems indicates a problem with the server's **nfsd** daemons. In this case, log in to the server and run the **ps** command to see if the **nfsd** daemon is running and accumulating CPU time. If not, you may be able to stop and then restart the **nfsd** daemon. If this does not work, you have to reboot the server.

Check your network connection and the connection of the server if other systems seem to be up and running.

Permissions and Authentication Schemes

Sometimes, after mounts have been successfully established, there are problems in reading, writing, or creating remote files or directories. Such difficulties are usually due to permissions or authentication problems. Permission and authentication problems can vary in cause depending on whether NIS is being used and secure mounts are specified.

The simplest case occurs when nonsecure mounts are specified and NIS is not used. In this case, user IDs (UIDs) and group IDs (GIDs) are mapped solely through the server and clients /etc/passwd and /etc/group files, respectively. In this scheme, for a user named john to be identified both on the client and on the server as john, the user john in the /etc/passwd file must have the same UID number. The following is an example of how this might cause problems:

```
User john is uid 200 on client foo. User john is uid 250 on server bar. User jane is uid 200 on server bar.
```

The /home/bar directory is mounted from server bar onto client foo. If user john is editing files on the /home/bar remote file system on client foo, confusion results when he saves files.

The server bar thinks the files belong to user jane, because jane is UID 200 on bar. If john logs on directly to bar by using the **rlogin** command, he may not be able to access the files he just created while working on the remotely mounted file system. jane, however, is able to do so because the machines arbitrate permissions by UID, not by name.

The only permanent solution to this is to reassign consistent UIDs on the two machines. For example, give john UID 200 on server bar or 250 on client foo. The files owned by john would then need to have the **chown** command run against them to make them match the new ID on the appropriate machine.

Because of the problems with maintaining consistent UID and GID mappings on all machines in a network, NIS or NIS+ is often used to perform the appropriate mappings so that this type of problem is avoided. See *AIX 4.3 NIS/NIS+ Guide* for more information.

Name Resolution on an NFS Server

When an NFS server services a mount request, it looks up the name of the client making the request. The server takes the client Internet Protocol (IP) address and looks up the corresponding host name that matches that address. Once the host name has been found, the server looks at the exports list for the requested directory and checks the existence of the client's name in the access list for the directory. If an entry exists for the client and the entry matches exactly what was returned for the name resolution, then that part of the mount authentication passes.

If the server is not able to perform the IP address—to—host—name resolution, the server denies the mount request. The server must be able to find some match for the client IP address making the mount request. If the directory is exported with the access being to all clients, the server still must be able to do the reverse name lookup to allow the mount request.

The server also must be able to look up the correct name for the client. For example, if there exists an entry in the /etc/exports file like the following:

```
/tmp -access=silly:funny
```

the following corresponding entries exist in the /etc/hosts file:

```
150.102.23.21 silly.domain.name.com
150.102.23.52 funny.domain.name.com
```

Notice that the names do not correspond exactly. When the server looks up the IP address—to—host—name matches for the hosts <code>silly</code> and <code>funny</code>, the string names do not match exactly with the entries in the access list of the export. This type of name resolution problem usually occurs when using the **named** daemon for name resolution. Most **named** daemon databases have aliases for the full domain names of hosts so that users do not have to enter full names when referring to hosts. Even though these host—name—to—IP address entries exist for the aliases, the reverse lookup may not exist. The database for reverse name lookup (IP address to host name) usually has entries containing the IP address and the full domain name (not the alias) of that host. Sometimes the export entries are created with the shorter alias name, causing problems when clients try to mount.

Limitations on the Number of Groups in the NFS Structure

On systems that use NFS Version 3.2, users cannot be a member of more than 16 groups without complications. (Groups are defined by the **groups** command.) If a user is a member of 17 or more groups, and the user tries to access files owned by the 17th (or greater) group, the system will not allow the file to be read or copied. To permit the user access to the files, rearrange the group order.

Mounting from NFS Servers That Have Earlier Version of NFS

When mounting a file system from a pre–Version 3 NFS server onto a Version 3 NFS client, a problem occurs when the user on the client executing the mount is a member of more than eight groups. Some servers are not able to deal correctly with this situation and deny the request for the mount. The solution is to change the user's group membership to a number less than eight and then retry the mount. The following error message is characteristic of this group problem:

RPC: Authentication error; why=Invalid client credential

Problems That Occur If the NFS Kernel Extension Is Not Loaded

Some NFS commands do not execute correctly if the NFS kernel extension is not loaded. Some commands with this dependency are: nfsstat, exportfs, mountd, nfsd, and biod. When NFS is installed on the system, the kernel extension is placed in the /usr/lib/drivers/nfs.ext file. This file is then loaded as the NFS kernel extension when the system is configured. The script that does this kernel extension loads the /etc/rc.net file. There are many other things done in this script, one of which is to load the NFS kernel extension. It is important to note that Transmission Control Protocol/Internet Protocol (TCP/IP) kernel extension should be loaded before the NFS kernel extension is loaded.

Note: The **gfsinstall** command is used to load the NFS kernel extension into the kernel when the system initially starts. This command can be run more than once per system boot and it will not cause a problem. The system is currently shipped with the **gfsinstall** command used in both the **/etc/rc.net** and **/etc/rc.nfs** files. This is correct. There is no need to remove either of these calls.

NFS Reference

List of Network File System (NFS) Files

bootparams
 exports
 Lists the directories that can be exported to NFS clients.
 networks
 Contains information about networks on the Internet network.
 pcnfsd.conf
 Provides configuration options for the rpc.pcnfsd daemon.

Configuration File

rpc Contains database information for Remote Procedure Call (RPC)

programs.

xtab Lists directories that are currently exported.

/etc/filesystems Lists all the file systems that are mounted at system restart.

List of NFS Commands

chnfs Starts a specified number of **biod** and **nfsd** daemons.

mknfs Configures the system to run NFS and starts NFS daemons.

nfso Configures NFS network options.

automount Mounts an NFS file system automatically.

chnfsexp
 chnfsmnt
 Changes the attributes of an NFS-exported directory.
 chnfsmnt
 changes the attributes of an NFS-mounted directory.
 exports
 Exports and unexports directories to NFS clients.

Isnfsexp Displays the characteristics of directories that are exported with NFS.

Isnfsmnt Displays the characteristics of mounted NFS systems.

mknfsexpmknfsmntmknfsmntmcontentmknfsmntmcontent<l

rmnfsexpRemoves NFS—exported directories from a server's list of exports.rmnfsmntRemoves NFS—mounted file systems from a client's list of mounts.

List of NFS Daemons

Locking Daemons

lockd Processes lock requests through the RPC package.

statd Provides crash–and–recovery functions for the locking services on NFS.

Network Service Daemons and Utilities

biod Sends the client's read and write requests to the server.

mountd Answers requests from clients for file system mounts.

nfsd Starts the daemons that handle a client's request for file system

operations.

pcnfsd Handles service requests from PC–NFS clients.

nfsstat Displays information about a machine's ability to receive calls.

on Executes commands on remote machines.

portmap Maps RPC program numbers to Internet port numbers.

rexd Accepts request to run programs from remote machines.

rpcgen Generates C code to implement an RPC protocol.

rpcinfo Reports the status of RPC servers.

rstatd Returns performance statistics obtained from the kernel.
rup Shows the status of a remote host on the local network.
rusers Reports a list of users logged on to the remote machines.

rusersdResponds to queries from the rusers command.rwallSends messages to all users on the network.rwalldHandles requests from the rwall command.

showmount Displays a list of all clients that have mounted remote file systems.

sprayspraydSends a specified number of packets to a host.spraydReceives packets sent by the spray command.

Secure Networking Daemons and Utilities

chkey Changes the user's encryption key.

keyenvoy Provides an intermediary between user processes and the key server.

keylogin Decrypts and stores the user's secret key.

keyserv Stores public and private keys.

mkkeyserv Starts the **keyserv** daemon and uncomments the appropriate entries in

the /etc/rc.nfs file.

newkey Creates a new key in the **publickey** file.

rmkeyserv Stops the keyserv daemon and comments the entry for the keyserv

daemon in the /etc/rc.nfs file.

ypupdated Updates information in Network Information Service (NIS) maps.

Sun Diskless Client Support

bootparamd Provides information necessary for booting to diskless clients.

NFS Subroutines

cbc_crypt, **des_setparity**, **or ecb_crypt** Implements Data Encryption Standard (DES) routines.

Chapter 11. AIX Fast Connect for Windows

AIX Fast Connect for Windows

AIX Fast Connect for Windows is server software that lets AIX servers and workstations share files and printers with personal computer clients running Windows NT, Windows 98, Windows 95, Windows For Workgroups, or OS/2 operating systems. This chapter discusses the following topics:

- AIX Fast Connect for Windows Overview, on page 11-2
- AIX Fast Connect for Windows Requirements, on page 11-7
- AIX Fast Connect for Windows Packaging and Installation, on page 11-8
- AIX Fast Connect for Windows Configuration, on page 11-10
- AIX Fast Connect for Windows Administration, on page 11-13
- Connecting Client PCs to AIX Fast Connect for Windows, on page 11-15
- AIX Fast Connect for Windows Problem Determination, on page 11-22

AIX Fast Connect for Windows Overview

Since AIX Fast Connect for Windows uses the Microsoft networking protocol, PC clients can access AIX files and printers using their native networking client software. PC users can use remote AIX file systems directly from their machines as if they were locally stored, and they can print jobs on printers using the AIX spool system, view available printers, and map a printer as a network printer.

Fast Connect provides these services by implementing the Server Message Block (SMB) networking protocol on top of the Network Basic Input/Output System (NetBIOS) over the Transmission Control Protocol/Internet Protocol (TCP/IP), based on the Internet Engineering Task Force requests for comments (RFCs) 1001 and 1002.

Important features of Fast Connect include:

- Tight integration with AIX and exploitation of features such as threads, kernel I/O, filesystem, and security
- SMB-based file and print services
- Login and passthrough authentication to NT domains
- Resource Browsing Protocol
- Opportunistic locking (oplock)
- WINS client and proxy
- B-node support
- Send File application programming interface (API) support
- · Traces and log capabilities
- Maintenance and administration using your user interface choice: commands,
 Web-based System Manager, or the System Management Interface Tool (SMIT)
- Streamlined configuration
- AIX long file name to DOS file name mapping support
- Unicode representation of share, user, file, and directory names.

Common Concepts and Terms

Explanation of some common networking terms is given below:

Broadcast (B-) Nodes

A Broadcast Node or B—node is a type of NetBIOS end node that supports NetBIOS service and contains applications. B—nodes communicate using a mix of UDP datagrams and TCP connections. B—nodes may freely interoperate with one another within a broadcast area. Other standard end nodes are Point—to—point nodes (P—nodes) and Mixed—mode nodes (M—nodes).

Browsing

Browsing refers to viewing the resources available on a network. The browse list on a Windows network is the list of other hosts and domains available on a network. Windows maintains the browse list to present other hosts offering network services through a point—and—click user interface rather than asking users to remember the names of remote hosts and services. Windows 95, Windows 98, and Windows NT use the browse list to construct the view of the network shown in the Network Neighborhood and Windows Explorer. The browse list is also accessible from the command line using the **NET VIEW** command.

Windows for Workgroups and Windows NT domains maintain the browse list on a computer called the Master Browser. Whenever a computer offers a network service for the first time,

it broadcasts a server announcement packet. The Master Browser receives this packet and adds the computer's name to its browse list. In response, the Master Browser transmits a list of backup browsers to the new computer.

Each domain or NT group contains at least one backup browser. A copy of the browse list is maintained on the backup browser to eliminate the need to rebuild the browse list if the Master Browser goes down.

CIFS

CIFS stands for Common Internet File System Protocol. CIFS provides an open cross-platform mechanism for client systems to request file services from server systems over a network. It is based on the Server Message Block protocol widely in use by PCs and workstations running a wide variety of operating systems. It is a draft submitted by Microsoft to the Internet Engineering Task Force for transparent file access across the Internet.

NetBIOS

NetBIOS, or Network Basic Input/Output System, is a vendor-independent network interface originally designed for IBM PC computer systems running PC-DOS or MS-DOS. NetBIOS is a software interface, not an actual networking protocol. It specifies the services that should be available without putting any restrictions on the protocol used to implement those services.

There is no officially defined NetBIOS standard. The original version, as described by IBM in 1984 in the IBM PC Network Technical Reference Manual, is treated as the de facto standard. Since its introduction, three main flavors of NetBIOS—each with its own transport protocol—have emerged: NetBEUI, NetBIOS over IPX, and NetBIOS over TCP/IP.

Fast Connect uses NetBIOS over TCP/IP.

NetBIOS Interface to Application Programs

On PCs, NetBIOS includes both a set of services and an exact program interface to those services. There are three types of NetBIOS services:

Name Service NetBIOS resources are referenced by name. Lower level addresses are not available to NetBIOS applications. An application representing a resource registers one or more names that it wishes to use.

> The name space is flat and not hierarchically organized. It uses 16 alphanumeric characters. Names may not start with an asterisk (*).

Registration implies bidding for use of a name. The bid may be for exclusive (unique) or shared (group) ownership. Each application contends with other applications in real time. No two applications on the NetBIOS network may use a unique name until the originating application requests that their name be deleted or the host is powered off or reset.

The three primitive operations provided by Name Service are Add Name, Add Group Name, and Delete Name.

Session Service

A session is a full-duplex, sequenced, and reliable message exchange conducted between a pair of NetBIOS applications. Data is organized into messages.

Multiple sessions may exist between any two applications. Both applications participating in the session have access to the name of the remote application. No specification is given for resolving session requests to a group name into a data connection. A service is provided for the detection of a session failure by an application.

The Session Service primitives are **Call**, **Listen**, **Hang Up**, **Send**, **Receive**, and **Session Status**.

Datagram Service

The Datagram Service is an unreliable, nonsequenced, and connectionless communication between two NetBIOS applications. It is analogous to UDP service under TCP/IP.

Datagrams are sent under cover of a name properly registered to the sender. Datagrams can be sent to a specific name or be explicitly broadcast.

Datagrams sent to an exclusive name are received, if at all, by the holder of that name. Datagrams sent to a group name are multicast to all holders of that name. The sending application can not distinguish between group and unique names and thus must act as if all nonbroadcast datagrams are multicast.

As with the Session Service, the receiver of the datagram is told the sending and receiving names.

The Datagram Service primitives are **Send Datagram**, **Send Broadcast Datagram**, **Receive Datagram**, and **Receive Broadcast Datagram**.

NetBIOS Name Resolution

Name Resolution refers to mapping a NetBIOS name to its corresponding IP address. The techniques commonly used for name resolution are the Windows Internet Name Service (WINS), the **LMHOSTS** file, and the domain name system (DNS). DNS is explained in TCP/IP Name Resolution, on page 3-97. The other techniques are explained below:

WINS

When a new service is made available on the network, such as when a Windows machine boots or when Fast Connect is started, the service must be registered with a WINS server before it can be available to clients located on other subnets. The WINS server records the name of the host, the NT domain the host is part of, and the IP address of the host. Whenever a machine attempts to resolve a host name, it first checks with the WINS server. If the host is not registered there, it will attempt to find the host using a broadcast. If the host is still not found, an error saying a computer or sharename could not be found is returned. Fast Connect properly registers itself with any WINS server.

WINS also includes a method for replicating its database of host names with other WINS servers to create a backup WINS server that can host queries if the primary WINS server is unavailable. It also allows large networks that are encumbered by slow links to distribute WINS servers closer to clients and provide faster name resolution. (WINS is a proprietary Microsoft protocol.)

LMHOSTS

LMHOSTS stands for LanManager Hosts and is analogous to the UNIX /etc/hosts file. The LMHOSTS file allows specific hostnames to be mapped to IP addresses. It also provides a syntax for defining the domain a host resides in as well as loading a LMHOSTS file from a shared directory on a server.

For more details on the **LMHOSTS** file, refer to the Windows NT Networking Guide or the Windows 95 Resource Kit.

NetBIOS over TCP/IP

NetBIOS over TCP/IP was first proposed in RFCs 1001 and 1002, which were submitted to the Internet Engineering Task Force in 1987. These RFCs describe an implementation of NetBIOS using Transmission Control Protocol (TCP) for connection-oriented session services and User Datagram Protocol (UDP) for datagram services.

This design has some significant advantages over NetBEUI and NetBIOS over IPX. First, it uses the existing TCP/IP protocols, so it can be routed across the global Internet and any other wide area networks. Secondly, software implementing the NetBIOS interface can be built using existing TCP/IP implementation without requiring any new network drivers. Since most operating systems already support TCP/IP, most are capable of supporting NetBIOS with minimal additional effort.

NetBIOS Scope

A NetBIOS Scope is the population of computers across which a registered NetBIOS name is known. NetBIOS broadcast and multicast datagram operations must reach the entire extent of the NetBIOS scope.

net Command The net command and its subcommands can be used to configure and administer the Fast Connect Server from the command line. Alternatively, Web-based System Manager and SMIT offer menu-driven interfaces for the same tasks. For detailed information about the net command, refer to the AIX Commands Reference, Volume 4.

Passthrough Authentication

Passthrough authentication is a mechanism employed by the Fast Connect server to validate user credentials with a domain controller and, if validated, to grant the user access to a resource on the Fast Connect server. For more details on passthrough authentication, refer to Internet Draft, CIFS Login and Passthrough Authentication, Prelim Draft, Jan 3, 1997 by Leach and Naik of Microsoft.

SMB

SMB stands for Server Message Block. It is the protocol used on top of NetBIOS to implement Windows file sharing and print services.

With this protocol, clients exchange messages (called Server Message Blocks) with a server to access resources on that server. Every SMB message has a common format: it consists of a fixed-sized header followed by a variable-sized parameter and data component.

SMB messages can be broken into four types:

- Session control messages start, authenticate, and terminate sessions.
- File and printer messages control file and printer access, respectively.
- Message commands allow an application to send or receive messages to or from another host.

When an SMB client negotiates a connection with an SMB server, the two parties decide on a common protocol to use for communication. This capability allows protocol extensions but makes SMB quite complex.

Shares Shares are resources shared by the Fast Connect server. The two types of shares supported by Fast Connect are AIX files and printers.

Workgroups

A workgroup is a logical collection of workstations and servers that do not belong to a domain. In a workgroup, each computer stores its own copy of user and group account information. Therefore, in workgroups, users can only log directly onto machines on which they have accounts. Workgroup members are able to view and use resources on other systems. To do this, resources are shared in the workgroup and network users are validated by the machine owning the resource.

AIX Fast Connect for Windows Limitations

The following limitations apply to AIX Version 4.3.3 and earlier versions of Fast Connect:

- Fast Connect does not allow multiple shares for the same print queue name. This means
 that if a printer share already exists for a given print queue, you can not create another
 one (having a different netname) for the same queue. If you do, you will get an
 Operation could not be performed error.
- Users who use OS/2 or other clients that do not support unicode must ensure client and server locales match.
- AIX ACLs are not recognized and not supported. Explicit user/group ACL permit/deny access behavior will be ignored.
- Security auditing is not accountable to any individual user actions through an SMB client. If auditing is turned on, some audit log event entries will have the root identity.
- Disk quota and ulimit of any user is not checked. Any user can fill up the shared file system.
- No support for DCE/DFS credentials or filesystems exists. DCE/DFS filesystems cannot be shared.
- There is no log information about which users are connected to the AIX server through AIX Fast Connect for Windows.
- Some AIX printer drivers backend add controls to the file that is being printed; others do
 not. But windows clients always send in a format that needs no controls. So, if the AIX
 printer driver is the former type, then set printer share options —o —dp when the printer
 share is created.

AIX Fast Connect for Windows Requirements

This section includes hardware and software requirements, both for the AIX server and for its PC clients.

Server Hardware Requirements

AIX Fast Connect for Windows runs on any machine that supports the AIX operating system with the exception of diskless or dataless machines. The machine must have a network adapter supported by TCP/IP and should be physically connected to a network. The system must have at least 32 MB of RAM (64 MB is preferred) and 50 MB of available disk space.

Server Software Requirements

The following are the server software requirements for AIX Fast Connect for Windows:

- AIX Version 4.3.2 or higher
- The size of /var should be large enough to temporarily store the largest file that can be printed by the print service
- Fileset bos.net.tcp.client version 4.3.2.0 or higher must be installed and configured.
- Fileset **bos.rte.loc** version 4.3.2.2 or higher must be installed and configured.
- AIX Version 4.3.2 APAR IX85388 is required for sendfile API support.

Client Hardware Requirements

Each client PC must have an installed LAN adapter and should be physically connected to a network.

Client Software Requirements

To use Fast Connect, all clients must have:

• A web browser with forms support (for example, Netscape) to be able to use the Web-based System Manager tool.

and one of the following operating systems:

- Windows NT 4.0
- Windows 98
- Windows 95 with service pack 1 or higher
- Windows for Workgroups 3.11 or higher
- OS/2 Warp 4.0 or higher

AIX Fast Connect for Windows Packaging and Installation

This section describes the AIX Fast Connect for Windows packaging images and installation requirements.

Packaging

Fast Connect packaging includes the following images:

Image Description

cifs.base Fast Connect Server Utilities
cifs.msg Fast Connect Server Messages
cifs.basic Fast Connect Server (Windows only)

or

cifs.advanced Fast Connect Advanced Server (Windows and OS2)

These images contain the following filesets:

Image	Fileset	Fileset Description
cifs.base	cifs.base.websm	Fast Connect Web-based System Manager
	cifs.base.smit	support
cif	cifs.base.cmd	Fast Connect SMIT support
		Fast Connect commands
cifs.msg	cifs.msg	Fast Connect Server messages (language is indicated by the fileset extension)
cifs.basic	cifs.basic.rte	Fast Connect Server files (Windows only)
cifs.advanced	cifs.advanced.rte	Fast Connect Advanced Server files (Windows and OS2)

Installation

AIX Fast Connect for Windows is installed on AIX Version 4.3.2 or higher. It requires that the following filesets are already installed and configured:

- bos.net.tcp.client version 4.3.2.0 or higher
- **bos.rte.loc** version 4.3.2.2. or higher

Also, AIX Version 4.3.2 IX85388 is required for **sendfile** API support.

In addition to these software requirements, Fast Connect requires approximately 50 MB of additional disk space. For further details, see the *AIX 4.3 Installation Guide*.

Once the installation is complete, the following files appear on the system:

File	Туре	Path	Description
cifsServer	binary	/usr/sbin	Fast Connect Server daemon
cifsServerAdv	binary	/usr/sbin	Fast Connect Advanced Server daemon
cifsPrintServer	binary	/usr/sbin	Fast Connect Print File Server daemon
net	binary	/usr/sbin	Fast Connect administration command

File	Туре	Path	Description
rc.cifs	Script	/etc	Fast Connect start/stop shell script for inittab
cifsConfig	ASCII	/etc/cifs	Fast Connect configuration file
cifsPasswd	ASCII	/etc/cifs	Fast Connect user configuration file
README	ASCII	/etc/cifs	Additional documentation for Fast Connect
cifsLog	ASCII	/var/cifs	Fast Connect log file
cifsTrace*	ASCII	/var/cifs	Fast Connect trace file
sm_smb.cat	Message catalog	/usr/lib/nls/msg	Fast Connect message catalog (language indicated in file name extension)

Note: The **cifsTrace** file does not appear on the system once the installation is completed.

AIX Fast Connect for Windows Configuration

You can use Web-based System Manager, SMIT, the net command, or a combination of these methods to configure the Fast Connect server for your site.

Configurable Parameters

Fast Connect is designed for ease of administration without eliminating required customization. Therefore, only a small set of configurable parameters is available for configuration.

Note: Most configurable parameters are dynamically configurable and do not require the server to be stopped and restarted for the changes to become effective.

A brief description of these parameters follows:

Parameter	Description	Type	Value	S/D ¹
autodisconnect	Timeout(minutes) for inactive sessions	int	(Default, Min, Max) 120, 0, 65535	D
backup_passthrough _authentication_server	IP address of the backup passthrough authentication server	String	n/a	D
comment	Comment of the server	String	n/a	D
domainname	Server domain	String	n/a	S
dosattrmapping	If set to 1, then Archive, System, and Hidden attributes are mapped to User, Group, and Other execute bits. Otherwise, these attributes are not supported.	int	(1, 0, 1)	D
dosfilenamemapchar	The character used to map long file names to 8.3 format. Valid values are tilde (~) and caret (^). Tilde (~) is the default.	char	~	D
dosfilenamemapping	If set to 1, long file names are mapped to 8.3 format. Otherwise, long file names are truncated. For more information, see Mapping AIX Long File Names to DOS File Names, on page 11-14	int	(1, 0, 1)	D
encrypt_passwords	If set to 0, plain text password. If set to a non-zero value, negotiates with client.	int	(1, 0, 2)	S
filenamematchinglevel	If set to 0, case—insensitive matching. If set to 1, case—sensitive matching.	int	(1, 0, 1)	S
guestlogonsupport	Whether or not guest access is allowed	int	(0, 0, 1)	D
guestname	Guest id	String	n/a	D
maxconnections ²	Maximum number of open connections on the server	int	0, 0, 1000	D

Parameter	Description	Туре	Value (Default, Min, Max)	S/D ¹
maxopens ²	Maximum number of open files on the server	int	0, 0, 1000	D
maxsearches ²	Maximum number of open searches on the server	int	0, 0, 1000	D
maxusers ²	Maximum number or user session permitted	int	0, 0, 1000	D
nbns	If set to 1, server acts like a NETBIOS name server	int	1, 0, 1	S
oplockfiles	Whether oplocks on files are supported	Y/N	yes, no, yes	S
oplocktimeout	Timeout in seconds for oplock	int	35, 35, 640	S
os2compatible	If set to 0, READONLY means all readable but not writable files. If set to 1, READONLY means all readable files.	int	1, 0, 1	D
passthrough_authentication _server	IP address of the passthrough authentication server	String	n/a	D
wins_proxy	Whether or not server has to act as WINS proxy	int	(0, 0, 1)	D
primary_wins_ipaddr	IP address of primary WINS address	String	n/a	D
secondary_wins_ipaddr	IP address of secondary WINS address	String	n/a	D
send_file_api	Turn on or off send file API	int	(1, 0, 1)	D
send_file_cache_size	If the send_file_api is 1 and the requested SMB read size is less than the value of this parameter, the send_file API will cache the file. The default value is zero, which means send_file API will not cache the file.	int	(0, 0, 4194304)	D
send_file_size	If the send_file_api is 1 and the requested SMB read size is greater than the value of this parameter, then send_file API will be used in the SMB operation	int	(4096, 1, 4194304)	D
servername	Name of the server	String	n/a	S
umask	AIX umask	octal	(022, 0, 0777)	D

Note:

- 1. S stands for static and D for dynamic. Any changes to static parameters require a Shutdown/Restart of the server before they take effect.
- 2. For maxusers, maxconnections, maxopens, and maxsearches, a default or minimum value of zero means unlimited (no restrictions).

File or Print Share Configuration

There are two types of shares that can be configured by Fast Connect: File Shares and Print Shares. When the server is first started, a file share with netname HOME is created and loaded by default. More file or print shares can be added by the administrator using Web-based System Manager, SMIT, or the net command.

Note: The default shares HOME, IBMLAN\$, and ADMIN\$ cannot be changed or deleted.

User Configuration

User accounts can be configured on the server using Web-based System Manager, SMIT, or the net command. Each defined Fast Connect user must also be a defined AIX user. Fast Connect supports user level authentication using plain text passwords. Resource access is permitted based on the authenticated user's AIX credentials.

Guest access is controlled with the help of **guest** and **guestname** configuration parameters. If enabled, Fast Connect allows guest access to server resources when the user ID received in the server session setup does not match any of the users defined on the server and is also different from the value of the **questname** parameter configured on the server.

Note: If encrypt passwords is set to 0, the user's password cannot be in mixed case. If encrypt_passwords is set to 1 or 2, add usernames to the Fast Connect Users database using Web-based System Manager, SMIT, or the net command. The Fast Connect Users database is a subset of the AIX Users database.

Configuration of Network Interfaces

Fast Connect complies with the SMB protocol requirement of using NetBIOS over TCP/IP protocol as described in IETF RFC1001 and RFC1002. It uses all the network interfaces that are defined and operational (are up) for AIX TCP/IP.

Initial Configuration

On a newly installed Fast Connect server, the initial configuration parameters are:

Server Name TCP/IP hostname of the server

Server Description Fast Connect server on hostname

Default Shares HOME, with the following attributes:

Network Name HOME

AIX directory User's home directory as defined in the user profile

Other server parameters are at default values.

AIX Fast Connect for Windows Administration

You can use Web-based System Manager, SMIT, or the net command to manage and administer the Fast Connect server. These tasks can fall in one of the three categories: controlling the server, changing server configuration, and getting information about the server.

	Administering AIX Fast Connect for Windows			
Web-based System Managers	wsm pc fast path			
(PC Services application)				
-OR-				
Task	SMIT Fast Path	Command or File		
Starting the Server	smit smbadminstart	net start		
Stopping the Server	smit smbadminstop	net stop		
Pausing the Server		net pause		
Resuming the Server		net resume		
Changing Parameters	smit smbcfghatt	net config		
Changing Resources	smit smbcfgresi	net config		
Adding Users	smit smbcfgusradd	net user		
Changing Users	smit smbchgusrlis	net user		
Changing a User Password	smit smbusrpwd	net user		
Removing a User	smit smbrmusrlis	net user		
Configuring nbns	smit smbwcfgn			
Adding a NetBIOS name	smit smbwcfgadd	net nbaddname or		
		net nbaddgroup or		
		net nbaddmulti		
Deleting a NetBIOS name	smit smbwcfgdel	net nbdelname		
Deleting a NetBIOS name	smit smbwcfdadd	net nbdeladdr		
by address and by name				
Backing Up NetBIOS name	smit smbwcfgbak	net nbbackup		
table	omit ombuvofaroo	net nbrestore		
Restoring NetBIOS name table	smit smbwcfgres	net hbrestore		
Listing All Shares	smit smbsrvlisall	net share		
Listing All File Shares	smit smbsrvfilist	net share		
Adding a File Share	smit smbsrvfiladd	net share		
Changing a File Share	smit smbsrvfilchg	net share		
Deleting a File Share	smit smbsrvfilrm	net share		
Adding Printer Share	smit smbsrvprtadd	net share		
Changing Printer Share	smit smbsrvprchg	net share		
Deleting Printer Share	smit ssrvprtrm	net share		
Showing Server Status	smit smbadminstatu	net status		
Showing the Configuration	smit smbcfg	net config		

Administering AIX Fast Connect for Windows			
Showing Statistics	smit smbadminstats	net statistics	
Showing Share	smit smbsrvlisall	net share	
Getting Help	smit then select Help	net help	

Mapping AIX Long File Names to DOS File Names

Older PC clients operating systems, such as Windows for Workgroups 3.11, do not support long file names. This restriction requires mapping long file names of AIX files to DOS file names. Simply truncating the name to the shorter name does not work for most situations because many file names will get mapped to the same names as long as the first 8 characters are same. Also, this feature is needed for many older (16-bit) applications running under Windows 95, Windows 98, and Windows NT.

Fast Connect maps AIX file names (AFN) to DOS File Names (DFN) ensuring file name uniqueness. It maps AFNs to DFNs using Microsoft Windows NT method for mapping names (that is, name conflicts are handled by using a delimiting character in in the short name followed by a unique numeric to make the name unique).

Fast Connect generates a mapped name whenever the AFN needs to be passed back to a DOS client.

DFNs generated by Fast Connect are not remembered across server restarts. However, mappings will stay consistent for the lifetime of the Fast Connect server.

It supports efficient lookups given DFN as well as AFN.

Fast Connect has a configuration option to turn off the mapping. When the mapping is turned off, the long names are simply truncated to conform to DOS file naming.

The character used as the delimiting character is configurable as well.

Note: AFN-to-DFN mapping may not map correctly if the server restarts. For example, consider two files in the root directory of an exported SMB share: LongFileName1.txt and LongFileName2.txt. Assume a Windows 3.11 client mounts this share and searches the directory. The resulting filenames will be:

```
LONGFI~1.TXT for LongFileName1.txt
LONGFI~2.TXT for LongFileName2.txt
```

The human user on this client notes this filename (LONGFI~1.TXT) and opens it for editing. The file is written, saving changes. Then the server shuts down. Someone then removes LongFileName1.txt from the server file system. Once the server is up and running, the user on the client again edits LONGFI~1. TXT. This time this 8.3 name maps to LongFileName2.txt, not the previously deleted file name! Thus, the client ends up editing the wrong file. Therefore, if the network drive is reconnected following server restart, new file lists must be obtained before accessing any mapped names.

Connecting Client PCs to AIX Fast Connect for Windows

The steps shown in the following sections are required to connect a PC client to the Fast Connect server.

TCP/IP Configuration

To access the Fast Connect Server, each client PC must be configured for NetBIOS over TCP/IP (RFC1001/1002). This can be accomplished for the various clients as shown in the following sections.

Windows 95, Windows 98 Clients

- 1. From the Start button, select **Settings->Control Panel->Network**.
- 2. On the Configuration tabbed panel (initially shown), verify that the following entries exist:
 - An entry for your networking—card (hardware driver)
 - TCP/IP (protocol)
 - Client for Microsoft Networks (client)

If any is missing, add it from your Windows 95 disks.

3. Click on the TCP/IP entry and select Properties. The TCP/IP Properties dialog box has several tabbed panels. Verify the following:

IP Address panel Configure as needed. (For initial testing, you may find it convenient

to manually specify unique IP addresses for each PC.)

Select Client for Microsoft Networks. Bindings panel

Additionally, you may wish to enable WINS support, DNS support, and/or GATEWAY support for each client. If so, configure each as needed.

4. TEST the client's TCP/IP configuration by ping-ing (by IP address) from the client PC's DOS prompt to the Fast Connect server, and vice versa.

Windows NT Clients

Note: You must be logged in as an Administrator.

- 1. From the Start button, select **Settings->Control Panel->Network**.
- 2. On the Services tabbed panel, verify that there are entries for the following services:
 - Computer Browser
 - NetBIOS Interface
 - Workstation

If any is missing, add it from your Windows NT CD.

3. On the Protocols panel, add TCP/IP (if missing), then select Properties.

The TCP/IP Properties dialog box has several tabbed panels. Verify the following:

Configure as needed. (For initial testing, you may find it convenient IP Address panel to manually specify unique IP addresses for each PC.)

You may also want to configure DNS, WINS Address, and Routing.

4. Test the client's TCP/IP configuration by ping-ing (by IP address) from the client PC's DOS prompt to the Fast Connect server and vice versa.

Windows For Workgroups (Windows 3.11) Clients

- 1. From group Network (within Program Manager), run Network Setup.
- 2. Verify that the following entries exist:

- Microsoft Windows Network (version 3.11) (network)
- An entry for your LAN adapter card (device-driver)
- Microsoft TCP/IP-32 3.11b (protocol)

You may need to install the TCP/IP protocol. TCP/IP is not included on the Windows 3.11 installation media. You can download a copy of Microsoft TCP/IP-32 3.11b from Microsoft's web site www.microsoft.com.)

To set up the TCP/IP configuration, double-click on Microsoft TCP/IP-32 3.11b. Configure the IP Address, Subnet Mask, Default Gateway, WINS Server(s), DNS, and other options as needed. (LMHOSTS and DNS enablement are available as Advanced options.)

3. Test the client's TCP/IP configuration by ping-ing (by IP address) from the client PC's DOS prompt to the Fast Connect server and vice versa.

OS/2 Clients

- 1. Install TCP/IP and NetBIOS support during OS/2 installation.
- 2. Use the TCP/IP Configuration program to verify and configure TCP/IP.
- 3. Use the Multi-Protocol Transport Services program (MPTS) to verify and configure the following protocols for your network adapter:
 - IBM OS/2 TCP/IP
 - IBM OS/2 NetBIOS OVER TCP/IP

These protocols should have the same LAN adapter number, which should match your TCP/IP interface.

Note: The default installation is IBM OS/2 NetBIOS. Be sure to add IBM OS/2 NetBIOS OVER TCP/IP if not already listed.)

4. Test the client's TCP/IP configuration by ping-ing (by IP address) from the client PC's DOS prompt to the Fast Connect server and vice versa.

User Authentication and Administration

supports the following methods for user authentication:

- AIX—based User Authentication using plain text CIFS passwords.
- CIFS LM (LAN Manager) password encrytion protocol
- NT Domain Passthrough authentication

User administration is dependent on the authentication method selected by Fast Connect administrator. Each type has its advantages and disadvantages. Which method you choose depends on your environment, your administration policy, and the type that you determine would be easiest to administer and use.

AIX-based User Authentication

AIX-based authentication uses AIX user definitions and passwords. Following session setup, a Fast Connect session gets the authenticated AIX user credentials (UID, GID and Secondary groups set).

Requirements

- 1. Clients must be able to negotiate plain text passwords. This may require enabling plain text passwords by updating required registry entries for Windows NT, 95, and 98 clients.
- 2. Fast Connect must be enabled for plain text passwords using SMIT, Web-based System Manager or the net command

Advantages

- 1. Low administrative overhead (uses existing AIX user information).
- AIX tools for managing users can be used.

Disadvantages

- 1. Windows registry update may be required.
- 2. Windows may require user ID and passwords to be typed again
- 3. Clear-text passwords are sent over the network.

CIFS LM (LAN Manager) Password Encryption Protocol

The CIFS LM protocol method uses Fast Connect user definitions and encrypted passwords for user authentication. Each user must be defined as an AIX user as well, and share the same user name. Fast Connect encrypts passwords and saves them in its files for use during session setup. Following session setup, a Fast Connect session gets the authenticated user's credentials (UID, GID and Secondary groups set).

Requirements

- 1. Users must be defined to Fast Connect using SMIT, Web-based System Manager or the net command.
- 2. AIX Fast Connect for Windows must be enabled for encrypted passwords using SMIT. Web-based System Manager or the **net** command.
- 3. Windows or OS/2 user logon passwords must be same as Fast Connect passwords. These passwords are not required to be same as AIX logon password.
- 4. Changing passwords requires root authority.

Advantages

- 1. No additional logon other than logging into the Windows or OS/2 workstation is required.
- 2. Clear text passwords are not sent over the network, which provides additional security.

Disadvantages

- Additional administrative tasks are needed for Fast Connect users.
- 2. Administrator intervention is needed for password update.

NT Domains Passthrough Authentication

This authentication method uses AIX user definitions and NT Server user authentication. Each user must be defined as an AIX user as well. Passthrough authentications is enabled using SMIT, Web-based System Manager or the net command by specifying Passthrough authentication NTserver IPaddress.

During session setup, Fast Connect forwards the session setup request to the NT server. If the NT server authenticates the user, Fast Connect grants access. Following session setup, a Fast Connect session gets the authenticated user's credentials (UID, GID and Secondary groups set).

Requirements

- 1. User must be defined on the Passthrough authentication server.
- 2. AIX Fast Connect for Windows must be enabled for Passthrough authentication using SMIT, Web-based System Manager or the net command to define IP address of the NT server.
- 3. NT user name must match AIX user name, although passwords can be different.

Advantages

1. No additional logon other than logging into the Windows or OS/2 workstation is required.

- 2. Clear text passwords are not sent over the network, which provides additional security.
- 3. Uses NT user definition, therefore less administrative overhead is needed.

Disadvantages

1. Requires an NT server.

NetBIOS Name Service (NBNS)

NetBIOS Name Service (NBNS) for AIX Fast Connect for Windows provides name resolution services. It also supports some functions of Windows Internet Name Service (WINS), such as registration of multihomed name and Internet group name.

To activate NBNS, 'nbns' in the cifsConfig file needs to be set to 1 (The default is 1) and restart the Fast Connect Server.

To turn off NBNS, set nbns = 0 in the cifsConfig file.

Administering NBNS Tasks		
Task	SMIT Fast Path	Command or File
List names in the NetBIOS Name Table	smit smb	net nblistnames
Add a NetBIOS Name	smit smb	et nbaddname /name:NetBIOS Name /ipaddress:ip address [/subcode:00-ff]*
Delete a NetBIOS name in Name Table	smit smb	net nbdelname /name:NetBIOS Name [/subcode:00-ff]
Delete by Address and by Name	smit smb	net nbdeladdr /name:NetBIOS Name /ipaddress:ipaddress*
Backup the NetBIOS Name Table to a File	smit smb	net nbbackup [/file:file name]
Restore a NetBIOS Name Table from Backup File	smit smb	net nbrestore [/file:file name]

Note: The value of *ip address* can be any number in IP address range.

Workgroups, Domains, and User Accounts

For ease of use, client PCs should be in the same Windows workgroup or NT domain as the Fast Connect server (or vice versa). Windows 3.11, Windows 95, and Windows NT all use WORKGROUP as a default workgroup name, and Fast Connect server initializes itself to use WORKGROUP, also. If your network uses NT domain login authentication, you can configure the Fast Connect server to verify Fast Connect access using the NT domain authentication servers.

Whether you use Workgroups or NT domains, access to Fast Connect is managed by user security. You must set up AIX user accounts for each Windows user who is accessing Fast Connect. It is easiest to use if the user accounts (and passwords) on AIX match the Windows or NT domain user accounts (and passwords).

On the Fast Connect server, use the SMIT fast path:

smit smbcfghatt

To use Workgroups, enter the workgroup name in the Domain Name field.

To use NT domain validation, enter the IP addresses for the NT domain authentication server(s) in the Passthrough Authentication Server and Backup Passthrough Authentication Server fields.

On PC clients running Windows 95 or Windows 98, select Start button -> Settings -> Control Panel -> Network. On the Identification panel, enter the computer name for that PC.

To use Workgroups, enter the workgroup name in the Workgroup field.

To use NT domain validation, go to the Configuration tabbed panel, select Client for Microsoft Networks, and click on Properties. Check the NT domain checkbox, and enter the NT domain name.

- On PC clients running Windows NT, make sure you are logged on as Administrator.
 - a. Select Start button -> Settings -> Control Panel -> Network.
 - b. On the Identification panel, select the **Change...** button.
 - c. Enter the Computer Name for that PC.
 - d. Enter the workgroup or domain name, as appropriate.
- On PC clients running Windows for Workgroups, configuration for the workgroup occurs during Windows 3.11 install/setup, but can be changed in the SYSTEM.INI file, if necessary.
- On PC clients running OS/2, configuration for the workgroup occurs during OS/2 installation, but can be changed in the DOMAIN parameter of IBMLAN.INI, if necessary.

Note: Use the OS/2 command LOGON to use NetBIOS services such as network browsing, NET VIEW, and NET USE.

Using Plain Text Passwords with Windows 98or Windows NT 4.0 (Service Pack 3)

For security reasons, Microsoft has disabled support for nonencrypted (plain text) network passwords in Windows 98 and Windows NT 4.0 (with Service Pack 3 installed). If you want to use plain text passwords on your network, these clients must be upgraded with the following Registry patches.

To install the Windows 98 Enable Plain Text Passwords patch:

1. Use **EDIT** or the NOTEPAD accessory to create the following text file, named **W98plain.reg**, as a local file on the Windows 98 machine:

REGEDIT4

; Registry file to allow plaintext passwords on Windows 98

[HKEY_LOCAL_MACHINE\System\CurrentControlSet\Services\VxD\VNETSUP] "EnablePlainTextPassword"=dword:0000001

- 2. Using Windows Explorer, double click on the **W98plain.reg** file namein the directory where you saved it. This action will update the WindowsRegistry for that client to allow plain text passwords.
- 3. Shutdown/Restart the Windows 98 machine. (Shutdown/Restart is requiredfor this patch to take effect.)

To install the Windows NT 4.0 Enable Plain Text Passwords patch,

1. Use **EDIT** or the NOTEPAD accessory to create the following text file, named NT4plain.reg, as a local file on the Windows NT machine:

REGEDIT4

; Registry file to allow plaintext passwords on Windows NT 4.0, SP3

[HKEY_LOCAL_MACHINE\SYSTEM\CurrentControlSet\Services\Rdr\Parameters] EnablePlainTextPassword=dword:00000001

- 2. Using Windows Explorer, double click on the NT4plain.reg file name in the directory where you saved it. This action will update the WindowsRegistry for that client to allow plain text passwords.
- 3. Shutdown/Restart the Windows NT machine. (Shutdown/Restart is required for this patch to take effect.)

Note: Even with the previous patch installed, all Windows NT 4.0 clients still require the user to type his/her password every time the user first connects to the Fast Connect server (by browsing, mapping drives, etc.). Once the user is successfully connected, additional browsing or drive mapping operations proceed without hindrance. The initial Password Invalid message is because Windows NT 4.0 attempts to use encrypted passwords while connecting to Fast Connect server, rather than plain text passwords.

Browsing the Network

AIX Fast Connect for Windows supports Browser operations such as Network Neighborhood and NET VIEW. These operations show the user a list of file and printer shares exported by each server.

Network Neighborhood can also be used as a convenient way to map drives. (Right-click on a file share name, then select Map Network Drive from the pop-up menu.)

However, note the following limitations on network browsing:

 To see the Fast Connect server in Network Neighborhood, a client PC needs to be able to see the Master Brower for the workgroup or domain for which that Fast Connect server is configured. (Fast Connect does not act as a browse master itself, but any Windows 95, Windows 98 or Windows NT clients in that workgroup/domain can perform the role.)

Network browsing generally works best if the client PC and the Fast Connect server are in the same workgroup/domain.

The browse list database that is maintained by the Master Browser is not always up-to-date. The list can show Fast Connect server names for servers that are currently down, off, physically disconnected, or otherwise unreachable. The Master Browser does not delete a server from the browse list until that servername's refresh timeout has expired, which could take several days. If a user tries to access that servername, though (by browsing share names, mapping drives, etc.), a disconnected Fast Connect server will be detected as unavailable.

Mapping Drives

Normally, PC clients will need to define drive mappings to use the Fast Connect-exported file shares. These drive mappings can be done from Windows or from the DOS command prompt.

You can use the following mechanisms to define/undefine mappings between PC drive letters and Fast Connect file shares. For the sake of the following examples, assume that the NetBIOS servername is cifs01, and that file shares apps and pcdata are defined.

From DOS:

```
DOS> net help
                                   (help info for DOS)
DOS> net use H: \\cifs01\home
                                  (pre-defined Fast Connect share)
DOS> net use F: \\cifs01\apps
DOS> copy F:\oldfile H:\newfile (uses previous drive-mappings)
DOS> net use F: /delete
                                  (delete the drive-mapping)
```

From Windows:

- 1. Find the Map Network Drive dialog box.
 - Select Windows Explorer -> Tools -> Map Network Drive.

or

- Right-click on Network Neighborhood and select Map Network Drive.
- 2. Select the desired drive (for example, F:) from the Drive: drop-down list, then
 - Enter in the path: (for example, \\cifs01\apps). or
 - Use the Shared Directories (browse tree) panel to select the network share desired.

Using Fast Connect Printers

For printing, DOS and Windows mappings are somewhat different. For the following examples, assume that Fast Connect server cifs01 has print shares netprint1 and pscolor defined.

For DOS applications, the following simple device—mappings can be used:

```
DOS> net use LPT1: \\cifs01\netprint1
DOS> net use LPT2: \\cifs01\pscolor
```

But to support printing from Windows applications, a Windows printer driver needs to be installed, and mapped to the network printer, as follows:

- 1. Select Start -> Settings -> Printers -> Add Printer.
- 2. Select Network Printer.
- 3. Enter the Fast Connect print share name (for example, \\cifs01\netprint1) or, use the browse list to select the desired print share.
- 4. Select the proper Windows printer driver for that network printer (for example, IBM 4039 Laser Printer PS), which will be installed from your Windows installation disks.

For Windows 3.11, install the desired printer driver through Control Panel, and use the Connect... button to map it to the Fast Connect print share.

AIX Fast Connect for Windows Problem Determination

Traces

The Fast Connect server comes with the ability to create AIX trace files to isolate problems. When a trace facility is active, information about selected events is recorded in the trace file. To obtain trace files, you must have the trace command installed on your machine. The trace command is in the **bos.sysmgt.trace** package.

The following trace hooks are being used by the Fast Connect server:

2EE	CIFS Enter
2EF	CIFS Exit
2F0	CIFS-FSS
2F1	CIFS-LOGON
2F2	CIFS-NET
2F3	CIFS-SMB PARSER
2F4	CIFS-PSS
2F5	CIFS-SMS

Trace files can be created by using either the SMIT interface or the command line interface.

From the command line:

1. Enter:

```
trace -a -j 2EE, 2EF, 2F0, 2F1, 2F2, 2F3, 2F4, 2F5 -o /tmp/cifs.trace
```

This creates a trace file named **cifs.trace** in the /tmp directory.

- 2. Recreate the problem.
- 3. Enter:

trcstop

4. Enter:

```
trcrpt -t /etc/trcfmt /tmp/cifs.trace
```

The trcrpt command formats the trace file into readable text and writes a report to standard output.

From the SMIT interface:

1. Enter the SMIT fast path:

```
smit trcstart
```

- 2. Select the CIFS hooks for ADDITIONAL event IDs to trace field, then exit SMIT. This creates a trace file named trcfile in the /var/adm/ras directory (default).
- 3. Recreate the problem.
- 4. Enter:

```
smit trcstop
```

then exit SMIT.

5. Enter:

```
smit trcrpt
```

and select the output format. It displays the trace file into readable format.

Logs

The Fast Connect server writes informational and error messages to a file in /var/cifs named cifsLog.

Common Problems and Their Solutions

Connection Problems

Cannot connect to server.

Check that the server has passthrough authentication server enabled.

access is denied password is invalid password is not correct not authorized to login

> When you get one of these error messages on the client PC, try the following:

- · Enter the correct password.
- · Check logon user ID and its password on clients that should have an account on the AIX server. Log off and log on clients with correct user ID/password.
- For clients with Window NT and Service Pack 3, the NET VIEW command returns access is denied. See Using Windows NT 4.0 with Service Pack 3, on page 11-19 for more information.

Note: Fast Connect does not support mixed case passwords.

System error 53 has occurred.

The network path was not found.

When you get this error message on the client PC, try the following:

- Check the NetBIOS name of the Fast Connect server.
- Check server status.
- See the Connection Checking Procedure, on page 11-23.

System error 51 has occurred.

The remote computer is not available.

When you get this error message on the client PC, check server status. It might be paused.

Connection Checking Procedure

- 1. **ping** the Fast Connect server by IP address. If timeout occurs, check:
 - cable for physical connection
 - status of the AIX machine
 - TCP/IP configuration on clients and on the AIX server
- 2. ping the Fast Connect server with its NetBIOS name. If it fails, refer to NetBIOS Name Resolution, on page 11-4 for more information.
- 3. Check server status on the AIX machine using **net config**, **net status**, and **net** statistics commands.

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